

CS M152A Project 4 Report

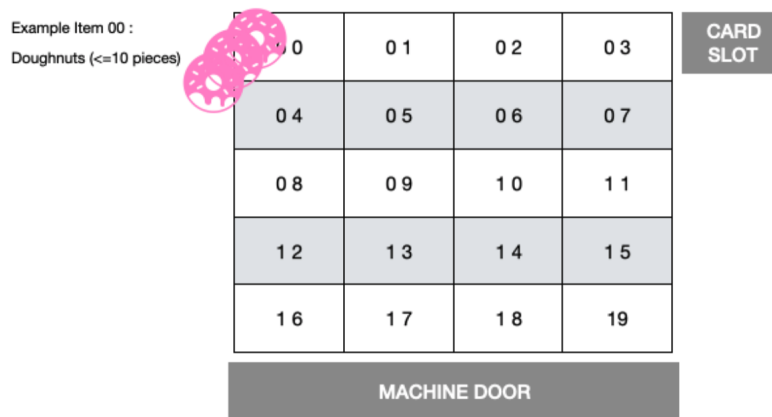
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1 Introduction and Requirement

The focus of this lab is for students to learn about how to design finite state machine that matches the specified behavior and to use the Xilinx ISE software to design and test state machines. Finite State Machines (FSMs) are used in many real world systems. An FSM has a finite number of states and can be in one state at a given time. There are two types of FSM machines: Moore Machine and Mealy Machine. A Moore machine's output only depends on which state it is in, but a Mealy machine's output depends on both the current state and input. For this assignment, we're tasked with designing a vending machine with the following characteristics:

1. Vending machine has 20 different snacks for sale. Each snack has two digit code (00 to 19).
2. Each snack is stored in separate slot. There can be up to 10 units of snack stored in 1 slot.
3. A buyer can purchase only 1 item at a time.
4. The machine only accepts payment by card.

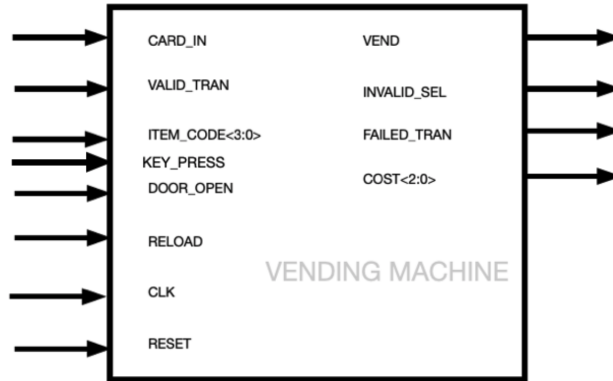


Vending Machine Figure from Manuscript

More details about specific behaviors of our vending machine FSM will be covered in the next sections.

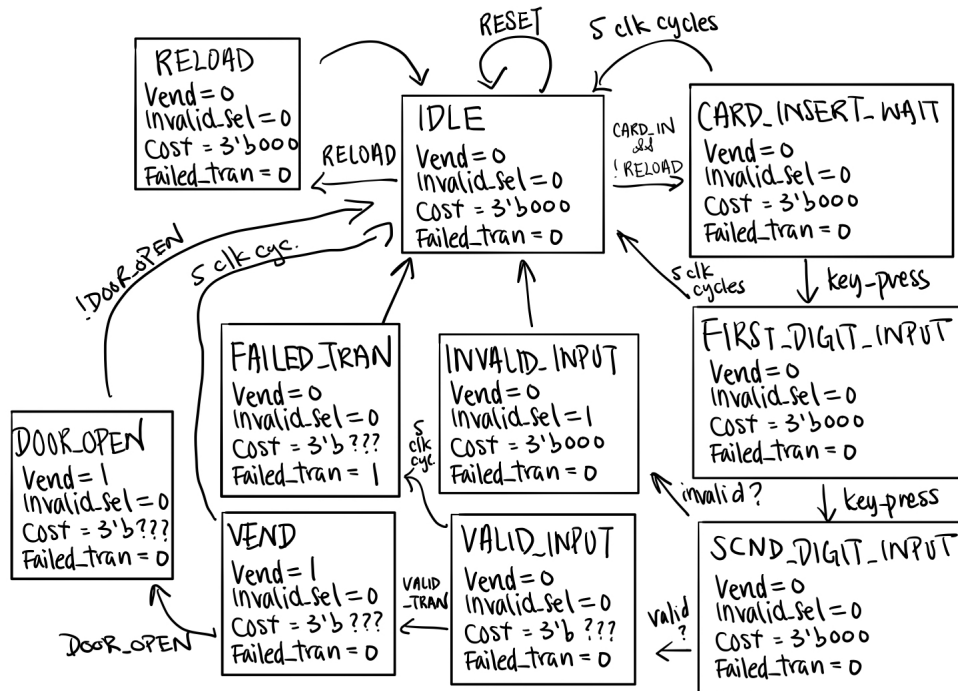
2 Design Description

For the design and implementation of the varying types of FSM, I followed guidance provided in the project specifications. I first consider the inputs provided and what my FSM will need to output. Our Vending Machine FSM will be given the following inputs, and we will need to output specific values for the outputs shown in the diagram below:



Vending Machine I/O from Manuscript

I then designed my FSM with around 10 states that correctly captures the behavior described in the manuscript. My FSM is a Moore machine, meaning that my output depends only on which state I am in, regardless of the inputs.



* if RESET signal is 1 in any state, return to IDLE state.
 All counters set to 0. All signal (arrow) above includes "and !RESET", except RESET arrow, not drawn to avoid repetition.

Moore FSM Design of Vending Machine

Descriptions of states and transitions of above FSM:

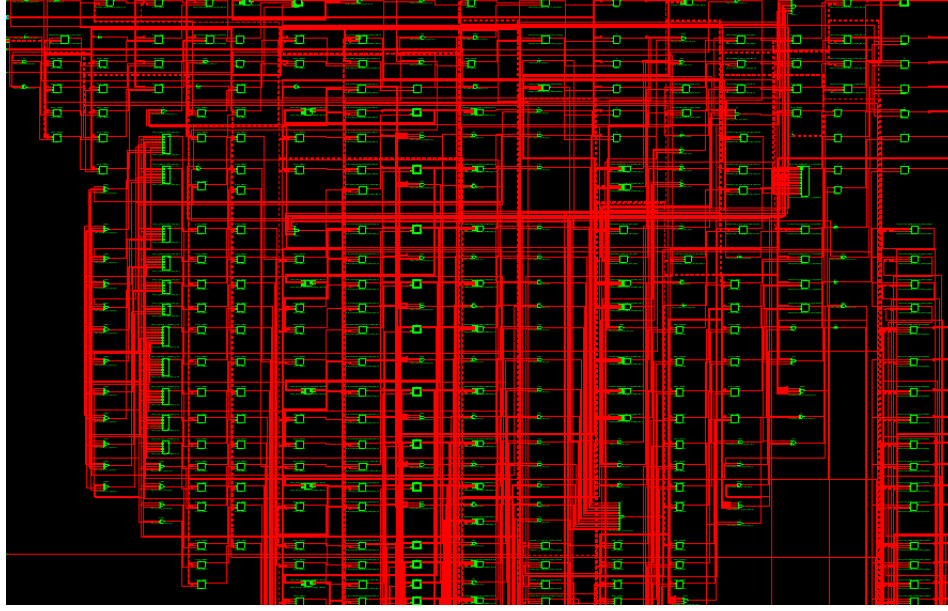
- Once RESET signal is 1, we will begin at the IDLE state, where our Vending Machine waits for the RELOAD or CARD_IN signal to be 1. All outputs should be 0, including count of items inside vending machine. This represents the state where no one is using our vending machine yet.
- Once RELOAD signal is 1, we transition to the RELOAD state where all our outputs remain unchanged, but the internal count for each slot in the vending machine should be set to 0. Once counts are set to 0, we directly return to IDLE state.

