

## Deriving formulas used in fastmath.c

This document derives the formulas to approximate floating-point operations, as well as their Newton–Raphson method.

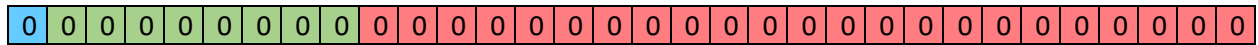
Based on the famous “Fast inverse square root” algorithm.

Additionally, the following resources helped me a great deal:

Lomont, Chris (February 2003). ["Fast Inverse Square Root"](#)

YouTube - @Nemean: ["Fast Inverse Square Root — A Quake III Algorithm"](#)

### Float 32



Bit representation =  $2^{M_{\text{bits}}} * \text{Exp} + \text{Man.}$

$$\text{Value} = \left(1 + \frac{\text{Man.}}{2^{M_{\text{bits}}}}\right) * 2^{\text{Exp} - E_{\text{bias}}}$$

Note: The mantissa and exponent are treated as unsigned integers.

The exponent needs to represent negative values for number less than one (e.g.,  $2^{-2} = 0.25$ ). Instead of using a signed bit/2's complement, a float's exponent is biased so the first half of values are negative.

### Taking the binary log of a float

$$\log_2(\text{float}) = \log_2 \left( \left(1 + \frac{\text{Man}}{2^{M_{\text{bits}}}}\right) * 2^{\text{Exp} - E_{\text{bias}}} \right)$$

$$= \log_2 \left( 1 + \frac{\text{Man}}{2^{M_{\text{bits}}}} \right) + \log_2 (2^{\text{Exp} - E_{\text{bias}}})$$

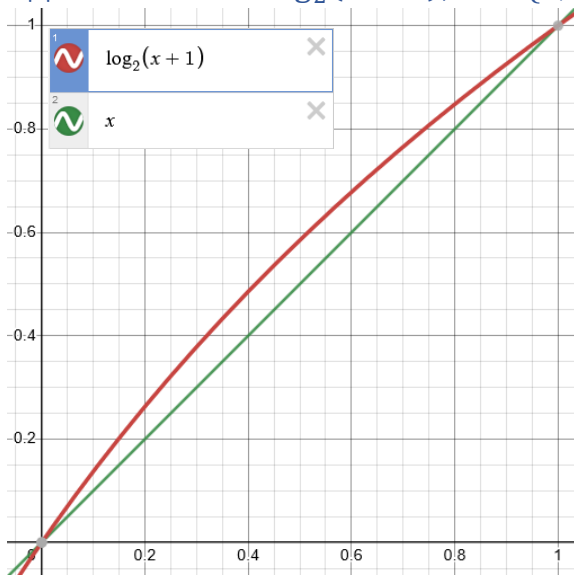
$$= \log_2 \left( 1 + \frac{\text{Man}}{2^{M_{\text{bits}}}} \right) + \text{Exp} - E_{\text{bias}}$$

$$\frac{\text{bit\_repr}}{2^{M_{\text{bits}}}} - E_{\text{bias}} = \text{Exp. Man} - E_{\text{bias}}$$

$$= \lfloor \log_2(x) \rfloor + \frac{1}{2} - \frac{1}{\pi} * \tan^{-1}(\cot(\pi * x * 2^{1 - \lfloor \log_2(x) \rfloor}))$$

('x' being the input float's value)

Approximation for  $\log_2(x + 1)$ ,  $x \in \{0,1\}$



As you can see,  $\log_2(x + 1) \approx x$ , notably they are the same at  $x=0$  and  $1$ . We can improve this approximation by reducing the absolute difference between them, so on average they are closer.

$$\int_0^1 (\log_2(x + 1) - x) dx \approx 0.0573$$

This is the term that will minimize the average difference between our two functions, so our better approximation will be:

$$\log_2(x + 1) \approx x + \mu = x + 0.0573$$

Using approximation to remove logarithm.

