



# Estruturas de Dados / Programação 2
## *Union-Find*

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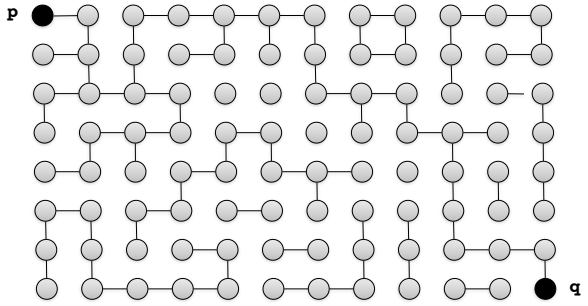
### Steps to develop an algorithm

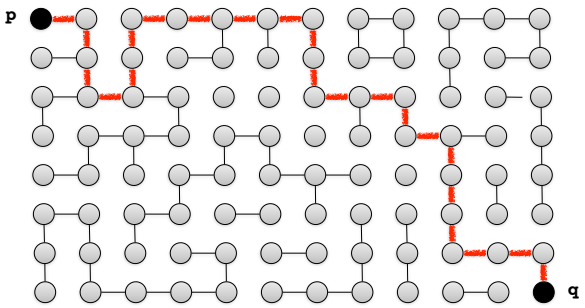
- Model the problem
- Find an algorithm to solve it
- Fast enough? Fits in memory?
- If not, figure out why
- Find a way to address the problem
- Iterate until satisfied

## Let's exercise these steps in this class!

### Network connectivity




### Network connectivity

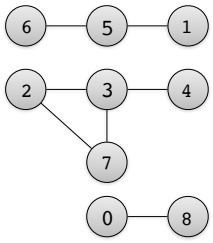



### Examples

- Variable names aliases
- Computers in a network
- Pixels in a digital photo

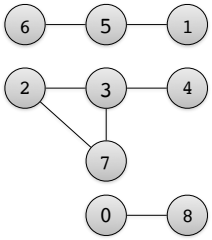
# Union-Find

## Connected nodes




{1 5 6} {2 3 4 7} {0 8}

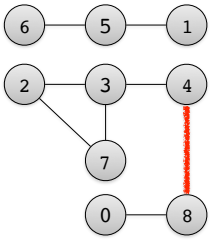
find(4, 8) = **false**




{1 5 6} {2 3 4 7} {0 8}

union(4, 8)

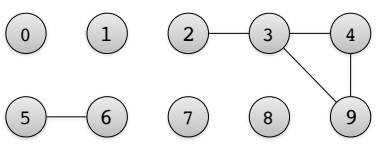
find(4, 8) = **true**




{1 5 6} {0 2 3 4 7 8}

# Quick find

## Data structure




| id[i] | 0 | 1 | 9 | 9 | 9 | 6 | 6 | 7 | 8 | 9 |
|---|---|---|---|---|---|---|---|---|---|---|
| i | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |

**p** and **q** are connected if they have the same id

## How does it work?

- Find: check if p and q have the same id
- Union: to merge sets containing p and q, change all entries with id[p] to id[q]

| id[i] | 0 | 1 | 9 | 9 | 9 | 6 | 6 | 7 | 8 | 9 |
|---|---|---|---|---|---|---|---|---|---|---|
| i | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |

Example: `union(3, 6)`

| id[i] | 0 | 1 | 6 | 6 | 6 | 6 | 6 | 7 | 8 | 6 |
|---|---|---|---|---|---|---|---|---|---|---|
| i | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |

# Abstract Data Type: Union-Find

## Union-Find ADT

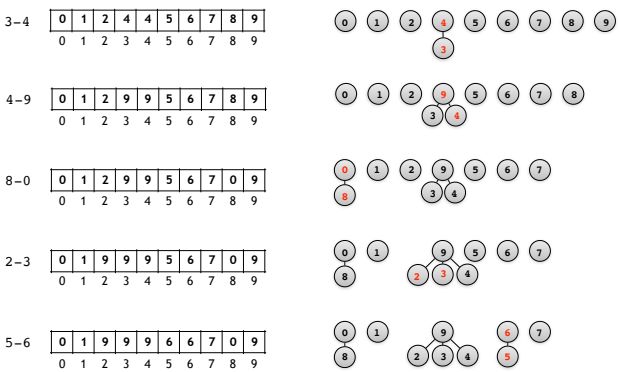
```
void initialize();

int find(int p, int q);

void union(int p, int q);
```

