

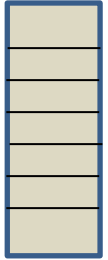
# Linked Lists



# Data Collection Structures

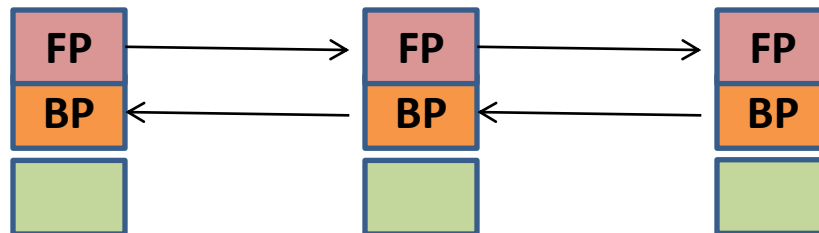
- Enhanced Arrays

- Elements consecutively stored
- Elements accessible by numeric key (Ex: `a[5]`)

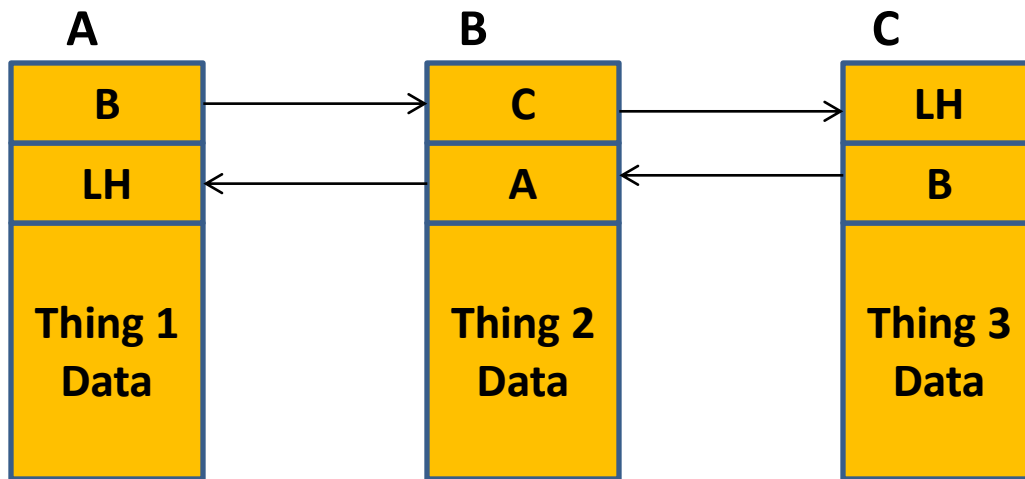


- **Linked Lists**

- Elements “daisy chained” together (hooks to neighbors)
- Access only by “walking the chain
- Why would anyone want to collect elements this way?



# Linked List: Implementation Issues



Class Thing

{

private:

Thing \* fp;

Thing \* bp;

// ?? Object must not  
assume its collection class!

public:

Thing \*getNext();

Thing \*getPrev();

bool insertNext (Thing \*);

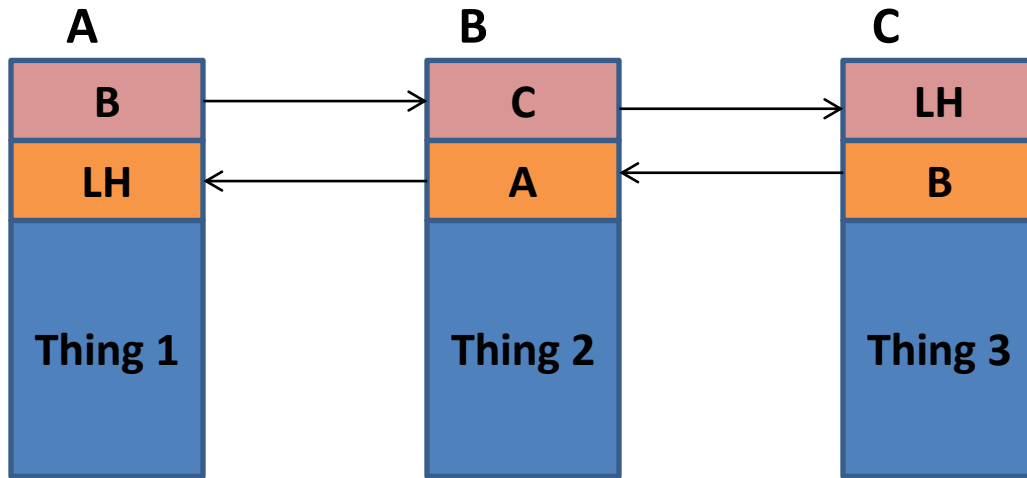
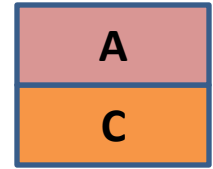
bool insertPrevious (Thing \*);

// ?? Who “*remembers*” the current position in the linked list to  
support these “next / previous” functions?

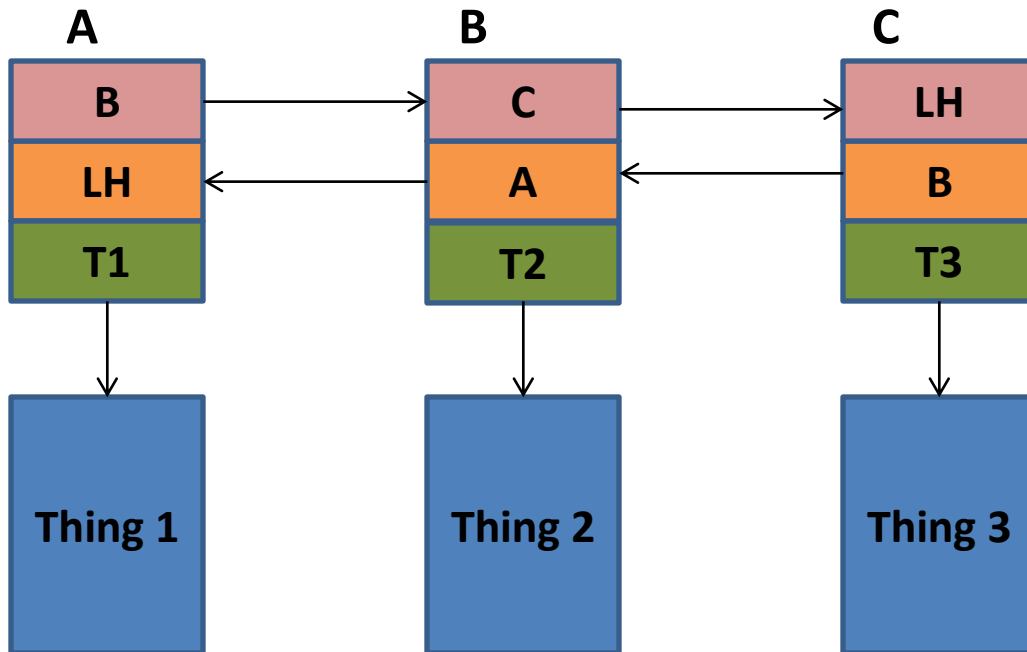
The Client? The List? The Element?