We can use that approximation to remove the logarithm in our function. Let the error correction term be,  $\mu = 0.0573$

$$= \log_2 \left( 1 + \frac{\text{Man}}{2^{M_{\text{bits}}}} \right) + \text{Exp} - E_{\text{bias}}$$

$$\approx \frac{\text{Man}}{2^{M_{\text{bits}}}} + \mu + \text{Exp} - E_{\text{bias}}$$

$$\approx \frac{\text{Man} + \text{Exp} * 2^{M_{\text{bits}}}}{2^{M_{\text{bits}}}} + \mu - E_{\text{bias}}$$

$$\approx \frac{1}{2^{M_{\text{bits}}}} * (\text{Man} + \text{Exp} * 2^{M_{\text{bits}}}) + \mu - E_{\text{bias}}$$

$$\approx \frac{1}{2^{M_{\text{bits}}}} * \text{bit\_repr} + \mu - E_{\text{bias}}$$

{Equal to our bit representation}

We can use this log2 approximation to calculate its inverse square root using some log laws.

$$\log\left(\frac{1}{\sqrt{n}}\right) = \log\left(n^{-\frac{1}{2}}\right) = -\frac{1}{2} * \log(n)$$

Let  $\Gamma$  be our solution.  $(1/\sqrt{y})$

$$\log_2(\Gamma) = -\frac{1}{2} \log_2(y)$$

Using our log2 approximation of a floating-point number, we get:

$$\frac{1}{2^{M_{\text{bits}}}} * \Gamma_{\text{bitrepr}} + \mu - E_{\text{bias}} = -\frac{1}{2} \left( \frac{1}{2^{M_{\text{bits}}}} * y_{\text{bitrepr}} + \mu - E_{\text{bias}} \right)$$

Solving for the bit representation of  $\Gamma$ .

$$\frac{1}{2^{M_{\text{bits}}}} * \Gamma_{\text{bitrepr}} = -\frac{1}{2} \left( \frac{1}{2^{M_{\text{bits}}}} * y_{\text{bitrepr}} + \mu - E_{\text{bias}} \right) + E_{\text{bias}} - \mu$$

$$\Gamma_{\text{bitrepr}} = \left( -\frac{1}{2} \left( \frac{1}{2^{M_{\text{bits}}}} * y_{\text{bitrepr}} + \mu - E_{\text{bias}} \right) + E_{\text{bias}} - \mu \right) * 2^{M_{\text{bits}}}$$

$$\Gamma_{\text{bitrepr}} = \left( -\frac{y_{\text{bitrepr}}}{2 * 2^{M_{\text{bits}}}} - \frac{\mu}{2} + \frac{E_{\text{bias}}}{2} + E_{\text{bias}} - \mu \right) * 2^{M_{\text{bits}}}$$

$$\Gamma_{\text{bitrepr}} = \left( -\frac{y_{\text{bitrepr}}}{2 * 2^{M_{\text{bits}}}} + \frac{E_{\text{bias}} + 2 * E_{\text{bias}} - \mu - 2\mu}{2} \right) * 2^{M_{\text{bits}}}$$

$$\Gamma_{\text{bitrepr}} = \left( -\frac{y_{\text{bitrepr}}}{2 * 2^{M_{\text{bits}}}} + \frac{3 * E_{\text{bias}} - 3\mu}{2} \right) * 2^{M_{\text{bits}}}$$

$$\Gamma_{\text{bitrepr}} = \left( -\frac{y_{\text{bitrepr}}}{2 * 2^{M_{\text{bits}}}} + \frac{3}{2} * (E_{\text{bias}} - \mu) \right) * 2^{M_{\text{bits}}}$$

$$\Gamma_{\text{bitrepr}} = -\frac{y_{\text{bitrepr}}}{2} + \frac{3}{2} * 2^{M_{\text{bits}}} * (E_{\text{bias}} - \mu)$$

$$\Gamma_{\text{bitrepr}} = \frac{3}{2} * 2^{M_{\text{bits}}} * (E_{\text{bias}} - \mu) - \frac{1}{2} * y_{\text{bitrepr}}$$

The less half of the equation is a constant, which in the original Quake 3 algorithm is the following:

$$0x5f3759df - (y \gg 1) \leftrightarrow \frac{3}{2} * 2^{M_{\text{bits}}} * (E_{\text{bias}} - \mu) - \frac{1}{2} * y_{\text{bitrepr}}$$

(Bit shifting a binary value to the right is equivalent to dividing by  $2^n$ ,  $n$  being the distance shifted.)

This approximation is pretty good, but we can improve it even further by using Newton's method.

$$x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}$$

This equation recursively finds an 'x', in which  $f(x) \rightarrow 0$ . In our case, we want to minimize the difference between our approximation and the actual value.

