

Review notes

Friday, February 10, 2023 4:29 PM

Group 1 Lecture 2 : Runtime Analysis |||

- Constant Time Operations

Hard:

- N vs. $2N+1$
- Asymptotic runtime analysis
 - How to tell runtime without the code
 - When functions count as constant time or not

tips:

- if it's not a loop, it's likely constant time (except for functions)
- $O(1)$ is constant time

2. Implementing binary search algorithm following above asymptotic analysis.

A:

- Counting Operations

Hard:

- asymptotic runtime vs. constant time

- Binary Search Case Study

Hard:

- best, avg, worst case analysis.

```
graph TD; A((A)) --- B((B)); A --- C((C)); B --- D((D)); B --- E((E)); C --- F((F)); C --- G((G)); D --- H((H)); D --- I((I)); E --- J((J)); E --- K((K)); F --- L((L)); F --- M((M)); G --- N((N)); G --- O((O))
```

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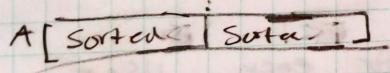
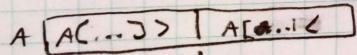
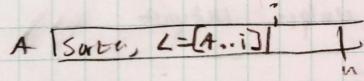
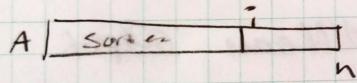
Abstract Data types
Interface vs Implementation
what it does vs how its done

Insertion Sort
Selection Sort

Asymptotic runtime analysis Big O
Common classes $O(\log n)$, $O(1)$, $O(n \log n)$, $O(n^2)$, $O(2^n)$, $O(n!)$

Box diagram of loop invariant

Which of the following represents the invariant for
Insertion Sort



- 5) D No, invalid
- 2) not ready, hidden output has no tangible connection to the input
- 3) No, its already bad

2) integers

3)

- 1) valid, an input will produce a different output
- 2) valid, unique input will have a unique output

c) return whom's accuracy unique for each user

Group 3 → lecture 4

Recursion

It's Easy!

Basics of Recursion

- -) precise specification
 -) base case
 -) recursive call makes progress
on subproblems
 -) replace recursive call with
spec and verify
overall behavior

Tips → Keep it simple

→ Don't try and write out
each call, just use the 4 steps!

Moses Miller 5 Anuk Centellas Wyatt Pratch Carver Rasmussen

- Topics:
 - MergeSort
 - Loop invariants
 - Divide & Conquer

- Hard Things:
 - Recursively creating subarrays
 - Runtime analysis
 - Divide & Conquer

- Tips:
 - The precondition for merge is that the arrays are already sorted
 - Review A1 (if done correctly...)

- Potential Questions:
 - Execute merge on this array:
[1, 4, ~~5~~, ~~7~~] [3, 6, ~~8~~, ~~9~~]

Answer: [1, 3, 4, 5, 6, 7, 8, 9]

- Do the subarrays need to be the same size?

Answer: No!

V ✓

Lecture 6

Topics - Quicksort

Hard things - Pivot not defined

- Can impact runtime

- Determining stability is tricky

- Conceptualizing mental partitions can be difficult to keep track of

Tips - Pivot should not be first or last

- Picture cards when running through QuickSort

Q's - What are the best, average, and worst-case runtimes for Quicksort?

- $n \log n$, $n \log n$, n^2

- When does the worst-case runtime occur?

- When pivot is on a minimum/maximum

Radix Sort

#7

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- Bucket Sort ; Counting Sort

Hard Things

- Counting sort ; Bucket sort both traverse backwards
- pseudo code

TIPS

- Stably sorted
- Sort from lowest to highest digit

Potential questions

Q: What is the best possible worst case runtime for a comparison-based sorting algorithm?

A: $O(n \log n)$

Q: What is the main difference between counting and bucket sort?

A: Bucket sort organizes numbers via digits
Counting sort organizes numbers by summing the count.

