Leonard's template

```
Leonard's template
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单链表

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```

```
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```

基础算法

快速排序

```
void qsort(int 1, int r) {
   int mid = a[1 + r >> 1], i = 1, j = r;
   while(i <= j) {
      while(a[i] < mid) i++;
      while(a[j] > mid) j--;
      if(i <= j) {
         swap(a[i], a[j]);
         i++;
         j--;
      }
   }
   if(1 < j) qsort(1, j);
   if(i < r) qsort(i, r);
}</pre>
```

归并排序

```
void merge_sort(int 1, int r) {
    int mid = 1 + r >> 1;
    if(1 != r) merge_sort(1, mid);
    if(1 != r) merge_sort(mid + 1, r);
    int k = 1, i = 1, j = mid + 1;
    while(i <= mid && j <= r) {
        if(a[i] <= a[j]) tmp[k++] = a[i++];
        else tmp[k++] = a[j++];
    }
    while(i <= mid) tmp[k++] = a[i++];
    while(j <= r) tmp[k++] = a[j++];
    for(int i = 1, j = 1; i <= r; i++) a[i] = tmp[j++];
}</pre>
```

逆序对的数量

```
i64 merge_sort(int 1, int r) {
    i64 res = 0;
    i64 \ mid = 1 + r >> 1;
    if(1 != r) res += merge_sort(1, mid);
    if(1 != r) res += merge\_sort(mid + 1, r);
    int k = 1, i = 1, j = mid + 1;
    while(i \leftarrow mid && j \leftarrow r) {
        if(a[i] \le a[j]) tmp[k++] = a[i++];
        else {
             tmp[k++] = a[j++];
             res += mid - i + 1;
    }
    while(i \leftarrow mid) tmp[k++] = a[i++];
    while(j \ll r) tmp[k++] = a[j++];
    for(int i = 1, j = 1; i \leftarrow r; i++) a[i] = tmp[j++];
    return res;
}
```

整数二分

```
bool check(int x)
// 区间[1, r]被划分成[1, mid]和[mid + 1, r]时使用:
int bsearch_1(int 1, int r) {
   while (1 < r) {
       int mid = 1 + r >> 1;
       if (check(mid)) r = mid;
       else l = mid + 1;
   }
   return 1;
// 区间[1, r]被划分成[1, mid - 1]和[mid, r]时使用:
int bsearch_2(int 1, int r) {
   while (1 < r) {
       int mid = 1 + r + 1 >> 1;
       if (check(mid)) 1 = mid;
       else r = mid - 1;
    }
    return 1;
```

浮点数二分

三分

三分法可以用来查找凸函数的最大(小)值。

```
lmid = left + (right - left >> 1);
rmid = lmid + (right - lmid >> 1); // 对右侧区间取半
if (cal(lmid) > cal(rmid)) right = rmid;
else left = lmid;
```

高精度加法

```
vector<int> add(vector<int> &A, vector<int> &B) {
    if (A.size() < B.size()) return add(B, A);</pre>
   vector<int> C;
   int t = 0;
    for (int i = 0; i < A.size(); i++) {
        t += A[i];
        if (i < B.size()) t += B[i];
        C.push_back(t % 10);
        t /= 10;
    }
    if (t) C.push_back(t);
    return C;
}
int main() {
    string a, b;
    vector<int> A, B;
    cin >> a >> b;
    for (int i = a.size() - 1; i \ge 0; i--) A.push_back(a[i] - '0');
    for (int i = b.size() - 1; i \ge 0; i--) B.push_back(b[i] - '0');
    auto C = add(A, B);
    for (int i = C.size() - 1; i >= 0; i--) cout << C[i];
    return 0;
}
```

高精度减法

```
bool cmp(vector<int> &A, vector<int> &B) {
   if (A.size() != B.size()) return A.size() > B.size();
    for (int i = A.size() - 1; i >= 0; i--)
        if (A[i] != B[i])
            return A[i] > B[i];
    return true;
}
vector<int> sub(vector<int> &A, vector<int> &B) {
    vector<int> C;
    for (int i = 0, t = 0; i < A.size(); i++) {
        t = A[i] - t;
        if (i < B.size()) t -= B[i];</pre>
        C.push_back((t + 10) % 10);
        if (t < 0) t = 1;
        else t = 0;
    }
    while (C.size() > 1 \& C.back() == 0) C.pop_back();
}
int main() {
    string a, b;
   vector<int> A, B;
    cin >> a >> b;
    for (int i = a.size() - 1; i >= 0; i--) A.push_back(a[i] - '0');
    for (int i = b.size() - 1; i \ge 0; i--) B.push_back(b[i] - '0');
    vector<int> C;
    if (cmp(A, B)) C = sub(A, B);
    else C = sub(B, A), cout << '-';
    for (int i = C.size() - 1; i >= 0; i--) cout << C[i];
    return 0;
}
```

高精度乘法

```
vector<int> mul(vector<int> &A, int b) {
    vector<int> C;
   int t = 0;
    for (int i = 0; i < A.size() || t; i++) {
        if (i < A.size()) t += A[i] * b;
        C.push_back(t % 10);
       t /= 10;
    while (C.size() > 1 \& C.back() == 0) C.pop_back();
    return C;
}
int main() {
   string a;
    int b;
    cin >> a >> b;
   vector<int> A;
    for (int i = a.size() - 1; i \ge 0; i--) A.push_back(a[i] - '0');
    auto C = mul(A, b);
    for (int i = C.size() - 1; i >= 0; i--) printf("%d", C[i]);
    return 0;
```

高精度除法

```
vector<int> div(vector<int> &A, int b, int &r) {
    vector<int> C;
    r = 0;
    for (int i = A.size() - 1; i >= 0; i--) {
        r = r * 10 + A[i];
        C.push_back(r / b);
        r %= b;
    }
    reverse(C.begin(), C.end());
    while (C.size() > 1 \& C.back() == 0) C.pop_back();
    return C;
int main() {
   string a;
    vector<int> A;
   int B;
    cin >> a >> B;
   for (int i = a.size() - 1; i >= 0; i--) A.push_back(a[i] - '0');
   int r;
    auto C = div(A, B, r);
    for (int i = C.size() - 1; i >= 0; i--) cout << C[i];
    cout << '\n' << r; //余数
    return 0;
}
```

一维前缀和

```
void init() {
    for (int i = 1; i <= n; i ++ ) scanf("%d", &a[i]);
    for (int i = 1; i <= n; i ++ ) s[i] = s[i - 1] + a[i];
}
int rangeSum(int l, int r) {
    return s[r] - s[l - 1];
}</pre>
```

二维前缀和

```
void init() {
    for (int i = 1; i <= n; i++)
        for (int j = 1; j <= m; j++)
            scanf("%d", &s[i][j]);
    for (int i = 1; i <= n; i++)
        for (int j = 1; j <= m; j++)
            s[i][j] += s[i - 1][j] + s[i][j - 1] - s[i - 1][j - 1];
}
int rangeSum(int l1, int r1, int l2, int r2) {
    return s[l2][r2] - s[l1 - 1][r2] - s[l2][r1 - 1] + s[l1 - 1][r1 - 1];
}</pre>
```

一维差分

```
void insert(int 1, int r, int c) {
    s[1] += c;
    s[r + 1] -= c;
}
void init() {
    for (int i = 1; i <= n; i++) s[i] += s[i - 1];
}
int rangeSum(int 1, int r) {
    return s[r] - s[1 - 1];
}</pre>
```

二维差分

```
void insert(int | 1, int r1, int | 12, int r2, int c) {
    s[11][r1] += c;
    s[12 + 1][r1] -= c;
    s[11][r2 + 1] -= c;
    s[12 + 1][r2 + 1] += c;
}

void init() {
    for (int i = 1; i <= n; i++)
        for (int j = 1; j <= m; j++)
            s[i][j] += s[i - 1][j] + s[i][j - 1] - s[i - 1][j - 1];
}
int rangeSum(int | 11, int r1, int | 12, int r2) {
    return s[12][r2] - s[11 - 1][r2] - s[12][r1 - 1] + s[11 - 1][r1 - 1];
}</pre>
```

搜索

全排列

```
#include <bits/stdc++.h>
using namespace std;
const int N = 10;
int n;
int path[N];
void dfs(int u, int state) {
    if (u == n) {
        for (int i = 0; i < n; i++) printf("%d ", path[i]);
        puts("");
        return;
    for (int i = 0; i < n; i++)
        if (!(state >> i & 1)) {
            path[u] = i + 1;
            dfs(u + 1, state + (1 << i));
        }
}
int main() {
    scanf("%d", &n);
    dfs(0, 0);
```

```
return 0;
}
```

n-皇后问题

```
#include <bits/stdc++.h>
using namespace std;
const int N = 20;
int n;
char g[N][N];
bool col[N], dg[N], udg[N];
void dfs(int u) {
   if (u == n) {
        for (int i = 0; i < n; i++) puts(g[i]);
        puts("");
        return;
    }
    for (int i = 0; i < n; i++)
        if (!col[i] \&\& !dg[u + i] \&\& !udg[n - u + i]) {
            g[u][i] = 'Q';
            col[i] = dg[u + i] = udg[n - u + i] = true;
            dfs(u + 1);
            col[i] = dg[u + i] = udg[n - u + i] = false;
            g[u][i] = '.';
        }
int main() {
   cin >> n;
    for (int i = 0; i < n; i++)
        for (int j = 0; j < n; j++)
            g[i][j] = '.';
   dfs(0);
    return 0;
}
```

图论

邻接表

```
int idx = 0;
vector<int> h(n + 1, -1), e(m * 2), ne(m * 2);
auto add = [&](int a, int b) {
    e[idx] = b, ne[idx] = h[a], h[a] = idx++;
};
```

拓扑排序

```
#include <bits/stdc++.h>
using namespace std;
const int N = 1e5 + 10, M = 1e5 + 10;
int n, m, e[M], ne[M], h[N], top[N], d[N], idx = 1, cnt;
void add(int a, int b){
    e[idx] = b, ne[idx] = h[a], h[a] = idx++;
}
bool topsort(){
    queue <int> q;
    for(int i = 1; i \le n; i++){
        if(d[i] == 0){
            q.push(i);
            top[++cnt] = i;
        }
    }
    while(q.size()){
        int t = q.front();
        q.pop();
        for(int i = h[t]; i != -1; i = ne[i]){
            int j = e[i];
            d[j]--;
            if(d[j] == 0){
                q.push(j);
                top[++cnt] = j;
            }
        }
    }
    if(cnt == n) return 1;
    else return 0;
int main(){
    scanf("%d%d", &n, &m);
    memset(h, -1, sizeof h);
    for(int i = 1; i \le m; i++){
        int a, b;
        scanf("%d%d", &a, &b);
        add(a, b);
        d[b]++;
    if(!topsort()) puts("-1");
    else{
        for(int i = 1; i \le n; i++){
            printf("%d ", top[i]);
        }
   return 0;
}
```

强连通分量

```
struct SCC{
  int n;
  vector<vector<int>> g;
  vector<int> stk;
```

```
vector<int> dfn, low, bel;
    int cur, cnt;
    SCC(int n, vector<vector<int>> &g){
        this->n = n;
        this->g = g;
        dfn.assign(n + 1, -1);
        low.resize(n + 1);
        bel.assign(n + 1, -1);
        stk.clear();
        cur = cnt = 0;
    }
    void dfs(int x) {
        dfn[x] = low[x] = cur++;
        stk.push_back(x);
        for (auto y : g[x]) {
            if (dfn[y] == -1) {
                dfs(y);
                low[x] = min(low[x], low[y]);
            } else if (bel[y] == -1) {
                low[x] = min(low[x], dfn[y]);
            }
        }
        if (dfn[x] == low[x]) {
            int y;
            cnt++;
            do {
                y = stk.back();
                bel[y] = cnt;
                stk.pop_back();
            } while (y != x);
       }
    }
    vector<int> work() {
        for (int i = 1; i \le n; i++) {
            if (dfn[i] == -1) {
                dfs(i);
        }
       return bel;
   }
    vector<vector<int>> get() {
        vector<vector<int>> v(cnt + 1);
        for(int i = 1; i <= n; i++) {
           v[bel[i]].push_back(i);
        }
        return v;
   }
};
    vector<vector<int>> g(n + 1);
    while(m--){
        int u, v;
```

```
cin >> u >> v;
    g[u].push_back(v);
    g[v].push_back(u);
}
SCC scc(n, g);
auto bel = scc.work();
auto v = scc.get();
int cnt = scc.cnt;
vector<> dp(cnt + 1);//int or 11
for(int i = 1; i <= cnt; i++){
    for(auto x : v[i]){
        for(auto j : g[x]){
            if (bel[j] != i){
            }
        }
    }
}
```

双连通分量

边双

```
void solve() {
   int n, m;
   cin >> n >> m;
   int idx = 0;
   vector<int> h(n + 1, -1), e(m * 2), ne(m * 2);
    auto add = [\&] (int a, int b) {
        e[idx] = b, ne[idx] = h[a], h[a] = idx++;
   };
   int timestamp = 0, top = 0, dcc_cnt = 0;
   vector<int> dfn(n + 1), low(n + 1);
   vector<int> stk(n + 1);
   vector<int> id(n + 1), is_bridge(m * 2);
    function<void(int, int)> dcc = [&](int u, int from) {
        dfn[u] = low[u] = ++timestamp;
        stk[++top] = u;
        for (int i = h[u]; ~i; i = ne[i]) {
            int j = e[i];
            if (!dfn[j]) {
                dcc(j, i);
                low[u] = min(low[u], low[j]);
                if (dfn[u] < low[j])</pre>
                    is_bridge[i] = is_bridge[i ^ 1] = true;
            else if (i != (from ^ 1))
                low[u] = min(low[u], dfn[j]);
        if (dfn[u] == low[u]) {
            ++dcc_cnt;
            int y;
            do {
```

```
y = stk[top--];
                id[y] = dcc_cnt;
            } while (y != u);
        }
   };
   while (m--) {
       int a, b;
        cin >> a >> b;
        add(a, b), add(b, a);
   dcc(1, -1);//因为边从0开始所以不能是0
   vector<vector<int>>> g(dcc_cnt + 1);
   for (int i = 1; i <= n; i++) {
        for (int j = h[i]; j != -1; j = ne[j]) {
            int v = e[j];
           if (is_bridge[j]) {
                g[id[i]].push_back(id[v]);
           }
       }
   }
}
```

最短路算法

朴素dijkstra算法

时间复杂度 $O(n^2 + m)$, n 表示点数, m 表示边数

```
int n, m, e[N][N], d[N];
bool st[N];
void dijkstra() {
    d[1] = 0;
    for(int i = 1; i \le n; i++) {
        int t = 0;
        for(int j = 1; j <= n; j++) {
            if(!st[j] \& (!t || d[j] < d[t])) t = j;
        st[t] = 1;
        for(int j = 2; j <= n; j++) {
            d[j] = min(d[j], d[t] + e[t][j]);
        }
    }
}
int main() {
    cin >> n >> m;
    memset(d, 0x3f3f3f3f, sizeof d);
    memset(e, 0x3f3f3f3f, sizeof e);
    for(int i = 1; i \le n; i++) e[i][i] = 0;
    while(m--) {
        int u, v, w;
        cin >> u >> v >> w;
        e[u][v] = min(e[u][v], w);
    dijkstra();
    if(d[n] == 0x3f3f3f3f) d[n] = -1;
```

```
printf("%d", d[n]);
return 0;
}
```

堆优化版dijkstra

时间复杂度O(mlogn), n 表示点数, m 表示边数。

```
auto dij = [\&](auto \&g, int s, int n) {
    vector<i64> d(n + 1, -1);
    vector<int> st(n + 1);
    priority_queue<array<i64, 2>, vector<array<i64, 2>>, greater<array<i64, 2>>>
pq;
    pq.push({0, s});
    while(!pq.empty()) {
        auto [dist, u] = pq.top();
        pq.pop();
        if(st[u]) continue;
        st[u] = 1;
        for(auto [v, w] : g[u]) {
            if(d[v] == -1 \mid \mid d[v] > dist + w) {
                pq.push({d[v] = dist + w, v});
            }
        }
    }
    return d;
};
```

bellman-ford

有边数限制的最短路(可能存在负权回路),时间复杂度O(nm)

```
const int N = 510, M = 10010;
int n, m, k, d[N], last[N];
struct Edge{
   int a, b, w;
}edge[M];
void bellman_ford(){
    memset(d, 0x3f, sizeof d);
    d[1] = 0;
    for(int i = 1; i \le k; i++){
        memcpy(last, d, sizeof d);
        for(int j = 1; j \le m; j++){
            auto e = edge[j];
            d[e.b] = min(d[e.b], last[e.a] + e.w);
    }
}
int main(){
    scanf("%d%d%d", &n, &m, &k);
    for(int i = 1; i \le m; i++){
        scanf("%d%d%d", &edge[i].a, &edge[i].b, &edge[i].w);
    }
    bellman_ford();
    if(d[n] > 0x3f3f3f3f / 2) puts("impossible");
    else printf("%d", d[n]);
```

```
return 0;
}
```

SPFA

```
const int N = 1e5 + 10, M = 1e5 + 10;
int n, m, e[M], ne[M], w[M], h[N], d[N], idx = 1;
bool st[N];
void add(int x, int y, int z){
    e[idx] = y, ne[idx] = h[x], w[idx] = z, h[x] = idx++;
}
int spfa(){
    memset(d, 0x3f, sizeof d);
    d[1] = 0;
    queue <int> q;
    q.push(1);
    st[1] = 1;
    while(q.size()){
        int t = q.front();
        st[t] = 0;
        q.pop();
        for(int i = h[t]; i != -1; i = ne[i]){
            int j = e[i];
            if(d[j] > d[t] + w[i]){
                d[j] = d[t] + w[i];
                if(!st[j]){
                    st[j] = 1;
                    q.push(j);
                }
            }
    }
    return d[n];
}
int main(){
    scanf("%d%d", &n, &m);
    memset(h, -1, sizeof h);
    for(int i = 1; i \le m; i++){
        int x, y, z;
        scanf("%d%d%d", &x, &y, &z);
        add(x, y, z);
    }
    int t = spfa();
    if(t == 0x3f3f3f3f) puts("impossible");
    else printf("%d", t);
    return 0;
}
```

SPFA求负环

```
const int N = 2010, M = 1e5 + 10;
int n, m, e[M], ne[M], w[M], h[N], cnt[N], d[N], idx = 1;
bool st[N];
void add(int x, int y, int z){
    e[idx] = y, ne[idx] = h[x], w[idx] = z, h[x] = idx++;
}
```

```
bool spfa(){
    queue <int> q;
    for(int i = 1; i \le n; i++){
        q.push(i);
        st[i] = 1;
    }
    while(q.size()){
        int t = q.front();
        q.pop();
        st[t] = 0;
        for(int i = h[t]; i != -1; i = ne[i]){
            int j = e[i];
            if(d[t] + w[i] < d[j]){
                d[j] = d[t] + w[i];
                cnt[j] = cnt[t] + 1;
                if(cnt[j] == n) return 1;
                if(!st[j]){
                    q.push(j);
                    st[j] = 1;
                }
           }
        }
    }
    return 0;
int main(){
    memset(h, -1, sizeof h);
    scanf("%d%d", &n, &m);
    for(int i = 1; i \le m; i++){
        int x, y, z;
        scanf("%d%d%d", &x, &y, &z);
        add(x, y, z);
    if(spfa()) puts("Yes");
    else puts("No");
    return 0;
}
```

Floyd求最短路

时间复杂度 $O(n^3)$

```
const int N = 210;
int n, m, k, d[N][N];
void floyd(){
    for(int k = 1; k <= n; k++){
        for(int j = 1; j <= n; j++){
            d[i][j] = min(d[i][j], d[i][k] + d[k][j]);
        }
    }
}
int main(){
    scanf("%d%d%d", &n, &m, &k);
    memset(d, 0x3f, sizeof d);
    for(int i = 1; i <= m; i++){</pre>
```

```
int x, y, z;
    scanf("%d%d%d", &x, &y, &z);
    d[x][y] = min(d[x][y], z);
}
for(int i = 1; i <= n; i++) d[i][i] = 0;
floyd();
while(k--){
    int x, y;
    scanf("%d%d", &x, &y);
    if(d[x][y] > 0x3f3f3f3f / 2) puts("impossible");
    else printf("%d\n", d[x][y]);
}
return 0;
}
```

最小生成树

Kruskal

```
int n, m, p[N];
struct edge {
   int u, v, w;
}e[N * 2];
bool cmp(edge a, edge b) {
    return a.w < b.w;</pre>
}
int find(int x) {
   if(p[x] != x) p[x] = find(p[x]);
    return p[x];
}
int kruskal() {
    int cnt = 0, ans = 0;
    sort(e + 1, e + 1 + m, cmp);
    for(int i = 1; i <= m; i++) {
        int fa = find(e[i].u), fb = find(e[i].v);
        if(fa != fb) {
            p[fa] = fb;
            cnt++;
            ans += e[i].w;
        }
    }
    if(cnt == n - 1) return ans;
    else return -1;
}
int main() {
    cin >> n >> m;
    for(int i = 1; i \le n; i++) p[i] = i;
    for(int i = 1; i \le m; i++) {
        int u, v, w;
        cin >> u >> v >> w;
        e[i] = \{u, v, w\};
    cout << kruskal();</pre>
    return 0;
}
```

Prim

```
#include <bits/stdc++.h>
using namespace std;
const int N = 510, INF = 0x3f3f3f3f3f;
int n, m, g[N][N], d[N];
bool st[N];
int prim(){
   memset(d, 0x3f, sizeof d);
    int res = 0;
    for(int i = 0; i < n; i++){
        int t = -1;
        for(int j = 1; j <= n; j++){
            if(!st[j] \& (t == -1 || d[t] > d[j])) t = j;
        if(i && d[t] == INF) return INF;
        if(i) res += d[t];
        st[t] = 1;
        for(int j = 1; j \le n; j++){
            d[j] = min(d[j], g[t][j]);
        }
    }
    return res;
}
int main(){
    memset(g, 0x3f, sizeof g);
    scanf("%d%d", &n, &m);
    for(int i = 1; i \le m; i++){
        int x, y, z;
        scanf("%d%d%d", &x, &y, &z);
        g[x][y] = g[y][x] = min(g[x][y], z);
   int t = prim();
    if(t == INF) puts("impossible");
    else printf("%d", t);
    return 0;
}
```

dsu on tree

```
sz[u] += sz[v];
        if(!heavy[u] || sz[v] > sz[heavy[u]]) heavy[u] = v;
    r[u] = tim;
};
function \langle void(int, int, bool) \rangle dfs_solve = [&](int u, int fa, bool keep) {
    for(auto v : g[u]) {
        if(v == fa \mid \mid v == heavy[u])
             continue;
        dfs_solve(v, u, 0);
    }
    if(heavy[u]) dfs_solve(heavy[u], u, 1);
    auto add = [\&](int x) {
    };
    auto query = [\&](int x) {
    };
    auto del = [\&](int x) {
    };
    for(auto v : g[u]) {
        if(v == fa \mid\mid v == heavy[u])
             continue;
        for(int i = 1[v]; i \leftarrow r[v]; i++)
             query(id[i]);
         for(int i = 1[v]; i \leftarrow r[v]; i++)
             add(id[i]);
    }
    query(u);
    add(u);
    if(!keep) {
        for(int i = l[u]; i \leftarrow r[u]; i++)
             del(id[i]);
    }
};
dfs_init(1, 0);
dfs_solve(1, 0, 0);
```