One way to measure how accurate an approximation is to a function, is to put that approximation through the inverse of that function and measure how well it matches the original input.

$$f(\text{approx}) = \frac{1}{\text{approx}^2} - \text{input}$$

Using this expression, we can create a Newton's method iteration. But first we need its derivative.

$$f'(\text{approx}) = \frac{dy}{d(\text{approx})} (\text{approx}^{-2} - \text{input})$$

$$f'(\text{approx}) = -2 * \text{approx}^{-3}$$

Plugging these into Newton's method, we get:

$$\text{approx}_{\text{new}} = \text{approx} - \frac{\frac{1}{\text{approx}^2} - \text{input}}{\frac{-2}{\text{approx}^3}}$$

$$\text{approx}_{\text{new}} = \text{approx} - \frac{\left(\frac{1}{\text{approx}^2} - \text{input}\right) * \text{approx}^3}{-2}$$

$$\text{approx}_{\text{new}} = \text{approx} - \frac{\left(\frac{1 - \text{approx}^2 * \text{input}}{\text{approx}^2}\right) * \text{approx}^3}{-2}$$

$$\text{approx}_{\text{new}} = \text{approx} - \frac{\text{approx} - \text{approx}^3 * \text{input}}{-2}$$

$$\text{approx}_{\text{new}} = \frac{2 * \text{approx} + \text{approx}}{2} - \frac{\text{input} * \text{approx}^3}{2}$$

$$\text{approx}_{\text{new}} = \text{approx} * \frac{3}{2} - \frac{1}{2} * \text{input} * \text{approx}^3$$

$$\text{approx}_{\text{new}} = \text{approx} * \left(\frac{3}{2} - \frac{1}{2} * \text{input} * \text{approx} * \text{approx}\right)$$

This is equivalent to the line:

$$y = y * (\text{threehalfs} - (0.5F * \text{number} * y * y))$$

### Exponentiation to any power

$$\frac{1}{2^{M_{\text{bits}}}} * \Gamma_{\text{bit\_repr}} = \text{pow} \left( \frac{1}{2^{M_{\text{bits}}}} * y_{\text{bit\_repr}} + \mu - E_{\text{bias}} \right) + E_{\text{bias}} - \mu$$

$$\Gamma_{\text{bit\_repr}} = \left( \text{pow} \left( \frac{1}{2^{M_{\text{bits}}}} * y_{\text{bit\_repr}} + \mu - E_{\text{bias}} \right) + E_{\text{bias}} - \mu \right) * 2^{M_{\text{bits}}}$$

$$\Gamma_{\text{bit\_repr}} = \text{pow} * 2^{M_{\text{bits}}} \left( \frac{1}{2^{M_{\text{bits}}}} * y_{\text{bit\_repr}} + \mu - E_{\text{bias}} \right) + 2^{M_{\text{bits}}} * E_{\text{bias}} - 2^{M_{\text{bits}}} * \mu$$

$$\Gamma_{\text{bit\_repr}} = \text{pow} * y_{\text{bit\_repr}} + 2^{M_{\text{bits}}} * \mu * \text{pow} - \text{pow} * 2^{M_{\text{bits}}} * E_{\text{bias}} + 2^{M_{\text{bits}}} * E_{\text{bias}} - 2^{M_{\text{bits}}} * \mu$$

$$\Gamma_{\text{bit\_repr}} = \text{pow} * y_{\text{bit\_repr}} + 2^{M_{\text{bits}}} * \mu * (\text{pow} - 1) - 2^{M_{\text{bits}}} * E_{\text{bias}} * (\text{pow} - 1)$$

$$\Gamma_{\text{bit\_repr}} = \text{pow} * y_{\text{bit\_repr}} + 2^{M_{\text{bits}}} * (\text{pow} - 1) * (\mu - E_{\text{bias}})$$

