# Q21: How do I efficiently collect marine environmental data?

Jie Wang, Xingjian Wang, Xuanchen Wu, Yanzuo Chen

#### **Table of Contents**

- Introduction
- 2 Problem Formulation
- Main Algorithm: Customized Block-Coordinate Descent
- Mumerical Simulations
- Conclusion

#### **Table of Contents**

- Introduction
- 2 Problem Formulation
- Main Algorithm: Customized Block-Coordinate Descent
- 4 Numerical Simulations
- Conclusion

#### Valuable Ocean Data

- The oceans: 71% of the Earth's surface
  - Vast unexplored areas
- Ocean temperatures determine climate and wind patterns
  - Affects life on land
- Marine pollution severely damages ecosystems

# Conventional Underwater Data Collection Methods

Technique	Limitations
Cable communication	<ul><li>High cost</li><li>Limited distance</li></ul>
Satellite communication with sea surface buoys	<ul><li>High cost</li><li>Low Rate</li></ul>
Multi-hop communication	<ul><li>Deployment overhead</li><li>Constant maintenance</li></ul>

# Efficient Underwater Sensor Network Data Collection Employing Unmanned Surface Vehicles

Jie Wang, Xingjian Wang, Xuanchen Wu, Yanzuo Chen

Network Coding Lab
The Chinese University of Hong Kong, Shenzhen



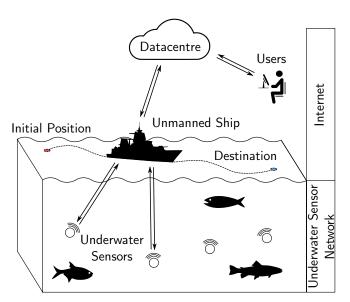
网络编码实验室 Network Coding Lab

December 12, 2019



Slides

## **Data Collection by Unmanned Ships**



#### Constraints to Consider

- Communication channel loss increases exponentially!
- Limited battery and transmission power for Underwater Sensor Nodes (USNs)

## As an Optimization Problem

Minimize the **maximum energy consumption** of all USNs by the joint design of...

- the path of the unmanned surface vehicle
- the wake-up schedule of the USNs

# As an Optimization Problem

#### Challenges

- Non-convexity
- Large problem sizes (Number of USNs, transmission time slots)

Not efficiently solved by existing algorithms and off-the-shelf tools!

#### Solution

Block-Coordinate Descent algorithm

#### **Table of Contents**

- Introduction
- Problem Formulation
- Main Algorithm: Customized Block-Coordinate Descent
- 4 Numerical Simulations
- Conclusion

### **Underwater Acoustic Channel Model**

Key assmptions<sup>1</sup>:

- Gaussian Noise;
- 2 The k-th node transmits with power  $p_k$ ;
- **3** Channel is separated into sub-channels, each with bandwidth  $\Delta f$  and frequency  $f_i$ .

## Transmission Rate Approximation

The transmission rate for the k-th node over distance d is approximated as

$$C(d,k) = \sum_{i} \log_{2} \left[ 1 + \frac{p_{k}/\Delta f}{N(f_{i}) \cdot A(d,f_{i})} \right] \Delta f$$

where  $A(d, f) \triangleq d^{\kappa}[\alpha(f)]^d$  denotes the attenuation factor; N(f) denotes noise p.s.d.

1. Milica Stojanovic. 2007. On the relationship between capacity and distance in an underwater acoustic communication channel.

## System Model

- An unmanned ship is to collect data from K USNs;
- Total time horizon is discretized into M time slots equally;
- Decision variable:

$\boldsymbol{q} := \{\boldsymbol{q}[m], 0 \le m \le M\}$	Path of unmanned ship
$\mathbf{x} := \{\mathbf{x}_k[m], 0 \le m \le M, 1 \le k \le K\}$	Wake-up schedule
$\boldsymbol{p} := \{p_k, 1 \leq k \leq K\}$	Transmission power of USNs

• Objective: minimize the maximum energy consumption for all USNs

$$\min_{\boldsymbol{p},\boldsymbol{q},\boldsymbol{x}} \max_{k} \sum_{m=0}^{M} x_{k}[m] p_{k}$$