二分图

判定二分图

```
#include <bits/stdc++.h>
using namespace std;
const int N = 1e5 + 10, M = 2e5 + 10;
int n, m, idx = 1, e[M], h[N], ne[M], color[N];
void add(int u, int v){
    e[idx] = v, ne[idx] = h[u], h[u] = idx++;
}
bool dfs(int u, int c){
    color[u] = c;
    for(int i = h[u]; i != -1; i = ne[i]){
        int j = e[i];
        if(!color[j]){
            if(!dfs(j, 3 - c)) return 0;
        }
}
```

```
else if(color[j] == c) return 0;
    return 1;
}
int main(){
    scanf("%d%d", &n, &m);
    memset(h, -1, sizeof h);
    for(int i = 1; i \le m; i++){
        int u, v;
        scanf("%d%d", &u, &v);
        add(u, v);
        add(v, u);
    }
    bool flag = 1;
    for(int i = 1; i \le n; i++){
        if(!color[i]){
            if(!dfs(i, 1)){
                flag = 0;
                break;
            }
        }
    }
    if(flag) puts("Yes");
    else puts("No");
    return 0;
}
```

增广路算法

二分图的匹配:给定一个二分图 G,在 G 的一个子图 M 中, M 的边集 { E } 中的任意两条边都不依附于同一个顶点,则称 M 是一个匹配。

二分图的最大匹配: 所有匹配中包含边数最多的一组匹配被称为二分图的最大匹配, 其边数即为最大匹配数。

时间复杂度: O(nm)

```
int n1, n2, m;
vector<int> e[N];
int match[N];
bool st[N];
bool find(int u) {
    for(int v : e[u]) {
        if(!st[v]) {
            st[v] = true;
            if(!match[v] || find(match[v])) {
                match[v] = u;
                return true;
            }
        }
    }
    return false;
}
void solve() {
    cin >> n1 >> n2 >> m;
```

```
for(int i = 1; i <= m; i++) {
    int u, v;
    cin >> u >> v;
    e[u].push_back(v);
}
int ans = 0;
for(int i = 1; i <= n1; i++) {
    memset(st, false, sizeof st);
    if(find(i)) ans++;
}
cout << ans;
}</pre>
```

网络流

最大流

dinic

时间复杂度 $O(n^2m)$ 二分图时间复杂度 $O(\sqrt{n}m)$

可求最大独立集。**独立集**是指在图中选取一个顶点子集,使得该子集中的任意两个顶点之间都不相邻(即没有边直接连接)。二分图的最大独立集为点数-最大匹配数 例题:CF2026E

```
template<class T>
struct Flow {
   const int n;
   struct Edge {
        int to;
       T cap;
        Edge(int to, T cap) : to(to), cap(cap) {}
   };
   vector<Edge> e;
   vector<vector<int>> g;
   vector<int> cur, h;
   Flow(int n) : n(n), g(n) {}
   bool bfs(int s, int t) {
        h.assign(n, -1);
        queue<int> que;
        h[s] = 0;
        que.push(s);
        while (!que.empty()) {
            const int u = que.front();
            que.pop();
            for (int i : g[u]) {
                auto [v, c] = e[i];
                if (c > 0 \& h[v] == -1) {
                    h[v] = h[u] + 1;
                    if (v == t) {
                        return true;
                    que.push(v);
                }
            }
        }
        return false;
```

```
T dfs(int u, int t, T f) {
        if (u == t) {
            return f;
        }
        auto r = f;
        for (int &i = cur[u]; i < int(g[u].size()); ++i) {</pre>
            const int j = g[u][i];
            auto [v, c] = e[j];
            if (c > 0 \& h[v] == h[u] + 1) {
                auto a = dfs(v, t, min(r, c));
                e[j].cap -= a;
                e[j \land 1].cap += a;
                r -= a;
                if (r == 0) {
                    return f;
                }
            }
        return f - r;
    void addEdge(int u, int v, T c) {
        g[u].push_back(e.size());
        e.emplace_back(v, c);
        g[v].push_back(e.size());
        e.emplace_back(u, 0);
    T maxFlow(int s, int t) {
        T ans = 0;
        while (bfs(s, t)) {
            cur.assign(n, 0);
            ans += dfs(s, t, numeric_limits<T>::max());
        }
        return ans;
    }
};
```

费用流

spfa (不一定能用)

```
const int N = 5010, M = 100010, INF = 1e8;
int n, m, S, T;
int h[N], e[M], f[M], w[M], ne[M], idx;
int q[N], d[N], pre[N], incf[N];
bool st[N];
void add(int a, int b, int c, int d) {
    e[idx] = b, f[idx] = c, w[idx] = d, ne[idx] = h[a], h[a] = idx++;
    e[idx] = a, f[idx] = 0, w[idx] = -d, ne[idx] = h[b], h[b] = idx++;
}
bool spfa() {
    int hh = 0, tt = 1;
    memset(d, 0x3f, sizeof d);
    memset(incf, 0, sizeof incf);
    q[0] = S, d[S] = 0, incf[S] = INF;
    while (hh != tt) {
```

```
int t = q[hh++];
        if (hh == N) hh = 0;
        st[t] = false;
        for (int i = h[t]; ~i; i = ne[i]) {
            int ver = e[i];
            if (f[i] && d[ver] > d[t] + w[i]) {
                d[ver] = d[t] + w[i];
                pre[ver] = i;
                incf[ver] = min(f[i], incf[t]);
                if (!st[ver]) {
                    q[tt++] = ver;
                    if (tt == N) tt = 0;
                    st[ver] = true;
                }
            }
        }
    }
   return incf[T] > 0;
}
void EK(int& flow, int& cost) {
   flow = cost = 0;
    while (spfa()) {
        int t = incf[T];
        flow += t, cost += t * d[T];
        for (int i = T; i != S; i = e[pre[i] \land 1]) {
            f[pre[i]] -= t;
            f[pre[i] ^ 1] += t;
        }
   }
}
int main() {
    scanf("%d%d%d%d", &n, &m, &S, &T);
    memset(h, -1, sizeof h);
    while (m--) {
        int a, b, c, d;
        scanf("%d%d%d%d", &a, &b, &c, &d);
        add(a, b, c, d);
    }
    int flow, cost;
    EK(flow, cost);
    printf("%d %d\n", flow, cost);
    return 0;
}
```

skip2004的板子

 $O(m^3)$

abc373G

```
namespace mcmf {
    using pr = std::pair<i64, int>;
    const int N = 10005, M = 1e6 + 10;
    struct edge {
        int to, nxt, v; i64 f;
    } e[M << 1];
    int h[N], num = 1;
    void link(int x, int y, int v, i64 f) {</pre>
```

```
e[++num] = \{y, h[x], v, f\}, h[x] = num;
        e[++num] = \{x, h[y], 0, -f\}, h[y] = num;
    i64 d[N], dis[N];
    int vis[N], fr[N];
    bool dijkstra(int s, int t) { // 正常题目不需要 dijk
        std::priority_queue<pr, std::vector<pr>, std::greater<pr>>> Q;
        for(int i = s; i \le t; ++i) dis[i] = d[i], d[i] = 1e18, vis[i] = fr[i] =
0; // CHECK
        for(Q.emplace(d[s] = 0, s); !Q.empty();) {
            int x = Q.top().second; Q.pop();
            if(vis[x]) continue;
            vis[x] = 1;
            for(int i = h[x];i;i = e[i].nxt) {
                const i64 v = e[i].f + dis[x] - dis[e[i].to];
                if(e[i].v \& d[e[i].to] > d[x] + v) {
                    fr[e[i].to] = i;
                    Q.emplace(d[e[i].to] = d[x] + v, e[i].to);
                }
            }
        }
        for(int i = s;i <= t;++i) d[i] += dis[i]; // CHECK
        return d[t] < 1e17;</pre>
    std::pair<i64, i64> EK(int s, int t) {
        i64 f = 0, c = 0;
        while(dijkstra(s, t)) { // 正常可以用 spfa
            i64 f1 = 1e18;
            for(int i = fr[t]; i; i = fr[e[i \land 1].to]) fl = std::min<i64>
(e[i].v, f1);
            for(int i = fr[t]; i; i = fr[e[i \land 1].to]) e[i].v -= fl, e[i \land 1].v
+= f1;
            f += f1, c += f1 * d[t];
        }
        return std::make_pair(f, c);
    }
}
void solve() {
    auto [x, y] = mcmf::EK(S, T); //y是最小费用
    for(int i = 1; i <= n; i++) {
        for(int j = mcmf::h[i]; j; j = mcmf::e[j].nxt) {
            if(mcmf::e[j].to > n \&\& mcmf::e[j].v == 0) {
                cout << mcmf::e[j].to - n << ' ';</pre>
                break;
        }
    }
    //求方案
    for(int i = 1; i \le n; i++) {
        using namespace mcmf;
        for(int j = h[i]; j; j = e[j].nxt) {
            if(e[j].to > n \& e[j].v == 0) {
                cout << e[j].to - n << ' ';</pre>
                break;
            }
        }
```

```
}
//直接using namespace mcmf 可以省很多代码 更方便但是要注意变量重名
}
```

LCA

倍增

```
struct LCA {
   int n, LOG;
    vector<vector<int>> fa;
    vector<int> d;
    LCA(int _n) : n(_n), LOG(__lg(n)), fa(n + 1, vector<int> (LOG + 1)), d(n + 1)
1){};
    void init(int s, auto &g) {
        d[s] = 1;
        queue <int> q;
        q.push(s);
        while (!q.empty()) {
            int u = q.front();
            q.pop();
            for (int v : g[u]) {
                if (!d[v]) {
                    q.push(v);
                    d[v] = d[u] + 1;
                    fa[v][0] = u;
                    for (int k = 1; k \le LOG; k++) fa[v][k] = fa[fa[v][k - 1]][k
- 1];
                }
            }
    };
    int anc(int x, int y) {
        if (d[x] > d[y]) swap(x, y);
        for (int k = LOG; k >= 0; k --) if (d[fa[y][k]] >= d[x]) y = fa[y][k];
        if (x == y) return x;
        for (int k = LOG; k >= 0; k --) if (fa[x][k] != fa[y][k]) x = fa[x][k], y
= fa[y][k];
        return fa[x][0];
    };
    int dist(int x, int y) {
        return d[x] + d[y] - 2 * d[anc(x, y)];
    };
};
```

树链剖分

```
vector<vector<int>>> g(n + 1);
for(int i = 1; i < n; i++) {
    int u, v;
    cin >> u >> v;
    g[u].push_back(v);
    g[v].push_back(u);
}
```

```
vector<int> heavy(n + 1), fa(n + 1), sz(n + 1, 1), dep(n + 1), top(n + 1);
function <void(int, int)> dfs_init = [&](int u, int f) {
    dep[u] = dep[f] + 1;
    for(int v : g[u]) {
        if(v == f) continue;
        fa[v] = u;
        dfs_init(v, u);
        sz[u] += sz[v];
        if(!heavy[u] || sz[v] > sz[heavy[u]]) heavy[u] = v;
    }
};
dfs_init(1, 0);
function <void(int, int)> dfs = [&](int u, int t) {
    top[u] = t;
    if(heavy[u]) dfs(heavy[u], t);
    for(int v : g[u]) {
        if(v == fa[u] || v == heavy[u]) continue;
        dfs(v, v);
    }
};
dfs(1, 1);
auto lca = [\&](int u, int v) {
    while(top[u] != top[v]) {
        if(dep[top[u]] < dep[top[v]]) v = fa[top[v]];</pre>
        else u = fa[top[u]];
    }
    return dep[u] < dep[v] ? u : v;</pre>
};
auto dist = [\&] (int u, int v) {
    return dep[u] + dep[v] - 2 * dep[lca(u, v)];
};
```

树链剖分(jiangly)

```
struct HLD {
   int n;
   vector<int> siz, top, dep, parent, in, out, seq;
   vector<vector<int>> g;
   //vector<vector<int>> c;//存每一条链
   int cur;
   HLD() {}
   HLD(int n) {
       init(n);
   }
   void init(int n) {
        this->n = n;
        siz.resize(n + 1);
        top.resize(n + 1);
        dep.resize(n + 1);
        parent.resize(n + 1);
        in.resize(n + 1);
        out.resize(n + 1);
        seq.resize(n + 1);
        cur = 0;
        g.assign(n + 1, \{\});
```

```
//c.assign(n + 1, {});
}
void addEdge(int u, int v) {
    g[u].push_back(v);
    g[v].push_back(u);
}
void work(int root = 1) {
    top[root] = root;
    dep[root] = 1;
    parent[root] = -1;
    dfs1(root);
    dfs2(root);
}
void dfs1(int u) {
    if (parent[u] != -1) {
        g[u].erase(find(g[u].begin(), g[u].end(), parent[u]));
    }
    siz[u] = 1;
    for (auto \&v : g[u]) {
        parent[v] = u;
        dep[v] = dep[u] + 1;
        dfs1(v);
        siz[u] += siz[v];
        if (siz[v] > siz[g[u][0]]) {
            swap(v, g[u][0]);
        }
    }
}
void dfs2(int u) {
    //c[top[u]].push_back(u); 速度会慢很多慎用
    in[u] = cur++;
    seq[in[u]] = u;
    for (auto v : g[u]) {
        top[v] = v == g[u][0] ? top[u] : v;
        dfs2(v);
    out[u] = cur;
}
int lca(int u, int v) {
    while (top[u] != top[v]) {
        if (dep[top[u]] > dep[top[v]]) {
            u = parent[top[u]];
        } else {
            v = parent[top[v]];
        }
    }
    return dep[u] < dep[v] ? u : v;</pre>
int dist(int u, int v) {
    return dep[u] + dep[v] - 2 * dep[lca(u, v)];
int jump(int u, int k) {
    if (dep[u] < k) {
        return -1;
    }
    int d = dep[u] - k;
    while (dep[top[u]] > d) {
        u = parent[top[u]];
```

```
return seq[in[u] - dep[u] + d];
    }
    bool isAncester(int u, int v) {
        return in[u] <= in[v] && in[v] < out[u];</pre>
    }
    int rootedParent(int u, int v) {
        swap(u, v);
        if (u == v) {
            return u;
        }
        if (!isAncester(u, v)) {
            return parent[u];
        }
        auto it = upper_bound(g[u].begin(), g[u].end(), v, [&](int x, int y) {
           return in[x] < in[y];</pre>
        }) - 1;
        return *it;
    }
    int rootedSize(int u, int v) {
        if (u == v) {
            return n;
        if (!isAncester(v, u)) {
            return siz[v];
        }
        return n - siz[rootedParent(u, v)];
    int rootedLca(int a, int b, int c) {
        return lca(a, b) \wedge lca(b, c) \wedge lca(c, a);
    }
};
```

O(nlogn) - O(1) LCA

```
vector<int> g[N];
int dep[N], dfn[N], a[N];
int anc[N][25];
int lg[N];
int chk(int x, int y) {
    return dep[x] < dep[y] ? x : y;
void init_lca(int n) {
    dep[1] = 1;
    int cur = 0;
    const int M = 2 * n;
    function<void(int)> dfs = [&](int u) {
        dfn[u] = ++cur;
        a[cur] = u;
        for(int v : g[u]) {
            dep[v] = dep[u] + 1;
            dfs(v);
            a[++cur] = u;
        }
    };
    dfs(1);
    lg[0] = -1;
```

```
for(int i = 1; i \le M; i++) lg[i] = lg[i / 2] + 1, anc[i][0] = a[i];
    for(int j = 1; j \leftarrow [M]; j++) {
        for(int i = 1; i + (1 << j) - 1 <= M; i++) {
            if(dep[anc[i][j - 1]] < dep[anc[i + (1 << j - 1)][j - 1]]) anc[i][j]
= anc[i][j - 1];
            else anc[i][j] = anc[i + (1 << j - 1)][j - 1];
        }
    }
}
int lca(int 1, int r) {
   1 = dfn[1], r = dfn[r];
    if(1 > r) swap(1, r);
    int k = \lg[r - 1 + 1];
    return chk(anc[1][k], anc[r - (1 << k) + 1][k]);
int dis(int x, int y){
    int 1 = lca(x, y);
    return dep[x] + dep[y] - 2 * dep[1];
}
```

虚树

```
void build() {
    sort(a + 1, a + 1 + k, [\&](int x, int y) {
        return in[x] < in[y];</pre>
   });
    stk[top = 1] = 1;
    for(int i = 1; i \le k; i++) {
        if(a[i] == 1) continue;
        int anc = lca(a[i], stk[top]);
        if(anc != stk[top]) {
            while(in[anc] < in[stk[top - 1]]) g[stk[top - 1]].push_back(stk[top-</pre>
-]);
            if(in[anc] > in[stk[top - 1]]) g[anc].push_back(stk[top]), stk[top]
= anc;
            else g[anc].push_back(stk[top--]);
        stk[++top] = a[i];
    for(int i = 1; i < top; i++) g[stk[i]].push_back(stk[i + 1]);
}//最后需要将虚树中的点记录下来清空边
void build() {
    sort(a + 1, a + 1 + k, [\&](int x, int y) {
        return in[x] < in[y];</pre>
    });
    stk[top = 1] = 1;
    for(int i = 1; i \le k; i++) {
        if(a[i] == 1) continue;
        int anc = lca(a[i], stk[top]);
        if(anc != stk[top]) {
            while(in[anc] < in[stk[top - 1]]) g[stk[top - 1]].push_back(stk[top-</pre>
-]);
            if(in[anc] > in[stk[top - 1]]) g[anc].push_back(stk[top]), stk[top]
= anc;
```

```
else g[anc].push_back(stk[top--]);
       }
       stk[++top] = a[i];
   }
   for(int i = 1; i < top; i++) g[stk[i]].push_back(stk[i + 1]);
}//最后需要将虚树中的点记录下来清空边
//----//
void build_virtual_tree() {
  sort(h + 1, h + m + 1, cmp); // 把关键点按照 dfn 序排序
  for (int i = 1; i < m; ++i) {
   a[++1en] = h[i];
   a[++len] = lca(h[i], h[i + 1]); // 插 \lambda lca
 a[++1en] = h[m];
  sort(a + 1, a + len + 1, cmp); // 把所有虚树上的点按照 dfn 序排序
 len = unique(a + 1, a + len + 1) - a - 1; // \pm 1
 for (int i = 1, lc; i < len; ++i) {
   lc = lca(a[i], a[i + 1]);
   conn(lc, a[i + 1]); // 连边, 如有边权 就是 distance(lc,a[i+1])
 }
}
```

三元环计数

时间复杂度 $m\sqrt{m}$

```
void solve() {
   int n, m;
    cin >> n >> m;
    vector<vector<int>> e(n + 1);
    vector<array<int, 2 >> g(m + 1);
    vector<int> d(n + 1);
    vector<int> st(n + 1);
    for(int i = 1; i \le m; i++) {
        cin >> g[i][0] >> g[i][1];
        d[g[i][0]]_{++}, d[g[i][1]]_{++};
    for(int i = 1; i <= m; i++) {
        auto [x, y] = g[i];
        if(d[x] < d[y] \mid | (d[x] == d[y] \& x < y)) e[y].push_back(x);
        else e[x].push_back(y);
    int res = 0;
    for(int i = 1; i <= n; i++) {
        for(int j : e[i]) st[j] = i;
        for(int j : e[i])
            for(int k : e[j])
                if(st[j] == st[k])
                     res++;
    }
    cout << res;</pre>
}
```

四元环计数

时间复杂度 $m\sqrt{m}$ 个数上界 $nm\sqrt{m}$

```
void solve() {
    int n, m;
    cin >> n >> m;
    vector<vector<int>>> e(n + 1);
    vector<int> a(n + 1), d(n + 1), rk(n + 1), st(n + 1);
    for(int i = 1; i <= m; i++) {
        int x, y;
        cin >> x >> y;
        d[x]++, d[y]++;
        e[x].push_back(y), e[y].push_back(x);
    }
    for(int i = 1; i \le n; i++) a[i] = i;
    sort(a.begin() + 1, a.end(), [\&](int x, int y) {
        return d[x] < d[y] \mid | (d[x] == d[y] && x < y);
    });
    for(int i = 1; i \le n; i++) rk[a[i]] = i;
    long long res = 0;
    for(int i = 1; i <= n; i++) {
        for(int j : e[i])
            if(rk[i] > rk[j])
                for(int k : e[j])
                     if(rk[i] > rk[k])
                         res += st[k]++;
        for(int j : e[i])
            if(rk[i] > rk[j])
                for(int k : e[j])
                     if(rk[i] > rk[k])
                         st[k] = 0;
    }
    cout << res;</pre>
}
```

点分治

一棵树,边带权,统计树上距离不超过k的点对数

```
void solve() {
   int n;
    cin >> n;
   vector<vector<array<int, 2>>> g(n + 1);
    for(int i = 1; i < n; i++) {
       int u, v, w;
        cin >> u >> v >> w;
        g[u].push_back({v, w});
        g[v].push_back({u, w});
   }
   int k;
   cin >> k;
   i64 \ ans = 0;
   vector<int> del(n + 1), maxs(n + 1), sz(n + 1);
    function<void(int, int)> treeDivide = [&](int u, int s) {
        int root = -1;
```

```
function<void(int, int)> center = [&](int u, int fa) {
            sz[u] = 1, maxs[u] = 0;
            for(auto [v, w] : g[u]) {
                if(del[v] || v == fa) continue;
                center(v, u);
                sz[u] += sz[v];
                maxs[u] = max(maxs[u], sz[v]);
            maxs[u] = max(maxs[u], s - sz[u]);
            if(root == -1 \mid | maxs[u] < maxs[root]) root = u;
        };
        center(u, 0);
        vector<int> d1, d2;
        d1.push_back(0);
        auto calc = [&](vector<int> &d) {
            sort(d.begin(), d.end());
            int m = d.size(), r = m - 1;
            i64 \ ans = 0;
            for(int i = 0; i < m; i++) {
                while(r >= 0 \&\& d[i] + d[r] > k) r--;
                if(i < r) ans += r - i;
            }
            return ans;
        };
        for(auto [v, w] : g[root]) {
            if(del[v]) continue;
            d2.clear();
            function<void(int, int, int)> dfs = [\&](int u, int fa, int dep) {
                sz[u] = 1;
                d1.push_back(dep);
                d2.push_back(dep);
                for(auto [v, w] : g[u]) {
                    if(del[v] || v == fa) continue;
                    dfs(v, u, dep + w);
                    sz[u] += sz[v];
                }
            };
            dfs(v, root, w);
            ans -= calc(d2);
        }
        ans += calc(d1);
        del[root] = true;
        for(auto [v, w] : g[root]) {
            if(!del[v]) treeDivide(v, sz[v]);
        }
   };
    treeDivide(1, n);
   cout << ans;</pre>
}
```