- From the IDLE state, when RELOAD signal is 0, and CARD_IN signal is 1, we transition into the CARD_INSERT_WAIT state. Outputs remain all zeroes.
- In CARD_INSERT_WAIT state, we wait for either 5 clock cycles to pass and transition back to the IDLE state or KEY_PRESS = 1 and transition to FIRST_DIGIT_INPUT state. After 5 cycles has passed, this represents our vending machine not receiving any code entered. When we detect a key press, we want to transition to next state where we prepare to read in a second key press.
- In FIRST_DIGIT_INPUT state, all our outputs remain 0. We check if 5 clock cycles has passed and return to the IDLE state, or we check KEY_PRESS = 1 and transition to SCND_DIGIT_INPUT state.
- In SCND_DIGIT_INPUT state, all our outputs remain 0. We have detected two different ITEM_CODE, so we have check whether the code represents a valid input. If the number falls between 00-19 and the item is still available, we transition to the VALID_INPUT state. Otherwise, we transition to the INVALID_INPUT state.
- In INVALID_INPUT state, we output INVALID_SEL = 1 and return to IDLE_STATE.
- In VALID_INPUT state, we output the COST corresponding to the valid input numbers, and rest of outputs remain 0. After 5 clock cycles, if VALID_TRAN remains 0, we transition to the FAILED_TRAN state. If VALID_TRAN = 1 within 5 cycles, we transition to the VEND state. This signifies the the card payment was accepted.
- In FAILED_TRAN state, we output FAILED_TRAN = 1, cost of item selected, and rest of outputs remain 0. We then return to the IDLE state in the next clock cycle.
- In VEND state, we output cost of item and VEND = 1. Rest of outputs remain 0. We wait 5 clock cycles for DOOR_OPEN to be 1 to transition into the DOOR_OPEN state. If door remains closed for 5 cycles, we return to IDLE state.
- In DOOR_OPEN state, we continue to output cost of item, VEND = 1, and rest of outputs are 0. We wait for DOOR_OPEN = 0 to transition to IDLE state, signifying end of our transaction. We remain in this state until DOOR_OPEN = 0 or we receive the reset signal.
- Note: If RESET = 1 at any point in our state diagram, it overrides other signal and arrows, reset all counts in vending machine to zero, and transitions to IDLE.

To implement the above FSM in Verilog, I created the `vending_machine` top module that takes in input listed in the Vending Machine I/O Diagram and outputs the outputs shown in the same diagram. Within my module, I initialized parameters representing all the states I have with a unique 4 bit binary number, two 4-bit registers `current_state` and `next_state`, and 20 4-bit registers corresponding to the count for each item in the vending machine. I have 3 main always block each with a specific task, similar to the sample code provided in the manuscript:

1. Sequential always block to update `current_state` to `next_state` at every `posedge` of the clock or update `current_state` to IDLE state when `RESET = 1`.
2. Combinational always block to decide what `next_state` should be based on the current state and current inputs.
3. Combinational always block to decide what outputs should be based on the current state and values stored in registers from inputs, such as which item is selected.

Aside from the blocks mentioned above, I have two more sequential always block that helps me with counting when 5 cycles has passed once we enter a particular state, such as `CARD_INPUT_WAIT` state. One always block uses a counter we implemented in previous projects and the other always block is in charge of setting a signal `clk_start` to signify that counter should start counting up.



RTL Schematic of Vending Machine FSM

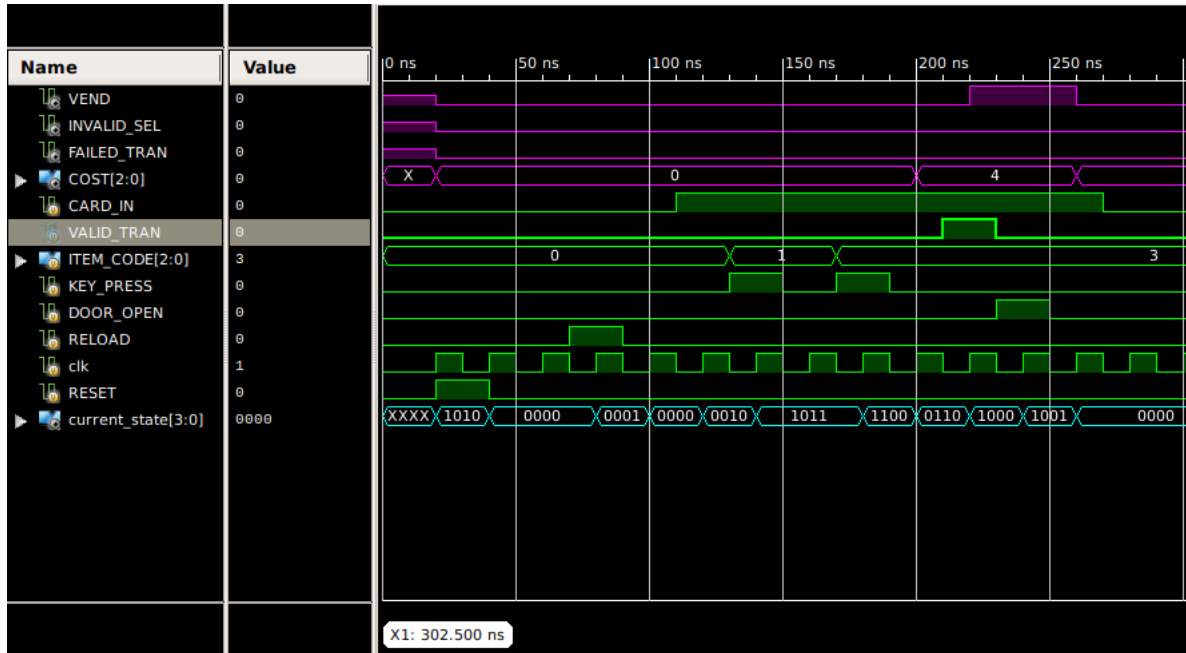
From first look, the RTL schematic above looks filled with registers, muxes, and counters. This makes sense as our FSM is fairly complicated having to keep track of 20 registers for each item in our vending machine on top of keeping track of the combinational logic to calculate the next state and the outputs. My Verilog code results in this schematic being generated as my code has 2 large combinational logic sections and requires many registers to keep track of internal signals to calculate the outputs and the next state.

3 Simulation Documentation

To test whether the FSM I implemented works as the manuscript described, I created my testbench that tests not only the successful runs of ordering an item from the vending machine, but all of the possible special cases, such as no key is pressed or an item has ran out. To simulate behavior of someone interacting with the vending machine, in my testbench, I created a system clock with $T=10\text{ns}$, and at each negative edge I change a particular input to observe how my FSM will deal with the change in input. Below are all the different cases I tested and their respective waveform.

