

**A GRADUATE COURSE  
IN  
CONVEX OPTIMIZATION**



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# A GRADUATE COURSE IN CONVEX OPTIMIZATION DDA6110 Notebook

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## Lecturer

*Prof. Yin Zhang*

*Institute for Data and Decision Analytics*

## Tex Written By

*Mr. Jie Wang*

*The Chinese University of Hong Kong, Shenzhen*



香港中文大學(深圳)

The Chinese University of Hong Kong, Shenzhen



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# Notations and Conventions

$\sup$	least upper bound
$\inf$	greatest lower bound
$\overline{E}$	closure of $E$
$f \circ g$	composition
$\limsup(\liminf)$	upper (lower) limit
$L(\mathcal{P}, f), U(\mathcal{P}, f)$	Riemann sums
$\mathcal{R}[a, b]$	classes of Riemann integrable functions on $[a, b]$
$\int_a^b f(x) \, dx, \overline{\int_a^b f(x) \, dx}$	Riemann integrals
$\langle \mathbf{x}, \mathbf{y} \rangle$	inner product
$\omega(f; E)$	oscillation of $f$ over set $E$
$\  \cdot \ $	norm
$\nabla f$	gradient
$\frac{\partial f}{\partial x_i}, f_{x_i}, f_i, \partial_i f, D_i f$	partial derivatives
$D_{\mathbf{v}} f$	directional derivative at direction $\mathbf{v}$
$\frac{\partial(y_1, \dots, y_m)}{\partial(x_1, \dots, x_n)}$	Jacobian
$\mathbb{S}^n$	set of real symmetric $n \times n$ matrices
$\succ (\succeq)$	positive (semi)-definite
$\mathcal{C}^m$	classes of $m$ -th order continuously differentiable functions
$\mathcal{C}(E; \mathbb{R}^m)$	set of $\mathcal{C}^1$ mapping from $E$ to $\mathbb{R}^m$
$(\mathcal{H}, d)$	metric space



# Chapter 1

## Week1

### 1.1. Monday

#### 1.1.1. Introduction to Convex Optimization

The basic optimization model is as follows:

$$\begin{array}{ll} \min & f(x) : \mathbb{R}^n \rightarrow \mathbb{R} \\ \text{s.t.} & x \in X \subseteq \mathbb{R}^n \end{array} \quad (1.1)$$

For instance, the constraint  $X$  can be union of some inequality constraints:

$$X = \left\{ x \in \mathbb{R}^n \mid f_i(x) \leq b_i, \quad i = 1, \dots, m \right\}$$

We only consider **convex** problems in this course, which means that in (1.1):

- the objective function  $f$  is convex;
- the constraint set  $X$  is a convex set.

The goal of optimization is to find  $x^* \in X$  such that

$$f(x^*) \leq f(x), \quad \forall x \in X.$$

or determine whether such an  $x^*$  exists or not.

■ **Example 1.1** The least squares problem is a convex problem:

$$\begin{aligned} \min \quad & \frac{1}{2} \| \mathbf{Ax} - \mathbf{b} \|_2^2 \\ \text{s.t.} \quad & \mathbf{x} \in \mathbb{R}^n \end{aligned}$$

The optimal solution can be compactly written as

$$\mathbf{x}^* = \mathbf{A}^* \mathbf{b}$$

where  $\mathbf{A}^*$  denotes the pseudo inverse of  $\mathbf{A}$ . ■

■ **Example 1.2** The linear programming is also a convex problem:

$$\begin{aligned} \min \quad & \mathbf{c}^T \mathbf{x} \\ \text{s.t.} \quad & \mathbf{a}_i^T \mathbf{x} \leq b_i, \quad i = 1, 2, \dots, m. \end{aligned}$$

The linear programming is a special case of the conic programming. We will cover conic programming in this course. ■

**Definition 1.1** [Line Segments] The line segment between  $\mathbf{x}_1, \mathbf{x}_2 \in \mathbb{R}^n$  is defined as

$$\{\theta \mathbf{x}_1 + (1 - \theta) \mathbf{x}_2 \mid \theta \in [0, 1]\} = \{\mathbf{x}_2 + \theta(\mathbf{x}_1 - \mathbf{x}_2) \mid \theta \in [0, 1]\}$$