### Newton's method for any power

$$x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}$$

$$f(\text{approx}) = \text{approx}^{\frac{1}{\text{pow}}} - \text{input}$$

$$f'(\text{approx}) = \frac{dy}{d(\text{approx})} \left( \text{approx}^{\frac{1}{\text{pow}}} - \text{input} \right)$$

$$f'(\text{approx}) = \frac{1}{\text{pow}} * \text{approx}^{\frac{1}{\text{pow}} - 1}$$

$$\text{approx}_{\text{new}} = \text{approx} - \frac{\text{approx}^{\frac{1}{\text{pow}}} - \text{input}}{\frac{1}{\text{pow}} * \text{approx}^{\frac{1}{\text{pow}} - 1}}$$

$$\text{approx}_{\text{new}} = \text{approx} - \text{pow} * \text{approx} * \frac{\text{approx}^{\frac{1}{\text{pow}}} - \text{input}}{\text{approx}^{\frac{1}{\text{pow}}}}$$

$$\text{approx}_{\text{new}} = \text{approx} * \left( 1 - \text{pow} * \left( 1 - \frac{\text{input}}{\text{approx}^{\frac{1}{\text{pow}}}} \right) \right)$$

This requires using an exponential, which is what this algorithm is trying to avoid.

General forms for  $\frac{1}{\sqrt[n]{x}}$

Float approximation

$$\Gamma_{\text{bitrepr}} = \frac{1}{n} \left( (n+1) * 2^{\text{Mbits}} * (\text{E}_{\text{bias}} - \mu) - \mathbf{y}_{\text{bitrepr}} \right)$$

Newton's method

$$\text{approx}_{\text{new}} = \text{approx} * \left( 1 + \frac{1}{n} * (1 - \text{input} * \text{approx}^n) \right)$$

Log2 Newton's method

$$f(\text{approx}) = 2^{\text{approx}} - \text{input}$$

$$f'(\text{approx}) = 2^{\text{approx}} * \ln(2)$$

$$\text{approx}_{\text{new}} = \text{approx} - \frac{2^{\text{approx}} - \text{input}}{2^{\text{approx}} * \ln(2)}$$

$$\text{approx}_{\text{new}} = \text{approx} - \frac{1}{\ln(2)} * \left( 1 - \frac{\text{input}}{2^{\text{approx}}} \right)$$

Reciprocal

Float approximation

$$\Gamma_{\text{bitrepr}} = (-1) * \mathbf{y}_{\text{bitrepr}} + 2^{\text{Mbits}} * (-1 - 1) * (\mu - \text{E}_{\text{bias}})$$

$$\Gamma_{\text{bitrepr}} = -2^{\text{Mbits}} * 2 * (\mu - \text{E}_{\text{bias}}) - \mathbf{y}_{\text{bitrepr}}$$

Newton's method

$$\text{approx}_{\text{new}} = \text{approx} * \left( 1 - (-1) * \left( 1 - \frac{\text{input}}{\text{approx}^{-1}} \right) \right)$$

$$\text{approx}_{\text{new}} = \text{approx} * (1 + (1 - \text{input} * \text{approx}))$$

$$\text{approx}_{\text{new}} = \text{approx} * (2 - \text{input} * \text{approx})$$

## Appendix

### Natural log

$$\ln(\text{float}) = \ln\left(1 + \frac{\text{Man}}{2^{\text{Mbits}}}\right) + \ln(2^{\text{Exp} - \text{E}_{\text{bias}}})$$

$$\ln(\text{float}) = \ln\left(1 + \frac{\text{Man}}{2^{\text{Mbits}}}\right) + (\text{Exp} - \text{E}_{\text{bias}}) * \ln(2)$$

$$\ln(\text{float}) \approx \ln(2) * \frac{\text{Man}}{2^{\text{Mbits}}} + \mu + (\text{Exp} - \text{E}_{\text{bias}}) * \ln(2) \quad \{\mu = 0.03972\}$$

$$\ln(\text{float}) \approx \ln(2) \left( \frac{\text{Man}}{2^{\text{Mbits}}} + \text{Exp} - \text{E}_{\text{bias}} \right) + \mu$$

$$\ln(\text{float}) \approx \ln(2) \left( \frac{\text{Man} + \text{Exp} * 2^{\text{Mbits}}}{2^{\text{Mbits}}} - \text{E}_{\text{bias}} \right) + \mu$$

$$\ln(\text{float}) \approx \ln(2) \left( \frac{\text{bit}_{\text{repr}}}{2^{\text{Mbits}}} - \text{E}_{\text{bias}} \right) + \mu$$

### Log gamma function approximations

Lanczos (very accurate)

$$\ln(\mathbf{x}!) \approx \frac{1}{2} \ln(2\pi) + \left(\mathbf{x} + \frac{1}{2}\right) * \ln\left(\mathbf{x} + A + \frac{1}{2}\right) - \left(\mathbf{x} + A + \frac{1}{2}\right) \quad A = 0.3$$