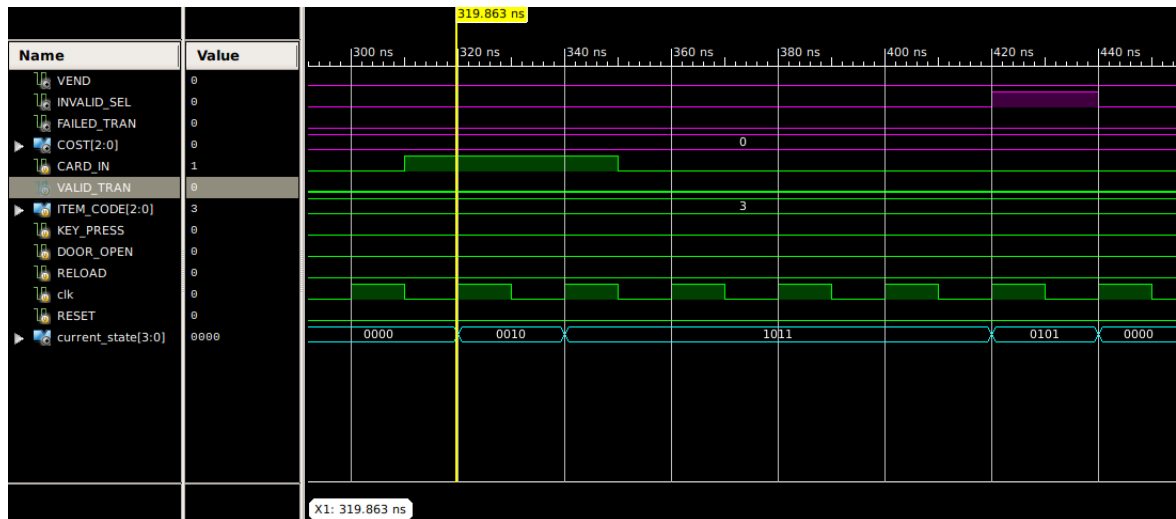
1. Successful Purchase of Valid In-Stock Item

This case represents a successful purchase from the vending machine machine where there were no errors all the way until vending machine door closes. The last light blue wave in the diagram below is an extra waveform I chose to display as output in order to show transition from one state to another. It is the binary representation of which state we're in. The way I simulated this behavior in the test bench is by first setting `RESET` as high to initialize all outputs to 0 and transition to the `IDLE` state. I then set `RELOAD = 1` in order to fill up the vending machine with stocks(shown at 70ns). Now to simulate someone using the machine, I set input `CARD_IN = 1` (shown at 110ns) and we're now in state `4'b0010 CARD_INSERT_WAIT` state. So far all output remains 0 as no mistakes has occurred and nothing has been inputted yet. We then set `KEY_PRESS` to high twice in the next 3 cycles in order to simulating pressing keys 1 and 3 into the machine(130ns-200ns). At 200ns, we transition into state `4'b0110 VALID_INPUT` as 13 corresponds to a valid in stock item in our machine. We correctly see that output `COST = 4` at 200ns and beyond, since from the manuscript cost of item 13 is 4. We then set `VALID_TRAN = 1` to represent the acceptance of our payment and once this signal is process we see that output `VEND = 1`. At 230ns, we set `DOOR_OPEN = 1` for one cycle to simulate item coming out of the machine. Once `DOOR_OPEN = 0` at 250ns, we see that `COST = 0` and `VEND = 0` as we have returned to `IDLE` state, which is the correct behavior.



2. Unsuccessful Purchase, No Key Press Detected

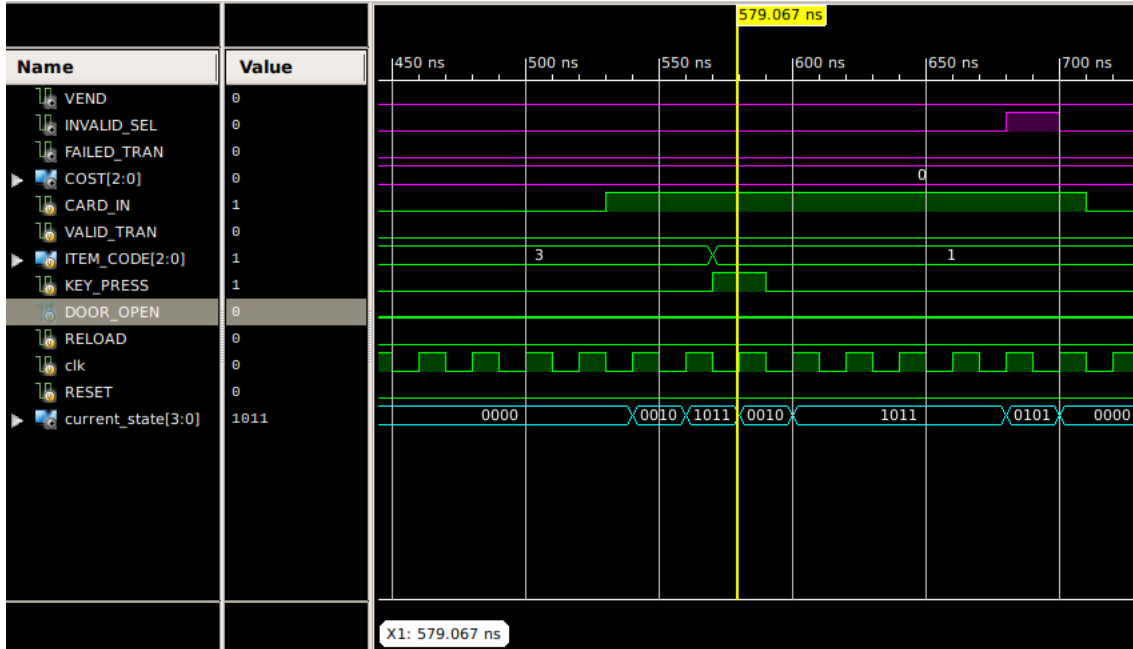
This case represents an unsuccessful purchase where we return to IDLE state as after 5 clock cycles, no key press is detected. To simulate this case, we simply set `CARD_IN = 1` to transition into state `4'b0010 CARD_INSERT_WAIT` and observe what happens after 5 clock cycles. At 320ns, we enter `CARD_INSERT_WAIT`, where all outputs are 0, and at 420ns, we see `INVALID_SEL = 1` for one cycle. One cycle is 20ns, thus our FSM correctly waits for 5 cycles before entering `4'b0101 INVALID_INPUT` state, and outputting `INVALID_SEL = 1`. After one cycle, we return to IDLE state where all output is 0. Thus, our FSM behaves as described by the manuscript in this edge case.