**Definition 1.2** [Convex Set in  $\mathbb{R}^n$ ] The set  $X \subseteq \mathbb{R}^n$  is convex if for any  $\mathbf{x}_1, \mathbf{x}_2 \in X$ ,

$$\theta \mathbf{x}_1 + (1 - \theta) \mathbf{x}_2 \in X, \quad \forall \theta \in [0, 1].$$

Here are two examples for convex sets over matrices:

$$\mathcal{S}^n = \{\mathbf{W} \in \mathbb{R}^n \times \mathbb{R}^n \mid \mathbf{W} = \mathbf{W}^T\}, \quad \mathcal{S}_+^n = \{\mathbf{W} \in \mathcal{S}^n \mid \mathbf{W} \succeq 0\},$$

- Ⓡ Sometimes the convex set could be a collection of matrices, but we can vectorize a matrix into a vector.

**Definition 1.3** [Convex Function] A function  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  is **convex** if

- $\text{dorm}(f)$  is a convex set, and
- $\forall x, y \in \text{dorm}(f)$ , it holds that for  $\theta \in [0, 1]$ ,

$$f(\theta x + (1 - \theta)y) \leq \theta f(x) + (1 - \theta)f(y),$$

which means that the secant line between any two points is above the function. ■

**Proposition 1.1**  $f(x) = \max_i f_i(x)$  is convex if  $f_i(x)$  is convex for each  $i$ .

*Proof.* For any  $\theta \in [0, 1]$ , the following inequality holds:

$$\begin{aligned} f(\theta x + (1 - \theta)y) &= \max_i f_i(\theta x + (1 - \theta)y) \\ &\leq \max_i \theta f_i(x) + (1 - \theta)f_i(y) \\ &\leq \max_i \theta f_i(x) + \max_i (1 - \theta)f_i(y) \\ &= \theta f(x) + (1 - \theta)f(y) \end{aligned}$$

■

■ **Example 1.3** Define the function

$$\begin{aligned} f : \mathbb{S}^n &\rightarrow \mathbb{R} \\ \text{with } f(\mathbf{X}) &= \lambda_{\max}(\mathbf{X}) \triangleq \max_{\|\mathbf{v}\|=1} \mathbf{v}^T \mathbf{X} \mathbf{v} \end{aligned}$$

This function is convex since  $f$  can be written as the maximization of a collection of affines

(in terms of  $\mathbf{X}$ ):

$$f(\mathbf{X}) = \max_{\|\mathbf{v}\|=1} f_{\mathbf{v}}(\mathbf{X}), \quad \text{with } f_{\mathbf{v}}(\mathbf{X}) = \mathbf{v}^T \mathbf{X} \mathbf{v}.$$

■

## 1.2. Wednesday

Convex Programming (CP):

$$\begin{array}{ll} \min & f_0(x) \\ \text{s.t.} & x \in X \subseteq \mathbb{R}^n \end{array}$$

with  $f$  and  $X$  being convex. The standard constraint set  $X$  can be written as:

$$X = \left\{ x \mid f_i(x) \leq b_i, \ i = 1, 2, 3, \dots, m, \ Ax = c \right\}$$

with  $A \in \mathbb{R}^{p \times n}$ ,  $p < n$ , and  $f_i$ 's are convex.

**Proposition 1.2** The constraint set

$$X = \left\{ x \mid f_i(x) \leq b_i, \ i = 1, \dots, m \right\}$$

is **convex** if all  $f_i$ 's are convex.

*Proof.* For any  $x_1, x_2 \in X$  and  $\theta \in [0, 1]$ , verify that

$$\begin{aligned} f_i(\theta x_1 + (1 - \theta)x_2) &\leq \theta f_i(x_1) + (1 - \theta)f_i(x_2) \\ &\leq \theta b_i + (1 - \theta)b_i = b_i \end{aligned}$$

where the first inequality is because of the convexity, and the second inequality is because that  $x_i \in X, i = 1, 2$ . As a result,  $\theta x_1 + (1 - \theta)x_2 \in X$ , i.e.,  $X$  is a convex set. ■

**Proposition 1.3** If  $\{X_i\}_{i \in \mathcal{I}}$  is a collection of convex sets, then the union  $X = \bigcap_{i \in \mathcal{I}} X_i$  is convex.