Stirling's

$$\ln(\mathbf{float}!) \approx \frac{1}{2} \ln(2\pi) + \left(\mathbf{x} + \frac{1}{2}\right) * \ln(\mathbf{Z}) - \mathbf{Z} \quad \mathbf{Z} = \mathbf{x} + A + \frac{1}{2}$$

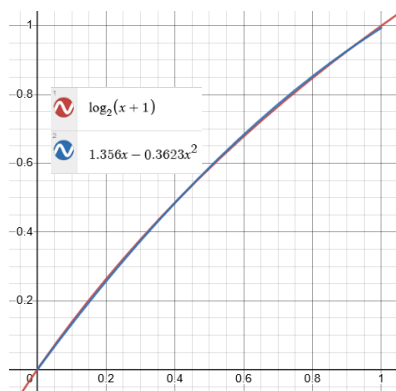
### Binary log for $x$ : (0,1)

Used for the fractional part in the *fastLog2Bits()* and in *fastLog2Alt()*.

$$\log_2(x + 1) \approx 1.356x - 0.3623x^2 \quad \{0 < x < 1\}$$

$$\log_2(x + 1) \approx 1.4088x - 0.49328x^2 + 0.086x^4 \quad \{0 < x < 1\}$$

$$\log_2(x + 1) \approx \frac{3.3739x}{x + 2.3791} \quad \{0 < x < 1\}$$





### Better general power approximation

$$\log_2(\mathbf{float}) = \log_2\left(1 + \frac{\mathbf{Man}}{2^{\mathbf{M}_{\text{bits}}}}\right) + \mathbf{Exp} - \mathbf{E}_{\text{bias}}$$

$$\log_2(\mathbf{float}^{\text{pow}}) = \text{pow}\left(\log_2\left(1 + \frac{\mathbf{Man}}{2^{\mathbf{M}_{\text{bits}}}}\right) + \mathbf{Exp} - \mathbf{E}_{\text{bias}}\right)$$

$$\frac{1}{2^{\mathbf{M}_{\text{bits}}}} * \mathbf{\Gamma}_{\text{bitrepr}} + \mu - \mathbf{E}_{\text{bias}} \approx \text{pow}\left(\log_2\left(1 + \frac{\mathbf{Man}}{2^{\mathbf{M}_{\text{bits}}}}\right) + \mathbf{Exp} - \mathbf{E}_{\text{bias}}\right)$$

$$\mathbf{\Gamma}_{\text{bitrepr}} \approx \left(\text{pow}\left(\log_2\left(1 + \frac{\mathbf{Man}}{2^{\mathbf{M}_{\text{bits}}}}\right) + \mathbf{Exp} - \mathbf{E}_{\text{bias}}\right) - \mu + \mathbf{E}_{\text{bias}}\right) * 2^{\mathbf{M}_{\text{bits}}}$$

$$\mathbf{\Gamma}_{\text{bitrepr}} \approx \text{pow}\left(2^{\mathbf{M}_{\text{bits}}} * \log_2\left(1 + \frac{\mathbf{Man}}{2^{\mathbf{M}_{\text{bits}}}}\right) + 2^{\mathbf{M}_{\text{bits}}} * \mathbf{Exp} - 2^{\mathbf{M}_{\text{bits}}} * \mathbf{E}_{\text{bias}}\right) - 2^{\mathbf{M}_{\text{bits}}} * \mu + 2^{\mathbf{M}_{\text{bits}}} * \mathbf{E}_{\text{bias}}$$

$$\mathbf{\Gamma}_{\text{bitrepr}} \approx \text{pow}\left(2^{\mathbf{M}_{\text{bits}}} * \log_2\left(1 + \frac{\mathbf{Man}}{2^{\mathbf{M}_{\text{bits}}}}\right) + 2^{\mathbf{M}_{\text{bits}}} * (\mathbf{Exp} - \mathbf{E}_{\text{bias}})\right) + 2^{\mathbf{M}_{\text{bits}}} * (\mathbf{E}_{\text{bias}} - \mu)$$