动态点分治

将每一层分治重心和上一层的连边得到分治树。这就是点分树。点分树是一棵重构树。

给一棵树,点有点权。 q 次询问给定 x , k , 求与 x 树上距离 $\leq k$ 的所有点的点权和。带修点权,强制在线。 n , $q \leq 10^5$ 。

```
template <typename T>
struct Fenwick {
    int n;
    vector<T> a;
    Fenwick(int _n = 0) : a(_n + 1), n(_n){}
    void add(int x, T v) {
        for (int i = x; i \le n; i += i \& -i) {
            a[i] += v;
        }
    }
    T sum(int x) {
        x = min(x, n);
        auto ans = T();
        for (int i = x; i >= 1; i -= i \& -i) {
            ans += a[i];
        }
        return ans;
    }
    T rangeSum(int 1, int r) {
        return sum(r) - sum(1 - 1);
   }
};
struct LCA {
   int n, LOG;
    vector<vector<int>> fa;
    vector<int> d;
    LCA(int _n) : n(_n), LOG(_lg(n)), fa(n + 1, vector<int> (_lg(n) + 1)), d(n)
+ 1){};
    void init(int s, auto &g) {
        d[s] = 1;
        queue <int> q;
        q.push(s);
        while (!q.empty()) {
            int u = q.front();
            q.pop();
            for (auto [v, w] : g[u]) {
                if (!d[v]) {
                    q.push(v);
                    d[v] = d[u] + 1;
                    fa[v][0] = u;
                    for (int k = 1; k \le LOG; k++) fa[v][k] = fa[fa[v][k - 1]][k
- 1];
                }
            }
        }
    }:
    int anc(int x, int y) {
        if (d[x] > d[y]) swap(x, y);
        for (int k = LOG; k >= 0; k--) if (d[fa[y][k]] >= d[x]) y = fa[y][k];
        if (x == y) return x;
        for (int k = LOG; k >= 0; k --) if (fa[x][k] != fa[y][k]) x = fa[x][k], y
= fa[y][k];
        return fa[x][0];
    };
    int dist(int x, int y) {
        return d[x] + d[y] - 2 * d[anc(x, y)];
    };
};
```

```
void solve() {
    int n, m;
    cin >> n >> m;
    vector<int> w(n + 1);
    vector<vector<array<int, 2>>> g(n + 1);
    vector<vector<int>>> e(n + 1);
    for(int i = 1; i <= n; i++) {
        cin >> w[i];
    }
    for(int i = 1; i < n; i++) {
        int u, v;
        cin >> u >> v;
        g[u].push_back({v, 1});
        g[v].push_back({u, 1});
    LCA lca(n);
    lca.init(1, g);
    vector<Fenwick<int>> fen1(n + 1), fen2(n + 1);
    vector<int> del(n + 1), maxs(n + 1), sz(n + 1), father(n + 1);
    function<void(int, int, int, int)> treeDivide = [&](int u, int s, int rt,
int lastmax) {
        int root = -1;
        function<void(int, int)> center = [&](int u, int fa) {
            sz[u] = 1, maxs[u] = 0;
            for(auto [v, w] : g[u]) {
                if(del[v] || v == fa) continue;
                center(v, u);
                sz[u] += sz[v];
                maxs[u] = max(maxs[u], sz[v]);
            }
            \max s[u] = \max(\max s[u], s - sz[u]);
            if(root == -1 || maxs[u] < maxs[root]) root = u;</pre>
        };
        center(u, 0);
        if(rt) e[rt].push_back(root), father[root] = rt;
        int mx = 0;
        for(auto [v, w] : g[root]) {
            if(del[v]) continue;
            function<void(int, int, int)> dfs = [&](int u, int fa, int dep) {
                sz[u] = 1;
                mx = max(mx, dep);
                for(auto [v, w] : g[u]) {
                    if(del[v] || v == fa) continue;
                    dfs(v, u, dep + w);
                    sz[u] += sz[v];
                }
            };
            dfs(v, root, w);
        }
        fen1[root] = Fenwick<int> (mx + 1);
        fen2[root] = Fenwick<int> (lastmax + 1);
        del[root] = true;
        for(auto [v, w] : g[root]) {
            if(!del[v]) treeDivide(v, sz[v], root, mx);
        }
    };
    treeDivide(1, n, 0, 0);
    for(int i = 1; i \le n; i++) {
```

```
int cur = i;
        while(cur) {
             fen1[cur].add(lca.dist(i, cur) + 1, w[i]);
            if(father[cur]) fen2[cur].add(lca.dist(i, father[cur]) + 1, w[i]);
            cur = father[cur];
        }
    }
    int lastans = 0;
    while(m--) {
        int op;
        cin >> op;
        int x, k;
        cin >> x >> k;
        x \wedge = lastans, k \wedge = lastans;
        if(!op) {
            int cur = x;
            int ans = 0;
            while(cur) {
                if(lca.dist(cur, x) <= k) ans += fen1[cur].sum(k - lca.dist(cur,</pre>
x) + 1);
                if(father[cur] && lca.dist(father[cur], x) <= k) ans -=</pre>
fen2[cur].sum(k - lca.dist(father[cur], x) + 1);
                cur = father[cur];
            }
            cout << ans << '\n';</pre>
            lastans = ans;
        }
        else {
            int cur = x;
            while(cur) {
                 fen1[cur].add(lca.dist(x, cur) + 1, -w[x]);
                 fen1[cur].add(lca.dist(x, cur) + 1, k);
                if(father[cur]) fen2[cur].add(lca.dist(x, father[cur]) + 1, -
w[x]), fen2[cur].add(lca.dist(x, father[cur]) + 1, k);
                cur = father[cur];
            }
            w[x] = k;
        }
    }
}
```

计算几何

板子

```
const double eps = 1e-8;
const double inf = 1e20;
const double pi = acos(-1.0);
const int maxp = 1010;

//判断正负
inline int sgn (double x) {
   if (fabs(x) < eps) return 0;
   if (x < 0) return -1;
   else return 1;
}</pre>
```

```
inline double sqr (double x) {return x * x;}
//精度高时用long long!
struct Point {
    double x, y;
   Point () {}
    Point (double _x, double _y) {
       x = _x;
       y = _y;
   void input () {scanf("%1f%1f", &x, &y);}
   void output () {printf("%.21f%.21f\n", x, y);}
   bool operator == (Point b) const {return sgn(x - b.x) == 0 \& sgn(y - b.y)
== 0;
    bool operator < (Point b) const {return sgn(x - b.x) == 0 ? sgn(y - b.y) < 0
: x < b.x;
    Point operator + (const Point &b) const {return Point (x + b.x, y + b.y);}
    Point operator - (const Point &b) const {return Point (x - b.x, y - b.y);}
    Point operator * (const double &k) const {return Point(x * k, y * k);}
    Point operator / (const double &k) const {return Point(x / k, y / k);}
   //叉积
   double operator ^ (const Point &b) const {return x * b.y - y * b.x;}
   //点乘
   double operator * (const Point &b) const {return x * b.x + y * b.y;}
    //距离原点长度
    double len () {return hypot(x, y);}
   //距离原点长度的平方
   double len2 () {return x * x + y * y;}
   //两点距离
    double distance (Point p) {return hypot(x - p.x, y - p.y);}
   //pa和pb所成的夹角
   double rad (Point a, Point b) {
       Point p = *this;
       return fabs(atan2(fabs((a - p) \land (b - p)), (a - p) * (b - p)));
   }
    //化为长度为r的向量
    Point trunc (double r) {
       double 1 = len();
       if (!sgn(1)) return *this;
       r /= 1;
       return Point (x * r, y * r);
   }
    //逆时针旋转90°
    Point rotleft () {return Point(-y, x);}
   //顺时针旋转90°
    Point rotright () {return Point(y, -x);}
    //绕着p点逆时针旋转angle
    Point rotate (Point p, double angle) {
```

```
Point v = (*this) - p;
        double c = cos(angle), s = sin(angle);
        return Point (p.x + v.x * c - v.y * s, p.y + v.x * s + v.y * c);
   }
};
//象限排序
inline int quadrant (Point a) {
   if (a.x > 0 \& a.y >= 0) return 1;
   if (a.x \le 0 \& a.y > 0) return 2;
   if (a.x < 0 \& a.y <= 0) return 3;
    return 4;
}
//极角排序
inline bool Pcmp (Point a, Point b) {
    double aa = a \wedge b;
    if (quadrant(a) == quadrant(b)) return aa > 0;
   return quadrant(a) < quadrant(b);</pre>
}
struct Line {
    Point s, e;
    Line () {}
    Line (Point _s, Point _e) {
        s = _S;
        e = _e;
    bool operator == (Line v) {return (s == v.s) && (e == v.e);}
    //根据一个点和倾斜角angle确定直线, 0 <= angle < pi
    Line (Point p, double angle) {
        s = p;
        if (sgn(angle - pi / 2) == 0) e = (s + Point(0, 1));
        else e = (s + Point(1, tan(angle)));
    Line (double a, double b, double c) {
        if (sgn(a) == 0) {
            s = Point(0, -c / b);
            e = Point(1, -c / b);
        } else if (sgn(b) == 0) {
            s = Point(-c / a, 0);
            e = Point(-c / a, 1);
        } else {
            s = Point(0, -c / b);
            e = Point(1, (-c - a) / b);
    }
    void input () {
        s.input();
        e.input();
    void adjust () {if (e < s) swap(s, e);}</pre>
    //求线段长度
    double length () {return s.distance(e);}
    //返回直线倾斜角
```

```
double angle () {
        double k = atan2(e.y - s.y, e.x - s.x);
        if (sgn(k) < 0) k += pi;
        if (sgn(k - pi) == 0) k -= pi;
        return k;
   }
   //点和直线的关系,1在左侧,2在右侧,3在直线上
   int relation (Point p) {
        int c = sgn((p - s) \land (e - s));
        if (c < 0) return 1;
        else if (c > 0) return 2;
        else return 3;
   }
    //点在直线上判断
   bool pointonseg (Point p) {return sgn((p - s) \land (e - s)) == 0 \& sgn((p - s)
(p - e) <= 0;
    //两向量平行(直线平行或重合)
   bool parallel (Line v) {return sgn((e - s) \land (v.e - v.s)) == 0;}
   //两线段相交判断, 2规范相交, 1非规范相交, 0不相交
   int segcrossseg (Line v) {
        int d1 = sgn((e - s) \land (v.s - s));
        int d2 = sgn((e - s) \land (v.e - s));
        int d3 = sgn((v.e - v.s) \land (s - v.s));
        int d4 = sgn((v.e - v.s) \land (e - v.s));
        if ((d1 \land d2) == -2 \&\& (d3 \land d4) == -2) return 2;
        return (d1 == 0 && sgn((v.s - s) * (v.s - e)) <= 0) ||
               (d2 == 0 \&\& sgn((v.e - s) * (v.e - e)) <= 0) ||
               (d3 == 0 \&\& sgn((s - v.s) * (s - v.e)) <= 0) ||
               (d4 == 0 \&\& sgn((e - v.s) * (e - v.e)) <= 0);
   }
   //直线(*this)与线段(v)相交判断,2规范相交,1非规范相交,0不相交
   int linecrossseg (Line v) {
        int d1 = sgn((e - s) \wedge (v.s - s));
        int d2 = sgn((e - s) \wedge (v.e - s));
        if ((d1 \land d2) == -2) return 2;
        return (d1 == 0 || d2 == 0);
   }
   //两直线关系,0平行,1重合,2相交
   int linecrossline (Line v) {
       if ((*this).parallel(v)) return v.relation(s) == 3;
        return 2;
   }
    //两直线交点
    Point crosspoint (Line v) {
        double a1 = (v.e - v.s) \land (s - v.s);
        double a2 = (v.e - v.s) \land (e - v.s);
        return Point((s.x * a2 - e.x * a1) / (a2 - a1), (s.y * a2 - e.y * a1) /
(a2 - a1));
   }
   //点到直线距离
```

```
double dispoint to line (Point p) {return fabs((p - s) \land (e - s) / length());}
    //点到线段距离
    double dispointtoseg (Point p) {
       if (sgn((p - s) * (e - s)) < 0 \mid | sgn((p - e) * (s - e)) < 0)
           return min(p.distance(s), p.distance(e));
       return dispointtoline(p);
   }
    //线段到线段的距离
   double dissegtosseg (Line v) {
       return min(min(dispointtoseg(v.s), dispointtoseg(v.e)),
min(v.dispointtoseg(s), v.dispointtoseg(e)));
   }
   //返回点p在直线上的投影
   Point lineprog (Point p) {
       return s + (((e - s) * ((e - s) * (p - s))) / ((e - s).len2()));
   }
   //返回点p关于直线的对称点
   Point symmetrypoint (Point p) {
       Point q = lineprog(p);
       return Point(2 * q.x - p.x, 2 * q.y - p.y);
   }
};
struct circle {
   Point p; //圆心
   double r; //半径
    circle() {}
    circle (Point _p, double _r) {
       p = _p;
       r = _r;
   }
    circle (double x, double y, double _r) {
       p = Point(x, y);
       r = _r;
   }
   //三角形的外接圆
   //需要 Point 的 + / rotate() 以及 Line 的 crosspoint()
    //利用两条边的中垂线得到圆心
   circle (Point a, Point b, Point c) {
       Line u = Line((a + b) / 2, ((a + b) / 2) + ((b - a).rotleft()));
       Line v = Line((b + c) / 2, ((b + c) / 2) + ((c - b).rotleft()));
       p = u.crosspoint(v);
       r = p.distance(a);
   }
   //三角形的内切圆
    //参数 bool t 没有作用,只是为了和外接圆函数区别
    circle (Point a, Point b, Point c, bool t) {
       double m = atan2(b.y - a.y, b.x - a.x), n = atan2(c.y - a.y, c.x - a.x);
       u.e = u.s + Point(cos((n + m) / 2), sin((n + m) / 2));
       v.s = b;
       m = atan2 (a.y - b.y, a.x - b.x), n = atan2(c.y - b.y, c.x - b.x);
```

```
v.e = v.s + Point(cos((n + m) / 2), sin((n + m) / 2));
    p = u.crosspoint(v);
    r = Line(a, b).dispointtoseg(p);
}
//输入
void input () {
    p.input();
    scanf("%1f", &r);
}
//输出
void output () {
    printf("%.21f %.21f %.21f\n", p.x, p.y, r);
bool operator == (circle v) {
    return (p == v.p) && sgn(r - v.r) == 0;
bool operator < (circle v) const {</pre>
    return ((p < v.p) \mid | ((p == v.p) \&\& sgn(r - v.r) < 0));
}
//面积
double area () {
   return pi * r * r;
}
//周长
double circumference () {
   return 2 * pi * r;
//点和圆的关系
//0 圆外
//1 圆上
//2 圆内
int relation (Point b) {
    double dst = b.distance(p);
    if (sgn(dst - r) < 0) return 2;
    else if (sgn(dst - r) == 0) return 1;
    return 0;
}
//线段和圆的关系
//比较的是圆心到线段的距离和半径的关系
int relationseg (Line v) {
    double dst = v.dispointtoseg(p);
    if (sgn(dst - r) < 0) return 2;
    else if (sgn(dst - r) == 0) return 1;
    return 0;
}
//直线和圆的关系
//比较的是圆心到直线的距离和半径的关系
int relationline (Line v) {
    double dst = v.dispointtoline(p);
    if (sgn(dst - r) < 0) return 2;
    else if (sgn(dst - r) == 0) return 1;
```

```
return 0;
}
//两圆的关系
//5 相离
//4 外切
//3 相交
//2 内切
//1 内含
//需要 Point 的 distance
int relationcircle (circle v) {
    double d = p.distance(v.p);
    if (sgn(d - r - v.r) > 0) return 5;
    if (sgn(d - r - v.r) == 0) return 4;
    double l = fabs(r - v.r);
    if (sgn(d - r - v.r) < 0 \& sgn(d - 1) > 0) return 3;
    if (sgn(d - 1) == 0) return 2;
    if (sgn(d - 1) < 0) return 1;
}
//求两个圆的交点,返回 0 表示没有交点,返回 1 是一个交点, 2 是两个交点
//需要 relationcircle
int pointcrosscircle (circle v, Point &p1, Point &p2) {
    int rel = relationcircle(v);
    if (rel == 1 || rel == 5) return 0;
    double d = p.distance(v.p);
    double l = (d * d + r * r - v.r * v.r) / (2 * d);
    double h = sqrt(r * r - 1 * 1);
    Point tmp = p + (v.p - p).trunc(1);
    p1 = tmp + ((v.p - p).rotleft().trunc(h));
    p2 = tmp + ((v.p - p).rotright().trunc(h));
    if (rel == 2 || rel == 4) return 1;
    return 2;
}
//求直线和圆的交点,返回交点个数
int pointcrossline (Line v, Point &p1, Point &p2) {
    if (!(*this).relationline(v)) return 0;
    Point a = v.lineprog(p);
    double d = v.dispointtoline(p);
    d = sqrt(r * r - d * d);
    if (sgn(d) == 0) {
        p1 = a;
       p2 = a;
       return 1;
    }
    p1 = a + (v.e - v.s).trunc(d);
    p2 = a - (v.e - v.s).trunc(d);
    return 2;
}
//得到过 a,b 两点,半径为 r1 的两个圆
int getcircle (Point a, Point b, double r1, circle &c1, circle &c2) {
    circle x(a, r1), y(b, r1);
    int t = x.pointcrosscircle(y, c1.p, c2.p);
    if (!t) return 0;
    c1.r = c2.r = r1;
    return t;
```

```
//得到与直线 u 相切,过点 q, 半径为 r1 的圆
        int getcircle (Line u, Point q, double r1, circle &c1, circle &c2) {
                 double dis = u.dispointtoline(q);
                 if (sgn(dis - r1 * 2) > 0) return 0;
                 if (sgn(dis) == 0) {
                         c1.p = q + ((u.e - u.s).rotleft().trunc(r1));
                         c2.p = q + ((u.e - u.s).rotright().trunc(r1));
                         c1.r = c2.r = r1;
                         return 2;
                 }
                 Line u1 = Line((u.s + (u.e - u.s).rotleft().trunc(r1)), (u.e + (u.e - u.e))
u.s).rotleft().trunc(r1)));
                 Line u2 = Line((u.s + (u.e - u.s).rotright().trunc(r1)), (u.e + (u.e - u.s).rotright().trunc(r1)), (u.e + (u.e - u.s).rotright().trunc(r1)), (u.e + u.e - u.e).rotright().trunc(r1)), (u.e - u.e - u.e).rotright().trunc(r1)), (u.e - u.e - u.e).rotright().trunc(r1)), (u.e - u.e - u.e)
u.s).rotright().trunc(r1)));
                 circle cc = circle(q, r1);
                 Point p1, p2;
                 if (!cc.pointcrossline(u1, p1, p2)) cc.pointcrossline(u2, p1, p2);
                 c1 = circle(p1, r1);
                 if (p1 == p2) {
                         c2 = c1;
                         return 1;
                 }
                 c2 = circle(p2, r1);
                 return 2;
        //同时与直线 u,v 相切,半径为 r1 的圆
        int getcircle (Line u, Line v, double r1, circle &c1, circle &c2, circle
&c3, circle &c4) {
                 if (u.parallel(v)) return 0; //两直线平行
                 Line u1 = Line(u.s + (u.e - u.s).rotleft().trunc(r1), u.e + (u.e - u.e)
u.s).rotleft().trunc(r1));
                 Line u2 = Line(u.s + (u.e - u.s).rotright().trunc(r1), u.e + (u.e - u.e)
u.s).rotright().trunc(r1));
                 Line v1 = Line(v.s + (v.e - v.s).rotleft().trunc(r1), v.e + (v.e - v.e)
v.s).rotleft().trunc(r1));
                 Line v2 = Line(v.s + (v.e - v.s).rotright().trunc(r1), v.e + (v.e - v.e)
v.s).rotright().trunc(r1));
                 c1.r = c2.r = c3.r = c4.r = r1;
                 c1.p = u1.crosspoint(v1);
                 c2.p = u1.crosspoint(v2);
                 c3.p = u2.crosspoint(v1);
                 c4.p = u2.crosspoint(v2);
                 return 4;
        }
        //同时与不相交圆 cx,cy 相切,半径为 r1 的圆
        int getcircle (circle cx, circle cy, double r1, circle &c1, circle &c2) {
                 circle x(cx.p, r1 + cx.r), y(cy.p, r1 + cy.r);
                 int t = x.pointcrosscircle(y, c1.p, c2.p);
                 if (!t) return 0;
                 c1.r = c2.r = r1;
                 return t;
        }
        //过一点作圆的切线 (先判断点和圆的关系)
        int tangentline (Point q, Line &u, Line &v) {
```

```
int x = relation(q);
        if (x == 2) return 0;
        if (x == 1) {
            u = Line(q, q + (q - p).rotleft());
            v = u;
            return 1;
        }
        double d = p.distance(q);
        double l = r * r / d;
        double h = sqrt(r * r - 1 * 1);
        u = Line(q, p + ((q - p).trunc(1) + (q - p).rotleft().trunc(h)));
        v = Line(q, p + ((q - p).trunc(1) + (q - p).rotright().trunc(h)));
        return 2;
    }
    //求两圆相交的面积
    double areacircle (circle v) {
        int rel = relationcircle(v);
        if (rel >= 4) return 0.0;
        if (rel <= 2) return min(area(), v.area());</pre>
        double d = p.distance(v.p);
        double hf = (r + v.r + d) / 2.0;
        double ss = 2 * sqrt(hf * (hf - r) * (hf - v.r) * (hf - d));
        double a1 = acos((r * r + d * d - v.r * v.r) / (2.0 * r * d));
        a1 = a1 * r * r;
        double a2 = acos((v.r * v.r + d * d - r * r) / (2.0 * v.r * d));
        a2 = a2 * v.r * v.r;
        return a1 + a2 - ss;
    }
    //求圆和三角形 pab 的相交面积
    double areatriangle (Point a, Point b) {
        if (sgn((p - a) \land (p - b)) == 0) return 0.0;
        Point q[5];
        int len = 0;
        q[len ++] = a;
        Line 1(a, b);
        Point p1, p2;
        if (pointcrossline(1, q[1], q[2]) == 2) {
            if (sgn((a - q[1]) * (b - q[1])) < 0) q[len ++] = q[1];
            if (sgn((a - q[2]) * (b - q[2])) < 0) q[len ++] = q[2];
        q[len ++] = b;
        if (len == 4 \& sgn((q[0] - q[1]) * (q[2] - q[1])) > 0) swap(q[1],
q[2]);
        double res = 0;
        for (int i = 0; i < len - 1; i ++) {
            if (relation(q[i]) == 0 \mid | relation(q[i + 1]) == 0)  {
                double arg = p.rad(q[i], q[i + 1]);
                res += r * r * arg / 2.0;
                res += fabs((q[i] - p) \land (q[i + 1] - p)) / 2.0;
        return res;
    }
};
```