*Proof.* Consider any  $x_1, x_2 \in X$ , i.e.,  $x_1, x_2 \in X_i, \forall i \in \mathcal{I}$ , then by the convexity of  $X_i$ s, the line segment

$$\theta x_1 + (1 - \theta)x_2 \in X_i, \forall i \in \mathcal{I},$$

i.e.,  $\theta x_1 + (1 - \theta)x_2 \in X$ . ■

**Definition 1.4** [Epigraph] Consider a function  $f : \mathbb{R}^n \rightarrow \mathbb{R}$ , the epigraph of  $f$  is a set defined as

$$\text{epi}(f) = \left\{ (x, t) \mid f(x) \leq t, x \in \text{dorm}(f) \right\} \in \mathbb{R}^{n+1}.$$

**Proposition 1.4** The function  $f$  is convex if and only if the epigraph  $\text{epi}(f)$  is convex.

*Proof.* We only talk about the proof on the forward direction. Consider any  $(x_1, t_1), (x_2, t_2) \in \text{epi}(f)$ , then we show that  $\theta(x_1, t_1) + (1 - \theta)(x_2, t_2) \in \text{epi}(f)$  for any  $\theta \in [0, 1]$ :

$$\begin{aligned} f(\theta x_1 + (1 - \theta)x_2) &\leq \theta f(x_1) + (1 - \theta)f(x_2) \\ &\leq \theta t_1 + (1 - \theta)t_2, \end{aligned}$$

which proves the desired result. ■

Let's start to introduce some useful convex sets.

**Definition 1.5** [Combination of two points] Consider the combination of two points  $x_1, x_2$ :

$$\theta_1 x_1 + \theta_2 x_2,$$

- When  $\theta_1, \theta_2 \in \mathbb{R}$ , it is called an affine combination; the resulted space is an **affine space**.
- When  $\theta_1, \theta_2 \in \mathbb{R}_+$ , it is called a non-negative combination; the resulted space is a **convex cone**.
- When  $\theta_2 = 1 - \theta_1, \theta_1 \in [0, 1]$ , it is called a convex combination; the resulted space is a **convex set**.

**Definition 1.6** [Affine Set] The standard form of an affine space is  $\{x \mid Ax = b\}$ . ■



**Definition 1.7** [Cone] A set is called a cone with a vertex at the origin if for any  $\mathbf{x} \in X$ ,  $a\mathbf{x} \in X$  for any  $a \geq 0$ . ■

The standard form of a convex conic programming is the following:

$$\begin{aligned} \min \quad & \langle \mathbf{C}, \mathbf{X} \rangle \\ \text{s.t.} \quad & \langle \mathbf{A}_i, \mathbf{X} \rangle = B_i, \quad i = 1, 2, \dots, m \\ & \mathbf{X} \in \mathcal{K} \end{aligned}$$

where  $\mathcal{K}$  is a convex cone.

One special form of conic programming is the linear programming:

$$\begin{aligned} \min \quad & \langle \mathbf{c}, \mathbf{x} \rangle \\ \text{s.t.} \quad & \mathbf{Ax} = \mathbf{b} \\ & \mathbf{x} \geq \mathbf{0} \end{aligned}$$

Another is the second-order cone programming:

$$\mathcal{K} = \{(\mathbf{x}, t) \mid \|\mathbf{x}\|_2 \leq t\}$$

The semidefinite programming also belongs to the case of conic programming:

$$\mathcal{K} = \mathcal{S}_+^n \triangleq \{\mathbf{X} \in \mathcal{S}^n \mid \mathbf{X} \succeq 0, \forall \mathbf{v} \in \mathbb{R}^n\}$$

Convex hull of a set  $S$ :

It is the smallest convex set containing  $S$ , called  $\text{conv}(S)$ .