# Linked List: The Internal “Node” <sup>LH</sup>



```
Class Thing {  
    private:  
        Thing *fp;  
        Thing *bp;  
  
    // Object must not assume  
    its collection class!
```



```
Class ThingNode  
{ // Internal List class to support  
  // a Thing object in a Linked List  
    private:
```

```
        ThingNode *fp;  
        ThingNode *bp;  
        Thing * tp;
```

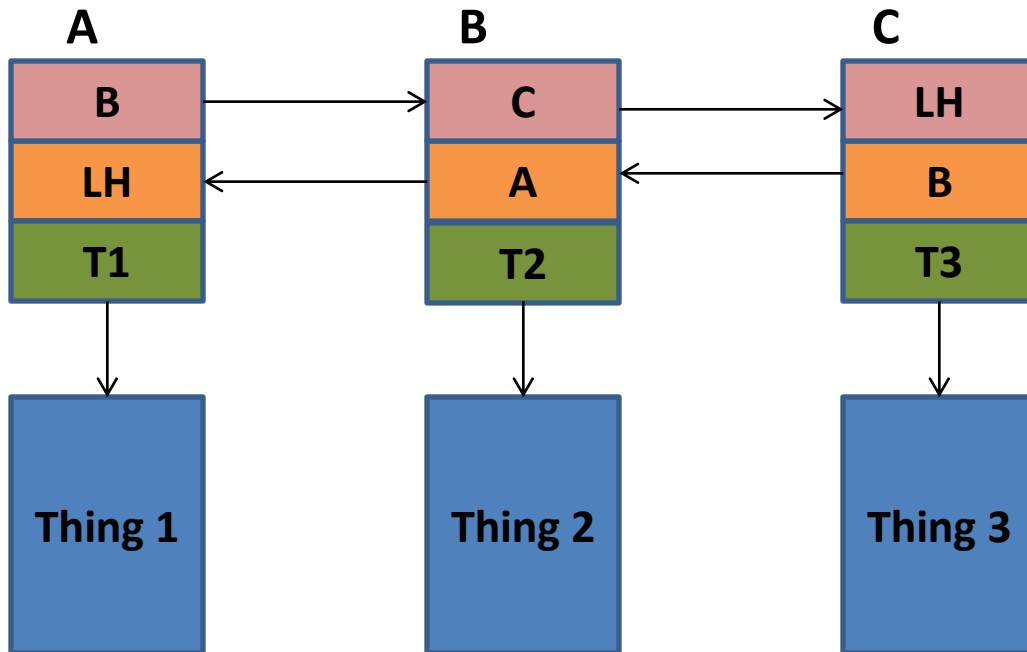
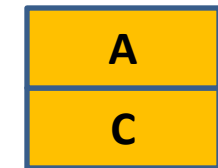
➔ Thing collected in a Linked List but doesn't know it!

# Linked List: The External “Iterator”

```
Class ThingLinkedList {  
    public:  
        bool isEmpty();  
        ThingIterator getIterator();  
}
```

```
Class ThingIterator  
{ // Works with LinkedList to “walk” element Nodes  
    public:  
        bool hasNext (); Thing& getNext();  
        bool hasPrevious (); Thing& getPrevious();  
        void remove (); // Delete current element  
        void set (Thing&); // Replace current element  
        void add (Thing&); // Insert element “here”  
}
```

LH



```
Class ThingNode  
{ // Internal List class to support  
  // a Thing object in a Linked List  
    private:
```

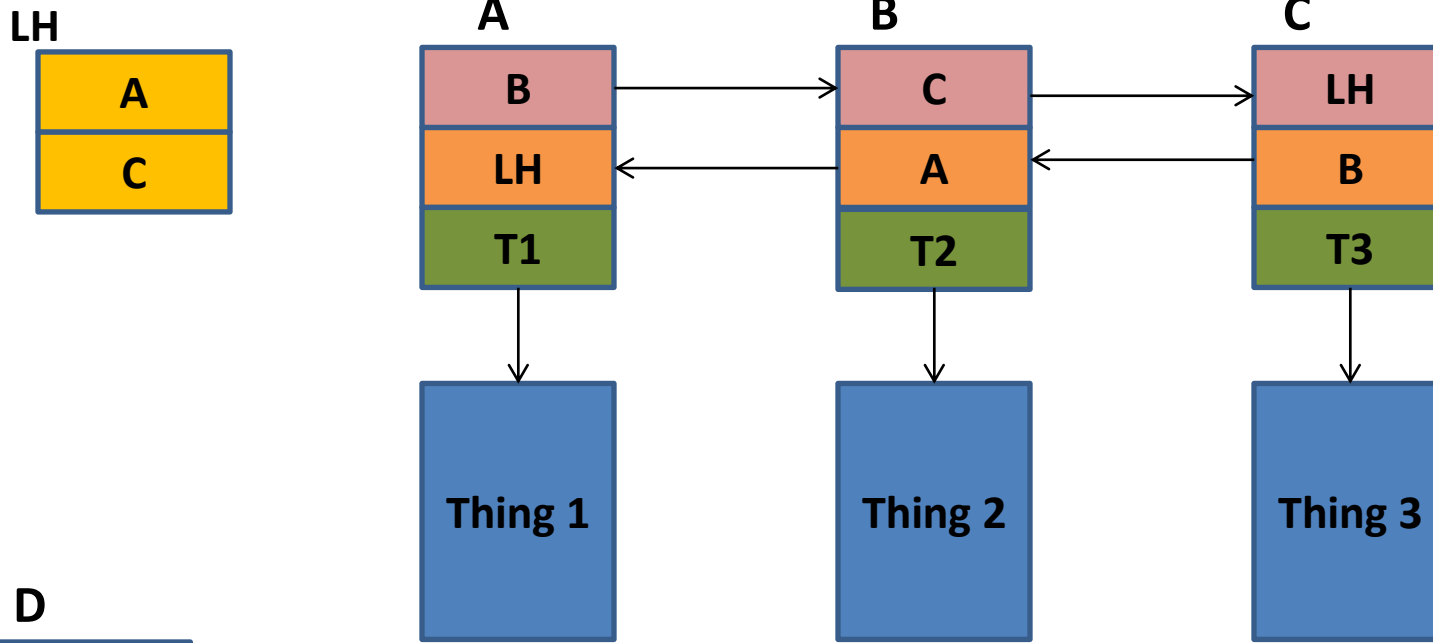
```
        ThingListEntry * fp;  
        ThingListEntry * bp;  
        Thing * tp;
```

.

.

➔ Thing collected in a Linked List but doesn't know it!