3. Unsuccessful Purchase, Only 1 Key Press Detected

This case represents an unsuccessful purchase where we return to IDLE state as we detect 1 key press, but after 5 cycles, we do not detect a second key press. To simulate this case, we simply set `CARD_IN = 1` to transition into state `4'b0010 CARD_INSERT_WAIT`, we then set `KEY_PRESS = 1` for one cycle,

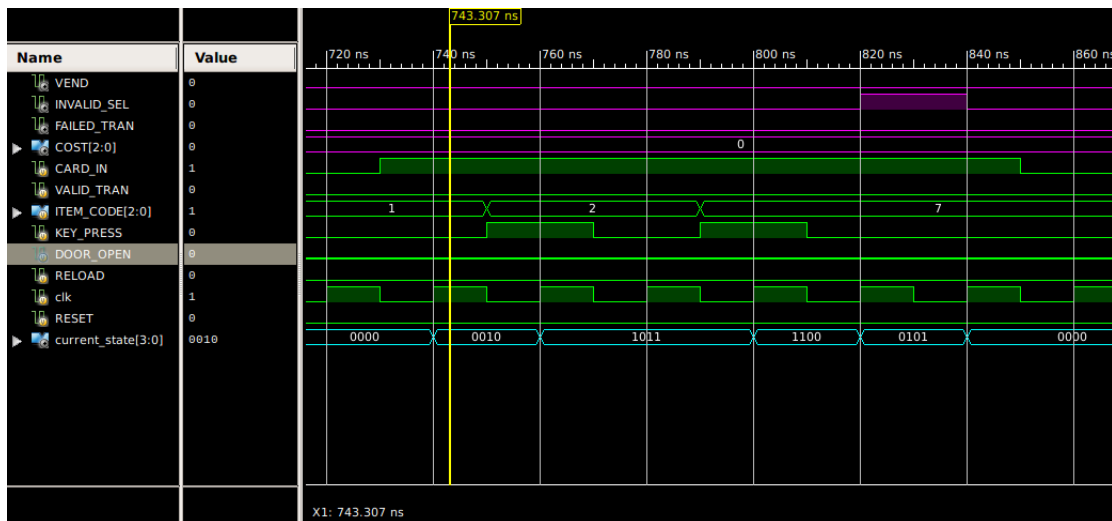
and observe what happens after 5 clock cycles. At 580ns, we detected and registered the first key press, all outputs remain 0 and we begin our counter. At 680ns, we see `INVALID_SEL = 1` for one cycle. One cycle is 20ns, thus our FSM correctly waits for 5 cycles before entering `4'b0101 INVALID_INPUT` state, and outputting `INVALID_SEL = 1`. After one cycle, we return to `IDLE` state where all output is 0. Thus, our FSM behaves as described by the manuscript in this edge case.



Simulation Waveform for Case 3

4. Unsuccessful Purchase, Out of Range Item Number Entered

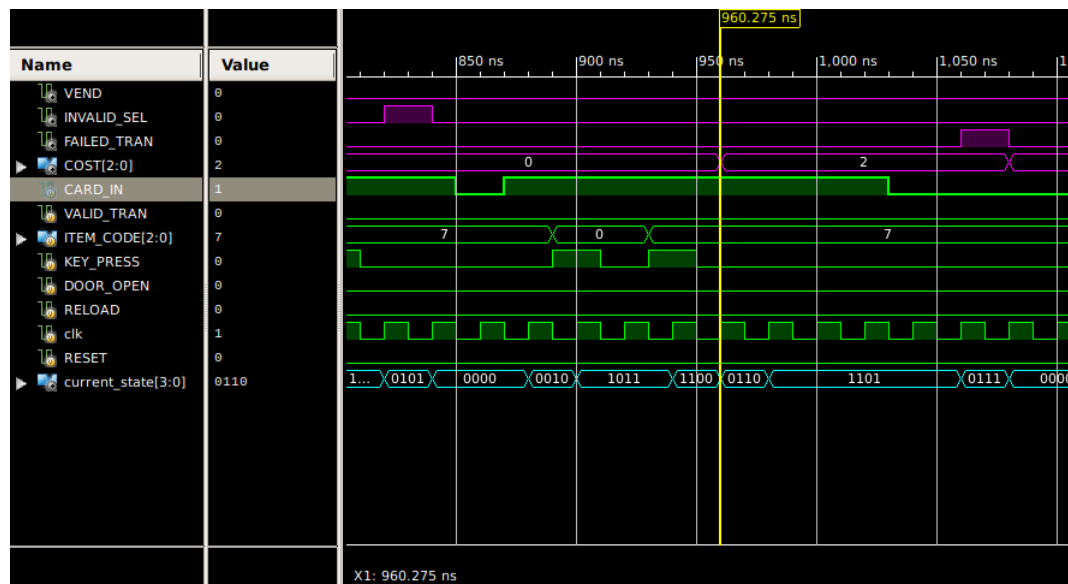
This case represents an unsuccessful purchase where two key press are detected, but the number they represent is out of range, so we want `INVALID_SEL = 1`. To simulate this case, we simply set `CARD_IN = 1` to transition into state `4'b0010 CARD_INSERT_WAIT`, we then set `KEY_PRESS = 1` for one cycle, and then set `KEY_PRESS = 1` for one cycle. The two keys we entered are 2, 7, representing an item that doesn't exist in our vending machine. At 760ns, we see that 2 is read by our FSM, and then at 800ns, 7 is read by our FSM. One cycle after 7 is read, we see that at 820ns, `INVALID_SEL = 1` for one cycle. This is the correct behavior as 27 is an invalid input.



Simulation Waveform for Case 4

5. Unsuccessful Purchase, Invalid Transaction

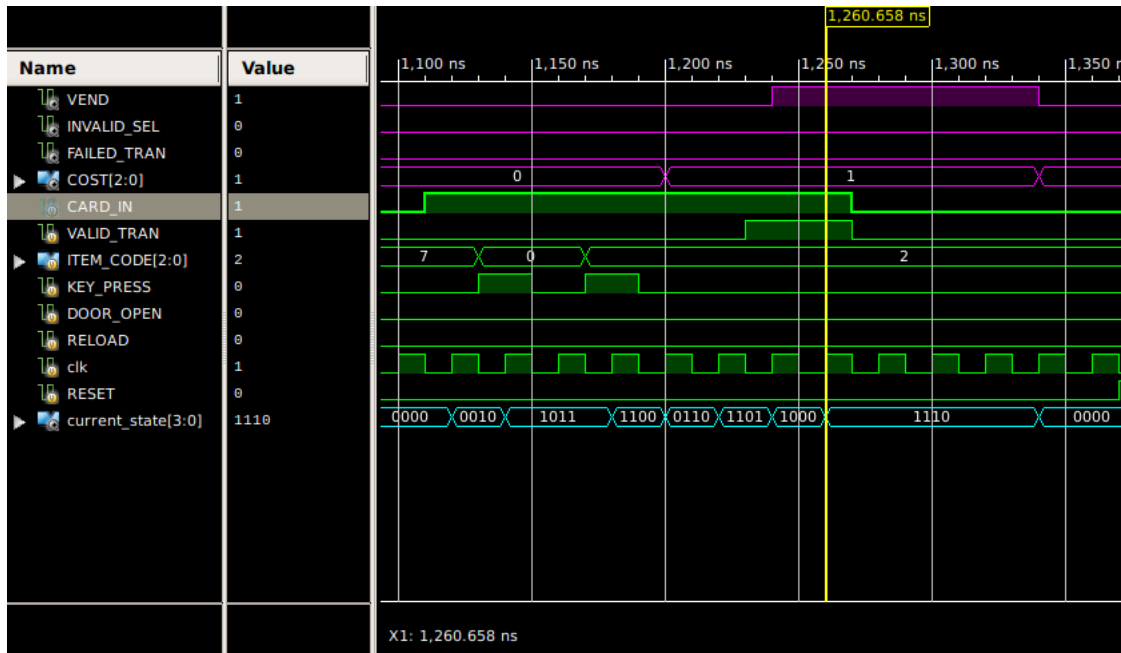
This case represents an unsuccessful purchase where valid keys are inputted and the item requested has stock, but our card transaction has failed. To simulate this case, we set **CARD_IN** = 1 to transition into state 4'b0010 **CARD_INSERT_WAIT**, we then set **KEY_PRESS** = 1 for one cycle, and then set **KEY_PRESS** = 1 for one cycle. From diagram below, we can see that we entered the key 0 and then the key 7, which is a valid item with stock as the machine was reloaded previously. At 960ns, we enter state 4'b0110 **VALID_INPUT** and wait for 5 cycles to pass. We see that output **COST** is correctly 2. At 1060ns, we see correctly observe **FAILED_TRAN** = 1.