Lecture 8

- Trees, Traversal, Revision On Trees

Hard things

- Tree Traversal ✓✓✓✓✓

TIPS

- Pre Order Parent is "Pre" PLR
- Post Order Parent is "Post" LPR
- In Order Parent is "in" between LPR

Write a base case for Tree traversal:

```
IF (T==NULL){  
    return False;}
```

What is the height of a tree?

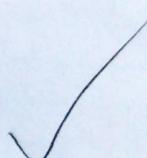
Insert elements into a tree in the order they are visited in Pre Order traversal

LECTURE 10

Topics: Generics, comparable interfaces, balance factors, & tree rotations
Hard things: Rotations - updating pointers when coding

Tips: When calculating balance factor, make sure to subtract the height of the node's children, not right & left subtrees.
Generics: a specialized class that is constructed to ignore types ~~types~~ to be able to be used across the board.

Questions:



Lecture 11 - AVL Trees

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TIPS

AVL invariant: $-1 \leq \text{balance}(n) \leq 1$

worst case for AVL insertion, search

removal is $O(\log n)$

know 2 cases for rebalancing

Hard Things

- Rebalancing
 - knowing which rotations are needed
- Time Complexity

Lecture 12

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Topics

- Priority Queue: Queue sorted by priority
- Heaps:
 - Implements Priority Queue
 - Each child node is \leq or \geq than its parent depending on if it's a min- or max-heap
 - No holes
 - Stored as an array

Hard Things

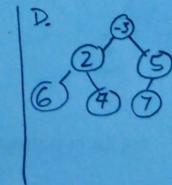
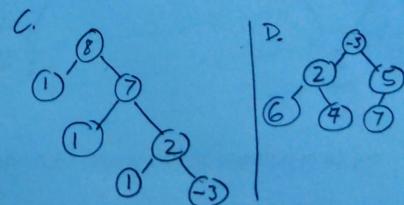
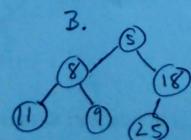
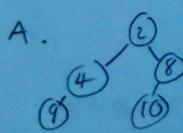
- Inserting nodes/values

Tips

- make sure to remember the formulae to find the parent node and child node

Potential Questions

Which of the trees
are valid min-heaps?



Answer: B, D

MAREN O'DAMN
KUSHA GOTHAM
COLIN SARGENT
FELIX KOMAR

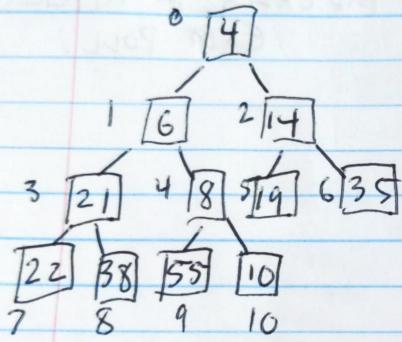
GROUP 12
LECTURE 13

HEAP SORT: (Ans: Priority queues)

- $\text{poll}()$
- level order w/ no holes

~~$O(n)$ construction~~ don't need to know

LEVEL ORDER TIPS:



TO FIND NODE k 's PARENT:

$$(k-1)/2$$

TO FIND NODE k 's CHILDREN:

$$2k+1 \quad \& \quad 2k+2$$

LEFT CHILD

RIGHT CHILD

poll() TIPS:

Hurt Things
 $O(n)$ construction

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HEAP SORT ?:

You HAVE A HEAP SORT THAT FAVORS THE MINIMUM VALUE. WRITE THE OUTPUT FOR THE FOLLOWING CALLS. WRITE THE FINAL TREE @ THE END OF THE CALLS & AFTER EACH POLL.

~~insert(12);~~ add(12);
~~insert(8);~~ add(8);
~~insert(13);~~ add(13);
~~peek();~~ peek();
poll();
add(4);
peek();
add(16);
poll();
add(1);