# Linked List: Insertion



Insert Thing 4 between Thing 1 and 2: → Adjust FP / BPs in adjacent nodes

1. Set D node's FP to B and its BP to A. It is ready to be inserted.

2. To actually insert it:

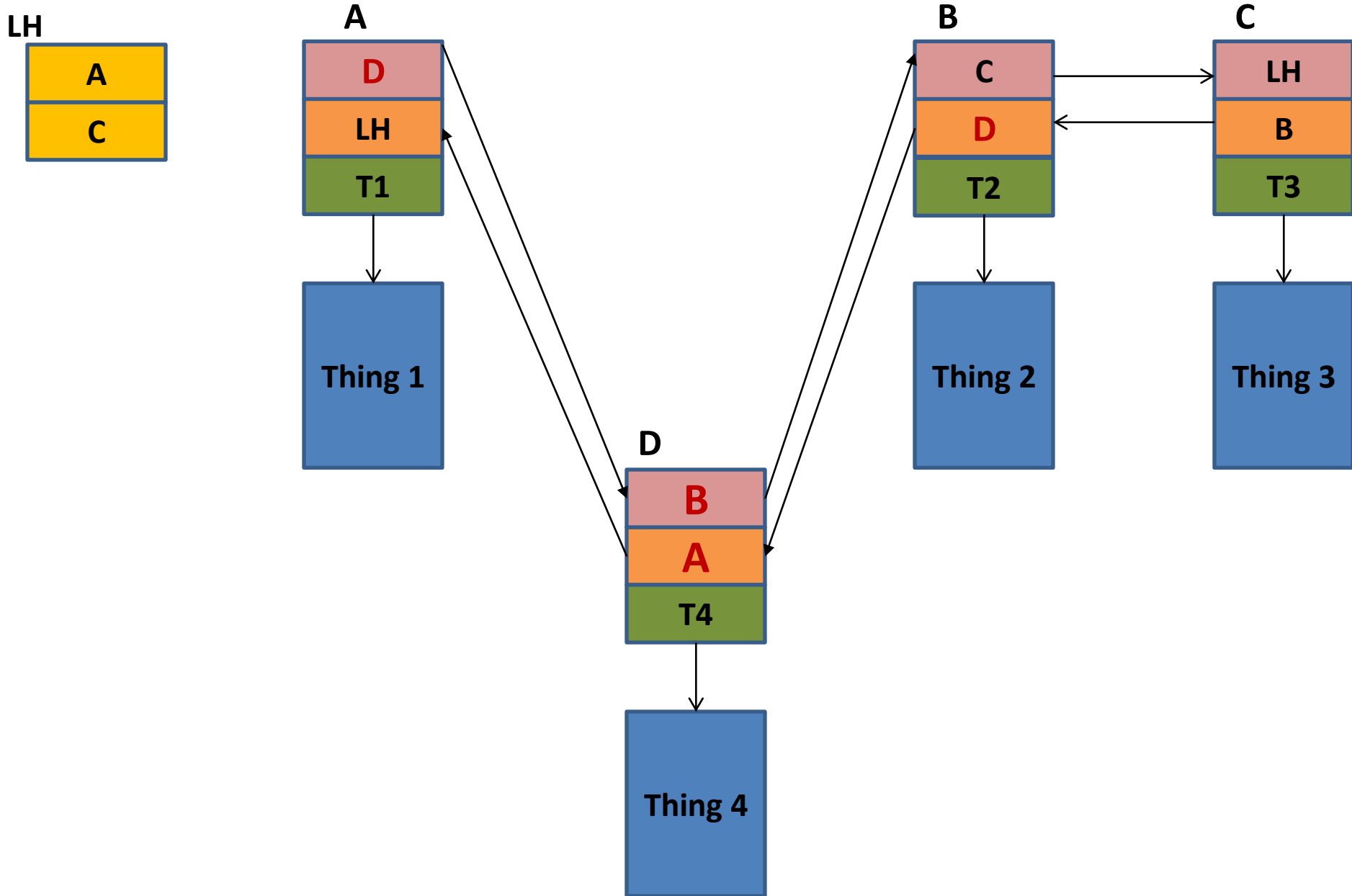
Set A's FP to D. // D now comes "after" A

Set B's BP to D // D now comes "before" B

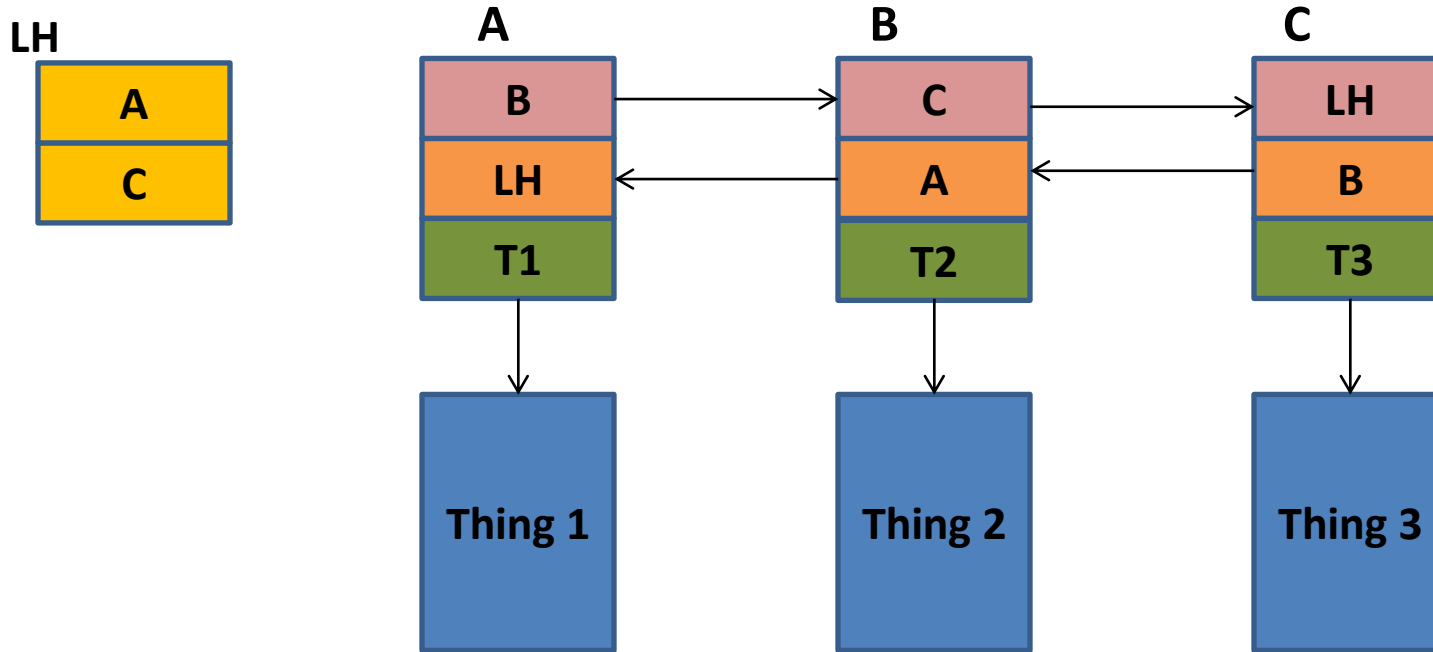
No matter how large the list, **this is all you have to do.**

**(Extended Array: Everything past insertion must be moved down 1 slot)**

# Linked List: Element Insertion



# Linked List: Element Removal



Delete Thing 2 from the List: → Adjust FP / BPs in adjacent nodes

Set A's FP to C. // B is off the forward chain

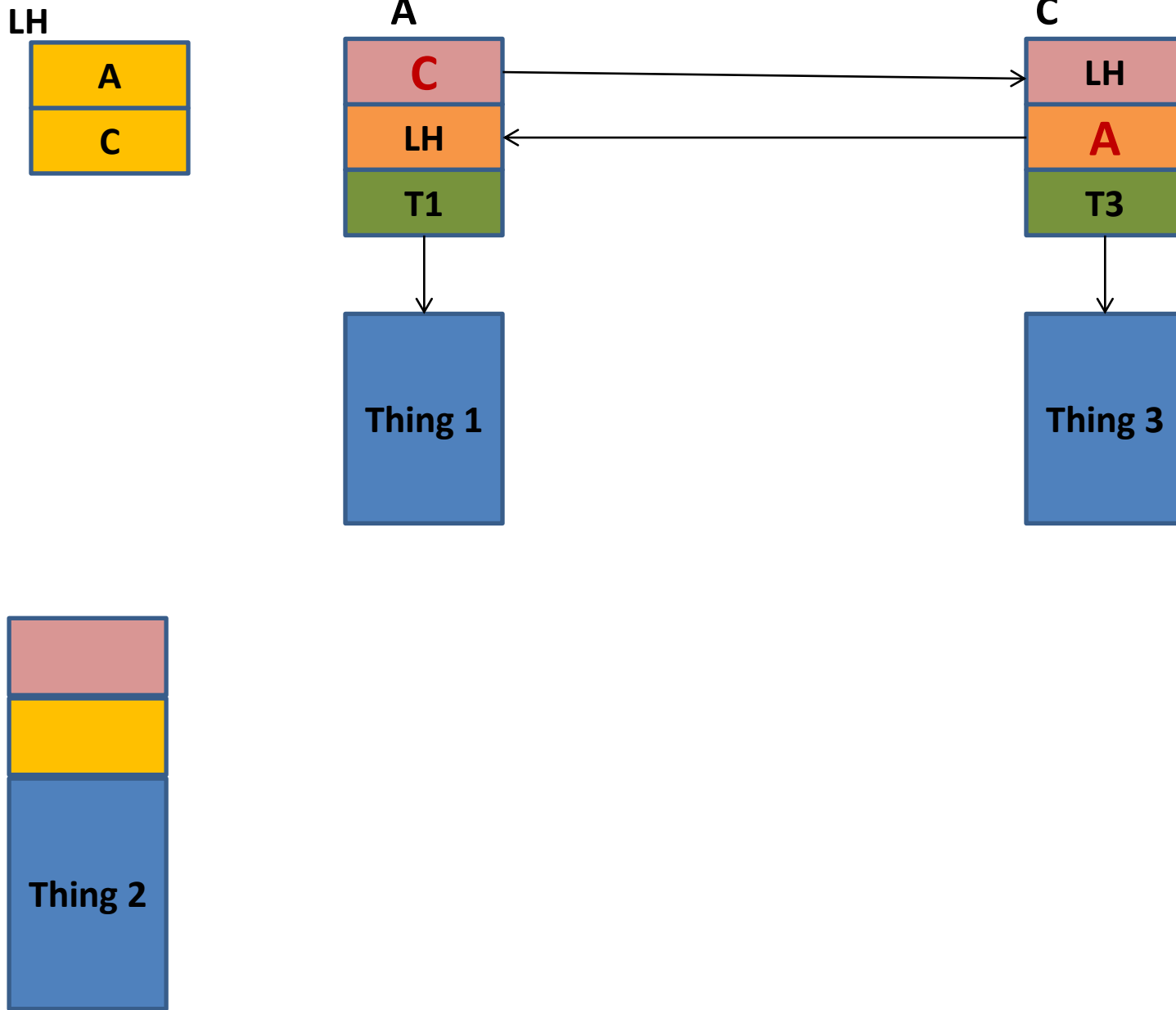
Set C's BP to A // B is off the backward chain