**Definition 1.8** [Polyhedron]

- A hyperplane in  $\mathbb{R}^n$  can be written as the form  $\{\mathbf{x} \mid \mathbf{a}^T \mathbf{x} = b\}$ .
- Half space:  $\{\mathbf{x} \mid \mathbf{a}^T \mathbf{x} \leq b\}$ .
- Polyhedron: a intersection of finite hyperplanes and half spaces.

- Ellipsoid:

$$\begin{aligned}\left\{ \mathbf{x} \mid (\mathbf{x} - \mathbf{x}_c)^T \mathbf{P}^{-1} (\mathbf{x} - \mathbf{x}_c) \leq 1, \mathbf{P} \succ 0 \right\} &= \{ \mathbf{x} \mid \|\mathbf{P}^{-1/2} (\mathbf{x} - \mathbf{x}_c)\|_2^2 \leq 1 \} \\ &= \{ \mathbf{x} \mid \mathbf{x} = \mathbf{x}_c + \mathbf{A}\mathbf{u}, \|\mathbf{u}\|_2 \leq 1 \}\end{aligned}$$



## Chapter 2

## Week2

### 2.1. Monday

Operations that preserve convexity:

- Intersection:

$$B = \bigcap_{\|a\|=1} \{\mathbf{x} \mid \mathbf{a}^T \mathbf{x} \leq 1\} = \{\mathbf{x} \mid \|\mathbf{x}\|_2 \leq 1\}.$$

- Affine function:  $f(\mathbf{x}) = \mathbf{A}\mathbf{x} + \mathbf{b}$  with  $\mathbf{A} \in \mathbb{R}^{m \times n}, \mathbf{b} \in \mathbb{R}^m$ . As long as  $\alpha + \beta = 1$  and  $\alpha, \beta \geq 0$ ,

$$f(\alpha \mathbf{x} + \beta \mathbf{y}) = \alpha f(\mathbf{x}) + \beta f(\mathbf{y}).$$

- When  $S \subseteq \mathbb{R}^n$  is convex,  $f(S) \subseteq \mathbb{R}^m$  is convex.
- When  $C \subseteq \mathbb{R}^m$  is convex, then  $f^{-1}(C) \subseteq \mathbb{R}^n$  is convex. Consider  $x_1, x_2 \in f^{-1}(C)$ , then  $f(x_1), f(x_2) \in C$ , and  $\theta f(x_1) + (1 - \theta)f(x_2) \in C$ , i.e.,  $f(\theta x_1 + (1 - \theta)x_2) \in C$ . As a result,  $\theta x_1 + (1 - \theta)x_2 \in f^{-1}(C)$ .

■ **Example 2.1** Consider the set

$$L = \left\{ \mathbf{x} \in \mathbb{R}^n \mid \sum_{i=1}^n x_i \mathbf{A}_i \preceq \mathbf{B} \right\}, \quad \text{with } \mathbf{A}_i, \mathbf{B} \in \mathbb{S}_+^n$$

Write  $f(\mathbf{x}) = \mathbf{B} - \sum_{i=1}^n x_i \mathbf{A}_i$ , which is an affine function. It follows that

$$L = \left\{ \mathbf{x} \in \mathbb{R}^n \mid f(\mathbf{x}) \preceq 0 \right\} = f^{-1}(\mathbb{S}_+^n)$$

Since  $S_+^n$  is convex,  $L$  is convex as well. ■

■ **Example 2.2** Consider the set

$$\mathcal{H} = \left\{ \mathbf{x} \mid \mathbf{x}^T \mathbf{P} \mathbf{x} \leq (\mathbf{c}^T \mathbf{x})^2, \quad \mathbf{c}^T \mathbf{x} \geq 0 \right\}, \quad \mathbf{P} \succeq 0.$$

Define

$$f(\mathbf{x}) = \begin{pmatrix} \mathbf{P}^{1/2} \\ \mathbf{c}^T \end{pmatrix} \mathbf{x} : \mathbb{R}^n \rightarrow \mathbb{R}^{n+1}$$

Therefore,  $\mathcal{H} = f^{-1}(\text{second order cone})$ , with

$$\text{second order cone} = \left\{ \begin{pmatrix} \mathbf{x} \\ t \end{pmatrix} \mid \|\mathbf{x}\| \leq t \right\}$$

Perspective transformation:

$$P(\mathbf{x}, t) = \frac{\mathbf{x}}{t}, \quad t > 0 : \mathbb{R}^{n+1} \rightarrow \mathbb{R}^n.$$

Consider  $S \subseteq \text{dorm}(P)$  and  $P(S) \subseteq \mathbb{R}^n$ .