链表

单链表

- 1. $H \times$, 表示向链表头插入一个数 x。
- 2. D k , 表示删除第 k 个插入的数后面的数 (当 k 为 00 时 , 表示删除头结点)。
- 3. I k x , 表示在第 k 个插入的数后面插入一个数 x (此操作中 k 均大于 0) 。

```
#include <bits/stdc++.h>
using namespace std;
const int N = 100010;
// head 表示头结点的下标
// e[i] 表示节点i的值
// ne[i] 表示节点i的next指针是多少
// idx 存储当前已经用到了哪个点
int head, e[N], ne[N], idx;
// 初始化
void init() {
   head = -1;
   idx = 0;
}
// 将x插到头结点
void add_to_head(int x) {
   e[idx] = x, ne[idx] = head, head = idx +++;
}
// 将x插到下标是k的点后面
void add(int k, int x) {
   e[idx] = x, ne[idx] = ne[k], ne[k] = idx ++;
// 将下标是k的点后面的点删掉
void remove(int k) {
   ne[k] = ne[ne[k]];
int main() {
   int m;
   cin >> m;
   init();
   while (m--) {
       int k, x;
       char op;
       cin >> op;
       if (op == 'H') {
           cin >> x;
           add_to_head(x);
       } else if (op == 'D') {
           cin >> k;
           if (!k) head = ne[head];
           else remove(k - 1);
       } else {
           cin >> k >> x;
           add(k - 1, x);
       }
   }
```

```
for (int i = head; i != -1; i = ne[i]) cout << e[i] << ' ';
return 0;
}</pre>
```

双链表

```
#include <bits/stdc++.h>
using namespace std;
const int N = 100010;
int m;
int e[N], 1[N], r[N], idx;
// 在节点a的右边插入一个数x
void insert(int a, int x) {
   e[idx] = x;
   l[idx] = a, r[idx] = r[a];
   l[r[a]] = idx, r[a] = idx ++ ;
}
// 删除节点a
void remove(int a) {
   l[r[a]] = l[a];
   r[1[a]] = r[a];
}
int main() {
   cin >> m;
    // 0是左端点,1是右端点
   r[0] = 1, \ l[1] = 0;
   idx = 2;
   while (m--) {
        string op;
        cin >> op;
        int k, x;
        if (op == "L") {
           cin >> x;
           insert(0, x);
        } else if (op == "R") {
           cin >> x;
            insert(1[1], x);
        } else if (op == "D") {
           cin >> k;
            remove(k + 1);
        } else if (op == "IL") {
            cin >> k >> x;
            insert(l[k + 1], x);
        } else {
           cin >> k >> x;
           insert(k + 1, x);
   }
   for (int i = r[0]; i != 1; i = r[i]) cout << e[i] << ' ';
   return 0;
}
```

模拟堆

维护一个集合, 初始时集合为空, 支持如下几种操作:

```
    Ix,插入一个数 x;
    PM,输出当前集合中的最小值;
    DM,删除当前集合中的最小值(数据保证此时的最小值唯一);
    Dk,删除第 k 个插入的数;
    Ckx,修改第 k 个插入的数,将其变为 x;
```

现在要进行 N 次操作,对于所有第2个操作,输出当前集合的最小值。

```
#include <bits/stdc++.h>
using namespace std;
const int N = 100010;
int h[N], ph[N], hp[N], cnt;
void heap_swap(int a, int b) {
    swap(ph[hp[a]], ph[hp[b]]);
    swap(hp[a], hp[b]);
    swap(h[a], h[b]);
void down(int u) {
   int t = u;
   if (u * 2 <= cnt && h[u * 2] < h[t]) t = u * 2;
   if (u * 2 + 1 \le cnt \& h[u * 2 + 1] < h[t]) t = u * 2 + 1;
    if (u != t) {
        heap_swap(u, t);
        down(t);
    }
}
void up(int u) {
    while (u / 2 \&\& h[u] < h[u / 2]) {
        heap_swap(u, u / 2);
        u >>= 1;
    }
int main() {
   int n, m = 0;
    scanf("%d", &n);
    while (n--) {
        char op[5];
        int k, x;
        scanf("%s", op);
        if (!strcmp(op, "I")) {
            scanf("%d", &x);
            cnt ++ ;
            m ++ ;
            ph[m] = cnt, hp[cnt] = m;
            h[cnt] = x;
            up(cnt);
        } else if (!strcmp(op, "PM")) printf("%d\n", h[1]);
        else if (!strcmp(op, "DM")) {
            heap_swap(1, cnt);
            cnt -- ;
            down(1);
```

```
} else if (!strcmp(op, "D")) {
            scanf("%d", &k);
            k = ph[k];
            heap_swap(k, cnt);
            cnt -- ;
            up(k);
            down(k);
        } else {
            scanf("%d%d", &k, &x);
            k = ph[k];
            h[k] = x;
            up(k);
            down(k);
        }
    }
   return 0;
}
```

字典树

字符串字典树

```
struct trie {
   vector<vector<int>> tr;
   vector<int> val;
   int idx = 0;
   vector<int> T;
    trie(int m) : val(1), T(m, 0){
        tr.emplace_back(T);
   };//m是字符集大小
    void insert(string &s) {
        int p = 0;
        for (int i = 0; i < s.size(); i++) {
            int u = s[i] - 'a';
            if (!tr[p][u]) {
                tr[p][u] = ++idx;
                tr.emplace_back(T);
                val.emplace_back(0);
            }
            p = tr[p][u];
        val[p]++;
    }
    int query(string &s) {
        int p = 0, ans = 0;
        for (int i = 0; i < s.size(); i++) {
            int u = s[i] - 'a';
            if (!tr[p][u]) return 0;
            p = tr[p][u];
        return val[p];
   }
};
```

01字典树

在给定的 N 个整数 A_1 , A_2 A_N 中选出两个进行 xor (异或) 运算,得到的结果最大/最小是多少?

```
int n, a[N], tr[N * 30][2], idx;
void insert(int x) {
   int p = 0;
    for (int i = 30; i >= 0; i--) {
        int \&s = tr[p][x >> i \& 1];
        if (!s) {
           s = ++idx;
        }
        p = s;
    }
}
int query_mx(int x) {
    int p = 0, res = 0;
    for (int i = 30; i >= 0; i--) {
        int s = x >> i & 1;
        if (tr[p][!s]) {
            res += 1 << i;
            p = tr[p][!s];
        } else p = tr[p][s];
    }
    return res;
}
int query_mn(int x) {
    int p = 0, res = 0;
    for (int i = 30; i >= 0; i--) {
        int s = x >> i & 1;
        if (tr[p][s]) {
            p = tr[p][s];
        } else p = tr[p][!s], res += 1 << i;
    }
   return res;
}
void del(int x) {
    for(int i = 0; i \le idx; i++) tr[i][0] = tr[i][1] = 0;
}
```

并查集

```
struct DSU {
    vector<int> f, siz;
    DSU(int n) : f(n + 1), siz(n + 1, 1) { iota(f.begin(), f.end(), 0); }
    int leader(int x) {
        while (x != f[x]) x = f[x] = f[f[x]];
        return x;
    }
    bool same(int x, int y) { return leader(x) == leader(y); }
    bool merge(int x, int y) {
        x = leader(x);
        y = leader(y);
        if (x == y) return false;
        siz[x] += siz[y];
    }
}
```

```
f[y] = x;
    return true;
}
int size(int x) { return siz[leader(x)]; }
};
```

可撤销并查集

```
struct DSU {
   int n;
   vector<int> fa, sz;
   stack<array<int, 2>> stk;
   void init(int n) {
        this->n = n;
        sz.assign(n + 1, 1);
        fa.resize(n + 1);
        iota(fa.begin(), fa.end(), 0);
   }
   int find(int x) {
       while(fa[x] != x) x = fa[x];
        return x;
   }
   void merge(int x, int y) {
        int fx = find(x), fy = find(y);
        if(fx != fy) {
            if(sz[fx] > sz[fy]) swap(fx, fy);
            sz[fy] += sz[fx];
           fa[fx] = fy;
           stk.push({fx, fy});
           //tag[fx] -= tag[fy];
   }
   void undo(int last) {
        while((int)stk.size() > last) {
            auto [x, y] = stk.top();
            stk.pop();
           sz[y] = sz[x];
            fa[x] = x;
           //tag[x] += tag[y];
        }
   }
//按秩合并并查集、可撤销并查集
DSU dsu;
```

带权并查集

```
struct DSU {
    struct T {
        int fa;
        i64 len;
    };
    vector<T> t;
    DSU(int n) : t(n + 1) {
        for (int i = 1; i <= n; i++) {
            t[i] = {i, 0};
    }
}</pre>
```

```
}
    int getfa(int v) {
        if (t[v].fa == v) {
            return v;
        }
        int f = getfa(t[v].fa);
        t[v].len += t[t[v].fa].len;
        t[v].fa = f;
        return f;
    }
    bool merge(int u, int v, int w) {
        int uf = getfa(u), vf = getfa(v);
        if (uf == vf) {
            return t[u].len - t[v].len == w;
        }
        t[uf].len = w + t[v].len - t[u].len;
        t[uf].fa = vf;
        return true;
   }
};
```

可持久化并查集

```
#include <bits/stdc++.h>
using namespace std;
using i64 = long long;
constexpr int N = 2e5 + 10;
struct DSU {
               int node = 0, n;
               vector<int> ls, rs, f, sz, r;
               DSU(int n, int m) : n(n), ls(n * 4 + 2 * (__lg(m) + 1) * m), rs(n * 4 + 2 * (__lg(m) + 1) * m)
 (_lg(m) + 1) * m), f(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * m), sz(n * 4 + 2 * (__lg(m) + 1) * 
1) * m), r(m + 1)\{\};
               void build(int 1, int r, int &x) {
                               x = ++node;
                               if(1 == r) {
                                               f[x] = 1;
                                               sz[x] = 1;
                               }
                               else {
                                               int mid = 1 + r \gg 1;
                                               build(1, mid, ls[x]);
                                               build(mid + 1, r, rs[x]);
                               }
               void modify(int pre, int &x, int 1, int r, int p, int v, bool op) {
                               x = ++node, ls[x] = ls[pre], rs[x] = rs[pre], f[x] = f[pre], sz[x] =
sz[pre];
                              if(1 == r) {
                                             if(op) f[x] = v;
                                               else sz[x] = v;
                               }
```

```
else {
            int mid = 1 + r \gg 1;
            if(p \leftarrow mid) modify(ls[pre], ls[x], l, mid, p, v, op);
            else modify(rs[pre], rs[x], mid + 1, r, p, v, op);
        }
    }
    int query(int 1, int r, int p, int x, bool op) {
        if(1 == r) {
           if(op) return f[x];
            else return sz[x];
        }
        else {
            int mid = 1 + r \gg 1;
            if(p <= mid) return query(1, mid, p, ls[x], op);</pre>
            else return query(mid + 1, r, p, rs[x], op);
        }
    }
    int find(int x, int id) {
        int fa = query(1, n, x, r[id], 1);
        if(fa == x) return x;
        else return find(fa, id);
    void merge(int x, int y, int id) {
        x = find(x, id), y = find(y, id);
        if(x == y) r[id + 1] = r[id];
        else {
            int sx = query(1, n, x, r[id], 0), sy = query(1, n, y, r[id], 0);
            if(sx < sy) swap(x, y);
            modify(r[id], r[id + 1], 1, n, y, x, 1);
            modify(r[id + 1], r[id + 1], 1, n, x, sx + sy, 0);
        }
    }
    bool check(int x, int y, int id) {
        return find(x, id) == find(y, id);
    }
};
void solve() {
   int n, m;
    cin >> n >> m;
    DSU dsu(n, m);
    dsu.build(1, n, dsu.r[0]);
    int id = 0;
    while(m--) {
        ++id;
        int op;
        cin >> op;
        if(op == 1) {
            int a, b;
            cin >> a >> b;
            dsu.merge(a, b, id - 1);
        }//合并a,b所在集合
        else if(op == 2) {
            int k;
            cin >> k;
            dsu.r[id] = dsu.r[k];
        }//回到第k次操作(执行三种操作中的任意一种都记为一次操作)之后的状态
        else {
```

```
int a, b;
            cin >> a >> b;
            dsu.r[id] = dsu.r[id - 1];
            cout << dsu.check(a, b, id) << '\n';</pre>
        }//询问a,b是否属于同一集合,如果是则输出1,否则输出0
   }
}
int main() {
   ios::sync_with_stdio(false);
   cin.tie(nullptr);
   int t;
   t = 1;
   while(t--) {
       solve();
   }
   return 0;
}
```

ST表

```
template <typename T>
struct ST_Table {
    vector<vector<T>> mx, mn, g;
    vector<int> lg;
   vector<T> a;
    ST_Table (int n) : a(n + 1), lg(n + 1) {
        this->n = n;
        lg[0] = -1;
        for(int i = 1; i \le n; i++) lg[i] = lg[i / 2] + 1;
    void buildmax() {
        mx.resize(n + 1);
        for(int i = 1; i \le n; i++) mx[i].resize(lg[n] + 1), mx[i][0] = a[i];
        for(int j = 1; j \leftarrow [n]; j++) {
            for(int i = 1; i + (1 << j) - 1 <= n; i++) {
                mx[i][j] = max(mx[i][j - 1], mx[i + (1 << j - 1)][j - 1]);
            }
        }
    }
    void buildmin() {
        mn.resize(n + 1);
        for(int i = 1; i \le n; i++) mn[i].resize(lg[n] + 1), <math>mn[i][0] = a[i];
        for(int j = 1; j \leftarrow [g[n]; j++) {
            for(int i = 1; i + (1 << j) - 1 <= n; i++) {
                mn[i][j] = min(mn[i][j - 1], mn[i + (1 << j - 1)][j - 1]);
            }
        }
    }
    void buildgcd() {
        g.resize(n + 1);
```

```
for(int i = 1; i \le n; i++) g[i].resize(lg[n] + 1), g[i][0] = a[i];
        for(int j = 1; j \leftarrow [n]; j++) {
            for(int i = 1; i + (1 << j) - 1 <= n; i++) {
                g[i][j] = \underline{gcd}(g[i][j-1], g[i+(1 << j-1)][j-1]);
            }
        }
    T querymax(int 1, int r) {
        int p = \lg[r - l + 1];
        return max(mx[1][p], mx[r - (1 << p) + 1][p]);
    }
    T querymin(int 1, int r) {
        int p = \lg[r - 1 + 1];
        return min(mn[1][p], mn[r - (1 << p) + 1][p]);
    T querygcd(int 1, int r) {
        int p = \lg[r - 1 + 1];
        return \_gcd(g[1][p], g[r - (1 << p) + 1][p]);
   }
};
```

笛卡尔树

```
struct Cartesian_tree {
   vector<int> a, 1, r;
    stack<int> stk;
   int root = 0, n;
   Cartesian_tree (int n) : a(n + 1), l(n + 1), r(n + 1) {this->n = n;};
   void build() {
        for(int i = 1; i <= n; i++) {
            int 1st = 0;
           while(!stk.empty() && a[stk.top()] > a[i]) {//求最大值时要改成小于号
                1st = stk.top();
                stk.pop();
            }
            if(!stk.empty()) r[stk.top()] = i;
            else root = i;
           l[i] = lst;
           stk.push(i);
        }
    }
};
```

普通莫队

```
int n, m, a[N], res[M], len, cnt[S], ans;
struct query {
    int l, r, id, b;
} q[M];
int get(int x) {
    return x / len;
}
bool cmp(query x, query y) {
    if (x.b != y.b) return x.b < y.b;</pre>
```

```
else {
        if (x.b \& 1) return x.r < y.r;
        else return x.r > y.r;
    }
}
void add(int x) {
    if (!cnt[x]) ans++;
    cnt[x]++;
}
void del(int x) {
    cnt[x]--;
    if (!cnt[x]) ans--;
}
int main() {
    cin >> n >> m;
    for (int i = 1; i \le n; i++) cin >> a[i];
    len = max(1, (int)sqrt((double)n * n / m));
    for (int i = 1; i \le m; i++) cin >> q[i].1 >> q[i].r, q[i].id = i, q[i].b =
get(q[i].1);
    sort(q + 1, q + 1 + m, cmp);
    int 1 = 1, r = 0;
    for (int i = 1; i <= m; i++) {
        while (r < q[i].r) add(a[++r]);
        while (1 > q[i].1) add(a[--1]);
        while (1 < q[i].1) del(a[1++]);
        while (r > q[i].r) del(a[r--]);
        res[q[i].id] = ans;
    for (int i = 1; i \leftarrow m; i++) cout \leftarrow res[i] \leftarrow '\n';
    return 0;
}
```

树状数组

单点修改+区间查询

```
template <typename T>
struct Fenwick {
   int n;
   vector<T> a;
   Fenwick(int _n = 0) : n(_n), a(_n + 1){}
   void init(int n) {
        a.assign(n + 1, T());
        this->n = n;
   void add(int x, T v) {
        for (int i = x; i \le n; i += i \& -i) {
            a[i] += v;
        }
   }
   T sum(int x) {
        auto ans = T();
        for (int i = x; i >= 1; i -= i \& -i) {
            ans += a[i];
        }
        return ans;
    }
```

```
T rangeSum(int 1, int r) {
        return sum(r) - sum(1 - 1);
    int kth(T k) {
        int x = 0;
        for (int i = 1 \ll std::_lg(n); i; i /= 2) {
            if (x + i \le n \&\& k >= a[x + i]) {
                x += i;
                k = a[x];
            }
        }
        return x;
    }
};
struct Max {
    int v;
    Max(int x = -1E9) : v{x} {}//注意范围
    Max &operator+=(Max a) {
       v = max(v, a.v);
        return *this;
    }
};
```

区间修改+区间查询

```
template <typename T>
struct Fenwick {
   int n;
   vector<T> a, b;
    Fenwick(int _n = 0) : a(n + 1, T()), b(n + 1, T()), n(_n){}
    void rangeAdd(int 1, int r, T v) {
        for (int i = 1; i \le n; i += i \& -i) {
            a[i] += v, b[i] += 1 * v;
        for (int i = r + 1; i \le n; i += i \& -i) {
            a[i] -= v, b[i] -= (r + 1) * v;
        }
    }
    T sum(int x) {
        auto ans = T();
        for (int i = x; i >= 1; i -= i \& -i) {
            ans += (x + 1) * a[i] - b[i];
        }
        return ans;
    T rangeSum(int 1, int r) {
        return sum(r) - sum(1 - 1);
   }
};
```

二维树状数组区间修改+区间查询

```
int tr[5][N][N];
int lowbit(int x) {
```

```
return x & -x;
}
void add(int x, int y, int v) {
    for (int i = x; i \leftarrow n; i \leftarrow lowbit(i)) {
        for (int j = y; j \leftarrow m; j \leftarrow lowbit(j)) {
             tr[0][i][j] += v;
            tr[1][i][j] += v * x;
            tr[2][i][j] += v * y;
            tr[3][i][j] += v * x * y;
        }
    }
void range_add(int x1, int y1, int x2, int y2) {
    add(x1, y1, 1);
    add(x2 + 1, y1, -1);
    add(x1, y2 + 1, -1);
    add(x2 + 1, y2 + 1, 1);
int ask(int x, int y) {
    int res = 0;
    for (int i = x; i \rightarrow 1; i \rightarrow 1) {
        for (int j = y; j \ge 1; j = lowbit(j)) {
             res += tr[0][i][j] * (x + 1) * (y + 1) - tr[1][i][j] * (y + 1) -
tr[2][i][j] * (x + 1) + tr[3][i][j];
        }
    }
    return res;
int range_sum(int x1, int y1, int x2, int y2) {
    return ask(x2, y2) - ask(x1 - 1, y2) - ask(x2, y1 - 1) + ask(x1 - 1, y1 - 1)
1);
}
```

线段树

单点修改区间查询(最大值)

```
struct node {
    int 1, r, v;
    #define ls u << 1
    #define rs u << 1 | 1
} tr[N << 2];</pre>
void pushup(int u) {
    tr[u].v = max(tr[ls].v, tr[rs].v);
}
void build(int u, int 1, int r) {
    tr[u] = \{1, r\};
    if(1 == r) {
        tr[u].v = a[1];
        return ;
    }
    int mid = 1 + r \gg 1;
    build(ls, l, mid), build(rs, mid + 1, r);
    pushup(u);
}
int query(int u, int 1, int r) {
    if(tr[u].l >= l \&\& tr[u].r <= r) return tr[u].v;
```

```
int mid = tr[u].l + tr[u].r >> 1;
    int v = 0;
    if(l <= mid) v = max(v, query(ls, l, r));
    if(r > mid) v = max(v, query(rs, l, r));
    return v;
}

void modify(int u, int x, int v) {
    if(tr[u].l == x && tr[u].r == x) tr[u].v = v;
    else {
        int mid = tr[u].l + tr[u].r >> 1;
        if(mid >= x) modify(ls, x, v);
        else modify(rs, x, v);
        pushup(u);
    }
}
```

```
struct Info {
    int mx;
};
Info operator+(const Info &a, const Info &b) {
    return {max(a.mx, b.mx)};
}
template <class T>
struct SegmentTree {
   int n;
    vector<T> info;
    vector<T> _init;
    SegmentTree() {}
    SegmentTree(int n) {
        _{\text{init.assign}}(n + 1, T());
        init(_init);
    SegmentTree(vector<T> &_init) {
        this->_init = _init;
        init(_init);
    }
    void init(vector<T> &_init) {
        n = (int)_init.size() - 1;
        info.assign(n * 4 + 1, T());
        build(1, 1, n);
    }
    void build(int p, int 1, int r) {
        if(1 == r) {
            info[p] = _init[1];
            return ;
        }
        int m = (1 + r) / 2;
        build(p * 2, 1, m);
        build(p * 2 + 1, m + 1, r);
        pull(p);
    };
    void pull(int p) {
        info[p] = info[p * 2] + info[p * 2 + 1];
    void modify(int p, int 1, int r, int x, T v) {
        if(1 == r) info[p] = v;
        else {
```

```
int m = (1 + r) / 2;
            if(x \le m) modify(p * 2, 1, m, x, v);
            else modify(p * 2 + 1, m + 1, r, x, v);
            pull(p);
        }
    }
    void modify(int x, T v) {
        modify(1, 1, n, x, v);
    }
    T query(int p, int 1, int r, int x, int y) {
        if(1 > y \mid \mid r < x) return T();
        if(x \le 1 \&\& y \ge r) return info[p];
        int m = (1 + r) / 2;
        return query(p * 2, 1, m, x, y) + query(p * 2 + 1, m + 1, r, x, y);
    T query(int 1, int r) {
        return query(1, 1, n, 1, r);
    }
};
```

ddp Info为矩阵时

```
struct Info {
   int a[2][2];
   void init() {
       for(int i = 0; i < 2; i++) {
           for(int j = 0; j < 2; j++) {
               //a[i][j] = inf - 1;
           }
       }
   }
   //array<array<int, 2>, 2> a {inf - 1, inf - 1, inf - 1};
};//数组最快 arrar次之 vector巨慢很容易超时
Info operator+(const Info &a, const Info &b) {
   Info c;
   c.init();
   for(int i = 0; i < 2; i++) {
       for(int j = 0; j < 2; j++) {
           for(int k = 0; k < 2; k++) {
               //c.a[i][k] = max(c.a[i][k], min(a.a[i][j], b.a[j][k]));
           }
       }
   return c;
}//广义矩阵乘法
```