$$\log_2(x + 1) \approx 1.356x - 0.3623x^2 \quad \{0 < x < 1\}$$

$$\mathbf{\Gamma}_{\text{bitrepr}} \approx \text{pow}\left(2^{\mathbf{M}_{\text{bits}}} * \frac{\mathbf{Man}}{2^{\mathbf{M}_{\text{bits}}}} \left(1.356 - 0.3623 \frac{\mathbf{Man}}{2^{\mathbf{M}_{\text{bits}}}}\right) + 2^{\mathbf{M}_{\text{bits}}} * (\mathbf{Exp} - \mathbf{E}_{\text{bias}})\right) + 2^{\mathbf{M}_{\text{bits}}} * (\mathbf{E}_{\text{bias}} - \mu)$$

$$\mathbf{\Gamma}_{\text{bitrepr}} \approx \text{pow}\left(\mathbf{Man} * \left(1.356 - 0.3623 \frac{\mathbf{Man}}{2^{\mathbf{M}_{\text{bits}}}}\right) + 2^{\mathbf{M}_{\text{bits}}} * (\mathbf{Exp} - \mathbf{E}_{\text{bias}})\right) + 2^{\mathbf{M}_{\text{bits}}} * (\mathbf{E}_{\text{bias}} - \mu)$$

$$\mathbf{\Gamma}_{\text{bitrepr}} \approx \text{pow}\left(\mathbf{Man} * \left(1.356 - \left(\frac{0.3623}{2^{\mathbf{M}_{\text{bits}}}}\right) * \mathbf{Man}\right) + 2^{\mathbf{M}_{\text{bits}}} * (\mathbf{Exp} - \mathbf{E}_{\text{bias}})\right) + 2^{\mathbf{M}_{\text{bits}}} * (\mathbf{E}_{\text{bias}} - \mu)$$

“Fast” multiplication float approximation

$$\mathbf{float}_{\text{value}} = \left(1 + \frac{\text{Man.}}{2^{\text{Mbits}}}\right) * 2^{\text{Exp} - \text{E}_{\text{bias}}}$$

$$\log_2(\mathbf{a} * \mathbf{b}) = \log_2 \left( \left(1 + \frac{\mathbf{a}_{\text{Man}}}{2^{\text{Mbits}}}\right) * 2^{\mathbf{a}_{\text{Exp}} - \text{E}_{\text{bias}}} * \left(1 + \frac{\mathbf{b}_{\text{Man}}}{2^{\text{Mbits}}}\right) * 2^{\mathbf{b}_{\text{Exp}} - \text{E}_{\text{bias}}} \right)$$

$$\log_2(\mathbf{a} * \mathbf{b}) = \log_2 \left( \left(1 + \frac{\mathbf{a}_{\text{Man}}}{2^{\text{Mbits}}}\right) * 2^{\mathbf{a}_{\text{Exp}} - \text{E}_{\text{bias}}} \right) + \log_2 \left( \left(1 + \frac{\mathbf{b}_{\text{Man}}}{2^{\text{Mbits}}}\right) * 2^{\mathbf{b}_{\text{Exp}} - \text{E}_{\text{bias}}} \right)$$

$$\log_2(\mathbf{a} * \mathbf{b}) = \log_2 \left( 1 + \frac{\mathbf{a}_{\text{Man}}}{2^{\text{Mbits}}} \right) + \mathbf{a}_{\text{Exp}} - \text{E}_{\text{bias}} + \log_2 \left( 1 + \frac{\mathbf{b}_{\text{Man}}}{2^{\text{Mbits}}} \right) + \mathbf{b}_{\text{Exp}} - \text{E}_{\text{bias}}$$

$$\log_2(\mathbf{a} * \mathbf{b}) \approx \frac{\mathbf{a}_{\text{Man}}}{2^{\text{Mbits}}} + \mu + \mathbf{a}_{\text{Exp}} - \text{E}_{\text{bias}} + \frac{\mathbf{b}_{\text{Man}}}{2^{\text{Mbits}}} + \mu + \mathbf{b}_{\text{Exp}} - \text{E}_{\text{bias}}$$

$$\log_2(\mathbf{a} * \mathbf{b}) \approx \frac{\mathbf{a}_{\text{bit\_repr}}}{2^{\text{Mbits}}} + \mu - \text{E}_{\text{bias}} + \frac{\mathbf{b}_{\text{bit\_repr}}}{2^{\text{Mbits}}} + \mu - \text{E}_{\text{bias}}$$

Let  $\Gamma$  be our solution.