Simulation Waveform for Case 5

6. Unsuccessful Purchase, Door Does Not Open

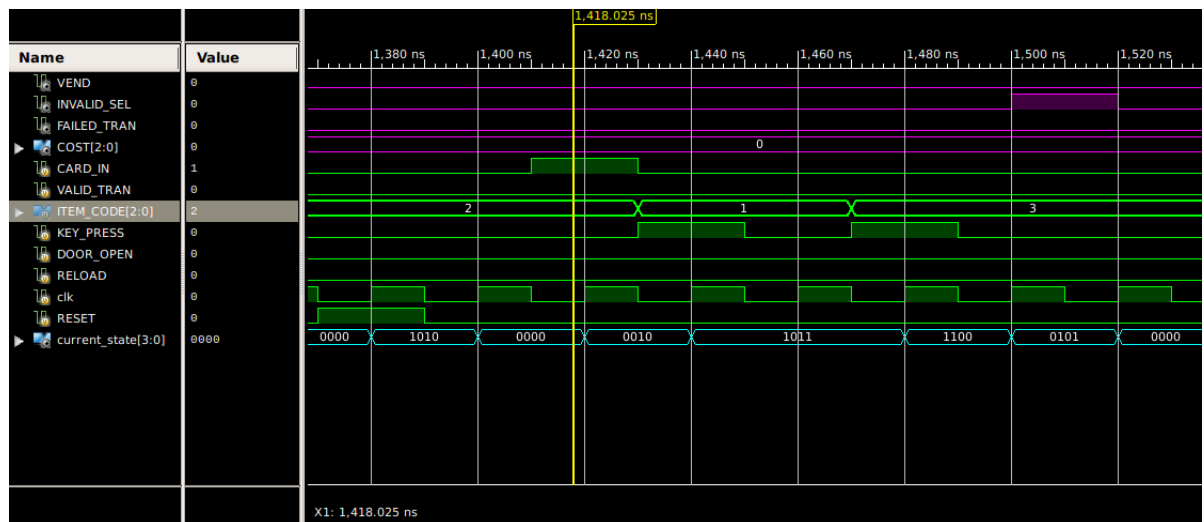
This case represents an unsuccessful purchase where valid keys are inputted, and the transaction is valid, but the door of vending machine does not open. To simulate this case, we set **CARD_IN** = 1 to transition into state 4'b0010 **CARD_INSERT_WAIT**, we then set **KEY_PRESS** = 1 for one cycle, and then set **KEY_PRESS** = 1 for one cycle. From diagram below, we can see that we entered the key 0 and then the key 2, which is a valid item with stock as the machine was reloaded previously. Thus, we see that output **COST** is correctly 2. At 1240ns, the **VALID_TRAN** signal is detected by our FSM, so output **VEND** = 1. From 1240ns on, we are in the state where we're waiting for either door to open or 5 cycles to pass. At 1340ns, 5 cycles has passed and door did not open, so we set all output to 0 and return to **IDLE** state.



Simulation Waveform for Case 6

7. Unsuccessful Purchase, Valid Input, No Stock

This case represents an unsuccessful purchase where valid keys are inputted, but because we have no more stock of the requested item left, we set `INVALID_SEL = 1` and return to IDLE state. To simulate this case, we have to first make sure there is no more stock of a particular item. To do this, we set `RESET = 1` for one cycle to make sure all counts of our item is 0. We then set `CARD_IN = 1` to transition into state 4'b0010 `CARD_INSERT_WAIT`, we then set `KEY_PRESS = 1` for one cycle, and then set `KEY_PRESS = 1` for one cycle. From diagram below, we can see that we entered the key 1 and then the key 3, which is a valid item number as it is within 00-19. But since our vending machine is out of stock, we can see that at 1500ns, `INVALID_SEL = 1` for one cycle after second key is entered. This is the correct expected behavior. All other outputs remain correctly at 0.



Simulation Waveform for Case 7

8. Case Where Signals Change

This case occurs throughout the waveform diagrams above as often the `CARD_IN` signal stays high throughout transaction or goes low in the middle of a transaction. We can see that we correctly handle this case as regardless of how long `CARD_IN` signal stays high, our module treats it correctly. If card is taken out earlier, it is okay, as our vending machine remembers the card info. So no odd behavior will happen when signals change arbitrarily.

One bug I found during simulation for this project was that my counter that keeps track of whether 5 cycle has passed was incorrect, as it continues to count up even in cases where the counter is not needed. I fixed this by adding an additional internal register to signal when count should be incremented.

4 Synthesis and Implementation Report

Sections of Synthesis and Implementation Report is attached at end of report. From the 'Design Summary' part of the synthesis report, we can conclude the different components that my implementation of the Vending Machine FSM uses. For example, we can see that we indeed do use a lot of registers from the synthesis report and we can also see the different clock signals needed for our module. From implementation report, we can see that there is no errors when trying to implement our module, but there are a few warnings due to truncation of numbers into registers. The warnings can be ignored as in my code I made sure that truncation will not occur for those registers.

5 Conclusion

In this project, I designed and implemented the Vending Machine FSM according to the behavior in the project manuscript. I designed the Vending Machine based on a Moore Machine, that has similar structure to the turnstile FSM example in the manuscript. To correctly implement this, I had to first draw the FSM diagram to allow me to clearly see what output should be in each state and how to transition to a different state. One of the major difficulty I encountered in this assignment was trying to incorporate a counter in order to keep track of whether 5 cycles has passed. I dealt with this by experimenting my code on a smaller, simpler module in order to see how to correctly use different signal within my module to signal counter to start counting and to detect output of the counter.

6 Reports

6.1 Synthesis Report

```
=====
*                               Design Summary                               *
=====

Top Level Output File Name      : vending_machine.ngc

Primitive and Black Box Usage:
-----
# BELS                          : 226
#     INV                      : 18
#     LUT2                     : 5
#     LUT3                     : 54
#     LUT4                     : 37
#     LUT5                     : 31
#     LUT6                     : 70
#     MUXF7                    : 11
# FlipFlops/Latches            : 89
#     FDR                     : 37
#     FDRE                    : 35
#     FDS                     : 4
#     LD                      : 13
# Clock Buffers                : 1
#     BUFGP                   : 1
# IO Buffers                   : 15
#     IBUF                    : 9
#     OBUF                    : 6

Device utilization summary:
-----

Selected Device : 6slx16csg324-3

Slice Logic Utilization:
Number of Slice Registers:      83 out of 18224    0%
Number of Slice LUTs:          215 out of 9112     2%
    Number used as Logic:      215 out of 9112     2%

Slice Logic Distribution:
Number of LUT Flip Flop pairs used: 220
    Number with an unused Flip Flop: 137 out of 220    62%
    Number with an unused LUT:      5 out of 220     2%
    Number of fully used LUT-FF pairs: 78 out of 220    35%
    Number of unique control sets: 27

IO Utilization:
Number of IOs:                  16
Number of bonded IOBs:          16 out of 232     6%
    IOB Flip Flops/Latches:      6

Specific Feature Utilization:
Number of BUFG/BUFGCTRLs:       1 out of 16        6%
-----
```

Partition Resource Summary:

No Partitions were found in this design.