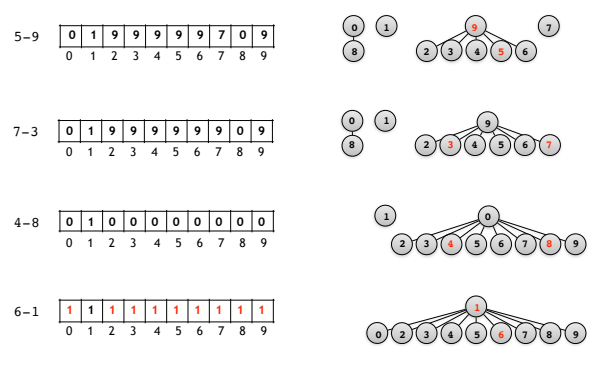
## Initialize and Find

```
void initialize()
{
  int i;
  for (i = 0; i < ARRAY_SIZE; i++) {
    id[i] = i;
  }
}

int find(int p, int q)
{
  return (id[p] == id[q]);
}
```

## Union

```
void union(int p, int q)
{
  int p_id = id[p];

  int i;
  for (i = 0; i < ARRAY_SIZE; i++) {
    if (id[i] == p_id) {
      id[i] = id[q];
    }
  }
}
```



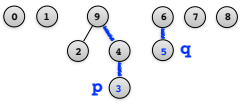




Problem: many values can change!

# Quick union

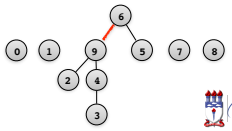
## Quick union

- Find: check if p and q have the same root
- Union: to merge sets containing p and q, set the id of q's root to the id of p's root
- id[i] is parent of i




| i | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
|---|---|---|---|---|---|---|---|---|---|---|
| id[i] | 0 | 1 | 9 | 4 | 9 | 6 | 6 | 7 | 8 | 9 |

| i | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
|---|---|---|---|---|---|---|---|---|---|---|
| id[i] | 0 | 1 | 9 | 4 | 9 | 6 | 6 | 7 | 8 | 6 |



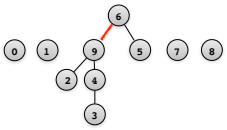

Only one value changes!

## Roots

```
int root(int i)
{
  while (i != id[i]) {
    i = id[i];
  }
  return i;
}

int find_with_roots(int p, int q)
{
  return (root(p) == root(q));
}

void union_with_roots(int p, int q)
{
  int p_root = root(p);
  int q_root = root(q);
  id[p_root] = q_root;
}
```

Any problem with this `union(3, 5)`?

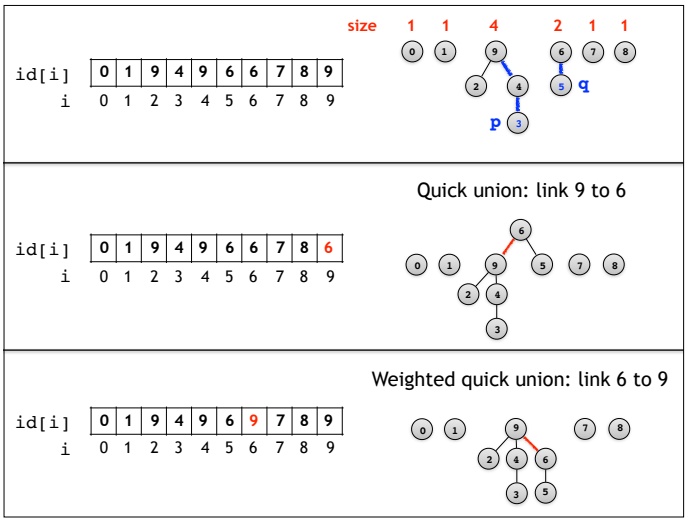
| i | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
|---|---|---|---|---|---|---|---|---|---|---|
| id[i] | 0 | 1 | 9 | 4 | 9 | 6 | 6 | 7 | 8 | 6 |



