



Estruturas de Dados / Programação 2 Heaps

Márcio Ribeiro
marcio@jc.ufal.br
twitter.com/marciomribeiro

Priority Queues

- We implemented using a naive approach...

- Enqueue: $O(n)$
- Dequeue: $O(1)$



```
enqueue(pq, "a", 17);    [<a, 17>]
enqueue(pq, "b", 12);    [<a, 17>, <b, 12>]
enqueue(pq, "c", 100);   [<c, 100>, <a, 17>, <b, 12>]
enqueue(pq, "d", 22);    [<c, 100>, <d, 22>, <a, 17>, <b, 12>]
dequeue(pq);              [<d, 22>, <a, 17>, <b, 12>]
dequeue(pq);              [<a, 17>, <b, 12>]
```

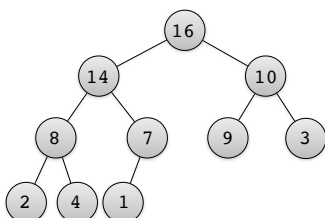


We can do much better!
 $O(\log n)$

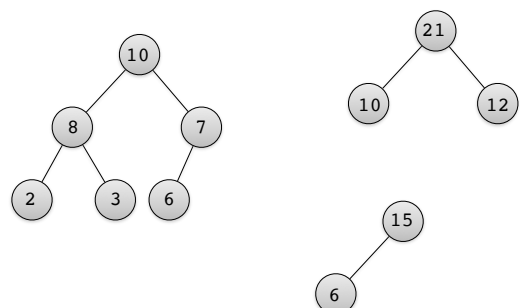
Heap

Definition

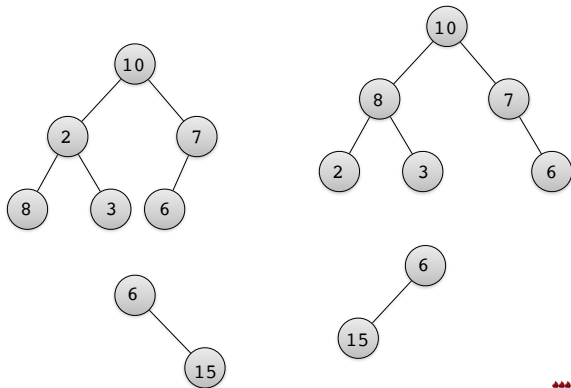
- For every node "i" other than the root, the value of a node is at most the value of its parent
- Completely filled on all levels except possibly the lowest, which is filled from left to right



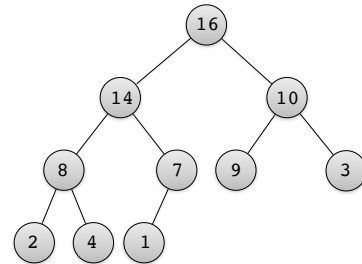
Heap examples



What about these ones?!



Array representation



$\text{Parent}(i) = i/2 = i \gg 1$

$\text{Left}(i) = 2i = i \ll 1$

$\text{Right}(i) = 2i + 1 = (i \ll 1) + 1$

$A[1 \dots A.\text{length}]$

16	14	10	8	7	9	3	2	4	1
1	2	3	4	5	6	7	8	9	10

Valid heap elements: $A[1 \dots A.\text{heap-size}]$



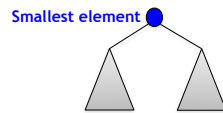
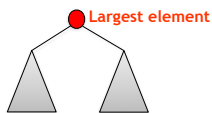
Max-Heaps and Min-Heaps

Max-heaps

$A[\text{Parent}(i)] \geq A[i]$

Min-heaps

$A[\text{Parent}(i)] \leq A[i]$



Abstract Data Type:
Heap

Heap ADT

```

heap* create_heap();

void enqueue(heap *heap, int item);

int dequeue(heap *heap);

int get_parent_index(heap *heap, int i);

int get_left_index(heap *heap, int i);

int get_right_index(heap *heap, int i);

void max_heapify(heap *heap, int i);

int item_of(heap *heap, int i);

void heapsort(heap *heap);

```



Struct

```

struct heap {
    int size;
    int data[MAX_HEAP_SIZE];
};