Timing Report

NOTE: THESE TIMING NUMBERS ARE ONLY A SYNTHESIS ESTIMATE.
FOR ACCURATE TIMING INFORMATION PLEASE REFER TO THE TRACE REPORT
GENERATED AFTER PLACE-and-ROUTE.

Clock Information:

Clock Signal	Clock buffer (FF name)	
_n1087(_n10871:0)	NONE(*) (INVALID_SEL)	1
_n1084(_n10841:0)	NONE(*) (FAILED_TRAN)	1
_n1079(_n10791:0)	NONE(*) (VEND)	1
CLK	BUFGP	7
current_state[3]_GND_6_o_Mux_472_o(Mmux_current_state[3]_GND_6_o_Mux_472_o11:0)	NONE(*) (code1_2)	3
current_state[3]_GND_2_o_Mux_464_o(Mmux_current_state[3]_GND_2_o_Mux_464_o11:0)	NONE(*) (code2_2)	3
current_state[3]_GND_10_o_Mux_480_o(Mmux_current_state[3]_GND_10_o_Mux_480_o1:0)	NONE(*) (code1inputted)	1
current_state[3]_PWR_14_o_Select_533_o(Mmux_current_state[3]_PWR_14_o_Select_533_o11:0)	NONE(*) (COST_2)	3

(*) These 7 clock signal(s) are generated by combinatorial logic,
and XST is not able to identify which are the primary clock signals.
Please use the CLOCK_SIGNAL constraint to specify the clock signal(s) generated by combinatorial logic.
INFO:Xst:2169 - HDL ADVISOR - Some clock signals were not automatically buffered by XST with BUFG/BUFR resources. P

Asynchronous Control Signals Information:

No asynchronous control signals found in this design

Timing Summary:

Speed Grade: -3

Minimum period: 4.511ns (Maximum Frequency: 221.683MHz)
Minimum input arrival time before clock: 5.316ns
Maximum output required time after clock: 3.648ns
Maximum combinational path delay: No path found

Timing Details:

All values displayed in nanoseconds (ns)

Timing constraint: Default period analysis for Clock 'CLK'
Clock period: 4.511ns (frequency: 221.683MHz)
Total number of paths / destination ports: 1267 / 179

Delay: 4.511ns (Levels of Logic = 3)
Source: count_2 (FF)
Destination: count_0 (FF)

Source Clock: CLK rising
Destination Clock: CLK rising

Data Path: count_2 to count_0

Cell:in->out	fanout	Gate		Net	Logical Name (Net Name)
		Delay	Delay		
FDRE:C->Q	4	0.447	0.912	count_2	count_2 (count_2)
LUT3:I0->O	11	0.205	0.883	GND_1_o_GND_1_o_equal_375_o<2>1	(GND_1_o_GND_1_o_equal_375_o)
LUT6:I5->O	1	0.205	0.684	Mmux_count_up22	(Mmux_count_up21)
LUT6:I4->O	3	0.203	0.650	Mmux_count_up25	(count_up)
FDRE:CE		0.322		count_0	
Total		4.511ns (1.382ns logic, 3.129ns route) (30.6% logic, 69.4% route)			

=====
Timing constraint: Default OFFSET IN BEFORE for Clock 'current_state[3]_GND_6_o_Mux_472_o'
Total number of paths / destination ports: 3 / 3
=====

Offset: 1.875ns (Levels of Logic = 1)
Source: ITEM_CODE<2> (PAD)
Destination: code1_2 (LATCH)
Destination Clock: current_state[3]_GND_6_o_Mux_472_o falling

Data Path: ITEM_CODE<2> to code1_2

Cell:in->out	fanout	Gate		Net	Logical Name (Net Name)
		Delay	Delay		
IBUF:I->O	2	1.222	0.616	ITEM_CODE_2_IBUF	(ITEM_CODE_2_IBUF)
LD:D		0.037		code1_2	
Total		1.875ns (1.259ns logic, 0.616ns route) (67.1% logic, 32.9% route)			

=====
Timing constraint: Default OFFSET IN BEFORE for Clock 'current_state[3]_GND_2_o_Mux_464_o'
Total number of paths / destination ports: 6 / 3
=====

Offset: 2.449ns (Levels of Logic = 2)
Source: KEY_PRESS (PAD)
Destination: code2_2 (LATCH)
Destination Clock: current_state[3]_GND_2_o_Mux_464_o falling

Data Path: KEY_PRESS to code2_2

Cell:in->out	fanout	Gate		Net	Logical Name (Net Name)
		Delay	Delay		
IBUF:I->O	11	1.222	0.987	KEY_PRESS_IBUF	(KEY_PRESS_IBUF)
LUT4:I2->O	1	0.203	0.000	current_state[3]_code2[2]_Mux_465_o1	(current_state[3]_code2[2]_Mux_465_o1)
LD:D		0.037		code2_2	
Total		2.449ns (1.462ns logic, 0.987ns route) (59.7% logic, 40.3% route)			

=====
Timing constraint: Default OFFSET IN BEFORE for Clock 'current_state[3]_PWR_14_o_Select_533_o'
Total number of paths / destination ports: 3 / 3
=====

Offset: 3.491ns (Levels of Logic = 3)
Source: RELOAD (PAD)
Destination: COST_2 (LATCH)
Destination Clock: current_state[3]_PWR_14_o_Select_533_o falling

Data Path: RELOAD to COST_2

Cell:in->out	fanout	Gate Delay	Net Delay	Logical Name (Net Name)
IBUF:I->0	6	1.222	0.745	RELOAD_IBUF (RELOAD_IBUF)
LUT5:I4->0	5	0.205	1.079	_n10871 (_n1087)
LUT6:IO->0	1	0.203	0.000	Mmux_current_state[3]_GND_1_o_Select_536_o11 (current_state[3]_GND_1_
LD:D		0.037		COST_0

Total		3.491ns (1.667ns logic, 1.824ns route)		
		(47.8% logic, 52.2% route)		

=====
Timing constraint: Default OFFSET IN BEFORE for Clock 'CLK'
Total number of paths / destination ports: 66 / 19

Offset: 5.316ns (Levels of Logic = 4)
Source: KEY_PRESS (PAD)
Destination: count_0 (FF)
Destination Clock: CLK rising

Data Path: KEY_PRESS to count_0

Cell:in->out	fanout	Gate Delay	Net Delay	Logical Name (Net Name)
IBUF:I->0	11	1.222	0.883	KEY_PRESS_IBUF (KEY_PRESS_IBUF)
LUT2:I1->0	1	0.205	0.944	Mmux_count_up21 (Mmux_count_up2)
LUT6:IO->0	1	0.203	0.684	Mmux_count_up22 (Mmux_count_up21)
LUT6:I4->0	3	0.203	0.650	Mmux_count_up25 (count_up)
FDRE:CE		0.322		count_0

