



**UTM**  
UNIVERSITI TEKNOLOGI MALAYSIA

# Chapter 1

Part 3

## Fundamental and Elements of Logic



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## Fundamental and Elements of Logic

# Why Are We Studying Logic?

## Some of the reasons:

- Logic is the foundation for computer operation
- Logical conditions are common in programs

### Example:

**Selection:** `if (score <= max) { ... }`

**Iteration:** `while (i < limit && list[i] != sentinel) ...`

- All manner of structures in computing have properties that need to be proven (and proofs that need to be understood).

**Example:** Trees, Graphs, Recursive Algorithms, ...

- Programs can be proven correct.
- Computational linguistics must represent and reason about human language, and language represents thought (and thus also logic).

# PROPOSITION

A **statement** or a **proposition**, is a declarative sentence that is **either TRUE or FALSE, but not both**.

## Example:

- 4 is less than 3. *false*
- 7 is an even integer. *false*
- Washington, DC, is the capital of United State. *True*

# Example

- i) Why do we study mathematics?
- ii) Study logic.
- iii) What is your name?
- iv) Quiet, please.

**The above sentences are not propositions. Why ?**

(i) & (iii) : is question, not a statement.  
(ii) & (iv) : is a command.

# Example

- i) The temperature on the surface of the planet Venus is 800 F.
- ii) The sun will come out tomorrow.

## Propositions? Why?

- Is a statement since it is either true or false, but not both.
- However, we do not know at this time to determine whether it is true or false.

# CONJUNCTIONS

- Compound propositions formed in English with the word “**and**”
- Formed in logic with the caret symbol (“ $\wedge$ ”) / & &
- **True** only when **both participating propositions** are true.

**TRUTH TABLE:** This tables aid in the evaluation of **compound propositions**.

p	q	$p \wedge q$
T	T	T
T	F	F
F	T	F
F	F	F

True (T), False (F)

# Example

$p$  : 2 is an even integer

$q$  : 3 is an odd number

propositions

$p \wedge q$

symbols

: 2 is an even integer and 3 is an odd number

statements

$p$  : today is Monday

$q$  : it is hot

$p \wedge q$ : today is Monday and it is hot



# Example

## Proposition

$p$  : 2 divides 4

$q$  : 2 divides 6

## Symbol & Statement

$p \wedge q$ : 2 divides 4 and 2 divides 6.

or,

$p \wedge q$ : 2 divides both 4 and 6.

# Example

## Proposition

$p$  : 5 is an integer

$q$  : 5 is not an odd integer

## Symbol & Statement

$p \wedge q$ : 5 is an integer and 5 is not an odd integer.

or,

$p \wedge q$ : 5 is an integer but 5 is not an odd integer.

# DISJUNCTION

- Compound propositions formed in English with the word “**or**”,
- Formed in logic with the caret symbol (“ $\vee$ ”) ✓ ||
- **True** when **one or both** participating propositions are true.

The **truth table** for  $p \vee q$

<b>p</b>	<b>q</b>	<b><math>p \vee q</math></b>
T	T	T
T	F	T
F	T	T
F	F	F

# Example

**$p$** : 2 is an integer

**$q$** : 3 is greater than 5

**$p \vee q$**  : 2 is an integer or 3 is greater than 5

**$p$**  :  $1+1=3$

**$q$**  : A decade is 10 years

**$p \vee q$**  :  $1+1=3$  or a decade is 10 years

# Example

**$p$**  : 3 is an even integer

**$q$**  : 3 is an odd integer

**$p \vee q$**

3 is an even integer or 3 is an odd integer  
or

3 is an even integer or an odd integer

# NEGATION

Negating a proposition simply flips its value. Symbols representing negation include:

$\neg x$ ,  $\bar{x}$ ,  $\sim x$ ,  $x'$  (NOT)

Let  $p$  be a proposition.  
The negation of  $p$ , written  $\neg p$  is the statement obtained by negating statement  $p$ .

The **truth table** of  $\neg p$

$p$	$\neg p$
T	F
F	T

# Example

$p$  : 2 is positive

$\neg p$  : 2 is not positive.