## **System Constraints**

• The path of the ship satisfies initial and final location constraints:

$$q[0] = q_0, \quad q[M] = q_f.$$

• The maximum speed constraints of the unmanned ship:

$$\|\boldsymbol{q}[m] - \boldsymbol{q}[m-1]\| \leq V_{\mathsf{max}}$$

Wake-up mechanism:

$$\begin{cases} \sum_{k=1}^{K} x_k[m] \le 1, & \forall m \\ x_k[m] \in \{0, 1\}, & \forall m, \forall k \end{cases}$$

• Data Load Constraint:

$$\sum_{m=1}^{M} x_k[m] R(p_k, \mathbf{q}[m]) \ge b_k, \quad \forall k$$

## Formulated Optimization Problem

The data collection scheme is formulated as the optimization problem<sup>1</sup>:

$$\begin{aligned} & \underset{\boldsymbol{p},\boldsymbol{q},\boldsymbol{x},\boldsymbol{\theta}}{\text{min}} & \boldsymbol{\theta} \\ & \text{s.t.} & & \sum_{m=1}^{M} x_k[m] p_k \delta \leq \boldsymbol{\theta}, \quad \forall k=1,\ldots,K \\ & \boldsymbol{q}[0] = \boldsymbol{q}_0, \quad \boldsymbol{q}[M] = \boldsymbol{q}_f \\ & & \| \boldsymbol{q}[m] - \boldsymbol{q}[m-1] \| \leq V_{\text{max}} \\ & & \sum_{k=1}^{K} x_k[m] \leq 1, \quad \forall m \\ & & \sum_{m=1}^{M} x_k[m] R(p_k, \boldsymbol{q}[m]) \geq b_k, \quad \forall k \\ & & x_k[m] \in \{0,1\}, \quad \forall m, \forall k \end{aligned}$$

1. Cheng Zhan, Yong Zeng, and Rui Zhang. 2018. Energy-efficient data collection in UAV enabled wireless sensor network

#### **Table of Contents**

- Introduction
- Problem Formulation
- Main Algorithm: Customized Block-Coordinate Descent
- 4 Numerical Simulations
- Conclusion

## Objective function

$$\min_{\boldsymbol{p},\boldsymbol{q},\boldsymbol{x},\boldsymbol{\theta}} \quad \boldsymbol{\theta}$$
s.t. 
$$\sum_{m=1}^{M} x_k[m] p_k \delta \leq \boldsymbol{\theta}, \quad \forall k = 1, \dots, K$$

$$\boldsymbol{q}[0] = \boldsymbol{q}_0, \quad \boldsymbol{q}[M] = \boldsymbol{q}_f$$

$$\|\boldsymbol{q}[m] - \boldsymbol{q}[m-1]\| \leq V_{\text{max}}$$

$$\sum_{k=1}^{K} x_k[m] \leq 1, \quad \forall m$$

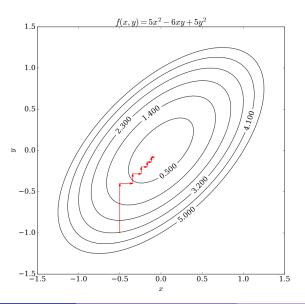
$$\sum_{m=1}^{M} x_k[m] R(p_k, \boldsymbol{q}[m]) \geq b_k, \quad \forall k$$

$$x_k[m] \in \{0, 1\}, \quad \forall m, \forall k$$

- x : The sensor wake-up scheduling policy
- *p* : The sensor power policy
- q : The path planning policy of the ship



## Coordinate Descent Algorithm



#### **Iteration**

Fixing  $\boldsymbol{p}$  (sensor power policy),  $\boldsymbol{q}$  (path planning), minimize over  $\boldsymbol{x}$  (wake up schedule)

$$\begin{aligned} & \underset{\boldsymbol{p},\boldsymbol{q},\boldsymbol{x},\boldsymbol{\theta}}{\min} & \boldsymbol{\theta} \\ & \text{s.t.} & & \sum_{m=1}^{M} x_k[m] p_k \boldsymbol{\delta} \leq \boldsymbol{\theta}, \quad \forall k \\ & \boldsymbol{q}[0] = \boldsymbol{q}_0, \quad \boldsymbol{q}[M] = \boldsymbol{q}_f \\ & & \| \boldsymbol{q}[m] - \boldsymbol{q}[m-1] \| \leq V_{\text{max}} \\ & & \sum_{k=1}^{K} x_k[m] \leq 1, \quad \forall m \\ & & \sum_{m=1}^{M} x_k[m] R(p_k, \boldsymbol{q}[m]) \geq b_k, \quad \forall k \\ & & x_k[m] \in \{0, 1\}, \quad \forall m, \forall k \end{aligned}$$