扫描线

```
struct node {
    int l, r;
    i64 sum, v;
    #define lc u << 1
    #define rc u << 1 | 1
}tr[N << 2];
void pushup(int u) {
    if(tr[u].sum) tr[u].v = tr[u].l + 1;</pre>
```

```
else if(tr[u].l!= tr[u].r) tr[u].v = tr[lc].v + tr[rc].v;
    else tr[u].v = 0;
void build(int u, int 1, int r) {
    if(1 == r) tr[u] = \{1, r, 0\};
    else {
        tr[u] = \{1, r\};
        int mid = 1 + r \gg 1;
        build(lc, l, mid), build(rc, mid + 1, r);
        pushup(u);
    }
void modify(int u, int 1, int r, i64 d) {
   if(tr[u].1 >= 1 && tr[u].r <= r) {
        tr[u].sum += d;
        pushup(u);
    }
    else {
        int mid = tr[u].l + tr[u].r >> 1;
        if(1 \le mid) modify(1c, 1, r, d);
        if(r > mid) modify(rc, 1, r, d);
        pushup(u);
    }
}
int query(int u, int 1, int r) {
    if(tr[u].1 >= 1 \&\& tr[u].r <= r) return tr[u].v;
    int mid = tr[u].l + tr[u].r >> 1;
    i64 sum = 0;
    if(1 \le mid) sum = sum + query(1c, 1, r);
    if(r > mid) sum = sum + query(rc, 1, r);
    return sum;
}
```

单点修改区间查询合并结果区间(最大连续子段和)

```
struct Node {
   int 1, r;
    int sum, 1max, rmax, tmax;
    #define lc u << 1
    #define rc u << 1 | 1
}tr[N * 4];
void pushup(Node &u, Node &1, Node &r) {
    u.sum = 1.sum + r.sum;
    u.lmax = max(1.lmax, l.sum + r.lmax);
    u.rmax = max(r.rmax, r.sum + 1.rmax);
    u.tmax = max(max(1.tmax, r.tmax), 1.rmax + r.lmax);
void pushup(int u) {
    pushup(tr[u], tr[lc], tr[rc]);
void build(int u, int 1, int r) {
    if(1 == r) tr[u] = \{1, r, a[1], a[1], a[1]\};
        tr[u].1 = 1, tr[u].r = r;
        int mid = (1 + r) >> 1;
        build(lc, l, mid), build(rc, mid + 1, r);
        pushup(u);
```

```
}
void modify(int u, int x, int v) {
    if(tr[u].] == x && tr[u].r == x) {
        tr[u] = \{x, x, v, v, v, v\};
    }
    else {
        int mid = (tr[u].l + tr[u].r) >> 1;
        if(x \le mid) modify(1c, x, v);
        else modify(rc, x, v);
        pushup(u);
    }
Node query(int u, int 1, int r) {
    if(tr[u].l >= l \& tr[u].r <= r) return tr[u];
    else {
        int mid = (tr[u].l + tr[u].r) >> 1;
        if(r <= mid) return query(lc, l, r);</pre>
        else if(l > mid) return query(rc, l, r);
        else {
            auto left = query(lc, l, r);
            auto right = query(rc, 1, r);
            Node res;
            pushup(res, left, right);
            return res;
        }
    }
}
```

懒标记线段树(区间乘加+区间求和) 洛谷模板

```
struct Tag {
    i64 \text{ add} = 0, \text{ mul} = 1;
    void apply(Tag &t) {
        add = (add * t.mul + t.add) % m;
        mul = mul * t.mul % m;
    }
};
struct Info {
    i64 \text{ sum} = 0, 1, r;
    void apply(Tag &t) {
        sum = (sum * t.mul + t.add * (r - l + 1) % m) % m;
    }
};
Info operator+(const Info &a, const Info &b) {
    return {(a.sum + b.sum) % m, a.1, b.r};
}
template <class Info, class Tag>
struct LazySegmentTree {
    int n;
    vector<Info> info;
    vector<Tag> tag;
    vector<Info> _init;
    LazySegmentTree() {}
    LazySegmentTree(int n) {
        _init.assign(n + 1, Info());
        init(_init);
```

```
LazySegmentTree(vector<Info> &_init) {
    this->_init = _init;
    init(_init);
}
void init(vector<Info> &_init) {
    n = (int)_init.size() - 1;
    info.assign(n * 4 + 1, Info());
    tag.assign(n * 4 + 1, Tag());
    build(1, 1, n);
}
void build(int p, int l, int r) {
    if(1 == r) {
        info[p] = _init[1];
        return ;
    }
    int m = (1 + r) / 2;
    build(p * 2, 1, m);
    build(p * 2 + 1, m + 1, r);
    pull(p);
};
void apply(int p, Tag &t) {
    info[p].apply(t);
    tag[p].apply(t);
void push(int p) {
    apply(p * 2, tag[p]);
    apply(p * 2 + 1, tag[p]);
    tag[p] = Tag();
void pull(int p) {
    info[p] = info[p * 2] + info[p * 2 + 1];
void modify(int p, int 1, int r, int x, int y, Tag &v) {
    if(1 > y \mid\mid r < x) return;
    if(1 >= x \& r <= y) {
        apply(p, v);
    }
    else {
        int m = (1 + r) / 2;
        push(p);
        modify(p * 2, 1, m, x, y, v);
        modify(p * 2 + 1, m + 1, r, x, y, v);
        pull(p);
    }
void modify(int x, int y, Tag &v) {
    modify(1, 1, n, x, y, v);
Info query(int p, int 1, int r, int x, int y) {
    if(1 > y \mid \mid r < x) return Info();
    if(x \ll 1 \& y \gg r)  {
        return info[p];
    }
    push(p);
    int m = (1 + r) / 2;
    return query(p * 2, 1, m, x, y) + query(p * 2 + 1, m + 1, r, x, y);
}
```

```
Info query(int 1, int r) {
    return query(1, 1, n, 1, r);
}
};
```

区间加+区间查询max abc397F

```
struct Tag {
   i64 \text{ add} = 0;
   void apply(Tag &t) {
        add += t.add;
   }
};
struct Info {
   i64 mx = 0, 1, r;
   void apply(Tag &t) {
        mx = mx + t.add;
    }
};
Info operator+(const Info &a, const Info &b) {
    return {max(a.mx, b.mx), a.1, b.r};
}
template <class Info, class Tag>
struct LazySegmentTree {
   int n;
    vector<Info> info;
    vector<Tag> tag;
    vector<Info> _init;
   LazySegmentTree() {}
    LazySegmentTree(int n) {
        _init.assign(n + 1, Info());
        init(_init);
    LazySegmentTree(vector<Info> &_init) {
        this->_init = _init;
        init(_init);
    void init(vector<Info> &_init) {
        n = (int)_init.size() - 1;
        info.assign(n * 4 + 1, Info());
        tag.assign(n * 4 + 1, Tag());
        build(1, 1, n);
    void build(int p, int l, int r) {
        if(1 == r) {
            info[p] = _init[1];
            return ;
        int m = (1 + r) / 2;
        build(p * 2, 1, m);
        build(p * 2 + 1, m + 1, r);
        pull(p);
    };
    void apply(int p, Tag &t) {
        info[p].apply(t);
        tag[p].apply(t);
    }
```

```
void push(int p) {
        apply(p * 2, tag[p]);
        apply(p * 2 + 1, tag[p]);
        tag[p] = Tag();
   }
   void pull(int p) {
        info[p] = info[p * 2] + info[p * 2 + 1];
   void modify(int p, int 1, int r, int x, int y, Tag &v) {
        if(1 > y \mid\mid r < x) return ;
        if(1 >= x \& r <= y) {
            apply(p, v);
        }
        else {
            int m = (1 + r) / 2;
            push(p);
            modify(p * 2, 1, m, x, y, v);
            modify(p * 2 + 1, m + 1, r, x, y, v);
            pull(p);
        }
   }
   void modify(int x, int y, Tag &v) {
        modify(1, 1, n, x, y, v);
   }
    Info query(int p, int 1, int r, int x, int y) {
        if(1 > y \mid \mid r < x) return Info();
        if(x \ll 1 \& y \gg r)  {
            return info[p];
        }
        push(p);
        int m = (1 + r) / 2;
        return query(p * 2, 1, m, x, y) + query(p * 2 + 1, m + 1, r, x, y);
   }
   Info query(int 1, int r) {
        return query(1, 1, n, 1, r);
   }
};
```

懒标记线段树(jiangly)

```
template<class Info, class Tag>
struct LazySegmentTree {
   int n;
   vector<Info> info;
   vector<Tag> tag;
   LazySegmentTree() : n(0) {}
   LazySegmentTree(int n_, Info v_ = Info()) {
      init(n_, v_);
   }
   template<class T>
   LazySegmentTree(vector<T> init_) {
      init(init_);
   }
   void init(int n_, Info v_ = Info()) {
      init(vector(n_, v_));
   }
   template<class T>
```

```
void init(vector<T> init_) {
        n = init_.size();
        info.assign(4 << __lg(n), Info());</pre>
        tag.assign(4 << lll lg(n), Tag());
        function<void(int, int, int)> build = [&](int p, int l, int r) {
            if (r - 1 == 1) {
                info[p] = init_[1];
                return;
            }
            int m = (1 + r) / 2;
            build(2 * p, 1, m);
            build(2 * p + 1, m, r);
            pull(p);
        };
        build(1, 0, n);
    }
    void pull(int p) {
        info[p] = info[2 * p] + info[2 * p + 1];
    }
    void apply(int p, const Tag &v) {
        info[p].apply(v);
        tag[p].apply(v);
    }
    void push(int p) {
        apply(2 * p, tag[p]);
        apply(2 * p + 1, tag[p]);
        tag[p] = Tag();
    void modify(int p, int 1, int r, int x, const Info \&v) {
        if (r - 1 == 1) {
            info[p] = v;
            return;
        int m = (1 + r) / 2;
        push(p);
        if (x < m) {
            modify(2 * p, 1, m, x, v);
        } else {
            modify(2 * p + 1, m, r, x, v);
        }
        pull(p);
    void modify(int p, const Info &v) {
        modify(1, 0, n, p, v);
    Info rangeQuery(int p, int 1, int r, int x, int y) {
        if (1 >= y || r <= x) {
            return Info();
        }
        if (1 >= x & r <= y) {
            return info[p];
        int m = (1 + r) / 2;
        push(p);
        return rangeQuery(2 * p, 1, m, x, y) + rangeQuery(2 * p + 1, m, r, x,
y);
    Info rangeQuery(int 1, int r) {
```

```
return rangeQuery(1, 0, n, 1, r);
}
void rangeApply(int p, int l, int r, int x, int y, const Tag &v) {
    if (1 >= y || r <= x) {
        return;
    }
    if (1 >= x & r <= y) {
        apply(p, v);
        return;
    int m = (1 + r) / 2;
    push(p);
    rangeApply(2 * p, 1, m, x, y, v);
    rangeApply(2 * p + 1, m, r, x, y, v);
    pull(p);
}
void rangeApply(int 1, int r, const Tag &v) {
    return rangeApply(1, 0, n, l, r, v);
}
template<class F>
int findFirst(int p, int 1, int r, int x, int y, F pred) {
    if (1 >= y || r <= x || !pred(info[p])) {
        return -1;
    }
    if (r - 1 == 1) {
        return 1;
    }
    int m = (1 + r) / 2;
    push(p);
    int res = findFirst(2 * p, 1, m, x, y, pred);
    if (res == -1) {
        res = findFirst(2 * p + 1, m, r, x, y, pred);
    }
    return res;
}
template<class F>
int findFirst(int 1, int r, F pred) {
    return findFirst(1, 0, n, 1, r, pred);
}
template<class F>
int findLast(int p, int 1, int r, int x, int y, F pred) {
    if (1 >= y || r <= x || !pred(info[p])) {
        return -1;
    }
    if (r - 1 == 1) {
        return 1;
    int m = (1 + r) / 2;
    push(p);
    int res = findLast(2 * p + 1, m, r, x, y, pred);
    if (res == -1) {
        res = findLast(2 * p, 1, m, x, y, pred);
    }
    return res;
}
template<class F>
int findLast(int 1, int r, F pred) {
    return findLast(1, 0, n, 1, r, pred);
```

```
};
struct Tag {
   int add = 0;
   void apply(Tag t) & {
        add += t.add;
   }
};
struct Info {
   int max = 0;
   void apply(Tag t) & {
       max += t.add;
   }
};
Info operator+(const Info &a, const Info &b) {
   return {max(a.max, b.max)};
}
```

区间修改区间查询

```
struct node {
   int 1, r;
   i64 sum, add, mul;
    #define lc u << 1
    #define rc u << 1 | 1
}tr[N << 2];</pre>
void pushup(int u) {
    tr[u].sum = tr[lc].sum + tr[rc].sum;
void eval(node &t, i64 add, i64 mul) {
    t.sum = t.sum * mul + (t.r - t.l + 1) * add;
    t.mul = t.mul * mul;
   t.add = t.add * mul + add;
void pushdown(int u) {
    eval(tr[lc], tr[u].add, tr[u].mul);
    eval(tr[rc], tr[u].add, tr[u].mul);
    tr[u].add = 0, tr[u].mul = 1;
void build(int u, int 1, int r) {
   if(1 == r) tr[u] = \{1, r, a[r], 0, 1\};
    else {
        tr[u] = \{1, r, 0, 0, 1\};
        int mid = 1 + r \gg 1;
        build(lc, l, mid), build(rc, mid + 1, r);
        pushup(u);
    }
void modify(int u, int 1, int r, i64 add, i64 mul) {
    if(tr[u].l >= l \&\& tr[u].r <= r) eval(tr[u], add, mul);
    else {
        pushdown(u);
        int mid = tr[u].l + tr[u].r >> 1;
        if(1 \le mid) modify(1c, 1, r, add, mul);
```

```
if(r > mid) modify(rc, l, r, add, mul);
    pushup(u);
}

int query(int u, int l, int r) {
    if(tr[u].l >= l && tr[u].r <= r) return tr[u].sum;
    pushdown(u);
    int mid = tr[u].l + tr[u].r >> 1;
    i64 sum = 0;
    if(l <= mid) sum = sum + query(lc, l, r);
    if(r > mid) sum = sum + query(rc, l, r);
    return sum;
}
```

将一段区间全部改为一个值

```
struct node {
    int 1, r, add, v;
#define lc u << 1
#define rc u << 1 | 1
} tr[N << 2];</pre>
void pushup(int u) {
    tr[u].v = (tr[]c].v + tr[rc].v);
void build(int u, int 1, int r) {
    tr[u] = \{1, r, -INT\_MAX, 0\};
    if(1 == r) {
        return ;
    }
    int mid = 1 + r \gg 1;
    build(lc, 1, mid);
    build(rc, mid + 1, r);
    pushup(u);
void pushdown(int u) {
    if(tr[u].add == -INT_MAX) return ;
    tr[lc].v = tr[u].add * (tr[lc].r - tr[lc].l + 1), tr[rc].v = tr[u].add *
(tr[rc].r - tr[rc].l + 1);
    tr[lc].add = tr[u].add, tr[rc].add = tr[u].add;
    tr[u].add = -INT_MAX;
void modify(int u, int 1, int r, int d) {
    if (tr[u].1 >= 1 && tr[u].r <= r) {
        tr[u].add = d;
        tr[u].v = d * (tr[u].r - tr[u].l + 1);
    } else {
        pushdown(u);
        int mid = tr[u].l + tr[u].r >> 1;
        if (1 \le mid) \mod ify(u \le 1, 1, r, d);
        if (r > mid) modify(u << 1 | 1, 1, r, d);
        pushup(u);
    }
int query(int u, int 1, int r) {
    if(tr[u].l >= l \&\& tr[u].r <= r) return tr[u].v;
    pushdown(u);
    int mid = tr[u].l + tr[u].r >> 1;
```

```
int res = 0;
if(1 <= mid) res += query(1c, 1, r);
if(r > mid) res += query(rc, 1, r);
pushup(u);
return res;
}
```

势能线段树

```
struct node {
    int 1, r;
    i64 sum;
    bool op;
    #define ls u << 1
    #define rs u << 1 | 1
}tr[N << 2];</pre>
void pushup(int u) {
    tr[u].sum = tr[ls].sum + tr[rs].sum;
    tr[u].op = (tr[ls].op & tr[rs].op);
void build(int u, int 1, int r) {
    if(1 == r) tr[u] = \{1, r, a[1]\};
    else {
        tr[u] = \{1, r, 0\};
        int mid = 1 + r \gg 1;
        build(ls, l, mid), build(rs, mid + 1, r);
        pushup(u);
    }
}
void modify(int u, int 1, int r) {
    if(tr[u].l >= l && tr[u].r <= r && tr[u].l == tr[u].r && !tr[u].op) {
    }
    else if(tr[u].op) {
        return ;
    }
    else {
        int mid = tr[u].l + tr[u].r >> 1;
        if(1 \le mid) modify(1s, 1, r);
        if(r > mid) modify(rs, 1, r);
        pushup(u);
    }
}
i64 query(int u, int 1, int r) {
    if(tr[u].l >= l \&\& tr[u].r <= r) return tr[u].sum;
    int mid = tr[u].l + tr[u].r >> 1;
    i64 \text{ sum} = 0;
    if(1 \le mid) sum = sum + query(1s, 1, r);
    if(r > mid) sum = sum + query(rs, 1, r);
    return sum;
}
```

线段树合并

```
int cur;
int root[N], ls[N], rs[N], sum[N], mxz[N];
void pushup(int u) {
   if(sum[]s[u]] >= sum[rs[u]])
        sum[u] = sum[]s[u]], mxz[u] = mxz[]s[u]];
    else
        sum[u] = sum[rs[u]], mxz[u] = mxz[rs[u]];
void update(int &u, int 1, int r, int z, int val) {
    if(!u) u = ++cur;
    if(1 == r) {
        sum[u] += val;
        mxz[u] = z;
        return ;
    int mid = 1 + r \gg 1;
    if(z \leftarrow mid) update(ls[u], l, mid, z, val);
    else update(rs[u], mid + 1, r, z, val);
    pushup(u);
int merge(int u, int v, int 1, int r) {
   if(!u || !v) return u + v;
    if(1 == r) {
        sum[u] += sum[v];
        return u;
    }
    int mid = 1 + r \gg 1;
    ls[u] = merge(ls[u], ls[v], l, mid);
    rs[u] = merge(rs[u], rs[v], mid + 1, r);
    pushup(u);
    return u;
}
```

线段树分治

CF1814F需要维护每个点的tag

```
int tag[N];
vector<array<int, 2>> sgt[N << 2];</pre>
struct DSU {
    int n;
    vector<int> fa, sz;
    stack<array<int, 2>> stk;
    void init(int n) {
        this->n = n;
        sz.assign(n + 1, 1);
        fa.resize(n + 1);
        iota(fa.begin(), fa.end(), 0);
    int find(int x) {
        while(fa[x] != x) x = fa[x];
        return x;
    }
    void merge(int x, int y) {
```

```
int fx = find(x), fy = find(y);
        if(fx != fy) {
            if(sz[fx] > sz[fy]) swap(fx, fy);
            sz[fy] += sz[fx];
            fa[fx] = fy;
            stk.push({fx, fy});
            //tag[fx] -= tag[fy];
        }
    }
    void undo(int last) {
        while((int)stk.size() > last) {
            auto [x, y] = stk.top();
            stk.pop();
            sz[y] = sz[x];
            fa[x] = x;
            //tag[x] += tag[y];
    }
};
//按秩合并并查集、可撤销并查集
DSU dsu;
void insert(int u, int 1, int r, int q1, int qr, int x, int y) {
    if(q1 \ll 1 \& qr >= r) {
        sgt[u].push_back({x, y});
        return ;
    }
    int m = 1 + r >> 1;
    if(ql \ll m) insert(u * 2, 1, m, ql, qr, x, y);
    if(qr > m) insert(u * 2 + 1, m + 1, r, ql, qr, x, y);
}
void dfs_sgt(int u, int 1, int r) {
    int last = dsu.stk.size();
    for(auto [x, y] : sgt[u]) {
        dsu.merge(x, y);
    if(1 == r) {
        //tag[dsu.find(1)] += 1;
        return ;
    }
    else {
        int m = 1 + r >> 1;
        dfs_sgt(u * 2, 1, m);
        dfs_sgt(u * 2 + 1, m + 1, r);
    }
    dsu.undo(last);
}
```

可持久化线段树

```
int cur, root[N], ls[N << 5], rs[N << 5];
struct Info {
   int v;
} sgt[N << 5];
Info operator+ (const Info &a, const Info &b) {</pre>
```

```
Info c;
    c.v = a.v + b.v;
    return c:
}
void pushup(int u) {
    sgt[u] = sgt[ls[u]] + sgt[rs[u]];
}
void build(int \&u, int 1, int r) {
    u = ++cur;
    if (1 == r) {
        sgt[u] = \{0\};
        return ;
    }
    int m = 1 + r >> 1;
    build(ls[u], 1, m);
    build(rs[u], m + 1, r);
    pushup(u);
void modify(int pr, int &u, int 1, int r, int p) {
    u = ++cur, ls[u] = ls[pr], rs[u] = rs[pr];
    if (1 == r) {
        sgt[u] = \{1\};
        return ;
    }
    int m = 1 + r >> 1;
    if (p <= m) modify(ls[pr], ls[u], l, m, p);
    else modify(rs[pr], rs[u], m + 1, r, p);
    pushup(u);
}
Info query(int u, int 1, int r, int q1, int qr) {
   if(qr < 1 || q1 > r) return Info();
    if(q1 \ll 1 \& qr >= r) {
        return sgt[u];
    }
    int m = 1 + r >> 1;
    Info ans = Info();
    if(q1 \ll m) {
        ans = query(ls[u], 1, m, q1, qr);
    }
    if(qr > m) {
        ans = ans + query(rs[u], m + 1, r, ql, qr);
    return ans;
}
int find_first(int u, int 1, int r, int q1, int qr, int k) {
   if(1 == r) {
        if(k == 1) return 1;
        else return -1;
    }
    int m = 1 + r >> 1;
    if(sgt[ls[u]].v >= k) return find_first(ls[u], 1, m, q1, qr, k);
    else return find_first(rs[u], m + 1, r, ql, qr, k - sgt[ls[u]].v);
//动态开点线段树上二分,全局从左到右第k个,q7等于1
void solve() {
    build(root[0], 1, n); //注意范围, 右边界可能不是n
    for (int i = 1; i <= n; i++) {
```

```
modify(root[i - 1], root[i], 1, n, p);
}
//同样需要注意是不是1到n
}
```