Total		5.316ns (2.155ns logic, 3.161ns route)		
		(40.5% logic, 59.5% route)		

=====
Timing constraint: Default OFFSET OUT AFTER for Clock 'current_state[3]_PWR_14_o_Select_533_o'
Total number of paths / destination ports: 3 / 3

Offset: 3.648ns (Levels of Logic = 1)
Source: COST_2 (LATCH)
Destination: COST<2> (PAD)
Source Clock: current_state[3]_PWR_14_o_Select_533_o falling

Data Path: COST_2 to COST<2>

Cell:in->out	fanout	Gate Delay	Net Delay	Logical Name (Net Name)
LD:G->Q	1	0.498	0.579	COST_2 (COST_2)
OBUF:I->0		2.571		COST_2_OBUF (COST<2>)

Total		3.648ns (3.069ns logic, 0.579ns route)		
		(84.1% logic, 15.9% route)		

Timing constraint: Default OFFSET OUT AFTER for Clock '_n1079'
 Total number of paths / destination ports: 1 / 1

Offset: 3.648ns (Levels of Logic = 1)
 Source: VEND (LATCH)
 Destination: VEND (PAD)
 Source Clock: _n1079 falling

Data Path: VEND to VEND

Cell:in->out	fanout	Gate Delay	Net Delay	Logical Name (Net Name)
LD:G->Q	1	0.498	0.579	VEND (VEND_OBUF)
OBUF:I->O		2.571		VEND_OBUF (VEND)
<hr/>				
Total		3.648ns (3.069ns logic, 0.579ns route)		
		(84.1% logic, 15.9% route)		

=====
 Timing constraint: Default OFFSET OUT AFTER for Clock '_n1087'
 Total number of paths / destination ports: 1 / 1

Offset: 3.648ns (Levels of Logic = 1)
 Source: INVALID_SEL (LATCH)
 Destination: INVALID_SEL (PAD)
 Source Clock: _n1087 falling

Data Path: INVALID_SEL to INVALID_SEL

Cell:in->out	fanout	Gate Delay	Net Delay	Logical Name (Net Name)
LD:G->Q	1	0.498	0.579	INVALID_SEL (INVALID_SEL_OBUF)
OBUF:I->O		2.571		INVALID_SEL_OBUF (INVALID_SEL)
<hr/>				
Total		3.648ns (3.069ns logic, 0.579ns route)		
		(84.1% logic, 15.9% route)		

=====
 Timing constraint: Default OFFSET OUT AFTER for Clock '_n1084'
 Total number of paths / destination ports: 1 / 1

Offset: 3.648ns (Levels of Logic = 1)
 Source: FAILED_TRAN (LATCH)
 Destination: FAILED_TRAN (PAD)
 Source Clock: _n1084 falling

Data Path: FAILED_TRAN to FAILED_TRAN

Cell:in->out	fanout	Gate Delay	Net Delay	Logical Name (Net Name)
LD:G->Q	1	0.498	0.579	FAILED_TRAN (FAILED_TRAN_OBUF)
OBUF:I->O		2.571		FAILED_TRAN_OBUF (FAILED_TRAN)
<hr/>				
Total		3.648ns (3.069ns logic, 0.579ns route)		
		(84.1% logic, 15.9% route)		

=====
 Cross Clock Domains Report:

Clock to Setup on destination clock CLK

	Src:Rise	Src:Fall	Src:Rise	Src:Fall
Source Clock	Dest:Rise	Dest:Rise	Dest:Fall	Dest:Fall
CLK	4.511			
current_state[3]_GND_10_o_Mux_480_o		4.668		
current_state[3]_GND_2_o_Mux_464_o		4.873		
current_state[3]_GND_6_o_Mux_472_o		4.456		

Clock to Setup on destination clock _n1079

	Src:Rise	Src:Fall	Src:Rise	Src:Fall
Source Clock	Dest:Rise	Dest:Rise	Dest:Fall	Dest:Fall
CLK			2.531	

Clock to Setup on destination clock _n1084

	Src:Rise	Src:Fall	Src:Rise	Src:Fall
Source Clock	Dest:Rise	Dest:Rise	Dest:Fall	Dest:Fall
CLK			2.531	

Clock to Setup on destination clock _n1087

	Src:Rise	Src:Fall	Src:Rise	Src:Fall
Source Clock	Dest:Rise	Dest:Rise	Dest:Fall	Dest:Fall
CLK			2.548	

Clock to Setup on destination clock current_state[3]_GND_10_o_Mux_480_o

	Src:Rise	Src:Fall	Src:Rise	Src:Fall
Source Clock	Dest:Rise	Dest:Rise	Dest:Fall	Dest:Fall
CLK			2.531	

Clock to Setup on destination clock current_state[3]_GND_2_o_Mux_464_o

	Src:Rise	Src:Fall	Src:Rise	Src:Fall
Source Clock	Dest:Rise	Dest:Rise	Dest:Fall	Dest:Fall
CLK			2.531	
current_state[3]_GND_10_o_Mux_480_o			1.597	

Clock to Setup on destination clock current_state[3]_PWR_14_o_Select_533_o

	Src:Rise	Src:Fall	Src:Rise	Src:Fall
Source Clock	Dest:Rise	Dest:Rise	Dest:Fall	Dest:Fall

CLK				3.813	
current_state[3]_GND_2_o_Mux_464_o				2.402	
current_state[3]_GND_6_o_Mux_472_o				2.445	
-----+-----+-----+-----+					

=====

Total REAL time to Xst completion: 6.00 secs
Total CPU time to Xst completion: 5.67 secs

-->

Total memory usage is 396832 kilobytes

Number of errors : 0 (0 filtered)
Number of warnings : 39 (0 filtered)
Number of infos : 1 (0 filtered)

6.2 Implementation Report

Release 14.7 Map P.20131013 (lin64)

Xilinx Mapping Report File for Design 'vending_machine'

Design Information

```
-----
Command Line   : map -intstyle ise -p xc6slx16-csg324-3 -w -logic_opt off -ol
high -t 1 -xt 0 -register_duplication off -r 4 -global_opt off -mt off -ir off
-pr off -lc off -power off -o vending_machine_map.ncd vending_machine.ngd
vending_machine.pcf
Target Device  : xc6slx16
Target Package : csg324
Target Speed   : -3
Mapper Version : spartan6 -- $Revision: 1.55 $
Mapped Date    : Sun May 24 23:44:11 2020
```

Design Summary

Number of errors: 0

Number of warnings: 7

Slice Logic Utilization:

Number of Slice Registers:	83 out of	18,224	1%
Number used as Flip Flops:	76		
Number used as Latches:	7		
Number used as Latch-thrus:	0		
Number used as AND/OR logics:	0		
Number of Slice LUTs:	156 out of	9,112	1%
Number used as logic:	156 out of	9,112	1%
Number using O6 output only:	129		
Number using O5 output only:	0		
Number using O5 and O6:	27		
Number used as ROM:	0		
Number used as Memory:	0 out of	2,176	0%

Slice Logic Distribution:

Number of occupied Slices:	55 out of	2,278	2%
Number of MUXCYs used:	0 out of	4,556	0%
Number of LUT Flip Flop pairs used:	164		
Number with an unused Flip Flop:	82 out of	164	50%
Number with an unused LUT:	8 out of	164	4%
Number of fully used LUT-FF pairs:	74 out of	164	45%
Number of unique control sets:	23		
Number of slice register sites lost to control set restrictions:	125 out of	18,224	1%

A LUT Flip Flop pair for this architecture represents one LUT paired with one Flip Flop within a slice. A control set is a unique combination of clock, reset, set, and enable signals for a registered element.