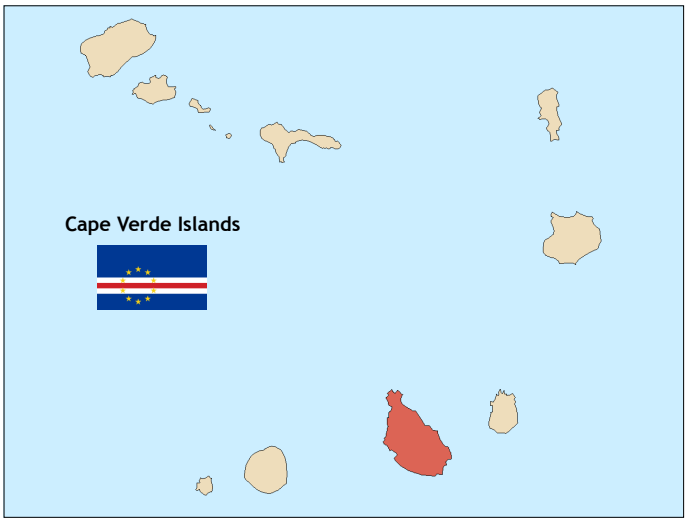

# Weighted roots

## Weighted roots

```
void union_with_weigthed_roots(int p, int q)
{
  int p_root = root(p);
  int q_root = root(q);

  if (size[p_root] < size[q_root]) {
    id[p_root] = q_root;
    size[q_root] = size[q_root] + size[p_root];
  } else {
    id[q_root] = p_root;
    size[p_root] = size[p_root] + size[q_root];
  }
}
```