指针版

主席树维护从root到各个节点的哈希值,然后在主席树上二分求出前k小

CF1957F2

```
struct Node {
    Node *1 = NULL;
    Node *r = NULL;
    i64 hash = 0;
} tr[N << 5];</pre>
int tot = 0;
Node *null = &tr[0];
Node *newNode() {
    return &tr[++tot];
}
const int B = 114514;
const i64 P = i64(6e17) + 7;
i64 pw[N];
void add(Node *&t, int 1, int r, int x) {
    t = new Node {*t};
    t->hash = (t->hash + pw[x]) % P;
    if (1 == r) {
        return;
    }
    int m = (1 + r) / 2;
    if (x \ll m) {
        add(t->1, 1, m, x);
    } else {
        add(t\rightarrow r, m + 1, r, x);
    }
void find(int 1, int r, Node *u1, Node *v1, Node *l1, Node *u2, Node *v2, Node
*12) {
    if(k == 0) return ;
    i64 \text{ hash1} = (((u1-)hash + v1-)hash - 2 * 11-)hash) % P + P) % P + (1 <= al1)
&& all \leftarrow r ? pw[all] : 0)) % P;
    i64 \text{ hash2} = (((u2-)\text{hash} + v2-)\text{hash} - 2 * 12-)\text{hash}) \% P + P) \% P + (1 <= a12)
&& al2 <= r ? pw[al2] : 0)) % P;
    if(hash1 == hash2) return ;
    if(1 == r) {
        ans.push_back(1);
        k--;
        return ;
    }
    int m = (1 + r) / 2;
    find(1, m, u1->1, v1->1, 11->1, u2->1, v2->1, 12->1);
    find(m + 1, r, u1->r, v1->r, l1->r, u2->r, v2->r, l2->r);
void solve() {
    null->l = null->r = null;
    vector<Node *> node(n + 1, null);
    function<void(int, int)> dfs = [&](int u, int fa) {
```

```
add(node[u], 1, N - 1, a[u]);
        for(int v : g[u]) {
            if(v == fa) continue;
            node[v] = node[u];
            dfs(v, u);
        }
    };
    while(q--) {
        cin >> u1 >> v1 >> u2 >> v2 >> k;
        11 = hld.lca(u1, v1), 12 = hld.lca(u2, v2);
        al1 = a[11], al2 = a[12];
        find(1, N - 1, node[u1], node[v1], node[l1], node[u2], node[v2],
node[12]);
        cout << ans.size() << " ";</pre>
        for(auto x : ans) cout << x << ' ';</pre>
        cout << '\n';</pre>
        ans.clear();
    }
}
```

treap

```
const int N = 100010, INF = 1e9;
int n;
struct Node {
   int 1, r;
   int key, val;
   int cnt, size;
} tr[N];
int root, idx;
void pushup(int p) {
    tr[p].size = tr[tr[p].1].size + tr[tr[p].r].size + tr[p].cnt;
}
int get_node(int key) {
    tr[ ++ idx].key = key;
    tr[idx].val = rand();
    tr[idx].cnt = tr[idx].size = 1;
    return idx;
}
void zig(int &p) { // 右旋
   int q = tr[p].1;
    tr[p].l = tr[q].r, tr[q].r = p, p = q;
    pushup(tr[p].r), pushup(p);
void zag(int &p) { // 左旋
   int q = tr[p].r;
    tr[p].r = tr[q].1, tr[q].1 = p, p = q;
    pushup(tr[p].1), pushup(p);
void build() {
    get_node(-INF), get_node(INF);
    root = 1, tr[1].r = 2;
    pushup(root);
    if (tr[1].val < tr[2].val) zag(root);</pre>
}
```

```
void insert(int &p, int key) {
    if (!p) p = get_node(key);
    else if (tr[p].key == key) tr[p].cnt++;
    else if (tr[p].key > key) {
        insert(tr[p].1, key);
        if (tr[tr[p].1].val > tr[p].val) zig(p);
    } else {
        insert(tr[p].r, key);
        if (tr[tr[p].r].val > tr[p].val) zag(p);
   pushup(p);
void remove(int &p, int key) {
   if (!p) return;
   if (tr[p].key == key) {
       if (tr[p].cnt > 1) tr[p].cnt--;
        else if (tr[p].1 || tr[p].r) {
            if (!tr[p].r || tr[tr[p].1].val > tr[tr[p].r].val) {
                remove(tr[p].r, key);
            } else {
               zag(p);
                remove(tr[p].1, key);
           }
        } else p = 0;
    } else if (tr[p].key > key) remove(tr[p].1, key);
    else remove(tr[p].r, key);
    pushup(p);
}
int get_rank_by_key(int p, int key) { // 通过数值找排名
   if (!p) return 0; // 本题中不会发生此情况
    if (tr[p].key == key) return tr[tr[p].l].size + 1;
   if (tr[p].key > key) return get_rank_by_key(tr[p].1, key);
    return tr[tr[p].1].size + tr[p].cnt + get_rank_by_key(tr[p].r, key);
int get_key_by_rank(int p, int rank) { // 通过排名找数值
   if (!p) return INF; // 本题中不会发生此情况
   if (tr[tr[p].1].size >= rank) return get_key_by_rank(tr[p].1, rank);
   if (tr[tr[p].1].size + tr[p].cnt >= rank) return tr[p].key;
    return get_key_by_rank(tr[p].r, rank - tr[tr[p].1].size - tr[p].cnt);
}
int get_prev(int p, int key) { // 找到严格小于key的最大数
   if (!p) return -INF;
    if (tr[p].key >= key) return get_prev(tr[p].1, key);
    return max(tr[p].key, get_prev(tr[p].r, key));
}
int get_next(int p, int key) { // 找到严格大于key的最小数
   if (!p) return INF;
    if (tr[p].key <= key) return get_next(tr[p].r, key);</pre>
    return min(tr[p].key, get_next(tr[p].1, key));
int main() {
   build();
    cin >> n;
   while (n--) {
       int op, x;
        cin >> op >> x;
        if (op == 1) insert(root, x);
```

```
else if (op == 2) remove(root, x);
else if (op == 3) cout << get_rank_by_key(root, x) - 1 << '\n';
else if (op == 4) cout << get_key_by_rank(root, x + 1) << '\n';
else if (op == 5) cout << get_prev(root, x) << '\n';
else cout << get_next(root, x) << '\n';
}
return 0;
}</pre>
```

CDQ分治

三位偏序, 陌上花开

有n朵花,每朵花有三个属性:花形(s),颜色(c),气味(m),用三个整数表示。

现在要对每朵花评级,一朵花的级别是它拥有的美丽能超过的花的数量。

定义一朵花 A 比另一朵花 B 要美丽,当且仅 $s_a \geq s_b, c_a \geq c_b, m_a \geq m_b$ 。

显然,两朵花可能有同样的属性。需要统计出评出每个等级的花的数量。

离线

```
constexpr int N = 2e5 + 10;
array<int, 5> a[N], tmp[N];
template <typename T>
struct Fenwick {
   int n:
    vector<T> a;
    Fenwick(int _n = 0) : a(_n + 1), n(_n){}
    void init(int n) {
        a.assign(n + 1, T());
        this->n = n;
    void add(int x, T v) {
        for (int i = x; i \le n; i += i \& -i) {
            a[i] += v;
        }
    }
    T sum(int x) {
        auto ans = T();
        for (int i = x; i >= 1; i -= i \& -i) {
            ans += a[i];
        }
        return ans;
    T rangeSum(int 1, int r) {
        return sum(r) - sum(1 - 1);
    }
};
void solve() {
   int n, k;
    cin >> n >> k;
    Fenwick<int> fen(k);
    for(int i = 1; i \le n; i++) {
        cin >> a[i][0] >> a[i][1] >> a[i][2];
        a[i][3] = 1;
```

```
sort(a + 1, a + 1 + n);
    int t = 0;
    for(int i = 1; i \le n; i++) {
        if(i == 1 || a[i][0] != a[t][0] || a[i][1] != a[t][1] || a[i][2] != a[t]
[2]) {
            a[++t] = a[i];
        }
        else {
            a[t][3]++;
        }
    }//例题题目要求去重
    function<void(int, int)> cdq = [\&](int 1, int r) {
        if(1 == r) return ;
        int mid = 1 + r \gg 1;
        cdq(1, mid), cdq(mid + 1, r);
        int p1 = 1, p2 = mid + 1;
        int p3 = 0;
        while(p1 <= mid \mid \mid p2 <= r) {
            if(p2 > r || (p1 <= mid && make_pair(a[p1][1], a[p1][2]) <=
make_pair(a[p2][1], a[p2][2]))) {
                fen.add(a[p1][2], a[p1][3]);
                tmp[++p3] = a[p1++];
            }//左半部分
            else {
                a[p2][4] += fen.sum(a[p2][2]);
                tmp[++p3] = a[p2++];
            }//右半部分
        }
        for(int i = 1; i \leftarrow mid; i++) fen.add(a[i][2], -a[i][3]);
        for(int i = 1; i \le p3; i++) a[1 + i - 1] = tmp[i];
    };
    cdq(1, t);
    vector<int> ans(n);
    for(int i = 1; i <= t; i++) {
        ans[a[i][3] + a[i][4] - 1] += a[i][3];
    }
   for(int i = 0; i < n; i++) cout << ans[i] << '\n';
}
```

字符串

KMP

```
vector<int> kmp(string a) {
    int n = a.size() - 1; //开头有空格
    vector<int> f(n + 1);
    for(int i = 2, j = 0; i <= n; i++) {
        while(j && a[j + 1] != a[i]) j = f[j];
        if(a[j + 1] == a[i]) f[i] = ++j;
    }
    return f;
}</pre>
```

Z函数 (拓展KMP)

约定:字符串下标以0为起点。

对于个长度为 n 的字符串 s 。定义函数 z[i] 表示 s 和 s[i,n-1] (即以 s[i] 开头的后缀)的最长公共前缀(LCP)的长度。 z 被称为 s 的 **Z 函数**。特别地,z[0]=0。1

```
vector<int> z_function(string &s) {
   int n = (int)s.length();
   vector<int> z(n);
   for (int i = 1, l = 0, r = 0; i < n; ++i) {
        if (i \le r \&\& z[i - 1] < r - i + 1) {
            z[i] = z[i - 1];
        } else {
            z[i] = max(0, r - i + 1);
            while (i + z[i] < n \& s[z[i]] == s[i + z[i]])  {
                ++z[i];
            }
        }
        if (i + z[i] - 1 > r) {
           l = i, r = i + z[i] - 1;
        }
   }
    return z;
}
```

字符串哈希

```
struct Hash {
   int Base, P;//常用模数1e9+7,998244353,常用base 131,13331
   vector<int> power, sum;
   void init(int base, int p, string &s) {
        //s需要从1开始
        power.resize(s.size()), sum.resize(s.size());
        P = p, Base = base, power[0] = 1;
        for(int i = 1; i < s.size(); i++) {
            power[i] = 1]] * power[i - 1] * Base % P;
            sum[i] = (111 * sum[i - 1] * Base + s[i]) % P;
        }
   }
   int gethash(int 1, int r) {
        return (sum[r] - 1]  * sum[l - 1]  * power[r - l + 1]  % P + P)  % P;
};
mt19937_64 rng(chrono::steady_clock::now().time_since_epoch().count());
int P1 = rng() % (int)1e9 + 1, P2 = rng() % (int)1e9 + 1;//可能会很慢
```

manacher

```
auto manacher(string &s) {
    string t = "$#";
    for(auto c : s) t += c, t += '#';
    t += '^';
    int n = t.size();
    vector<int> p(n);
```

```
int r = 0, mid;
for(int i = 1; i < n; i++) {
    if(i <= r) p[i] = min(p[2 * mid - i], r - i + 1);
    else p[i] = 1;
    while(t[i - p[i]] == t[i + p[i]]) p[i] += 1;
    if(i + p[i] - 1 > r) r = i + p[i] - 1, mid = i;
}
return p;
}
int res = 0;
for(int i = 1; i < p.size(); i++) {
    res = max(res, p[i] - 1);
    int l = i - p[i] + 2, r = i + p[i] - 2;//以第i个字符为中心的最长回文边界if(l > r) continue;
}//最长回文串长度
```

AC自动机

```
struct AC_automation {
   vector<vector<int>> tr, g;
   vector<int> val, fail;
   int idx = 0;
   vector<int> T;
   AC_automation(int m) : val(1), fail(1), T(m, 0){
        tr.emplace_back(T);
   };
   void insert(string &s, int v) {
        int p = 0;
        for (int i = 0; i < s.size(); i++) {
           int u = s[i] - 'a';
            if (!tr[p][u]) {
                tr[p][u] = ++idx;
                tr.emplace_back(T);
                val.emplace_back(0);
                fail.emplace_back(0);
           p = tr[p][u];
        }
       val[p] += v;
   }
   void build() {
        queue<int> q;
        for(int i = 0; i < 26; i++)
            if(tr[0][i])
                q.push(tr[0][i]);
        while(!q.empty()) {
           int u = q.front();
            q.pop();
            for(int i = 0; i < 26; i++) {
                int v = tr[u][i];
                if(!v)
                    tr[u][i] = tr[fail[u]][i];
                else {
                    fail[v] = tr[fail[u]][i];
                    q.push(v);
                }
```

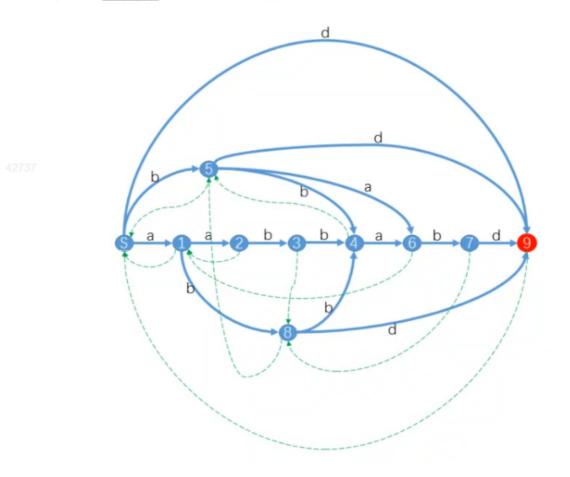
```
}
    int find(string &s) {
        int ans = 0;
        for(int i = 0, j = 0; i < s.size(); i++) {
            int u = s[i] - 'a';
            j = tr[j][u];
            int p = j;
            while(p \&\& val[p] >= 0) {
                ans += val[p];
                val[p] = -1;//优化为trie图
                p = fail[p];
            }
        }
        return ans;
    }
    void addfail() {
        g.resize(idx + 1);
        for(int i = 1; i \le idx; i++) {
            g[fail[i]].push_back(i);
        }
    }
};
```

后缀数组

```
int t, n, m = 122, a[N], sa[N], c[N], x[N], y[N], rk[N], height[N];
string s;
void get_sa() {
    for(int i = 1; i \le n; i++) c[x[i] = s[i]]++;
    for(int i = 2; i \le m; i++) c[i] += c[i - 1];
    for(int i = n; i >= 1; i--) sa[c[x[i]]--] = i;
    for(int k = 1; k <= n; k <<= 1) {
        int num = 0;
        for(int i = n - k + 1; i \le n; i++) y[++num] = i;
        for(int i = 1; i \le n; i++) if(sa[i] > k) y[++num] = sa[i] - k;
        for(int i = 1; i \le m; i++) c[i] = 0;
        for(int i = 1; i \le n; i++) c[x[i]]++;
        for(int i = 2; i \le m; i++) c[i] += c[i - 1];
        for(int i = n; i >= 1; i--) sa[c[x[y[i]]]--] = y[i], y[i] = 0;
        for(int i = 1; i \le n; i++) swap(x[i], y[i]);
        x[sa[1]] = 1, num = 1;
        for(int i = 2; i \le n; i++) x[sa[i]] = (y[sa[i]] == y[sa[i - 1]] &&
y[sa[i] + k] == y[sa[i - 1] + k])? num : ++num;
        if(num == n) break;
        m = num;
    }
}
void get_height() {
    for(int i = 1; i \le n; i++) rk[sa[i]] = i;
    for(int i = 1, k = 0; i \le n; i++) {
        if(rk[i] == 1) continue;
        if(k) k--;
        int j = sa[rk[i] - 1];
        while(i + k <= n && j + k <= n && s[i + k] == s[j + k]) k++;
        height[rk[i]] = k;
    }
```

后缀自动机

比如对于字符串 S = aabbabd, 它的后缀自动机是



```
const int N = 1e6 + 10; //endpoint的最大个数是2*n级别, N要开两倍
string s;
int tot = 1, last = 1;
int cnt[N], len[N], p[N], sum[N]; //cnt是这个子串的数量, sum是以这个子串为前缀的数量
struct Node {
    int len, fa;
   int ch[26];
} node[N];
void extend(int c) {
   int p = last, np = last = ++tot;
   cnt[tot] = 1;
   node[np].len = node[p].len + 1;
    for(; p && !node[p].ch[c]; p = node[p].fa) node[p].ch[c] = np;
   if (!p) node[np].fa = 1;
   else {
       int q = node[p].ch[c];
       if (node[q].len == node[p].len + 1) node[np].fa = q;
       else {
           int nq = ++tot;
           node[nq] = node[q], node[nq].len = node[p].len + 1;
           node[q].fa = node[np].fa = nq;
           for(; p \& node[p].ch[c] == q; p = node[p].fa) node[p].ch[c] = nq;
       }
   }
}
```

```
int h[N], e[N], ne[N], idx;
void add(int a, int b) {
    e[idx] = b, ne[idx] = h[a], h[a] = idx++;
}
void build() {
    for(char c : s) extend(c - 'a');
    for(int i = 1; i <= tot; i++) len[node[i].len]++;</pre>
    for(int i = 1; i <= s.size(); i++) len[i] += len[i - 1];
    for(int i = 1; i <= tot; i++) p[len[node[i].len]--] = i;
    for(int i = tot; i >= 1; i--) cnt[node[p[i]].fa] += cnt[p[i]];
    sum[1] = cnt[1] = 0;
    for(int i = 1; i \leftarrow tot; i++)
        sum[i] = 1;
    for(int i = tot; i >= 1; i--) {
        for(int j = 0; j < 26; j++) {
            if(node[p[i]].ch[j]) {
                sum[p[i]] += sum[node[p[i]].ch[j]];
            }
        }
    for(int i = 1; i <= tot; i++) {
        node[i].len = node[i].fa = cnt[i] = p[i] = len[i] = 0;
        for(int j = 0; j < 26; j++)
            node[i].ch[j] = 0;
    }//清空
}
```

回文自动机

```
string s;
int len[N], fail[N], tr[N][26], tot, cur, cnt[N];
void init() {
    for(int i = 0; i <= tot; i++)
        for(int j = 0; j < 26; j++)
            tr[i][j] = 0;
    fail[0] = 1, len[1] = -1, tot = 1, cur = 0;
}
int getfail(int x, int i) {
    while(s[i - len[x] - 1] != s[i])
        x = fail[x];
    return x;
}
void insert() {
    for(int i = 0; i < s.size(); i++) {
        int p = getfail(cur, i), u = s[i] - 'a';
        if(!tr[p][u]) {
            fail[++tot] = tr[getfail(fail[p], i)][u];
            tr[p][u] = tot;
            len[tot] = len[p] + 2;
        cnt[cur = tr[p][u]]++;
    for(int i = tot; i >= 2; i--)
        cnt[fail[i]] += cnt[i];
}
void solve() {
```

```
cin >> s;
init();
insert();
}
```

数学知识

线性筛质数

```
int p[N], minp[N], pcnt;
void get_prime(int n) {
    for(int i = 2; i <= n; i++) {
        if(!minp[i]) {
            minp[i] = i;
            p[++pcnt] = i;
        }
    for(int j = 1; p[j] <= n / i && j <= pcnt; j++) {
            minp[i * p[j]] = p[j];
            if(i % p[j] == 0) break;
        }
    }
}</pre>
```

快速幂

```
i64 power(i64 a, i64 b) {
    a %= P;
    i64 res = 1;
    while (b) {
        if (b & 1) res = res * a % P;
        a = a * a % P;
        b >>= 1;
    }
    return res % P;
}
```

最大公约数

```
int gcd(int a, int b) {
    return b ? gcd(b, a % b) : a;
}
```

拓展欧几里得

解出来的 x 和 y 是特解,通解为 $x'=x+k\times (b/gcd(a,b)), y'=y-k\times (a/gcd(a,b)),$ k是整数

```
i64 exgcd(i64 a, i64 b, i64 &x, i64 &y) {
    if(!b) {
        x = 1;
        y = 0;
        return a;
    }
    else {
        i64 g = exgcd(b, a % b, y, x);
        y -= a / b * x;
        return g;
    }
}
```

中国剩余定理

```
i64 exgcd(i64 a, i64 b, i64 &x, i64 &y) {
    if(!b) {
        x = 1;
        y = 0;
        return a;
   }
   else {
        i64 g = exgcd(b, a \% b, y, x);
        y -= a / b * x;
        return g;
    }
int m[N], a[N];
i64 CRT(){
   int n;
   cin >> n;
   i64 \, mul = 1;
    for(int i = 1; i <= n; i++) {
        cin >> m[i] >> a[i];
        mul *= m[i];
    }
    i64 res = 0;
    for(int i = 1; i <= n; i++) {
        i64 M = mul / m[i];
        i64 x, y;
        exgcd(M, m[i], x, y);
        res = (res + a[i] * M * ((x % m[i] + m[i]) % m[i]));
   return res % mul;
}
```

拓展中国剩余定理

```
else {
        i64 g = exgcd(b, a \% b, y, x);
       y -= a / b * x;
        return g;
   }
}
i64 m[N], a[N];
i64 excRT() {
   int n;
    cin >> n;
    for(int i = 1; i \le n; i++) cin >> m[i] >> a[i];
    i64 res = a[1], p = m[1];
    for(int i = 2; i <= n; i++) {
        i64 x, y;
        \_int128 b = (a[i] - res % m[i] + m[i]) % m[i];
        i64 g = exgcd(p, m[i], x, y);
        if(b % g) return -1;//无解
        i64 k = m[i] / g;
        x = (x \% k * (b / g \% k) \% k + k) % k; // k 很大时会爆11, 所以b需要开__int128
        res += x * p;
        p *= k;
    }
   return res;
}
```