The Slice Logic Distribution report is not meaningful if the design is over-mapped for a non-slice resource or if Placement fails.

IO Utilization:

Number of bonded IOBs:	16 out of	232	6%
IOB Latches:	6		

Specific Feature Utilization:

Number of RAMB16BWERs:	0 out of	32	0%
Number of RAMB8BWERs:	0 out of	64	0%

Number of BUFIO2/BUFIO2_2CLKs:	0 out of	32	0%
Number of BUFIO2FB/BUFIO2FB_2CLKs:	0 out of	32	0%
Number of BUFG/BUFGMUXs:	1 out of	16	6%
Number used as BUFGs:	1		
Number used as BUFGMUX:	0		
Number of DCM/DCM_CLKGENs:	0 out of	4	0%
Number of ILOGIC2/ISERDES2s:	0 out of	248	0%
Number of IODELAY2/IODRP2/IODRP2_MCBs:	0 out of	248	0%
Number of OLOGIC2/OSERDES2s:	6 out of	248	2%
Number used as OLOGIC2s:	6		
Number used as OSERDES2s:	0		
Number of BSCANs:	0 out of	4	0%
Number of BUFHs:	0 out of	128	0%
Number of BUFPLLs:	0 out of	8	0%
Number of BUFPLL_MCBs:	0 out of	4	0%
Number of DSP48A1s:	0 out of	32	0%
Number of ICAPs:	0 out of	1	0%
Number of MCBs:	0 out of	2	0%
Number of PCILOGICSEs:	0 out of	2	0%
Number of PLL_ADVs:	0 out of	2	0%
Number of PMVs:	0 out of	1	0%
Number of STARTUPs:	0 out of	1	0%
Number of SUSPEND_SYNCs:	0 out of	1	0%

Average Fanout of Non-Clock Nets: 4.56

Peak Memory Usage: 676 MB

Total REAL time to MAP completion: 10 secs

Total CPU time to MAP completion: 7 secs

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- Section 9 - Area Group and Partition Summary
- Section 10 - Timing Report
- Section 11 - Configuration String Information
- Section 12 - Control Set Information
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Section 1 - Errors

Section 2 - Warnings

WARNING:PhysDesignRules:372 - Gated clock. Clock net
current_state[3]_GND_6_o_Mux_472_o is sourced by a combinatorial pin. This is
not good design practice. Use the CE pin to control the loading of data into
the flip-flop.

WARNING:PhysDesignRules:372 - Gated clock. Clock net _n1087 is sourced by a
combinatorial pin. This is not good design practice. Use the CE pin to
control the loading of data into the flip-flop.

WARNING:PhysDesignRules:372 - Gated clock. Clock net

current_state[3]_PWR_14_o_Select_533_o is sourced by a combinatorial pin.
This is not good design practice. Use the CE pin to control the loading of data into the flip-flop.

WARNING:PhysDesignRules:372 - Gated clock. Clock net
current_state[3]_GND_10_o_Mux_480_o is sourced by a combinatorial pin. This is not good design practice. Use the CE pin to control the loading of data into the flip-flop.

WARNING:PhysDesignRules:372 - Gated clock. Clock net
current_state[3]_GND_2_o_Mux_464_o is sourced by a combinatorial pin. This is not good design practice. Use the CE pin to control the loading of data into the flip-flop.

WARNING:PhysDesignRules:372 - Gated clock. Clock net _n1084 is sourced by a combinatorial pin. This is not good design practice. Use the CE pin to control the loading of data into the flip-flop.

WARNING:PhysDesignRules:372 - Gated clock. Clock net _n1079 is sourced by a combinatorial pin. This is not good design practice. Use the CE pin to control the loading of data into the flip-flop.

Section 3 - Informational

INFO:MapLib:562 - No environment variables are currently set.

INFO:LIT:244 - All of the single ended outputs in this design are using slew rate limited output drivers. The delay on speed critical single ended outputs can be dramatically reduced by designating them as fast outputs.

INFO:Pack:1716 - Initializing temperature to 85.000 Celsius. (default - Range: 0.000 to 85.000 Celsius)

INFO:Pack:1720 - Initializing voltage to 1.140 Volts. (default - Range: 1.140 to 1.260 Volts)

INFO:Map:215 - The Interim Design Summary has been generated in the MAP Report (.mrp).

INFO:Pack:1650 - Map created a placed design.

Section 4 - Removed Logic Summary

Section 5 - Removed Logic

To enable printing of redundant blocks removed and signals merged, set the detailed map report option and rerun map.

Section 6 - IOB Properties

IOB Name	Type	Direction	IO Standard	Diff Term	Drive Strength	Slew Rate
CARD_IN	IOB	INPUT	LVC MOS25			
CLK	IOB	INPUT	LVC MOS25			
COST<0>	IOB	OUTPUT	LVC MOS25		12	SLOW
COST<1>	IOB	OUTPUT	LVC MOS25		12	SLOW
COST<2>	IOB	OUTPUT	LVC MOS25		12	SLOW
DOOR_OPEN	IOB	INPUT	LVC MOS25			
FAILED_TRAN	IOB	OUTPUT	LVC MOS25		12	SLOW
INVALID_SEL	IOB	OUTPUT	LVC MOS25		12	SLOW
ITEM_CODE<0>	IOB	INPUT	LVC MOS25			
ITEM_CODE<1>	IOB	INPUT	LVC MOS25			
ITEM_CODE<2>	IOB	INPUT	LVC MOS25			

KEY_PRESS	IOB	INPUT	LVCMOS25			
RELOAD	IOB	INPUT	LVCMOS25			
RESET	IOB	INPUT	LVCMOS25			
VALID_TRAN	IOB	INPUT	LVCMOS25			
VEND	IOB	OUTPUT	LVCMOS25		12	SLC

Section 7 - RPMs

Section 8 - Guide Report

Guide not run on this design.

Section 9 - Area Group and Partition Summary

Partition Implementation Status

No Partitions were found in this design.

Area Group Information

No area groups were found in this design.

Section 10 - Timing Report

A logic-level (pre-route) timing report can be generated by using Xilinx static timing analysis tools, Timing Analyzer (GUI) or TRCE (command line), with the mapped NCD and PCF files. Please note that this timing report will be generated using estimated delay information. For accurate numbers, please generate a timing report with the post Place and Route NCD file.

For more information about the Timing Analyzer, consult the Xilinx Timing Analyzer Reference Manual; for more information about TRCE, consult the Xilinx Command Line Tools User Guide "TRACE" chapter.

Section 11 - Configuration String Details

Use the "-detail" map option to print out Configuration Strings

Section 12 - Control Set Information

Use the "-detail" map option to print out Control Set Information.

Section 13 - Utilization by Hierarchy

Use the "-detail" map option to print out the Utilization by Hierarchy section.