No matter how large the list, **this is all you have to do**

**(Extended Array: Everything past insertion must be moved up 1 slot).**



# Linked List: Element Removal



# Linked List Advantages / Disadvantages

- **Advantage**

- Lists perform generally better in inserting & extracting elements in any position within the container **to which an iterator is already positioned.**

- **Disadvantages**

- Requires 3 classes to implement (List, Node, Iterator)
- Lists lack direct access to the elements by their position, so the list must be walked to get into position.

➔ **Ex: To access the  $n$ th element in a list, the client first has to iterate from a known position (like the beginning or the end) to that position, which takes linear time.**

# Linked List (LL) Usage Guidelines

1. An LL should never be considered when the contained elements do not need to be ordered.
2. An LL might be considered if element ordering is required, access to the “sorted collection” is frequent, and the number of elements changes frequently:  $(\text{inserts} + \text{deletes}) > \# \text{ retrieves}$ .
3. An LL should seriously be considered if the elements being collected are objects and not object pointers (avoids copy constructor calls on each element move).

# Linked List vs. Ordered Array

Comparisons (N elements)	Linked List	Ordered
Element “Find”	Scan all elements until match $O(N/2)$	Binary search for specific element $O(\log N)$
New Element Addition	Find $O(N/2)$ + Add $O(1)$ = $O(N/2)$	Find $O(\log N)$ + push every element past insertion point back one = $\sim O(N/2)$
Element Deletion (once there)	Find $O(N/2)$ + Remove $O(1)$ = $O(N/2)$	Find $O(\log N)$ + push every element past remove point forward one = $\sim O(N/2)$
Element “Ordering”	Ordered at Element insertion. $O(1)$	Ordered at Element insertion. $O(1)$
Reporting	All elements are present in order $O(N)$	All elements are present in order $O(N)$

# “Ordered” Containers Comparison: Usage

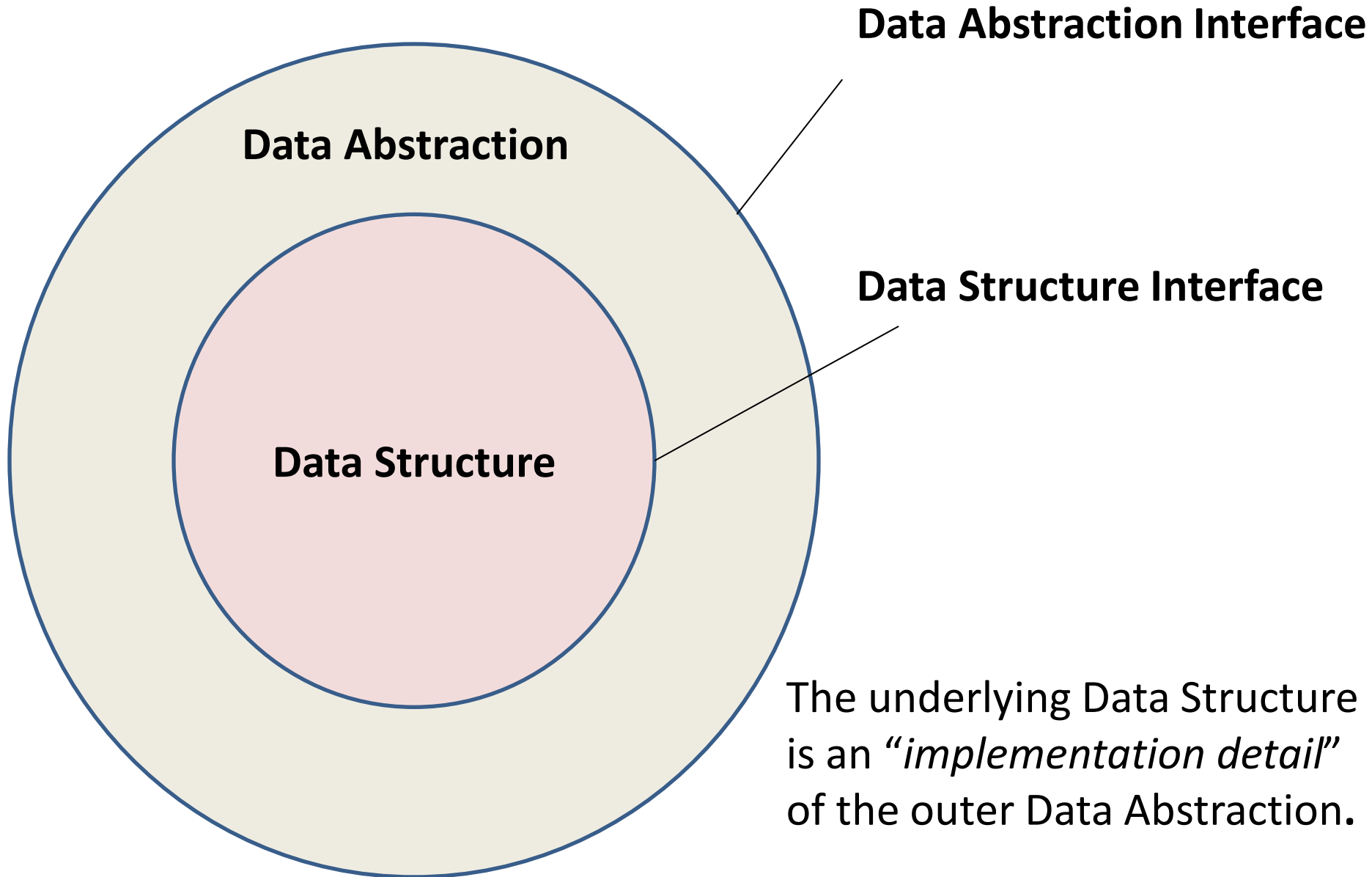
- **Number of elements in Collection  $< 100$** 
  - Preallocate Array and self-manage
- **Collecting objects rather than object PTRS**
  - Use LL to avoid object “moves” on insert/delete
- **Integer UID range  $< 10$  times number of elements**
  - Use Sparse array (ID = Index)
- **Ratio of Retrievals / Updates (Inserts & Deletes)**
  - Low ( $R \ll U$ )  $\rightarrow$  LL or OA.
  - High ( $R \gg U$ )  $\rightarrow$  OA ... or Unordered Array + Sorting
    - LL “Find” times much worse