$p$  : 4 is less than 3

$\neg p$  : 4 is not less than 3.

# CONDITIONAL PROPOSITIONS

Let *p* and *q* be propositions.

*“if **p**, then **q**”*

is a statement called a **conditional proposition**,  
written as

$$p \rightarrow q$$



# CONDITIONAL PROPOSITIONS

The **truth table** of  $p \rightarrow q$   
 $\Rightarrow$  *Cause and effect relationship*

FALSE if  
 $p = \text{True}$   
 and  $q$   
 $= \text{false}$

$p \wedge q$                        $q$                        $(p \wedge q)$

$p$	$q$	$p \rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

TRUE if  
 both  
 true or  
 $p = \text{false}$   
 for any  
 value of  
 $q$

# Example

***p*** : today is Sunday

***q*** : I will go for a walk

***p* → *q*** : If today is Sunday, then I will go for a walk.

***p*** : I get a bonus

***q*** : I will buy a new car

***p* → *q***: If I get a bonus, then I will buy a new car

# Example

***p*** :  $x/2$  is an integer.

***q*** :  $x$  is an even integer.

$p \rightarrow q$  : if  $x/2$  is an integer, then  $x$  is an even integer.

# BICONDITIONAL

Let  $p$  and  $q$  be propositions.

*“ $p$  if and only if  $q$ ”*

is a statement called a **biconditional proposition**,  
written as

$$p \leftrightarrow q \quad \text{or} \quad p \text{ iff } q$$

# BICONDITIONAL

If both  $p$  and  $q$  are true or false  
then it is true

The **truth table** of  $p \leftrightarrow q$ :

$p$	$q$	$p \leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

# Example

$p$  : my program will compile

$q$  : it has no syntax error.

$p \leftrightarrow q$  : My program will compile if and only if it has no syntax error.

$p$  :  $x$  is divisible by 3

$q$  :  $x$  is divisible by 9

$p \leftrightarrow q$  :  $x$  is divisible by 3 if and only if  $x$  is divisible by 9.

- The compound propositions  **$Q$**  and  **$R$**  are made up of the propositions  $p_1, \dots, p_n$ .
- **$Q$**  and  **$R$**  are logically equivalent and write,  
$$Q \equiv R$$
provided that given any truth values of  $p_1, \dots, p_n$ , either  **$Q$**  and  **$R$**  are **both true** or  **$Q$**  and  **$R$**  are **both false**.

# Example

$$Q = p \rightarrow q \quad R = \neg q \rightarrow \neg p$$

Show that,  $Q \equiv R$

The **truth table** shows that,  $Q \equiv R$

Shown

$p$	$q$	$p \rightarrow q$	$\neg q \rightarrow \neg p$	$\neg q$	$\neg p$
T	T	T	T	F	F
T	F	F	F	T	F
F	T	T	T	F	T
F	F	T	T	T	T



# Example

Show that,  $\neg(p \rightarrow q) \equiv p \wedge \neg q$

The **truth table** shows that,  $\neg(p \rightarrow q) \equiv p \wedge \neg q$

$p$	$q$	$\neg(p \rightarrow q)$	$p \wedge \neg q$	$\neg q$
T	T	F ✓	F ✓	F
T	F	T ✓	T ✓	T
F	T	F ✓	F ✓	F
F	F	F ✓	F ✓	T

# PRECEDENCE OF LOGICAL CONNECTIVES

*Keutamaan*

**Precedence of logical connectives  
is as follows:**

not	$\neg$	↑	Highest
and	$\wedge$		
or	$\vee$		
If...then	$\rightarrow$		
If and only if	$\leftrightarrow$		Lowest

# Example

Construct the truth table for,  $A = \neg(p \vee q) \rightarrow (q \wedge p)$

## Solution

$p$	$q$	$(p \vee q)$	$\neg(p \vee q)$	$(q \wedge p)$	$A$
T	T	T	F	T	T
T	F	T	F	F	T
F	T	T	F	F	T
F	F	F	T	F	F

Logic and set theory go very well together. The previous definitions can be made very succinct: <sup>ringkas</sup>

