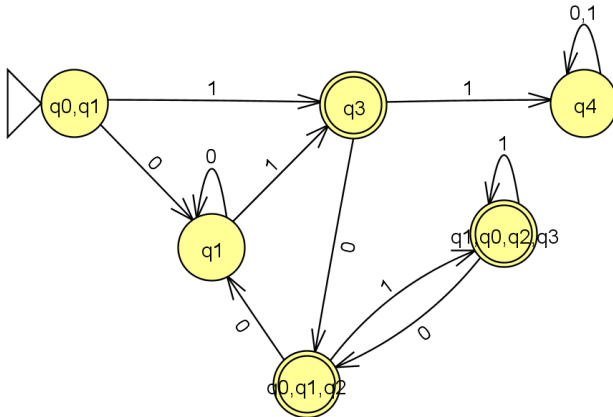


## CS 321 HW - 2

Assignment:

1) (4 pts) Convert the following NFA into an equivalent DFA

Answer:



2) (3 pts) Show that the language  $L = \{vwv : v, w \in \{a,b\}^*, |v| = 2\}$  is a regular language.

Answer:

Regular expression:

$$L(M) = ab(a+b)^*ab + bb(a+b)^*bb + ba(a+b)^*ba + aa(a+b)^*aa;$$

In result we are proof that L is regular language.

3) (4 pts) Prove that if L is regular language then  $L^R$  is a regular language.

Answer:

[Definition of reverse]

Let  $M = (Q_M, \Sigma_M, \delta, q_M, F_M)$  be a NFA and  $L = L(M)$ .

The definition of  $L^R$  will be:

$$1. M^R = \{M_R \cup P\{q_s\}, \Sigma_M, \delta_{M^R}, q_s, \{q_M\}\} \text{ and } q_s \notin Q_M$$

2.  $p \in \delta_M(q, a) \Leftrightarrow p \in \delta_{M^R}(p, a)$ , where  $a \in \Sigma_M$  and  $q, p \in Q_M$

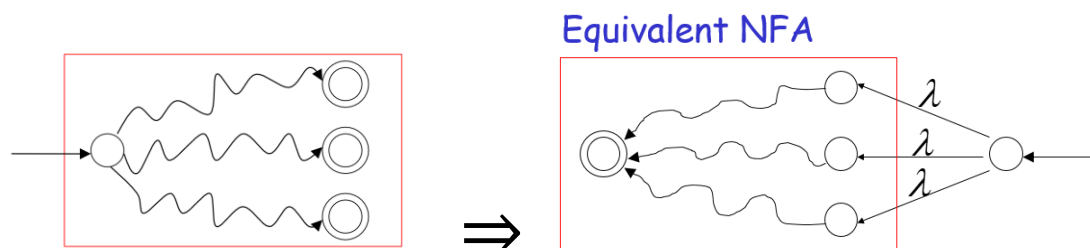
3. if  $q_M$  equal to  $F_M$  in  $L$ , then  $q_s$ , we will be  $F_M$  in  $L^R$

[Inductive hypothesis]

Assume for regular expressions  $r_1$ , that  $L(r_1)$  and  $L^R(r_1)$  are regular languages.

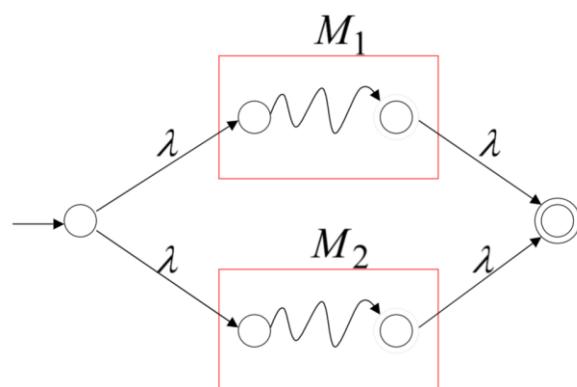
[Inductive step]

If there are multiple final state in  $L$  regular language and we need to reverse this. We have to create an equivalent of NFA, where states will connect to new node, which is a initial state, by  $\lambda$ .