约数个数

```
egin{aligned} N &= p_1{}^{a_1} * p_2{}^{a_2} \cdots * p_k{}^{a_k} \ \\ ans &= (a_1+1)(a_2+1) \cdots (a_k+1) \end{aligned}
```

```
using i64 = long long;
const int P = 1e9 + 7;
int n, a[110];
unordered_map <int, int> primes;
int main(){
    scanf("%d", &n);
    for(int i = 1; i <= n; i++) {
        scanf("%d", &a[i]);
        for(int j = 2; j \le a[i] / j; j++) {
            while(a[i] \% j == 0) {
                a[i] /= j;
                primes[j] ++;
            }
        if(a[i] != 1) primes[a[i]] ++;
    i64 \text{ ans} = 1;
    for(auto prime: primes) {
        ans = ans * (prime.second + 1) % P;
    printf("%11d", ans);
    return 0;
}
```

约数之和

```
(p_1^0 + p_1^1 + \ldots + p_1^{c1}) * \ldots * (p_k^0 + p_k^1 + \ldots + p_k^{ck})
```

```
using i64 = long long;
const int mod = 1e9 + 7;
int n;
unordered_map <int, int> primes;
int main(){
    scanf("%d", &n);
    while(n--) {
        int x;
        scanf("%d", &x);
        for(int i = 2; i \le x / i; i++) {
            while(x \% i == 0) {
                primes[i]++;
                x /= i;
        }
        if(x > 1) primes[x]++;
    }
    i64 res = 1;
    for(auto i: primes) {
        i64 t = 1;
        i64 a = i.first, b = i.second;
        while(b--) {
            t = (t * a + 1) \% mod;
        res = res * t % mod;
    }
    printf("%11d", res);
    return 0;
}
```

dfs求约数 (利用大数约数数量大约为 $n^{1/3}$ 的结论)

```
function<void(int, i64)> dfs = [&](int u, i64 num) {
    if(u == (int)p.size()) {

        return ;
    }
    i64 res = num;
    for(int i = 0; i <= p[u].second; i++) {
        dfs(u + 1, res);
        res *= p[u].first;
    }
};</pre>
```

欧拉函数

```
int prime[N], phi[N];
bool st[N];
void get_euler(int n) {
    phi[1] = 1;
    for(int i = 2; i <= n; i++) {
        if(!st[i]) {
            prime[++m] = i;
            phi[i] = i - 1;
        for(int j = 1; prime[j] <= n / i; j++) {
            st[i * prime[j]] = 1;
            if(i % prime[j] == 0) {
                phi[i * prime[j]] = prime[j] * phi[i];
                break;
            phi[i * prime[j]] = phi[i] * (prime[j] - 1);
        }
   }
}
```

莫比乌斯反演

莫比乌斯反演是数论中的重要内容。对于一些函数 f(n),如果很难直接求出它的值,而容易求出其倍数和或约数和 g(n),那么可以通过莫比乌斯反演简化运算,求得 f(n) 的值。

性质

莫比乌斯函数不仅是积性函数,还有如下性质:

$$\sum_{d|n} \mu(d) = \begin{cases} 1 & n=1 \\ 0 & n \neq 1 \end{cases}$$

即
$$\sum_{d|n} \mu(d) = \varepsilon(n)$$
, $\mu*1 = \varepsilon$

证明

设
$$n = \prod_{i=1}^k p_i^{\ c_i}, n' = \prod_{i=1}^k p_i$$

那么
$$\sum_{d|n} \mu(d) = \sum_{d|n'} \mu(d) = \sum_{i=0}^k \binom{k}{i} \cdot (-1)^i = (1+(-1))^k$$

根据二项式定理,易知该式子的值在 k=0 即 n=1 时值为 1 否则为 0 ,这也同时证明了 $\sum_{d|n}\mu(d)=[n=1]=\varepsilon(n)$ 以及 $\mu*1=\varepsilon$

如果有
$$f(n) = \sum_{d|n} g(d)$$
, 那么有 $g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d})$

如果有
$$f(n) = \sum_{n|d} g(d)$$
, 那么有 $g(n) = \sum_{n|d} \mu(\frac{d}{n}) f(d)$

```
void get_mu(int n) {
    mu[1] = 1;
```

```
for(int i = 2; i <= n; i++) {
        if(!np[i]) {
            p[++m] = i;
            mu[i] = -1;
        for(int j = 1; j \le m \&\& p[j] \le n / i; j++) {
            np[i * p[j]] = 1;
            if(i \% p[j] == 0) {
                mu[i * p[j]] = 0;
                break;
            mu[i * p[j]] = -mu[i];
        }
   }
   for(int i = 1; i <= n; i++) {
       sum[i] = sum[i - 1] + mu[i];
    }
}
```

逆元

线性求逆元

```
inv[1] = 1;
for (int i = 2; i <= n; ++i) {
   inv[i] = (11)(P - P / i) * inv[P % i] % P;
}</pre>
```

费马小定理求逆元

```
i64 power(i64 a, i64 b) {
    a %= P;
    i64 res = 1;
    while (b) {
        if (b & 1) res = res * a % P;
        a = a * a % P;
        b >>= 1;
    }
    return res % P;
}
i64 inv(i64 x) {
    return power(x, P - 2);
}
```

组合数

```
constexpr int P = 1e9 + 7;
i64 power(i64 a, i64 b) {
    a %= P;
    i64 res = 1 % P;
    while (b) {
        if (b & 1) {
            res = res * a % P;
        }
        a = a * a % P;
```

```
b >>= 1;
    return res;
}
vector<i64> fact, inv_fact, _inv;
i64 c(int n, int m) {
    if(n < m || m < 0) return 0;
    return fact[n] * inv_fact[m] % P * inv_fact[n - m] % P;
}
i64 inv(i64 x) {
   return power(x, P - 2);
}
void init(int n) {
   fact.resize(n + 1);
    inv_fact.resize(n + 1);
    _{inv.resize(n + 1);}
   fact[0] = _inv[1] = 1;
    for (int i = 1; i \le n; i++) fact[i] = fact[i - 1] * i \% P;
    inv_fact[n] = inv(fact[n]);
    for (int i = n - 1; i >= 0; i--) inv_fact[i] = inv_fact[i + 1] * (i + 1) %
Р;
    for(int i = 2; i <= n; i++) _inv[i] = fact[i - 1] * inv_fact[i] % P;
}
```

矩阵快速幂

```
struct Matrix {
   vector<vector<i64>> a;
    int n, m;
    Matrix(int _n, int _m) : n(_n), m(_m), a(_n, vector < i64 > (_m)){}
};
Matrix operator * (const Matrix &A, const Matrix &B) {
   Matrix C(A.n, B.m);
    for (int i = 0; i < A.n; i++) {
        for (int j = 0; j < A.m; j++) {
            for (int k = 0; k < B.m; k++) {
                C.a[i][k] = C.a[i][k] + A.a[i][j] * B.a[j][k];
        }
    }
    return C;
Matrix operator * (const Matrix &A, const Matrix &B) {
    Matrix C(A.n, B.m);
    for (int i = 0; i < A.n; i++) {
        for (int j = 0; j < A.m; j++) {
            for (int k = 0; k < B.m; k++) {
                C.a[i][k] = max(C.a[i][k], A.a[i][j] + B.a[j][k]);
            }
        }
    }
    return C;
}//广义矩阵快速幂
    auto power = [&](Matrix a, Matrix b, i64 c) {
        C--;
        Matrix ans = a;
        while(c) {
```

```
if(c & 1) ans = ans * b;
b = b * b;
c >>= 1;
}
return ans;
};
```

高斯消元

```
const int N = 110;
const double eps = 1e-8;
int n;
double a[N][N];
int gauss() { // 高斯消元, 答案存于a[i][n]中, 0 <= i < n
   int c, r;
    for (c = 0, r = 0; c < n; c++) {
       int t = r;
       for (int i = r; i < n; i++) if(fabs(a[i][c]) > fabs(a[t][c])) t = i; //
找绝对值最大的行
       if (fabs(a[t][c]) < eps) continue;</pre>
       for (int i = c; i <= n; i++) swap(a[t][i], a[r][i]); // 将绝对值最大的行换
到最顶端
       for (int i = n; i >= c; i--) a[r][i] /= a[r][c]; // 将当前行的首位变成1
       for (int i = r + 1; i < n; i++) // 用当前行将下面所有的列消成0
           if (fabs(a[i][c]) > eps)
               for (int j = n; j >= c; j--) a[i][j] -= a[r][j] * a[i][c];
       r++;
   }
    if (r < n) {
       for (int i = r; i < n; i++)
           if (fabs(a[i][n]) > eps) return 2; // 无解
       return 1; // 有无穷多组解
   }
    for (int i = n - 1; i >= 0; i -- )
       for (int j = i + 1; j < n; j ++ )
           a[i][n] -= a[i][j] * a[j][n];
    return 0; // 有唯一解
}
int main() {
   scanf("%d", &n);
    for (int i = 0; i < n; i ++)
       for (int j = 0; j < n + 1; j ++)
           scanf("%1f", &a[i][j]);
   int t = gauss();
   if (t == 2) puts("No solution");
    else if (t == 1) puts("Infinite group solutions");
    else {
       for (int i = 0; i < n; i ++) {
           if (fabs(a[i][n]) < eps) a[i][n] = 0; // 去掉输出 -0.00 的情况
           printf("%.21f\n", a[i][n]);
       }
   return 0;
}
```

BSGS

给定正整数 a , p , b 数据保证 a 和 p 互质。

求满足 $a^x \equiv b \pmod{p}$ 的最小非负整数 x。

```
int bsgs(int a, int b, int p) {
    if (1 % p == b % p) return 0;
    int k = sqrt(p) + 1;
    unordered_map<int, int> hash;
    for (int i = 0, j = b % p; i < k; i++) {
        hash[j] = i;
        j = 1i64 * j * a % p;
    }
    int ak = 1;
    for (int i = 0; i < k; i ++ ) ak = 1i64 * ak * a % p;
    for (int i = 1, j = ak; i <= k; i++) {
        if (hash.count(j)) return 1i64 * i * k - hash[j];
        j = 1i64 * j * ak % p;
    }
    return -1;
}</pre>
```

拓展BSGS

给定整数a,p,b。

求满足 $a^x \equiv b \pmod{p}$ 的最小非负整数 x。

```
const int INF = 1e8;
int exgcd(int a, int b, int& x, int& y) {
    if (!b) {
        x = 1, y = 0;
        return a;
    }
    int d = exgcd(b, a \% b, y, x);
    y -= a / b * x;
    return d;
int bsgs(int a, int b, int p) {
    if (1 % p == b % p) return 0;
    int k = sqrt(p) + 1;
    unordered_map<int, int> hash;
    for (int i = 0, j = b \% p; i < k; i++) {
        hash[j] = i;
        j = 1i64 * j * a % p;
    }
    int ak = 1;
    for (int i = 0; i < k; i ++) ak = 1i64 * ak * a % p;
    for (int i = 1, j = ak; i \le k; i++) {
        if (hash.count(j)) return i * k - hash[j];
        j = 1i64 * j * ak % p;
    }
    return -INF;
int exbsgs(int a, int b, int p) {
    b = (b \% p + p) \% p;
```

```
if (1 % p == b % p) return 0;
int x, y;
int d = exgcd(a, p, x, y);
if (d > 1) {
    if (b % d) return -INF;
    exgcd(a / d, p / d, x, y);
    return exbsgs(a, 1i64 * b / d * x % (p / d), p / d) + 1;
}
return bsgs(a, b, p);
}
```

FFT

```
constexpr double PI = acos(-1);
struct Complex {
    double x, y;
    Complex operator+ (const Complex& t) const {
        return \{x + t.x, y + t.y\};
    Complex operator- (const Complex& t) const {
        return \{x - t.x, y - t.y\};
    Complex operator* (const Complex& t) const {
        return \{x * t.x - y * t.y, x * t.y + y * t.x\};
    }
auto fft = [&](auto &a, int inv, int tot, auto &rev) {
    for (int i = 0; i < tot; i++) {
        if (i < rev[i]) swap(a[i], a[rev[i]]);</pre>
    for (int mid = 1; mid < tot; mid <<=1) {
        auto w1 = Complex({cos(PI / mid), inv * sin(PI / mid)});
        for (int i = 0; i < tot; i += mid * 2) {
            auto wk = Complex(\{1, 0\});
            for (int j = 0; j < mid; j++, wk = wk * w1) {
                auto x = a[i + j], y = wk * a[i + j + mid];
                a[i + j] = x + y, a[i + j + mid] = x - y;
        }
    }
};
void solve() {
   int n, m;
    cin >> n >> m;
    int bit = 0;
    while((1 << bit) < n + m + 1) bit++;
    int tot = 1 << bit;</pre>
    vector<Complex> a(2 * (n + m + 1) + 1), b(2 * (n + m + 1) + 1);
    for(int i = 0; i \le n; i++) cin >> a[i].x;
    for(int i = 0; i \le m; i++) cin >> b[i].x;
    for(int i = 0; i < tot; i++) rev[i] = (rev[i >> 1] >> 1) | ((i & 1) << (bit
- 1));
    fft(a, 1, tot, rev), fft(b, 1, tot, rev);
    for (int i = 0; i < tot; i++) a[i] = a[i] * b[i];
    fft(a, -1, tot, rev);
    for(int i = 0; i \le n + m; i++) cout << (long long)(a[i].x / tot + 0.5) << '
'://注意值域可能超出int范围
```

三次变两次

```
constexpr int N = 2e6 + 10; //N要大于(n+m+1)*2
constexpr double PI = acos(-1);
int n, m;
struct Complex {
    double x, y;
    Complex operator+ (const Complex& t) const {
        return \{x + t.x, y + t.y\};
    }
    Complex operator- (const Complex& t) const {
        return \{x - t.x, y - t.y\};
    }
    Complex operator* (const Complex& t) const {
        return \{x * t.x - y * t.y, x * t.y + y * t.x\};
    }
} a[N];
int rev[N], bit, tot;
void fft(Complex a[], int inv) {
    for(int i = 0; i < tot; i++) {
        if(i < rev[i]) {
            swap(a[i], a[rev[i]]);
        }
    for(int mid = 1; mid < tot; mid <<= 1) {</pre>
        auto w1 = Complex({cos(PI / mid), inv * sin(PI / mid)});
        for(int i = 0; i < tot; i += mid * 2) {
            auto wk = Complex(\{1, 0\});
            for(int j = 0; j < mid; j++, wk = wk * w1) {
                auto x = a[i + j], y = wk * a[i + j + mid];
                a[i + j] = x + y, a[i + j + mid] = x - y;
            }
        }
    }
    if(inv != 1) {
        for(int i = 0; i < tot; i++) {
            a[i].x /= tot, a[i].y /= tot;
        }
    }
void solve() {
    cin >> n >> m;
    for(int i = 0; i <= n; i++) {
        cin \gg a[i].x;
    for(int i = 0; i \le m; i++) {
        cin \gg a[i].y;
    while((1 << bit) < n + m + 1) {
        bit++;
    tot = 1 << bit;
    for(int i = 0; i < tot; i++) {
        rev[i] = (rev[i >> 1] >> 1) | ((i & 1) << (bit - 1));
    }
```

```
fft(a, 1);
for(int i = 0; i < tot; i++) {
    a[i] = a[i] * a[i];
}
fft(a, -1);
for(int i = 0; i <= n + m; i++) cout << (long long)(a[i].y / 2 + 0.5) << '
';
}</pre>
```

分治FFT

```
int rev[N];
vector<Complex> cal(int 1, int r) {
    if(1 == r) {
       //vector<Complex> a(2);
        //a[0].x = 1 - p[1];
        //a[1].x = p[1];
       return a;
    }
   int mid = (1 + r) / 2;
    auto a = cal(1, mid);
    auto b = cal(mid + 1, r);
   int bit = 0;
    int n = (int)a.size(), m = (int)b.size();
    a.resize(2 * (n + m + 1) + 1);
    b.resize(2 * (n + m + 1) + 1);
    while((1 << bit) < n + m + 1) bit++;
    int tot = 1 << bit;</pre>
    for(int i = 0; i < tot; i++) {
        rev[i] = (rev[i >> 1] >> 1) | ((i & 1) << (bit - 1));
    fft(a, 1, tot, rev), fft(b, 1, tot, rev);
    for (int i = 0; i < tot; i++) a[i] = a[i] * b[i];
    fft(a, -1, tot, rev);
    vector<Complex> c(n + m - 1);
    for(int i = 0; i \le n + m - 2; i++) {
        c[i].x = a[i].x / tot;
    }
    return c;
void solve() {
    cin >> n;
    //for(int i = 1; i <= n; i++) {
   // cin >> p[i];
    //}
    auto ans = cal(1, n);
   double sum = 0;
    //for(int i = n / 2 + 1; i <= n; i++) {
        sum += ans[i].x;
   //
   //}
   cout << fixed << setprecision(10) << sum;</pre>
}
```

```
constexpr int N = 4e6 + 10; //(n+m+1)*2
constexpr i64 P = 998244353;
constexpr int g = 3;
i64 a[N], b[N], rev[N], gn[N];
i64 power(i64 a, i64 b) {
    a %= P;
    i64 res = 1;
    while (b) {
        if (b & 1) res = (__int128)res * a % P;
        a = (\underline{\quad}int128)a * a % P;
        b >>= 1;
    }
    return res % P;
void ntt(auto a[], int len, auto opt) {
    for (int i = 0; i < len; i++)
        if (i < rev[i])</pre>
            swap(a[i], a[rev[i]]);
    for (int i = 2; i <= len; i <<= 1) {
        int m = i \gg 1;
        i64 gw = power(g, (opt == 1) ? ((P - 1) / (m << 1)) : (P - 1 - (P - 1) /
(m << 1));
        qn[0] = 1;
        for (int j = 1; j < m; j++)
            gn[j] = (\_int128)gn[j - 1] * gw % P;
        for (int j = 0; j < len; <math>j += i) {
            for (int k = j; k < j + m; k++) {
                i64 x = a[k], y = (_int128)gn[k - j] * a[k + m] % P;
                a[k] = (x + y) \% P;
                a[k + m] = (x - y + P) \% P;
            }
        }
    }
void solve() {
    int n, m;
    cin >> n >> m;
    for(int i = 0; i \le n; i++) cin >> a[i];
    for(int i = 0; i <= m; i++) cin >> b[i];
    int bit = 0;
    while((1 << bit) < n + m + 1) bit++;
    int tot = 1 << bit;</pre>
    for(int i = 0; i < tot; i++) rev[i] = (rev[i >> 1] >> 1) | ((i & 1) << (bit
- 1));
    ntt(a, tot, 1), ntt(b, tot, 1);
    for (int i = 0; i < tot; i++)
        a[i] = (\underline{\quad}int128)a[i] * b[i] % P;
    ntt(a, tot, -1);
    i64 inv = power(tot, P - 2);
    for (int i = 0; i \le n + m; i++) a[i] = (_int128)a[i] * inv % P;
    for (int i = 0; i \le n + m; i++) cout << a[i] << ' ';
}
```

分治NTT

```
constexpr int N = 1e6 + 10;
constexpr i64 P = 998244353;
constexpr int g = 3;
int n;
using Poly = vector<i64>;
i64 gn[N], rev[N];
i64 power(i64 a, i64 b) {
    a %= P;
    i64 res = 1;
    while (b) {
        if (b & 1) res = res * a % P;
        a = a * a % P;
        b >>= 1;
    }
    return res % P;
}
i64 inv(i64 x) {
    return power(x, P - 2);
}
int gw[2][N];
void init(int n) {
    for(int i = 1; i <= n; i *= 2) {
        gw[0][i] = power(g, (P - 1 - (P - 1) / (i << 1)));
        gw[1][i] = power(g, ((P - 1) / (i << 1)));
    }
}
void ntt(Poly &a, int len, auto opt) {
    for (int i = 0; i < len; i++)
        if (i < rev[i])
            swap(a[i], a[rev[i]]);
    for (int i = 2; i \le len; i \le 1) {
        int m = i \gg 1;
        i64 w = gw[opt == 1][m];
        //i64 \text{ gw} = \text{power}(g, (opt == 1) ? ((P - 1) / (m << 1)) : (P - 1 - (P - 1)
/ (m << 1)));//预处理gw会快很多
        gn[0] = 1;
        for (int j = 1; j < m; j++)
            gn[j] = gn[j - 1] * w % P;
        for (int j = 0; j < len; <math>j += i) {
            for (int k = j; k < j + m; k++) {
                i64 x = a[k], y = gn[k - j] * a[k + m] % P;
                a[k] = (x + y) \% P;
                a[k + m] = (x - y + P) \% P;
            }
        }
    }
Poly operator*(Poly a, Poly b) {
    int n = a.size() - 1, m = b.size() - 1;
    int bit = 0;
    while((1 \ll bit) < n + m + 1) bit++;
    int tot = 1 << bit;</pre>
    for(int i = 0; i < tot; i++) rev[i] = (rev[i >> 1] >> 1) | ((i & 1) << (bit
- 1));
    a.resize(tot);
```

```
b.resize(tot);
    \mathsf{ntt}(\mathsf{a},\ \mathsf{tot},\ 1)\,,\ \mathsf{ntt}(\mathsf{b},\ \mathsf{tot},\ 1)\,;
    for (int i = 0; i < tot; i++)
         a[i] = a[i] * b[i] % P;
    ntt(a, tot, -1);
    i64 inv = power(tot, P - 2);
    for (int i = 0; i \le n + m; i++) a[i] = a[i] * inv % P;
    a.resize(n + m + 1);
    return a;
}
Poly MulAll(int 1, int r, vector<Poly> &a) {
    if(1 == r) return a[1];
    int m = 1 + r >> 1;
    auto v = MulAll(l, m, a) * MulAll(m + 1, r, a);
    return v;
void solve() {
    init(n); //预处理gw, n的数量级为答案的最高幂次
    vector<Poly> a(n + 1);
    for(int i = 1; i \le n; i++) {
        //int x;
        //cin >> x;
        //a[i] = \{1, x\};
    auto v = Mulall(1, n, a);
}
```