```



```

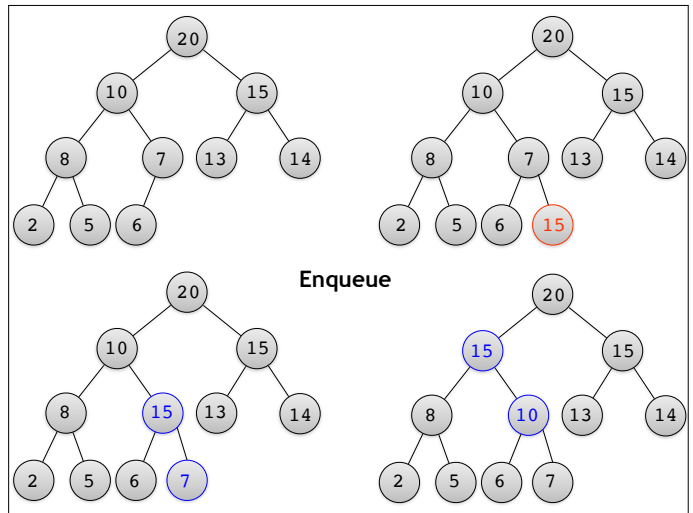
int get_parent_index(heap *heap, int i)
{
    return i/2;
}

int get_left_index(heap *heap, int i)
{
    return 2*i;
}

int get_right_index(heap *heap, int i)
{
    return 2*i + 1;
}

int item_of(heap *heap, int i)
{
    return heap->data[i];
}

```



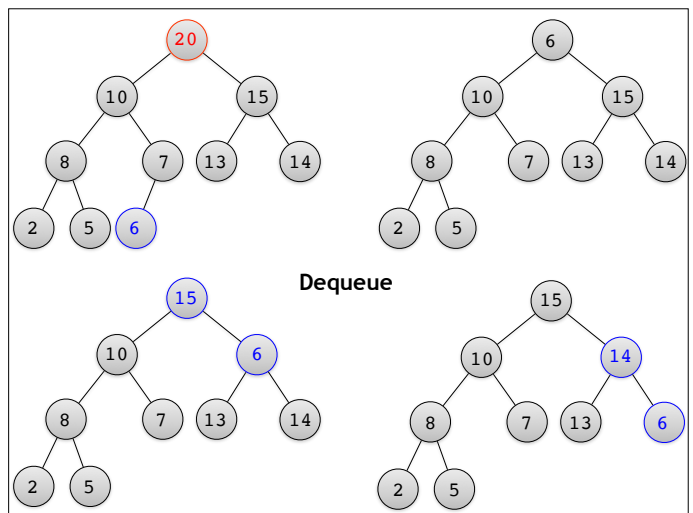
```

void enqueue(heap *heap, int item)
{
    if (heap->size >= MAX_HEAP_SIZE) {
        printf("Heap overflow");
    } else {
        heap->data[++heap->size] = item;

        int key_index = heap->size;
        int parent_index = get_parent_index(heap, heap->size);

        while (parent_index >= 1 &&
            heap->data[key_index] > heap->data[parent_index]) {
            swap(&heap->data[key_index], &heap->data[parent_index]);
            key_index = parent_index;
            parent_index = get_parent_index(heap, key_index);
        }
    }
}

```



```

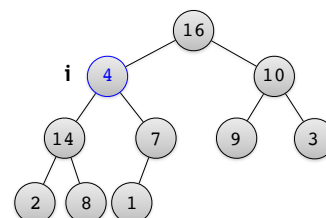
int dequeue(heap *heap)
{
    if (is_empty(heap)) {
        printf("Heap underflow");
        return -1;
    } else {
        int item = heap->data[1];
        heap->data[1] = heap->data[heap->size];
        heap->size--;
        max_heapify(heap, 1);
        return item;
    }
}

```

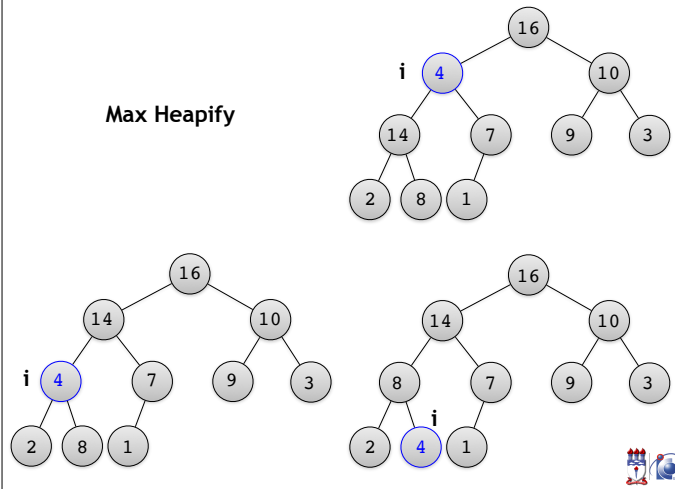


Maintaining the Heap property

- MAX-Heapify
- When called, it assumes that the binary trees rooted at LEFT(i) and RIGHT(i) are max-heaps