# Data Abstractions

(creating a “false front” for a Data Structure)



# Data Abstraction “wraps” Data Structure



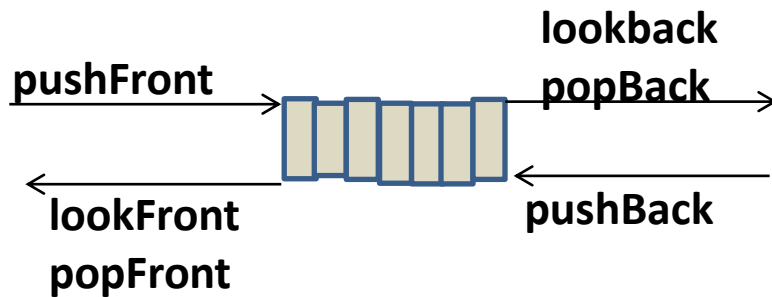
# Some Data Collection Abstractions

- **Deque (double Q)**
  - **FIFO Queue**
  - **LIFO Stack**
  - **Priority Queue**
- 
- Hash Tables
  - Trees
  - Graphs
  - Heaps



# DeQueue (DQ) Abstraction (underlying Data Structure TBD)

## Deque for “Thing”



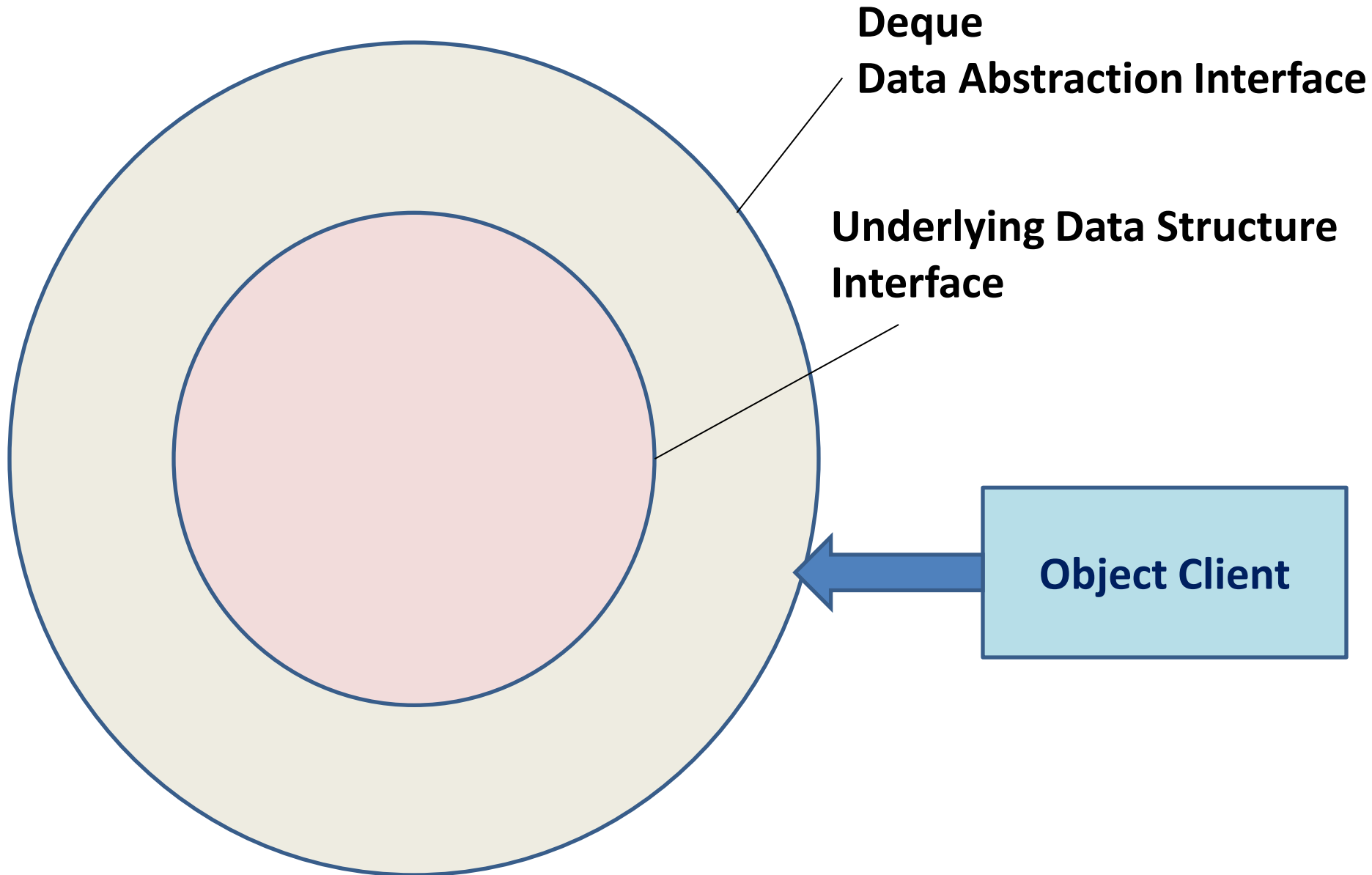
```
class DQ
{ public:
    DQ (unsigned int maxSize);
    DQ (DQ&); ~DQ();

    void pushFront (Thing &);
    Thing& lookback ();
    Thing& popBack ();

    void pushBack (Thing &);
    Thing & lookFront ();
    Thing & popFront ();

    bool isEmpty ();
    unsigned int getNumElements();
};
```

# “Data Abstraction “wraps” Data Structure



# Wrapper Design Pattern

**Object Type A** “contains” **Object Type B** (or a Ptr to one) in its private data and uses it to deliver most / all of its functionality. In this “Wrapper” case:

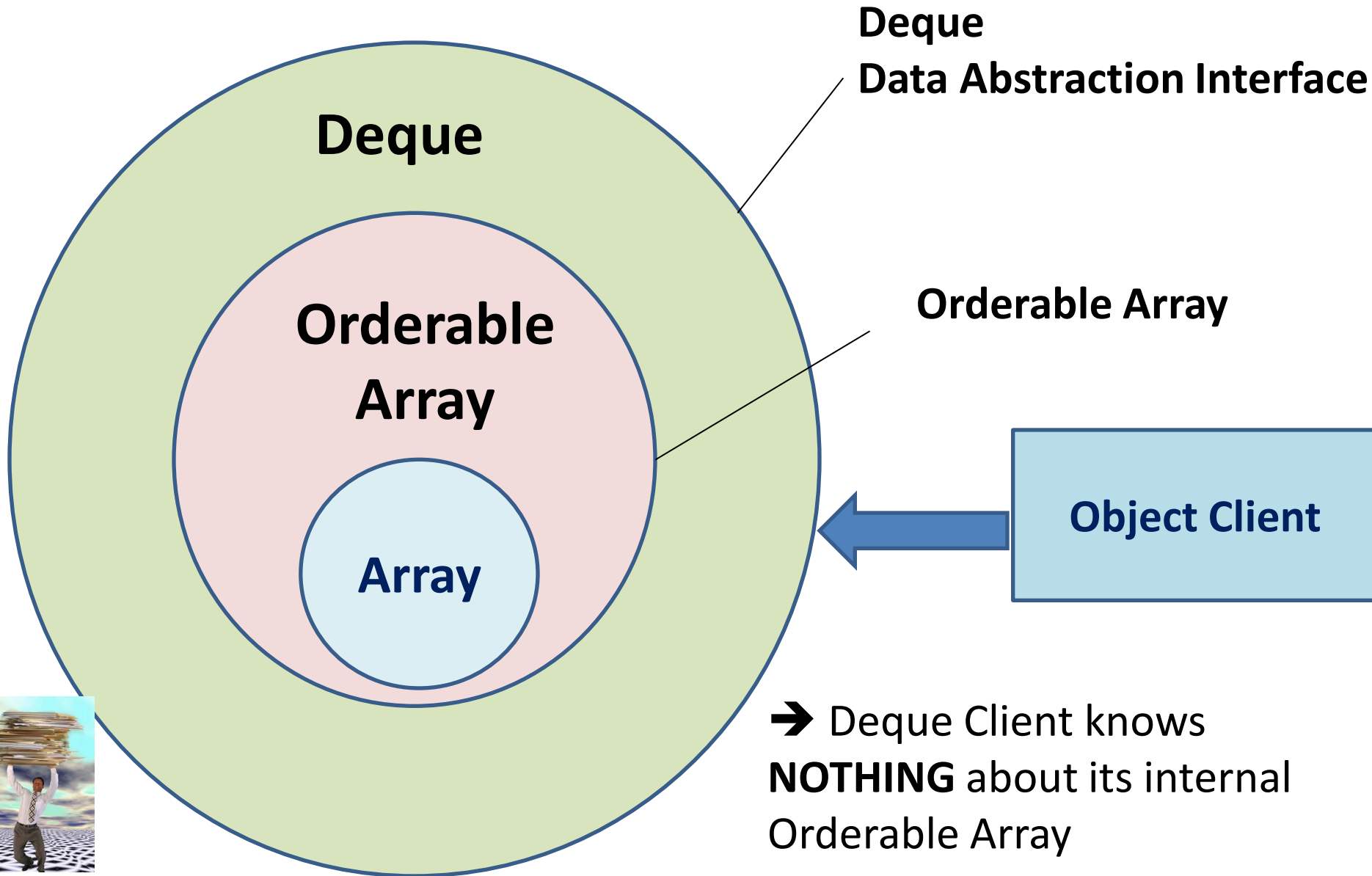
1. Does the Client see both interfaces? If not, which interface does the Client see?
2. What should the underlying Data Structure for Deque Data Abstraction be? **Why?**
3. How can the underlying Data Structure be optimally used in supporting the Deque Data Abstraction?



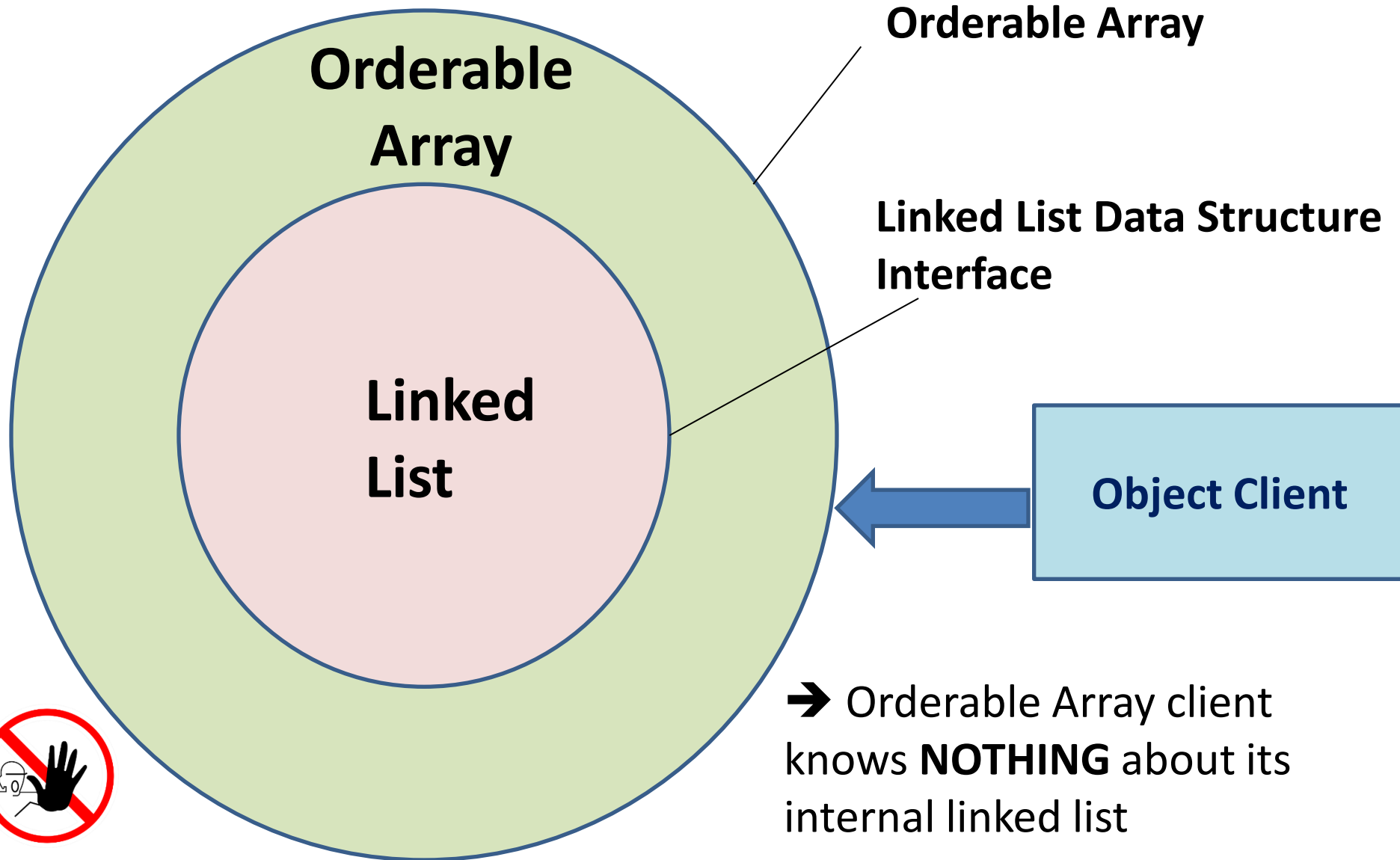
# Data Abstraction Wrapper

- Which interface does the Client see?
  - The Data Abstraction (“wraps” inner Data Structure)
- What should the underlying Data Structure for Deque be? Why? **Orderable Array**
  - Element insertion / deletion is either at front or back
  - Strength of LL is insertion / deletion at middle
- How can the underlying Data Structure be optimally used in supporting the Deque?
  - DQ Constructor creates internal OA

# Deque “wraps” OA



# Could an Orderable Array Wrap a LL?



# Could an OA Wrap a LL?

## Link List Iterator API

Thing& getNext();

Thing& getPrevious();

void remove ();

void set (Thing&);

void add (Thing&);

## Orderable Array API

Thing& get(int index);

int remove(int index);

int set (Thing&, int index);

int insert (Thing&, int index);

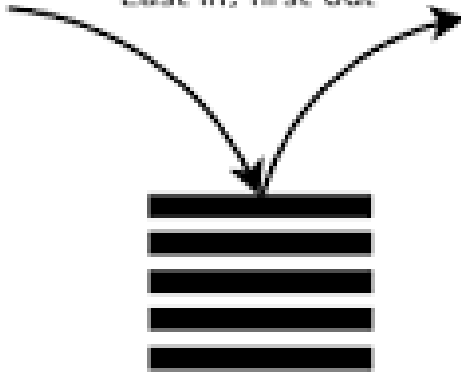
int append (Thing&);