$x \notin A$  if and only if  $\neg(x \in A)$

$A \subseteq B$  if and only if  $(x \in A \rightarrow x \in B)$  is True

$x \in (A \cap B)$  if and only if  $(x \in A \wedge x \in B)$

$x \in (A \cup B)$  if and only if  $(x \in A \vee x \in B)$

$x \in A - B$  if and only if  $(x \in A \wedge x \notin B)$

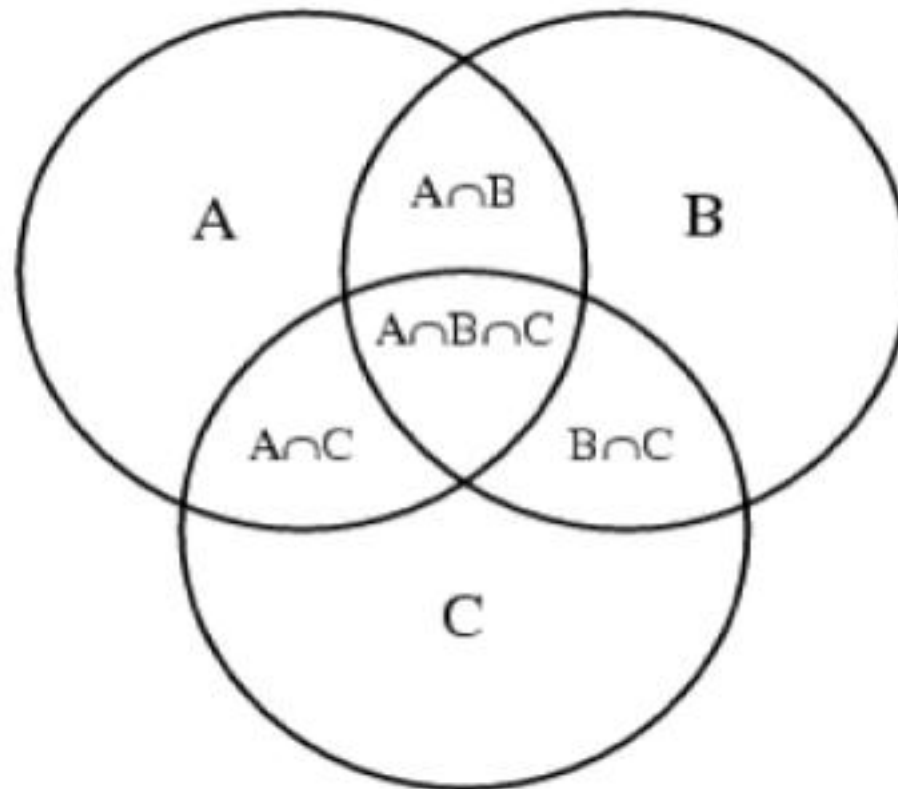
$x \in A \Delta B$  if and only if  $(x \in A \wedge x \notin B) \vee (x \in B \wedge x \notin A)$

$x \in A'$  if and only if  $\neg(x \in A)$

$X \in P(A)$  if and only if  $X \subseteq A$

# Venn Diagrams

**Venn Diagrams** are used to depict the various unions, subsets, complements, intersections etc. of sets.



# Logic and Sets are closely related

## Tautology

$$p \vee q \leftrightarrow q \vee p$$

$$p \wedge q \leftrightarrow q \wedge p$$

$$p \vee (q \vee r) \leftrightarrow (p \vee q) \vee r$$

$$p \wedge (q \wedge r) \leftrightarrow (p \wedge q) \wedge r$$

$$p \vee (q \wedge r) \leftrightarrow (p \vee q) \wedge (p \vee r)$$

$$p \wedge (q \vee r) \leftrightarrow (p \wedge q) \vee (p \wedge r)$$

$$p \wedge \neg q \leftrightarrow p \wedge \neg(p \wedge q)$$

$$p \wedge \neg(q \vee r) \leftrightarrow (p \wedge \neg q) \wedge (p \wedge \neg r)$$

$$p \wedge \neg(q \wedge r) \leftrightarrow (p \wedge \neg q) \vee (p \wedge \neg r)$$

$$p \wedge (q \wedge \neg r) \leftrightarrow (p \wedge q) \wedge \neg(p \wedge \neg r)$$

$$p \vee (q \wedge \neg r) \leftrightarrow (p \vee q) \wedge \neg(r \wedge \neg p)$$

$$p \wedge \neg \vee (q \wedge \neg r) \leftrightarrow (p \wedge \neg q) \vee (p \wedge r)$$