We show that line segments map to line segments:

*Proof.* Let  $(x_1, t_1), (x_2, t_2) \in S$  with  $t_1, t_2 > 0$ . Consider

$$\begin{aligned} \begin{pmatrix} \theta x_1 + (1 - \theta)x_2 \\ \theta t_1 + (1 - \theta)t_2 \end{pmatrix} &\rightarrow \frac{\theta x_1 + (1 - \theta)x_2}{\theta t_1 + (1 - \theta)t_2} \\ &= \frac{\theta t_1}{\theta t_1 + (1 - \theta)t_2} \frac{x_1}{t_1} + \frac{(1 - \theta)t_2}{\theta t_1 + (1 - \theta)t_2} \frac{x_1}{t_2} \\ &= \mu \frac{x_1}{t} + (1 - \mu) \frac{x_2}{t} \end{aligned}$$

■

Linear fractional:

$$f(\mathbf{x}) = \frac{\mathbf{Ax} + \mathbf{b}}{\mathbf{c}^T \mathbf{x} + \mathbf{d}}, \quad \mathbf{c}^T \mathbf{x} + \mathbf{d} > 0.$$

Proper cone  $\mathcal{K}$ :

- convex;
- closed;
- solid;
- pointed, i.e., contain no line.

Generalized inequality:

$$\mathbf{x} \preceq_{\mathcal{K}} \mathbf{y} \implies \mathbf{y} - \mathbf{x} \in \mathcal{K}$$

Define minimum element v.s. minimal element.

## 2.2. Wednesday

$K$  proper cone: convex, closed, solid, pointed.

**Definition 2.1** [Generalized Inequality] Given a set  $K$ , we define  $x \leq_K y$  if  $y - x \in K$ . ■

Pointed cone:

$$x \leq_K y, y \leq_K x \implies x = y.$$

**Definition 2.2** [Dual Cone] Given a cone  $K \subseteq \mathbb{R}^n$ , the dual cone is defined as

$$K^* = \{y \in \mathbb{R}^n \mid y^T x \geq 0, \forall x \in K\}.$$

**R** The dual cone  $K^*$  is convex, closed, and pointed. The convexity is because that it is an intersection of many half-spaces.

Self-dual cones:

$$(\mathbb{R}_+^n)^* = \mathbb{R}_+^n, \quad (\mathbb{S}_+^n)^* = \mathbb{S}_+^n.$$

*Proof.* Let  $X, Y \in \mathbb{S}_+^n$ , then

$$\langle X, Y \rangle = \sum_i \lambda_i q_i^T Y q_i \geq 0.$$

Therefore,  $\mathbb{S}_+^n \subseteq (\mathbb{S}_+^n)^*$ . When  $Y \notin (\mathbb{S}_+^n)^*$ , there exists  $u$  such that

$$u^T Y u < 0 \implies \langle Y, uu^T \rangle < 0,$$

which implies  $Y \notin (\mathbb{S}_+^n)^*$ . ■

### 2.2.1. Separation Theorems

If  $C$  and  $D$  are disjoint convex sets, there exists affine function such that

$$a^T x \leq b, \forall x \in C, \quad a^T x \geq b, \forall x \in D.$$

**Theorem 2.1** If  $C$  is convex, then for any  $x_0 \in \partial C$ , there exists a hyper-plane  $\{\mathbf{x} \mid \mathbf{a}^T \mathbf{x} = b\}$  with  $\mathbf{a} \neq \mathbf{0}$  such that

$$\mathbf{a}^T \mathbf{x}_0,$$

and  $\mathbf{a}^T \mathbf{x} \leq b, \forall \mathbf{x} \in C$ .

The converse is true if  $C$  is closed.





## **Chapter 3**

### **Week3**

#### **3.1. Monday**

##### **3.1.1. Convex Functions**