**➔ Yes! Process the index argument by traversing that many nodes in the LL. But why??**

# Stacks & Queues

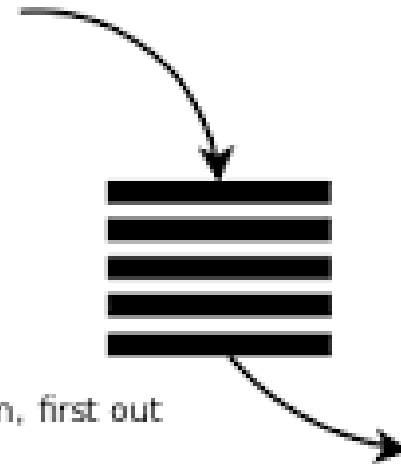
**Stack:**

Last in, first out



**Queue:**

First in, first out



**Real World Examples??**

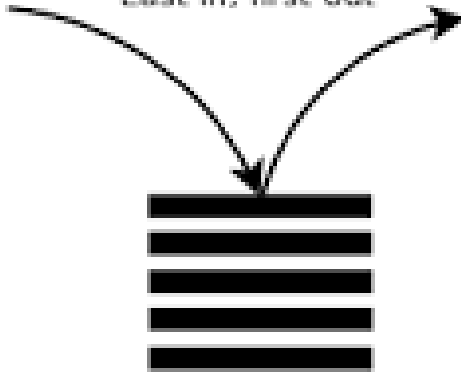




# Stacks & Queues

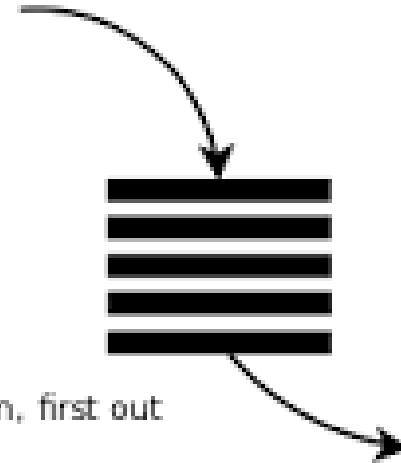
**Stack:**

Last in, first out



**Queue:**

First in, first out



**Employed Teachers  
when layoffs happen.**

**Employed Teachers when  
assigning available Sections.**

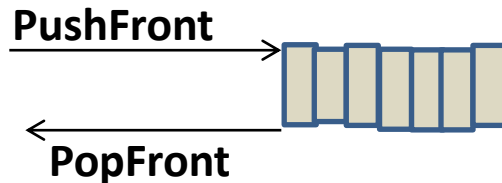
# Creating FIFO Queues and LIFO Stacks

## FIFO Queue



```
class FQ
{ public:
    void pushFront (Thing &);
    Thing & popBack ();
};
```

## LIFO Stack



```
class LS
{ public:
    void pushFront (Thing &);
    Thing & popFront ();
};
```

# Relationship: FQ with DQ

## 1. Wrapper: DQ “wraps” 2 FQ’s

- One supports pushBack(), popFront()
- One supports pushFront(), popBack()

## 2. Inheritance (DQ is a child of FQ)?

- FQ has 2 of the 4 public functions of DQ
- DQ inherits the FQ interf / impl ops, adds the other 2
  - Why couldn’t FQ be a child of DQ?

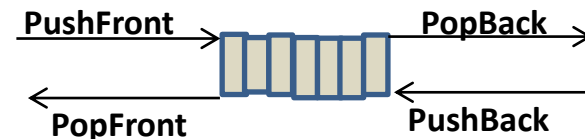
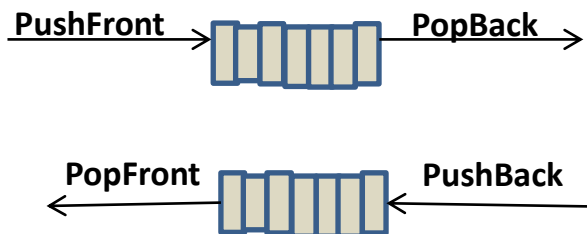
## 3. Wrapper: FQ “wraps” DQ

- pushFront() & popBack() invoke DQ equivalents
- Remaining functions unsupported in public interface



# 1. DQ wraps 2 FQs: Fails

2 Single FIFO Queues **Are NOT equivalent to** 1 Deque



**\*\* Problem:** A Thing inserted via **PushFront (or PushBack)** can be popped off **EITHER** the front **or** the back. If DQ was a collection of 2 FQ's this would not be possible!!

➔ DQ cannot wrap 2 separate FQ's because even taken together, they cannot implement the public function requirements of a DQ.

## 2. DQ inherits from FQ: Fails



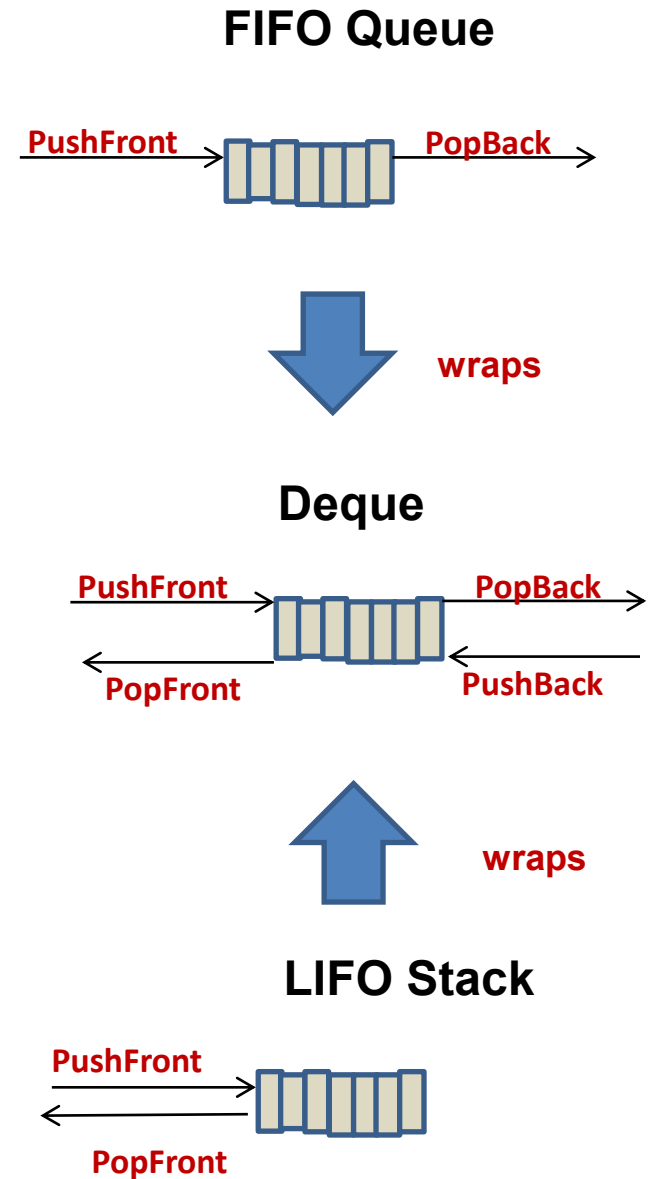
**\*\* Problem:** How could the DQ implementations of PushBack and PopFront learn where the array used internally by FQ methods PushFront and PopBack was, how many entries were in it, what was the current front and back ... ??

➔ As a child of a FIFO Queue, a DQ will inherit the implementation of 2 of the 4 functions it needs, but it cannot use them to add the other 2 functions it must support!