#### **Iteration**

Fixing p (sensor power policy), q (path planning), minimize over x (wake up schedule)

$$\min_{\mathbf{x},\theta} \quad \theta$$
s.t. 
$$\sum_{m=1}^{M} x_k[m] p_k \delta \leq \theta, \quad \forall k$$

$$\sum_{k=1}^{K} x_k[m] \le 1, \quad \forall m$$

$$\sum_{m=1}^{M} x_k[m] R(p_k, \mathbf{q}[m]) \ge b_k, \quad \forall k$$

$$x_k[m] \in \{0, 1\}, \quad \forall m, \forall k$$

## **Advantages**

- Computational cost is usually less than other methods
- Sub-problems are usually easier to solve
- Ease of implementation

## **Trajectory Optimization**

For fixed wake-up schedule x and transmission power policy p,

$$\begin{array}{ll} \max\limits_{\boldsymbol{q},\eta} & \eta \\ \text{s.t.} & \|\boldsymbol{q}[m] - \boldsymbol{q}[m-1]\| \leq V_{\text{max}}, \ \forall m=1,\ldots,M \\ & \boldsymbol{q}[0] = \boldsymbol{q}_0, \ \boldsymbol{q}[M] = \boldsymbol{q}_f \\ & \frac{1}{b_k} \sum_{m=1}^M x_k[m] \frac{R(p_k,\boldsymbol{q}[m])}{R(p_k,\boldsymbol{q}[m])} \geq \eta, \ \forall k=1,\ldots,K \end{array}$$

where  $R(p_k, \mathbf{q}[m])$  is the *channel gain* at time slot m for sensor k:

$$R(p_k, \mathbf{q}[m]) \triangleq \sum_{i} \Delta f \cdot \log_2 \left( 1 + \frac{A_{k,m}^{(i)}}{\|\mathbf{q}[m] - \mathbf{I}[k]\|^{\kappa} \cdot \alpha^{\|\mathbf{q}[m] - \mathbf{I}[k]\|}} \right)$$

## **Successive Convex Approximation**

canonical form 
$$\max_{m{q}} \quad f_0(m{q}) \\ \text{s.t.} \quad f_i(m{q}) \geq 0, \quad i=1,\dots,I$$

Traditional method is by applying  $f_i(\mathbf{q}) \geq f_{i,\text{lb}}^{(\ell)}(\mathbf{q}), \forall \mathbf{q}$ :

```
convex relaxation \max_{m{q}} \quad f_0(m{q}) \\ \text{s.t.} \quad f_{i,\text{lb}}^{(\ell)}(m{q}) \geq 0, \quad i=1,\dots,I
```

- Easy to solve the relaxation problem
- Solution is feasible to the nominal problem.

# Standard SCA is not Applicable!

$$\begin{array}{ll} \max\limits_{\boldsymbol{q},\eta} & \eta \\ \text{s.t.} & \|\boldsymbol{q}[m] - \boldsymbol{q}[m-1]\| \leq V_{\text{max}}, \ \forall m=1,\ldots,M \\ & \boldsymbol{q}[0] = \boldsymbol{q}_0, \ \boldsymbol{q}[M] = \boldsymbol{q}_f \\ & \frac{1}{b_k} \sum_{m=1}^M x_k[m] R(p_k, \boldsymbol{q}[m]) \geq \eta, \ \forall k=1,\ldots,K \end{array}$$

#### Global concave lower bound is not feasible to find

For the channel gain function

$$R(p_k, \mathbf{q}[m]) \triangleq \sum_{i} \Delta f \cdot \log_2 \left( 1 + \frac{A_{k,m}^{(i)}}{\|\mathbf{q}[m] - \mathbf{I}[k]\|^{\kappa} \cdot \alpha^{\|\mathbf{q}[m] - \mathbf{I}[k]\|}} \right),$$

as long as  $\alpha \neq 1$ , the first-order taylor expansion is not its global concave lower bound.

