Cryptology Hw3

到6024

(a) If c=0, choose  $M_1,M_2=0$  $EL(K,0)=EL(K,0)=EL(K,0)\oplus EL(K,0)=0$ 

面心で有的大是し

把 chosen ciphertext 六要ith是 1的就製造1個ci (其中ith bit=1,其餘為0)則cipher c 可以被唯一的一組 Ci set Ci E [1,256]),這些 Ci XOR可以組成 c 写這些i被收集在Ic E [1,23",256]

 $C = \bigoplus Ci = \bigoplus E(k, mi) = E(k, \bigoplus mi)$ ielo
ielo
ielo

EX. 11010=10000 \$01000 \$ 00010

## ui) chall zero string

- :: linear A ELCK, O) = ELCK, O) = ELCK, O) = O
- : If c=0 then its plaintext m=0.

## 作法

Decrypt\_without\_key ()

l If Cis zerostring, then return M=0 必至多256次
else ask for the plaintext of Ci (i & I c)
(mi, i & I c)
return M=⊕ Mi
i & I c

#2

井う

 $\begin{array}{c|c} \bigcirc & \text{Lis} \| R_{15} \xrightarrow{T_{16}} & \text{Ris} \| f(R_{15}) \oplus L_{15} \xrightarrow{T_{17}} f(R_{15}) \oplus L_{15} \| R_{15} \\ \xrightarrow{\text{IP}} \xrightarrow{\text{IP}} \xrightarrow{\text{TDI}} & \text{Ris} \| f(R_{15}) \oplus f(R_{15}) \oplus L_{15} \rightarrow R_{15} \| o \oplus L_{15} \\ & \rightarrow & \text{Ris} \| L_{15} \\ & & \end{array}$ 

(b)  $L_{15} \| R_{15} \xrightarrow{T_{16}} R_{15} \| f(R_{15}) \oplus L_{15} \xrightarrow{IP} IP$   $f(R_{15}) \oplus L_{15} \| f(f(R_{15}) \oplus L_{15}) \oplus R_{15}$ 

If we want to hold the above equals to LIS || RIS. then  $f(RIS) = 0 \Rightarrow$  the righthand side become  $f(LIS) \oplus RIS$  again, f(LIS) = 0, the condition is impossible because function f is 1-to-1 (so that it exists inverse f')  $\Rightarrow$  1+'s impossible for f(LIS) = f(RIS) = 0

Hence, we disprove the equality.