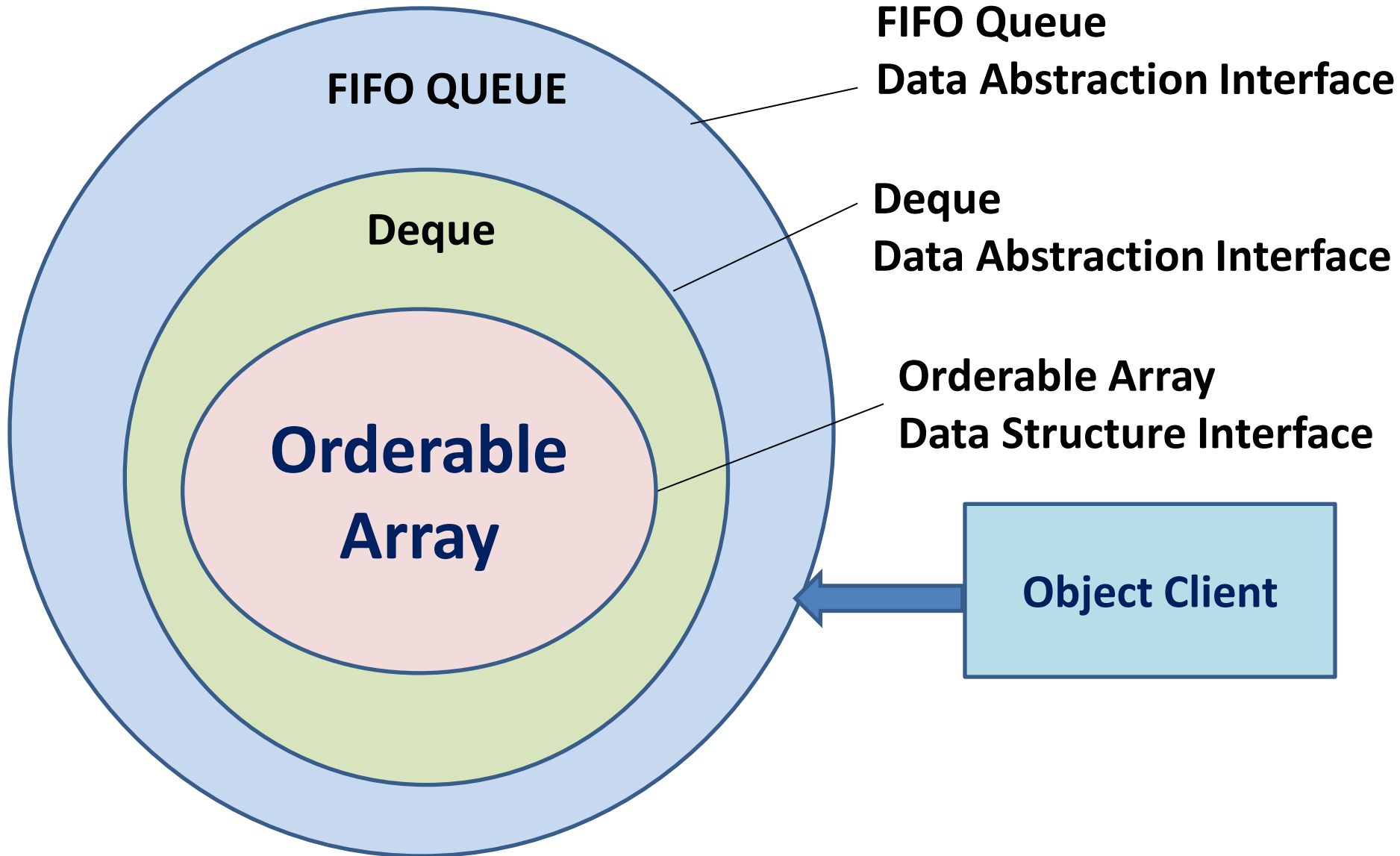
### 3. FQ (and Lstack) each wrap DQ: Success!

Once a Deque is implemented, it can be wrapped by FIFO Queue (and LIFO Stack) to do all their required functions!

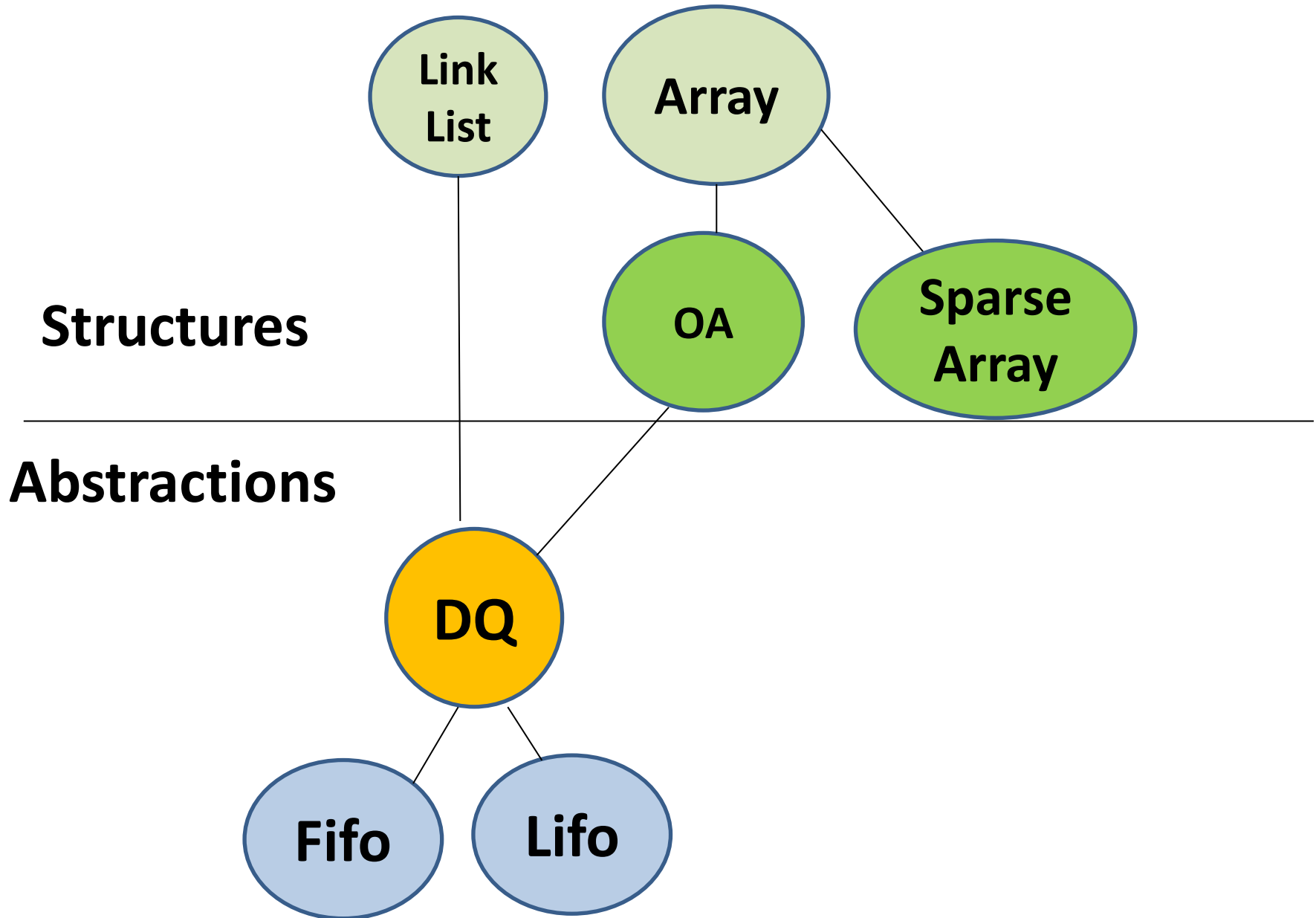
```
class Fifo {  
    private  
        DQ dq; // Created at construction  
    public:  
        void pushFront (Thing& x)  
        { return (dq.pushFront (x)); }  
  
        Thing & popBack()  
        { return (dq.popBack()); }  
};
```



# FQ / LS “wrap” DQ



# Relationship of Collections





# Questions?