FWT

```
int m;
cin >> m;
int n = 1 \ll m;
vector<i64> a(n), b(n);
for(int i = 0; i < n; i++) {
    cin \gg a[i];
for(int i = 0; i < n; i++) {
    cin >> b[i];
auto fwtor = [&](vector<i64> &a, int op) {
    for(int i = 1; i < n; i *= 2) {
        for(int p = i \ll 1, j = 0; j < n; j += p) {
            for(int k = 0; k < i; k++) {
                a[i + j + k] = (a[i + j + k] + a[j + k] * op + P) % P;
            }
        }
    }
};
auto fwt_or = [\&] (vector<i64> &a, vector<i64> &b) {
    auto A = a, B = b;
    fwtor(A, 1), fwtor(B, 1);
    for(int i = 0; i < n; i++) A[i] = A[i] * B[i] % P;
    fwtor(A, -1);
    return A;
};
auto A = fwt_or(a, b);
```

```
auto fwtand = [&](vector<i64> &a, int op) {
    for(int i = 1; i < n; i *= 2) {
        for(int p = i \ll 1, j = 0; j < n; j += p) {
            for(int k = 0; k < i; k++) {
                a[j + k] = (a[j + k] + a[i + j + k] * op + P) % P;
            }
        }
    }
};
auto fwt_and = [&](vector<i64> &a, vector<i64> &b) {
    auto A = a, B = b;
    fwtand(A, 1), fwtand(B, 1);
    for(int i = 0; i < n; i++) A[i] = A[i] * B[i] % P;
    fwtand(A, -1);
    return A;
};
auto B = fwt_and(a, b);
auto fwtxor = [&](vector<i64> &a, int op) {
    for(int i = 1; i < n; i *= 2) {
        for(int p = i \ll 1, j = 0; j \ll n; j \neq p) {
            for(int k = 0; k < i; k++) {
                i64 x = a[j + k], y = a[i + j + k];
                a[j + k] = (x + y) \% P;
                a[i + j + k] = (x - y + P) \% P;
                if(op == -1) {
                    a[j + k] = a[j + k] * inv2 % P;
                    a[i + j + k] = a[i + j + k] * inv2 % P;
                }
            }
        }
    }
};
auto fwt_xor = [\&] (vector<i64> &a, vector<i64> &b) {
    auto A = a, B = b;
    fwtxor(A, 1), fwtxor(B, 1);
    for(int i = 0; i < n; i++) A[i] = A[i] * B[i] % P;
    fwtxor(A, -1);
    return A;
};
auto C = fwt\_xor(a, b);
for(int i = 0; i < n; i++) cout << A[i] << " ";
cout << '\n';</pre>
for(int i = 0; i < n; i++) cout << B[i] << " ";
cout << '\n';
for(int i = 0; i < n; i++) cout << C[i] << " ";
```

Pollard-Rho

```
i64 mul(i64 a, i64 b, i64 m) {
    return static_cast<__int128>(a) * b % m;
}
i64 power(i64 a, i64 b, i64 m) {
    i64 res = 1 % m;
    for (; b; b >>= 1, a = mul(a, a, m))
        if (b & 1)
```

```
res = mul(res, a, m);
    return res;
}
bool isprime(i64 n) {
   if (n < 2)
        return false;
    static constexpr int A[] = \{2, 3, 5, 7, 11, 13, 17, 19, 23\};
    int s = __builtin_ctzll(n - 1);
    i64 d = (n - 1) >> s;
    for (auto a : A) {
        if (a == n)
            return true;
        i64 x = power(a, d, n);
        if (x == 1 || x == n - 1)
            continue;
        bool ok = false;
        for (int i = 0; i < s - 1; ++i) {
            x = mul(x, x, n);
            if (x == n - 1) {
                ok = true;
                break;
            }
        }
        if (!ok)
            return false;
    }
    return true;
vector<i64> factorize(i64 n) {
    vector<i64> p;
    function<void(i64)> f = [&](i64 n) {
        if (n <= 10000) {
            for (int i = 2; i * i <= n; ++i)
                for (; n \% i == 0; n /= i)
                    p.push_back(i);
            if (n > 1)
                p.push_back(n);
            return;
        }
        if (isprime(n)) {
            p.push_back(n);
            return;
        }
        auto g = [\&](i64 x) {
            return (mul(x, x, n) + 1) \% n;
        };
        i64 x0 = 2;
        while (true) {
            i64 x = x0;
            i64 y = x0;
            i64 d = 1;
            i64 power = 1, lam = 0;
            i64 v = 1;
            while (d == 1) {
                y = g(y);
                ++1am;
                v = mul(v, abs(x - y), n);
                if (lam % 127 == 0) {
```

```
d = gcd(v, n);
                    v = 1;
                }
                if (power == lam) {
                    x = y;
                    power *= 2;
                    lam = 0;
                    d = gcd(v, n);
                    v = 1;
                }
            }
            if (d != n) {
                f(d);
                f(n / d);
                return;
            }
            ++x0;
        }
   };
    sort(p.begin(), p.end());
    return p;
}
```

类欧几里得算法

分数类

```
struct Frac {
   i128 x, y;
   Frac(i128 \ a = 0, i128 \ b = 1)  {
        if(b < 0) a = -a, b = -b; //保证分母是正数
        x = a, y = b;
   }
};
bool operator < (const Frac &a, const Frac &b) {</pre>
   return a.x * b.y < a.y * b.x;
bool operator <= (const Frac &a, const Frac &b) {</pre>
   return a.x * b.y <= a.y * b.x;
}
bool operator == (const Frac &a, const Frac &b) {
   return a.x * b.y == a.y * b.x;
}
bool operator > (const Frac &a, const Frac &b) {
   return a.x * b.y > a.y * b.x;
bool operator >= (const Frac &a, const Frac &b) {
   return a.x * b.y >= a.y * b.x;
Frac operator + (const Frac &a, const Frac &b) {
   return Frac(a.x * b.y + a.y * b.x, a.y * b.y);
Frac operator - (const Frac &a, const Frac &b) {
   return Frac(a.x * b.y - a.y * b.x, a.y * b.y);
```

```
Frac operator * (const Frac &a, const Frac &b) {
    return Frac(a.x * b.x, a.y * b.y);
}

struct Line {
    Frac k, b;
} line[N];

pair<Frac, Frac> LineIntersection(Line &a, Line &b) {
    Frac frac1(b.b - a.b), frac2(a.k - b.k);
    return make_pair(frac1, frac2);
    //return make_pair(b.b - a.b, a.k - b.k);
} //求线段交点的横坐标
//用的过程中小心爆int128, 在abc372G中因为求线段交点的大小关系时爆int128debug很久
```

动态规划

换根dp

有一棵n个节点的树,节点编号从1到n。请求出每个节点到其他所有节点的距离的和。 定义两个节点的距离为它们之间的简单路径上经过了多少条边。

http://oj.daimayuan.top/course/5/problem/224

```
vector<int> g[N];
int n, sz[N];
i64 dn[N], up[N];
void dfs1(int u, int fa) {
    sz[u] = 1;
    for(int v : g[u]) {
        if(v == fa) continue;
        dfs1(v, u);
        sz[u] += sz[v];
        dn[u] += dn[v] + sz[v];
    }
void dfs2(int u, int fa) {
   for(int v : g[u]) {
        if(v == fa) continue;
        up[v] = up[u] + dn[u] - (dn[v] + sz[v]) + n - sz[v];
        dfs2(v, u);
    }
}
void solve() {
    cin >> n;
    for(int i = 1; i < n; i++) {
        int u, v;
        cin >> u >> v;
        g[u].push_back(v);
        g[v].push_back(u);
    }
    dfs1(1, 0);
    dfs2(1, 0);
```

```
for(int i = 1; i <= n; i++) {
    cout << up[i] + dn[i] << '\n';
}
//dn是节点下面的点到该节点的距离和, up是节点上面的点到该节点的距离和
}
```

最长上升子序列

```
template <typename T>
struct Fenwick {
    int n;
   vector<T> a;
    Fenwick(int _n = 0) : a(_n + 1), n(_n){}
    void init(int n) {
        a.assign(n + 1, T());
        this->n = n;
    }
    void add(int x, T v) {
        for (int i = x; i \le n; i += i \& -i) {
            a[i] = max(a[i], v);
    }
    T sum(int x) {
        int ans = 0;
        for (int i = x; i >= 1; i -= i \& -i) {
            ans = max(ans, a[i]);
        }
        return ans;
    }
    T rangeSum(int 1, int r) {
        return sum(r) - sum(1 - 1);
    }
    int kth(T k) {
        int x = 0;
        for (int i = 1 \ll std::_lg(n); i; i /= 2) {
            if (x + i \le n \&\& k >= a[x + i]) {
                x += i;
                k = a[x];
            }
        }
        return x;
    }
};
    vector<int> b = a;
    sort(b.begin(), b.end());
    auto find(int x) = [\&](int x) {
        return lower_bound(b + 1, b + 1 + n, x) - b;
    };
    Fenwick<int> fen1(n), fen2(n);
    vector<int> 1(n + 1), r(n + 1);
    for(int i = 1; i <= n; i++) {
        int pos = find(a[i]);
        int v = fen1.sum(pos - 1);
        l[i] = v + 1;
        fen1.add(pos, v + 1);
    for(int i = n; i >= 1; i--) {
```

```
int pos = n - find(a[i]) + 1;
int v = fen2.sum(pos - 1);
r[i] = v + 1;
fen2.add(pos, v + 1);
}
```

单调队列优化多重背包

时间复杂度 O(nm)

```
void solve() {
    int n, m;
    cin >> n >> m;
    vector<int> dp(m + 1), ndp(m + 1);
    vector<int> q(m + 1);
    for(int i = 1; i <= n; i++) {
        int v, w, s;
        cin >> v >> w >> s;
        for(int r = 0; r < v; r++) {//r是模数
             int hh = 0, tt = -1;
             for(int j = r; j \leftarrow m; j \leftarrow v) {
                 while(hh \leftarrow tt && q[hh] < j - s * v) hh++;
                 while(hh <= tt \&\& dp[q[tt]] + (j - q[tt]) / v * w <= dp[j]) tt-
-;
                 q[++tt] = j;
                 ndp[j] = dp[q[hh]] + (j - q[hh]) / v * w;
            }
        }
        swap(dp, ndp);
    cout << dp[m];</pre>
}
```

斜率优化dp

```
void solve() {
   int n;
   i64 L;
   cin >> n >> L;
   L++;
   vector<i64> dp(n + 1), s(n + 1), q(n + 1);
   for(int i = 1; i <= n; i++) {
        cin >> s[i];
        s[i] += s[i - 1] + 1;
   }
   auto sq = [\&](int x) {
       return 1i64 * x * x;
   };
   auto Y = [\&](int i) {
        return dp[i] + sq(s[i]);
   };
   auto X = [\&](int i) {
        return s[i];
   };
```

```
auto slope = [&](int i, int j) {
    return 1. * (Y(j) - Y(i)) / (X(j) - X(i));
};
int hh = 0, tt = 0;
for(int i = 1; i <= n; i++) {
    while(hh < tt && slope(q[hh], q[hh + 1]) <= 2 * (s[i] - L)) hh++;
    int j = q[hh];
    dp[i] = dp[j] + sq(s[i] - s[j] - L);
    while(hh < tt && slope(q[tt - 1], q[tt]) >= slope(q[tt - 1], i)) tt--;
    q[++tt] = i;
}
cout << dp[n];
}</pre>
```

数位dp

统计0~9 在 [a,b] 中出现了多少次

```
int num[20];
i64 dp[20][20];
i64 dfs(int pos, bool limit, bool lead, int sum, int d) {
    i64 ans = 0;
    if(!pos) return sum;
    if(!limit && lead && dp[pos][sum] != -1) return dp[pos][sum];
    int up = 9;
    if(limit) up = num[pos];
    for(int j = 0; j \leftarrow up; j \leftrightarrow ds) ans += dfs(pos - 1, (j == up) & limit, lead ||
j, sum + ((j \mid | lead) && (j == d)), d);
    if(!limit && lead) dp[pos][sum] = ans;
    return ans;
}//lead=0表示存在前导0
i64 get(i64 x, int d) {
    int len = 0;
    while(x) {
        num[++1en] = x \% 10;
        x /= 10;
    memset(dp, -1, sizeof dp);
    return dfs(len, 1, 0, 0, d);
}
```

不同子序列的个数

有一个长度为 N 的数列 A_i 。

你可以进行若干次,最多 N-1 次操作,选择相邻的两个数,删去他们,并在原位置放上他们两个的 和。

现在你需要求出可能产生的序列个数。

做法:一种合法的序列就是由前缀和中按顺序任选其中几个确定,那本质不同的序列也就是有多少种不同的选择的方法,也就是子序列的数量。

```
const int P = 998244353;
int n;
i64 s[N];
map<i64, i64> cnt;
```

```
i64 dp[N];
void solve() {
   int n;
   cin >> n;
    for(int i = 1; i <= n; i++) {
       i64 x;
       cin >> x;
       s[i] = s[i - 1] + x;
   }
   for(int i = 1; i < n; i++) {
       if(cnt.count(s[i])) {
           dp[i] = (dp[i - 1] * 2 - cnt[s[i]] + P) % P;
       }
       else {
           dp[i] = (dp[i - 1] * 2 + 1) % P;
       cnt[s[i]] = dp[i - 1];
   cout << (dp[n - 1] + 1) % P;//该题中 s[n] 必须要选择,遇到其他题时需注意
}
```

其他

树的拓扑序计数

有根树: $n! \times \prod_{1}^{n} \frac{1}{sz[i]}$

无根树:换根dp处理有根树的式子

切比雪夫距离转换曼哈顿距离

```
int n, x[N], y[N];
vector<array<int, 2>> v1, v2; //v1存x和y的和为偶数的点, v2存为奇数的点
i64 \text{ ans} = 0;
void cal(auto a) {
   vector<11> x, y;
    for(auto [u, v] : a) {
        x.push_back(u + v);
       y.push_back(u - v);
    }
    sort(x.begin(), x.end());
    sort(y.begin(), y.end());
    i64 s = 0;
    for(int i = 0; i < x.size(); i++) {
        ans += (x[i] * i) - s;
        s += x[i];
    }
    s = 0;
    for(int i = 0; i < y.size(); i++) {
        ans += (y[i] * i) - s;
        s += y[i];
void solve() {
```

```
cin >> n;
for(int i = 1; i <= n; i++) {
    cin >> x[i] >> y[i];
    if((x[i] + y[i]) % 2 == 0) v1.push_back({x[i], y[i]});
    else v2.push_back({x[i], y[i]});
}
cal(v1), cal(v2);
cout << ans / 2;
}</pre>
```

杂项

随机数

```
mt19937_64 rng(chrono::steady_clock::now().time_since_epoch().count());//使用64位时,存储要用unsigned long long mt19937 rng(time(0));
```

unordered_map 防 TLE

```
struct custom_hash {
    static uint64_t splitmix64(uint64_t x) {
        // http://xorshift.di.unimi.it/splitmix64.c
        x += 0x9e3779b97f4a7c15;
        x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
        x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
        return x ^ (x >> 31);
    }

    size_t operator()(uint64_t x) const {
        static const uint64_t FIXED_RANDOM =
    chrono::steady_clock::now().time_since_epoch().count();
        return splitmix64(x + FIXED_RANDOM);
    }
};
unordered_map <int, int, custom_hash> mx;
```

初始代码模板

```
#include<bits/stdc++.h>

using namespace std;
using i64 = long long;

constexpr int N = 2e5 + 10;

void solve() {
}

int main() {
```

```
ios::sync_with_stdio(false);
    cin.tie(nullptr);

int t;
    t = 1;

while (t--) {
        solve();
    }

    return 0;
}
```

int128读入输出

```
inline __int128 read(){
   _{\text{int}128} x = 0, f = 1;
   char ch = getchar();
   while(ch < '0' || ch > '9'){
        if(ch == '-')
            f = -1;
        ch = getchar();
   }
    while(ch >= '0' && ch <= '9'){
       x = x * 10 + ch - '0';
        ch = getchar();
   }
    return x * f;
}
inline void print(__int128 x){
   if(x < 0){
        putchar('-');
        x = -x;
   }
   if(x > 9)
        print(x / 10);
   putchar(x \% 10 + '0');
}
```

取余模板

```
constexpr int P = 998244353;
using i64 = long long;
// assume -P <= x < 2P
int norm(int x) {
    if (x < 0) {
        x += P;
    }
    if (x >= P) {
        x -= P;
    }
    return x;
}
```

```
T power(T a, i64 b) {
    T res = 1;
    for (; b; b /= 2, a *= a) \{
        if (b % 2) {
            res *= a;
        }
    }
   return res;
}
struct Z {
   int x;
    Z(int x = 0) : x(norm(x)) \{\}
    Z(i64 x) : x(norm(x \% P)) \{ \}
    int val() const {
        return x;
    }
    Z operator-() const {
        return Z(norm(P - x));
    }
    Z inv() const {
        assert(x != 0);
        return power(*this, P - 2);
    Z &operator*=(const Z &rhs) {
        x = i64(x) * rhs.x % P;
        return *this;
    Z &operator+=(const Z &rhs) {
        x = norm(x + rhs.x);
        return *this;
    Z &operator-=(const Z &rhs) {
        x = norm(x - rhs.x);
        return *this;
    Z &operator/=(const Z &rhs) {
        return *this *= rhs.inv();
    friend Z operator*(const Z &lhs, const Z &rhs) {
        z res = 1hs;
        res *= rhs;
        return res;
    friend Z operator+(const Z &lhs, const Z &rhs) {
        z res = 1hs;
        res += rhs;
        return res;
    friend Z operator-(const Z &lhs, const Z &rhs) {
        z res = 1hs;
        res -= rhs;
        return res;
    friend Z operator/(const Z &lhs, const Z &rhs) {
        z res = 1hs;
        res /= rhs;
        return res;
    }
```

```
friend istream &operator>>(istream &is, Z &a) {
    i64 v;
    is >> v;
    a = Z(v);
    return is;
}
friend ostream &operator<<(ostream &os, const Z &a) {
    return os << a.val();
}
};</pre>
```

lambda函数

```
实现dfs:
function<void(int, int, int)> dfs = [&](int x, int y, int z) {
};
其他:
auto fun = [&](int x, int y, int z) {
};
sort(a.begin(), a.end(), [&](int x, int y) {
    return ***;
});
```

开O2

```
#pragma GCC optimize(2)
```

windows系统下的对拍

```
#include <bits/stdc++.h>
using namespace std;
using i64 = long long;
constexpr int N = 2e5 + 10;
void solve() {
    for(int i = 1; i \le 100000; i++) {
        system("datamaker.exe > data.txt");
        system("A.exe < data.txt > A.txt");
        double start = clock();
        system("test.exe < data.txt > test.txt");
        double end = clock();
        if(system("fc A.txt test.txt")) {cout << "WA\n"; break;}</pre>
        else cout << "AC TIME : " << end - start << '\n';</pre>
    }
}
int main() {
    int t = 1;
    while(t--) {
```

```
solve();
}
return 0;
}
```

linux系统下的对拍

```
需要在两个代码中读写文件,格式:
freopen("data.in","r",stdin);
freopen("std.out", "w", stdout);
#include<bits/stdc++.h>
using namespace std;
int main() {
   for (int i = 1; i \le 100000; i++){
       printf("The result of No. %d Case is: ", i);
       system("./datamaker");
       system("./std");
       system("./test");
       if (system("diff std.out test.out")) {
            printf("Wrong Answer\n");
           break;
       }
       else printf("Accepted\n");
   }
   return 0;
}
```

codeblocks终端修改

由于codeblocks默认的终端字体不易观看,更改步骤如下:

- 打开codeblocks
- Settings -> Environment
- 将Terminal to launch console programs栏中改为:

gnome-terminal --geometry 80x20+100+100 --hide-menubar -t \$TITLE -x

原根表

```
prime
               r k
                     g
               1 1
3
                     2
5
               1 2
                     2
17
               1 4 3
               3 5 5
97
               3 6 5
193
               1 8 3
257
               15 9 17
7681
12289
               3 12 11
               5 13 3
40961
              1 16 3
65537
              3 18 10
786433
              11 19 3
5767169
              7 20 3
7340033
               11 21 3
23068673
               25 22 3
104857601
```

```
167772161 5 25 3
 469762049
                   479 21 3
15 27 3
                       7 26 3
 1004535809
 2013265921 15 27 31

2281701377 17 27 3

3221225473 3 30 5

75161927681 35 31 3

77309411329 9 33 7

206158430209 3 36 22
3 36 22

2748779069441 5 39 3

6597069766657 3 41 5

39582418599937 9 42 5

79164837199873 9 43 5

263882790666241
 1231453023109121 35 45 3
 3799912185593857 27 47 5
 4222124650659841 15 48 19
 7881299347898369 7 50 6
 31525197391593473 7 52 3
 180143985094819841 5 55 6
 1945555039024054273 27 56 5
 4179340454199820289 29 57 3
```

C++ STL简介

```
vector, 变长数组, 倍增的思想
   size() 返回元素个数
   empty() 返回是否为空
   clear() 清空
   front()/back()
   push_back()/pop_back()
   begin()/end()
   Г٦
   支持比较运算,按字典序
pair<int, int>
   first,第一个元素
   second, 第二个元素
   支持比较运算,以first为第一关键字,以second为第二关键字(字典序)
string,字符串
   size()/length() 返回字符串长度
   empty()
   clear()
   substr(起始下标,(子串长度)) 返回子串
   c_str() 返回字符串所在字符数组的起始地址
queue, 队列
   size()
   empty()
   push() 向队尾插入一个元素
   front() 返回队头元素
   back() 返回队尾元素
   pop() 弹出队头元素
priority_queue, 优先队列,默认是大根堆
   size()
```

```
empty()
   push() 插入一个元素
   top() 返回堆顶元素
   pop() 弹出堆顶元素
   定义成小根堆的方式: priority_queue<int, vector<int>, greater<int>> q;
stack, 栈
   size()
   empty()
   push() 向栈顶插入一个元素
   top() 返回栈顶元素
   pop() 弹出栈顶元素
deque, 双端队列
   size()
   empty()
   clear()
   front()/back()
   push_back()/pop_back()
   push_front()/pop_front()
   begin()/end()
set, map, multiset, multimap, 基于平衡二叉树(红黑树), 动态维护有序序列
   size()
   empty()
   clear()
   begin()/end()
   ++, -- 返回前驱和后继, 时间复杂度 O(logn)
   set/multiset
      insert() 插入一个数
      find() 查找一个数
      count() 返回某一个数的个数
      erase()
          (1) 输入是一个数x, 删除所有x O(k + logn)
          (2) 输入一个迭代器, 删除这个迭代器
      lower_bound()/upper_bound()
          lower_bound(x) 返回大于等于x的最小的数的迭代器
          upper_bound(x) 返回大于x的最小的数的迭代器
   map/multimap
      insert() 插入的数是一个pair
      erase() 输入的参数是pair或者迭代器
      find()
      [] 注意multimap不支持此操作。 时间复杂度是 O(logn)
      lower_bound()/upper_bound()
unordered_set, unordered_map, unordered_multiset, unordered_multimap, 哈希表
   和上面类似,增删改查的时间复杂度是 O(1)
   不支持 lower_bound()/upper_bound(), 迭代器的++, --
bitset, 圧位
   bitset<10000> s;
   ~, &, |, ^
   >>, <<
   ==, !=
   Г٦
   count() 返回有多少个1
   any() 判断是否至少有一个1
   none() 判断是否全为0
   set() 把所有位置成1
   set(k, v) 将第k位变成v
   reset() 把所有位变成0
   flip() 等价于~
```

flip(k) 把第k位取反

$n \le$	10 ¹	10^{2}	10^{3}	10^{4}	10^{5}	10^{6}	10^{7}	10^{8}	10^{9}
$\max\{\omega(n)\}$	2	3	4	5	6	7	8	8	9
$\max\{d(n)\}$	4	12	32	64	128	240	448	768	1344
$n \le$	10 ¹⁰	10^{11}	10^{12}	10^{13}	10^{14}	10^{15}	10^{16}	10^{17}	10^{18}
$\max\{\omega(n)\}$	10	10	11	12	12	13	13	14	15
$\max\{d(n)\}$	2304	4032	6720	10752	17280	26880	41472	64512	103680