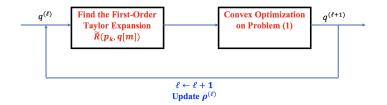
## **SCA** with Trust Region Heuristic

$$\max_{\boldsymbol{q},\eta} \quad \eta$$
s.t. 
$$\|\boldsymbol{q}[m] - \boldsymbol{q}[m-1]\| \leq V_{\text{max}}, \ \forall m = 1, \dots, M$$

$$\boldsymbol{q}[0] = \boldsymbol{q}_0, \ \boldsymbol{q}[M] = \boldsymbol{q}_f \qquad (1)$$

$$\frac{1}{b_k} \sum_{m=1}^{M} x_k[m] \tilde{\boldsymbol{R}}(\boldsymbol{p}_k, \boldsymbol{q}[m]) \geq \eta, \ \forall k = 1, \dots, K$$

$$\boldsymbol{q} \in \mathcal{T}^{(\ell)} \triangleq \{\boldsymbol{q} \mid \|\boldsymbol{q} - \boldsymbol{q}^{(\ell)}\| \leq \rho^{(\ell)}\}$$



## **Techniques Summarization**

$$\min_{\boldsymbol{p},\boldsymbol{q},\boldsymbol{x}} \quad \max_{k} \sum_{m=1}^{M} x_k[m] p_k \delta \tag{12a}$$

s.t. 
$$\|q[m] - q[m-1]\| \le V_{\text{max}}, \ \forall m$$
 (12b)

$$\mathbf{q}[0] = \mathbf{q}_0, \mathbf{q}[M] = \mathbf{q}_f \tag{12c}$$

$$\sum_{k=1}^{K} x_k[m] \le 1, \quad \forall m \tag{12d}$$

$$\sum_{m=1}^{M} x_k[m] R(p_k, \boldsymbol{q}[m]) \ge b_k, \quad \forall k$$
 (12e)

$$x_k[m] \in \{0,1\}, \quad \forall k, \forall m. \tag{12f}$$

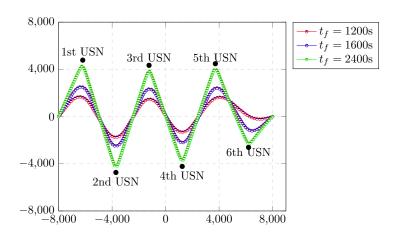
#### Convergence Comments

Our customized algorithm is guaranteed to converge.

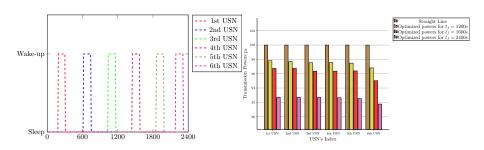
#### **Table of Contents**

- Introduction
- Problem Formulation
- Main Algorithm: Customized Block-Coordinate Descent
- 4 Numerical Simulations
- Conclusion

### **Numerical Simulation**



## Transmission Scheduling and Power Control



#### **Table of Contents**

- Introduction
- Problem Formulation
- Main Algorithm: Customized Block-Coordinate Descent
- 4 Numerical Simulations
- Conclusion

#### **Conclusion**

- Data collection task by employing unmanned surface vechicles
- Jointly optimize the transmission scheduling, powers, and trajectory.
- Solving the non-convex optimization by:
  - block-coordinate descent
  - successive convex approximation with trust region heuristic
- Other useful techniques:
  - Approximate Dynamic Programming
  - Machine learning for online trajectory optimization

#### References

- J. Wang, J. Ma, J. Yang, and S. Yang, "Efficient underwater sensor network data collection employing unmanned ships," in Proceedings of the Thirteenth ACM International Conference on Underwater Networks & Systems, ser. WUWNet'19, 2019.
- J. Borden and J. DeArruda, "Long range acoustic underwater communication with a compact auv," in 2012 Oceans, Oct 2012, pp. 1-5.
- Y. Su, Y. Zuo, Y. Li, Z. Jin, and X. Fu, "An underwater data acquisition and transmission testbed based on beidou satellite system (bds) and underwater acoustic communication technology," in 2018 OCEANS - MTS/IEEE Kobe Techno-Oceans (OTO), May 2018, pp. 1-4.
- S.Yang, J.Ma, and X.Huang, "Multi-hop underwater acoustic networks based on bats codes," in Proceedings of the Thirteenth ACM Interna- tional Conference on Underwater Networks & Systems, ser. WUWNet '18. New York, NY, USA: ACM, 2018, pp. 30:1-30:5.