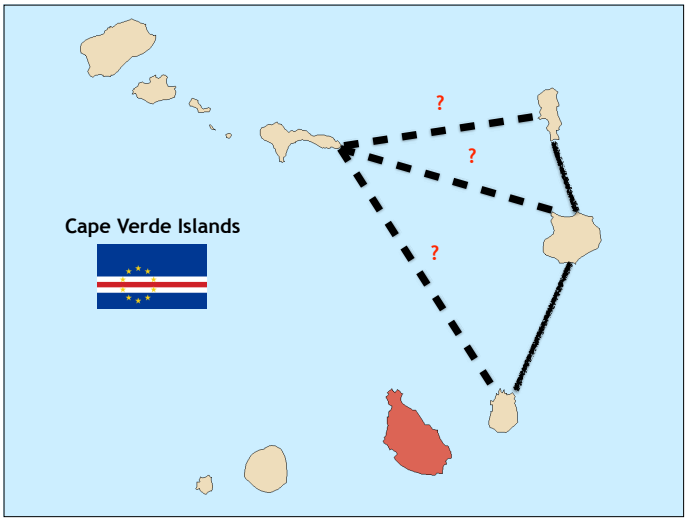


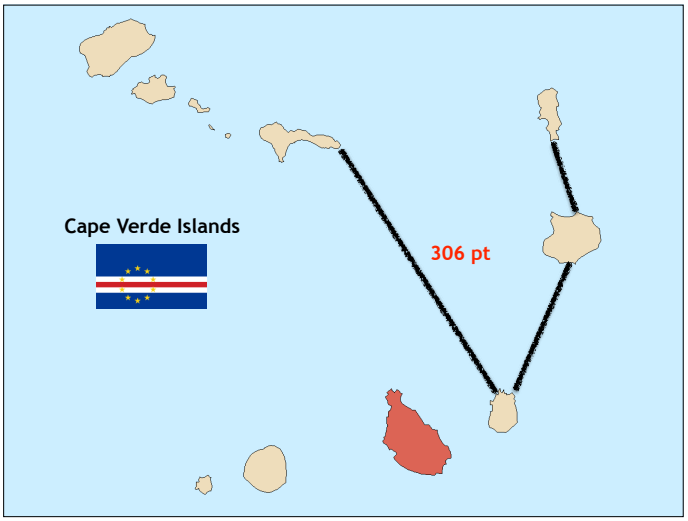

# Application

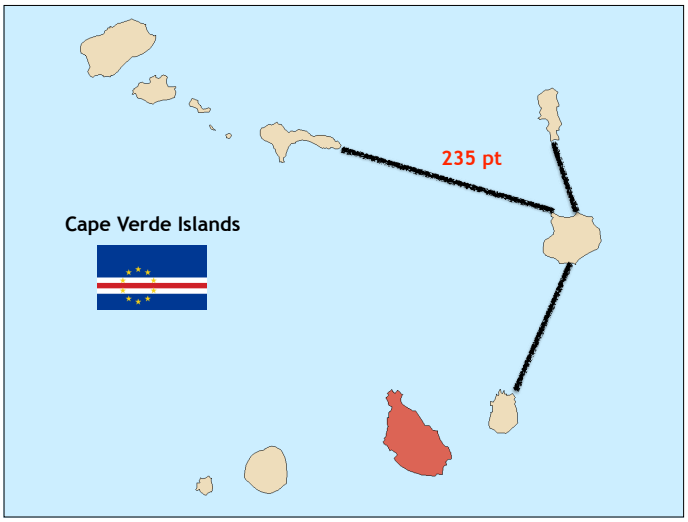


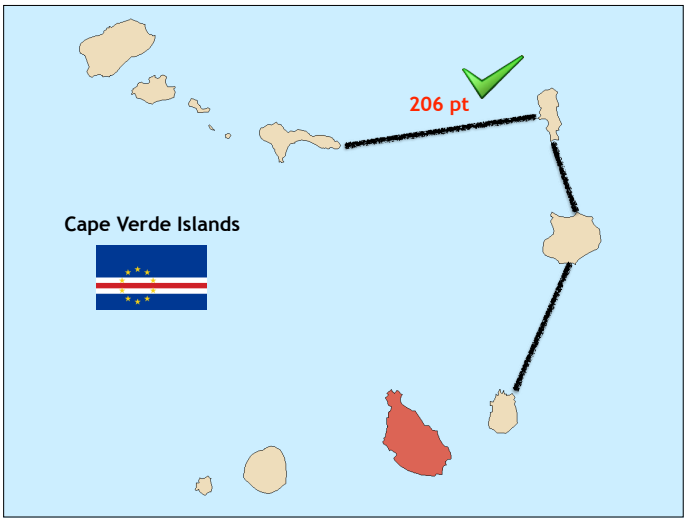

**Building bridges**

- The government wants to connect all islands
- The cost to build a bridge is proportional to the bridge's length
- How to connect all islands with the minimum cost?



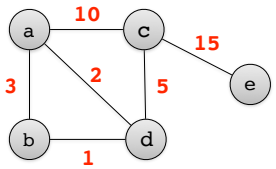






## Graph representation

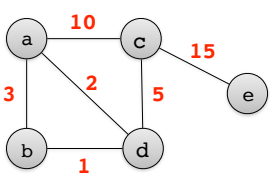
- 5 islands: the government wants to connect all
- The cost to build a bridge is proportional to the bridge's length

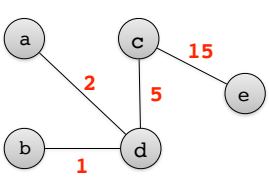



# Minimum Spanning Trees

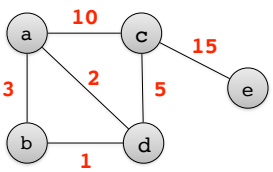
## Building bridges

- 6 islands: the government wants to connect all
- The cost to build a bridge is proportional to the bridge's length




## Kruskal Algorithm




V = {{a}, {b}, {c}, {d}, {e}}

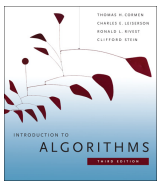
E = {{b,d}, {a,d}, {a,b}, {c,d}, {a,c}, {c,e}}

```
FIND-SET(b, d) = false
A = {} U {(b,d)} = {(b,d)}
V = {{a},{b,d},{c},{e}}

FIND-SET(a, d) = false
A = {(a,d)} U {(b,d)} = {(b,d),(a,d)}
V = {{a,b,d},{c},{e}}

FIND-SET(a, b) = true

FIND-SET(c, d) = false
A = {(c,d)} U {(b,d),(a,d)} = {(b,d),(a,d),(c,d)}
V = {{a,b,d,c},{e}}

FIND-SET(a, c) = true

FIND-SET(c, e) = false
A = {(c,e)} U {(b,d),(a,d),(c,d)} = {(b,d),(a,d),(c,d),(c,e)}
V = {{a,b,d,c,e}}
```

## Kruskal Algorithm

```
kruskal(G)

  A = ∅
  for each vertex v ∈ G.V
    MAKE-SET(v)

  sort(G.E)

  for each edge (p, q) ∈ G.E
    if !FIND-SET(p, q)
      A = A U {(p, q)}
      UNION(p, q)

  return A
```

## References




Chapter 21


Chapter 8