As a result, we can prove that the states in  $L^R$  will be the union of  $M^R$  and  $q_s$ .

If we can make some operation like Union, Concatenation, Star or at least one of them, with a reverse language that language is regular.



4) (9 pts) Give regular expressions for the following languages on  $\Sigma = \{a, b\}$

a)  $L_1 = \{ w : n_a(w) \bmod 3 = 1 \}$ .

Answer:

$b^*a b^* (a b^*a b^*a b^*)^*$

b)  $L_2 = \{ w : w \text{ ends in } aa \}$ .

Answer:

$(a+b)^* aa$

c)  $L_3 =$  all strings containing no more than three a's.

Answer:

$b^*ab^*ab^*ab^* + b^* + b^*ab^* + b^*ab^*ab^*$

5) (4 pts) Consider a type of scientific notation for real numbers with the following rules:

a. A number can be preceded by a "+" or "-" sign or the sign may be absent.

b. Numeric values must be of the form  $cb_1b_2\dots b_n$  where  $b_i$  is any digit, but  $c$  must be nonzero.

c. The number may be followed by an exponent field of the form  $e^{+}y_1y_2$  or  $e^{-}y_1y_2$ , where  $y_i$  can be any digit  $\{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$ .

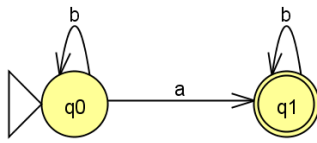
For example the strings -123e+10 and 257 represent real number in this scientific format. Give a regular expression for this scientific notation. Let  $\Sigma = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9, "+", "-", e\}$ . (Note: With this convention "+" is the sign associated with the scientific number and + the operator of the regular expression.)

Answer:

$( "+" + "-" + \lambda ) (1+2+3+4+5+6+7+8+9) (0+1+2+3+4+5+6+7+8+9)^* + ( "+" + "-" + \lambda )$   
 $(1+2+3+4+5+6+7+8+9) (0+1+2+3+4+5+6+7+8+9)^* ( "e" ) ( "+" + "-" )$   
 $(0+1+2+3+4+5+6+7+8+9) (0+1+2+3+4+5+6+7+8+9)$

6) (6 pts) Find a regular grammars for the following languages on  $\Sigma=\{a, b\}$ :

a)  $L_0$  is all strings with exactly one a



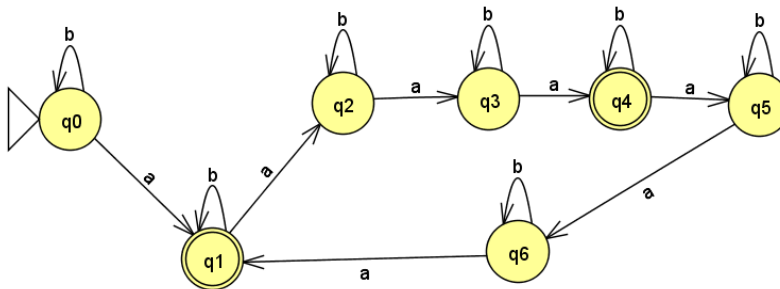
$q_0 \rightarrow S$  ;  $q_1 \rightarrow A$ ;

Grammar G:

$S \rightarrow bS \mid aA$

$A \rightarrow bA \mid \lambda$

b)  $L_1 = \{ w : n_a(w) \bmod 3 = 1 \}$ .



$q_0 \rightarrow S$  ;  $q_1 \rightarrow A$ ;  $q_2 \rightarrow B$ ;  $q_3 \rightarrow C$  ;  $q_4 \rightarrow D$ ;  $q_5 \rightarrow E$ ;  $q_6 \rightarrow F$ ;

Grammar G:

$S \rightarrow bS \mid aA$

$A \rightarrow bA \mid aB \mid \lambda$

$B \rightarrow bB \mid aC$

$C \rightarrow bC \mid aD$

$D \rightarrow bD \mid aE \mid \lambda$

$E \rightarrow bE \mid aF$

$F \rightarrow bF \mid aA$