$$\frac{1}{2^{\text{Mbits}}} * \Gamma_{\text{bit\_repr}} + \mu - \text{E}_{\text{bias}} \approx \frac{\mathbf{a}_{\text{bit\_repr}}}{2^{\text{Mbits}}} + \mu - \text{E}_{\text{bias}} + \frac{\mathbf{b}_{\text{bit\_repr}}}{2^{\text{Mbits}}} + \mu - \text{E}_{\text{bias}}$$

$$\Gamma_{\text{bit\_repr}} \approx \left( \frac{\mathbf{a}_{\text{bit\_repr}}}{2^{\text{Mbits}}} + \mu - \text{E}_{\text{bias}} + \frac{\mathbf{b}_{\text{bit\_repr}}}{2^{\text{Mbits}}} + \mu - \text{E}_{\text{bias}} - \mu + \text{E}_{\text{bias}} \right) * 2^{\text{Mbits}}$$

$$\Gamma_{\text{bit\_repr}} \approx \mathbf{a}_{\text{bit\_repr}} + \mathbf{b}_{\text{bit\_repr}} + 2^{\text{Mbits}} * (\mu - \text{E}_{\text{bias}})$$

Division

$$\Gamma_{\text{bit\_repr}} \approx \left( \frac{\mathbf{a}_{\text{bit\_repr}}}{2^{\text{Mbits}}} + \mu - \text{E}_{\text{bias}} - \frac{\mathbf{b}_{\text{bit\_repr}}}{2^{\text{Mbits}}} - \mu + \text{E}_{\text{bias}} - \mu + \text{E}_{\text{bias}} \right) * 2^{\text{Mbits}}$$

$$\Gamma_{\text{bit\_repr}} \approx \left( \frac{\mathbf{a}_{\text{bit\_repr}}}{2^{\text{Mbits}}} - \frac{\mathbf{b}_{\text{bit\_repr}}}{2^{\text{Mbits}}} - \mu + \text{E}_{\text{bias}} \right) * 2^{\text{Mbits}}$$

$$\Gamma_{\text{bit\_repr}} \approx \mathbf{a}_{\text{bit\_repr}} - \mathbf{b}_{\text{bit\_repr}} - 2^{\text{Mbits}} * (\mu - \text{E}_{\text{bias}})$$

Constant raised to a variable power

Float approximation for  $2^n$

$$\log_2(\Gamma) = \log_2(2^{\text{pow}})$$

$$\frac{1}{2^{M_{\text{bits}}}} * \Gamma_{\text{bitrepr}} + \mu - E_{\text{bias}} \approx \text{pow}$$

$$\Gamma_{\text{bitrepr}} \approx (\text{pow} - \mu + E_{\text{bias}}) * 2^{M_{\text{bits}}}$$

Float approximation for  $e^n$

$$\frac{1}{2^{M_{\text{bits}}}} * \Gamma_{\text{bitrepr}} + \mu - E_{\text{bias}} \approx \text{pow} * \log_2(e)$$

$$\Gamma_{\text{bitrepr}} \approx (\text{pow} * \log_2(e) - \mu + E_{\text{bias}}) * 2^{M_{\text{bits}}}$$

$$\Gamma_{\text{bitrepr}} \approx \text{pow} * \log_2(e) * 2^{M_{\text{bits}}} + 2^{M_{\text{bits}}} * (E_{\text{bias}} - \mu)$$

Use 15 for  $x$ :  $(0, 1)$

$$\ln(x + 1) = 2 * \sum_{n=1}^{\infty} \left( \frac{1}{2n-1} * \left( \frac{x}{x+2} \right)^{2n-1} \right)$$

$$2 \left( \frac{1}{1} * \frac{x}{x+2} + \frac{1}{3} * \frac{x}{x+2} * \frac{x}{x+2} * \frac{x}{x+2} + \frac{1}{5} * \frac{x}{x+2} * \frac{x}{x+2} * \frac{x}{x+2} * \frac{x}{x+2} * \frac{x}{x+2} \dots \right)$$

$$= \frac{1}{1} * c \quad c = \frac{x}{x+2}$$

$$+= \frac{1}{3} * c \quad c *= \frac{x}{x+2} * \frac{x}{x+2}$$

$$+= \frac{1}{5} * c \quad c *= \frac{x}{x+2} * \frac{x}{x+2}$$

$$+= \frac{1}{7} * c \quad c *= \frac{x}{x+2} * \frac{x}{x+2}$$