## Set Operation Identity

$$A \cup B = B \cup A$$

$$A \cap B = B \cap A$$

$$A \cup (B \cup C) = (A \cup B) \cup C$$

$$A \cap (B \cap C) = (A \cap B) \cap C$$

$$A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$$

$$A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$$

$$A - B = A - (A \cap B)$$

$$A - (B \cap C) = (A - B) \cup (A - C)$$

$$A - (B \cup C) = (A - B) \cap (A - C)$$

$$A \cap (B - C) = (A \cap B) - (A \cap C)$$

$$A \cup (B - C) = (A \cup B) - (C - A)$$

$$A - (B - C) = (A - B) \cup (A \cap C)$$

The above identities serve as the basis for an "algebra of sets".

# Logic and Sets are closely related

## Tautology

$$p \wedge p \leftrightarrow p$$

$$p \vee p \leftrightarrow p$$

$$p \wedge \neg(q \wedge \neg q) \leftrightarrow p$$

$$p \vee \neg(q \wedge \neg q) \leftrightarrow p$$

## Contradiction

$$p \wedge \neg p$$

$$p \wedge (q \wedge \neg q)$$

$$p \wedge \neg p$$

## Set Operation Identity

$$A \cap A = A$$

$$A \cup A = A$$

$$A - \emptyset = A$$

$$A \cup \emptyset = A$$

## Set Operation Identity

$$A - A = \emptyset$$

$$A \cap \emptyset = \emptyset$$

$$A - A = \emptyset$$

The above identities serve as the basis for an "algebra of sets".

# Theorem for Logic

Let  $p$ ,  $q$  and  $r$  be propositions.

## Idempotent laws:

$$p \wedge p \equiv p$$

$$p \vee p \equiv p$$

## Truth table

$p$	$p \wedge p$	$p \vee p$
T	T	T
F	F	F



# Theorem for Logic

**Double negation law:**

$$\neg \neg p \equiv p$$

**Commutative laws:**

$$p \wedge q \equiv q \wedge p$$

$$p \vee q \equiv q \vee p$$

**Associative laws:**

$$(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$$

$$(p \vee q) \vee r \equiv p \vee (q \vee r)$$

# Theorem for Logic

**Distributive laws:**

$$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$$

$$p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$$

$p$	$q$	$r$	$p \vee (q \wedge r)$	$(p \vee q) \wedge (p \vee r)$
T	T	T	T	T
T	T	F	T	T
T	F	T	T	T
T	F	F	T	T
F	T	T	T	T
F	T	F	F	F
F	F	T	F	F
F	F	F	F	F

**PROVE**

# Theorem for Logic

**Absorption laws:**

$$p \wedge (p \vee q) \equiv p$$

$$p \vee (p \wedge q) \equiv p$$

**PROVE**

$p$	$q$	$p \wedge (p \vee q)$	$p \vee (p \wedge q)$	$p \wedge q$	$p \vee q$
T	T	T	T	T	T
T	F	T	T	F	T
F	T	F	F	F	T
F	F	F	F	F	F

Shown

# Theorem for Logic

**De Morgan's laws:**

$$\neg(p \wedge q) \equiv (\neg p) \vee (\neg q)$$

$$\neg(p \vee q) \equiv (\neg p) \wedge (\neg q)$$

The **truth table** for  $\neg(p \vee q) \equiv (\neg p) \wedge (\neg q)$

$p$	$q$	$\neg(p \vee q)$	$\neg p \wedge \neg q$
T	T	F	F
T	F	F	F
F	T	F	F
F	F	T	T

# Thank You

# Thank You