Max Heapify



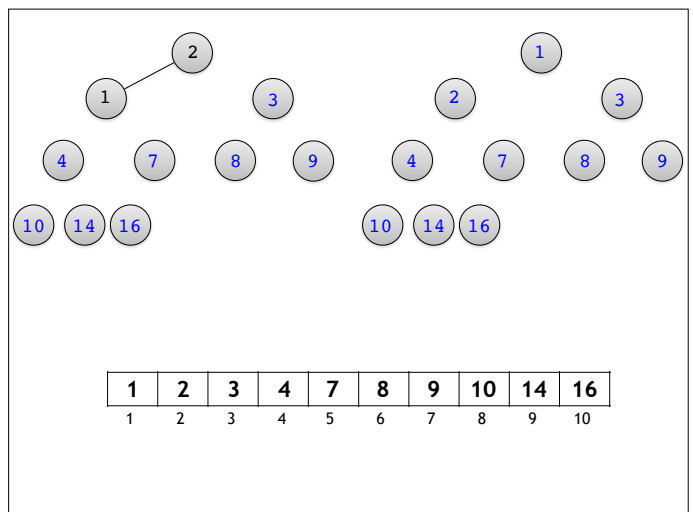
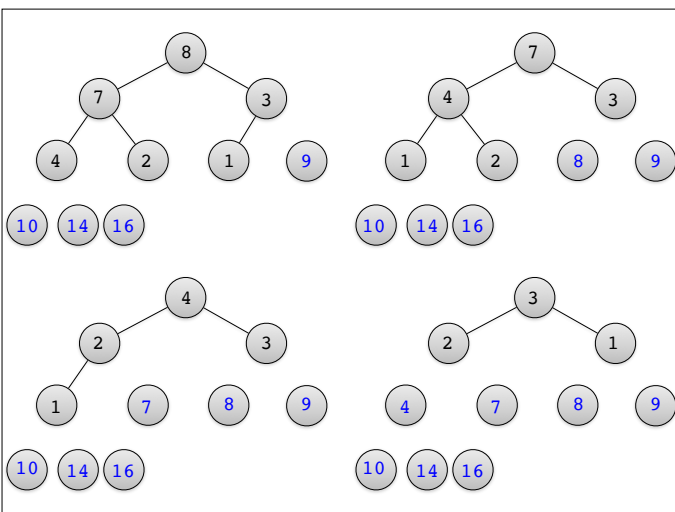
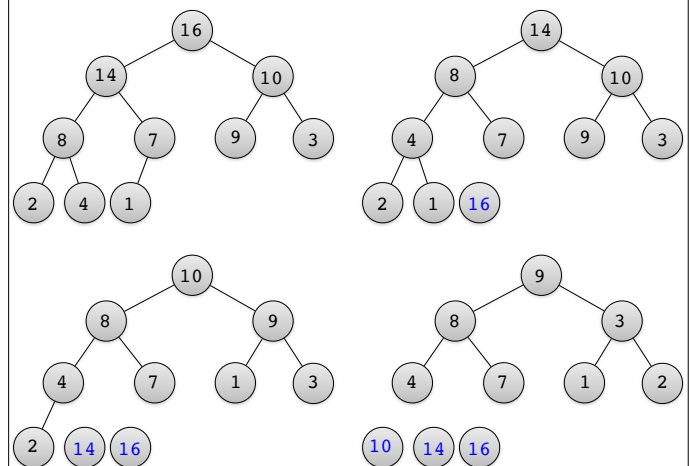
```
void max_heapify(heap *heap, int i)
{
    int largest;
    int left_index = get_left_index(heap, i);
    int right_index = get_right_index(heap, i);

    if (left_index <= heap->size &&
        heap->data[left_index] > heap->data[i]) {
        largest = left_index;
    } else {
        largest = i;
    }

    if (right_index <= heap->size &&
        heap->data[right_index] > heap->data[largest]) {
        largest = right_index;
    }

    if (heap->data[i] != heap->data[largest]) {
        swap(&heap->data[i], &heap->data[largest]);
        max_heapify(heap, largest);
    }
}
```

Heapsort



1	2	3	4	7	8	9	10	14	16
1	2	3	4	5	6	7	8	9	10

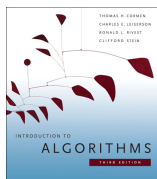
```
void heapsort(heap *heap)
{
    int i;
    for (i = heap->size; i >= 2; i--) {
        swap(&heap->data[1], &heap->data[i]);
        heap->size--;
        max_heapify(heap, 1);
    }
}
```



Execution: Heap Sort

6 5 3 1 8 7 2 4

References



Chapter 6



Chapters 6 and 9

