## ElGanal Public-Key Cryptographic System

Global Public Elements a prime number 2 & its primitive root a.

key generation by Alice

1. Generate a private key XA.,

2. Compute  $\underline{Y}_A = \alpha^{X_A} \mod 2$ . Public ky (2, a, TA)

Encryption by Bob with Atres Public tey

1. massage M: 0≤M≤2-1

2. choose a random roteger k: 15 k = 7-1. Compute a cone-time pad  $K = (Y_A) \mod g$ . 3. Encrypt M as  $(C_1, C_2)$ , where  $\chi^{K_{OKA}}$ .

CI= of mod g, Cz= KM mod g

Decryption by Alice with her private key

1. Recover  $K'=C_1^{XA} \mod 2 = \chi^{KXA} \mod 2$ . 2. compute  $M=(K)^T C_2 \mod 2$ .

Afre 
$$X_1 = 0$$
.

Afre  $X_2 = 0$ .

 $X_1 = 0$  A mod  $Y_2 = 0$  mod  $Y_3 = 0$ .

Private key  $Y_1 = 0$ .

 $Y_2 = 0$  A mod  $Y_3 = 0$ .

 $Y_4 = 0$  A mod  $Y_5 = 0$ .

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A unique k should be used for each block of message

If  $C_{i}^{(0)} = x^{k} \mod 2$ ,  $C_{2}^{(0)} = K^{i} M_{i} \mod 2$   $C_{i}^{(0)} = \alpha^{k} \mod 2$ ,  $C_{2}^{(0)} = K^{i} M_{2} \mod 2$ .  $C_{2}^{(0)} = \alpha^{k} \mod 2$ ,  $C_{2}^{(0)} = K^{i} M_{2} \mod 2$ .  $C_{2}^{(0)} = M_{i} \mod 2$ 

## Cubic function fix= $ax^3+bx^2+cx+d$ . $0 = f(x) = 30x^2+2bx+c$ . $\Rightarrow x = \frac{-b^{\frac{1}{2}}b^{\frac{1}{2}}3x}{3a}$ $\Delta_0 = b^{\frac{1}{2}}3ac$ $0 \Rightarrow 0$ $\Rightarrow 0$ $\Rightarrow 0$ $\Rightarrow 0$ $\Rightarrow 0$ $\Rightarrow 0$

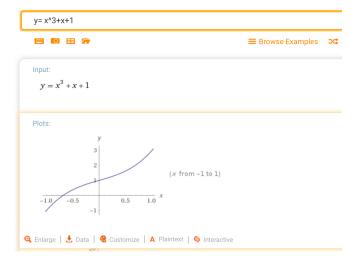
Consider 
$$f(x) = x^3 + px + q$$
. Pr  $q \in \mathbb{R}$ 

$$\Delta = -4p^3 - 2\eta q^2$$

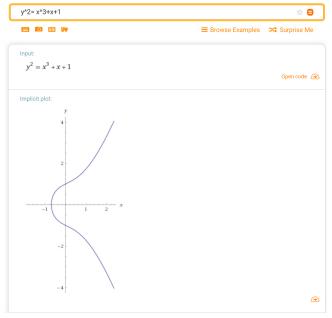
$$\Delta_0 = -3p$$

DD>0, foo) has three distinct voots. DD=0, foo) has a multiple voot.

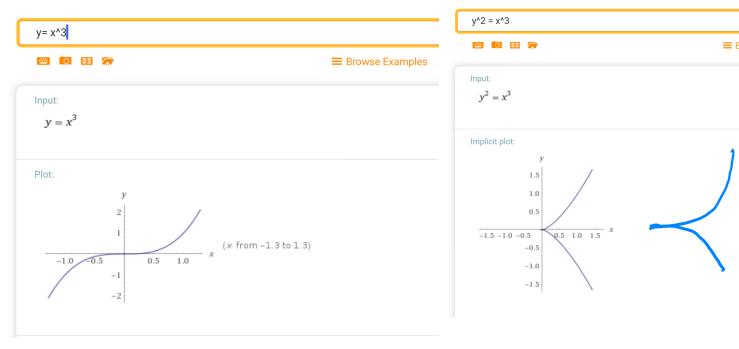
## \*WolframAlpha computational intelligence.



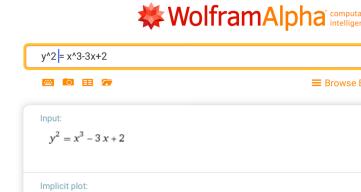


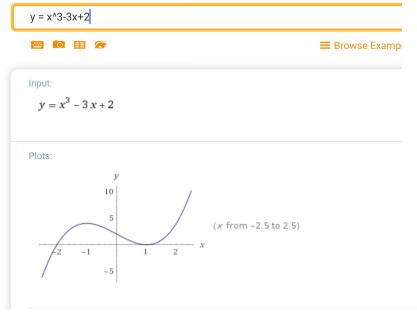


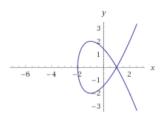
## **WolframAlpha** computate intelligence











y= x3+ px+9, p, & =1R.

the curve is symmetric about y=0.

Let O be the point at infinity.

Let E(p,q)= 5 (x, y) (R": y=x3+px+2) U SO}

Suppose 4p3+2792 = 0.

It can be shown that every line intersecting E(P, g)

intersects it in exactly three points

where a point P is counted twice it the true is tangent to the cure at P and also the point at offinity is also counted (where the line is vertical.)