# Security Engineering 3rd Problem Set

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# Organization

### Mini projects

If you are presenting a mini project, please submit it till Thursday 13:30 before the Problem Session.

# Miniprojects

#### Recap

- Get familiar with the basics of Ada: √
- Get familiar with unit testing: √
- Get familiar with at least one testing tool: √
- Organize sources and testing code: √

#### 4th Problem Set

#### 4th Problem Set: Next Steps

- Reflect the quality of your tests
- The formal backgrounds of how-to-prove the correctness of program code
- Proofing the correctness of your code

#### How to Measure Test Quality?

- Analyzing covered lines of source code
- Analyzing covered branches
- Rule-of-thumb:
  - All lines covered is the wish (loops?)
  - ightharpoonup > 95% of lines are desirable

#### gcov and lcov

- gcov
  - Many of you already know (comes with gcc) gcc.gnu.org/onlinedocs/gcc/Gcov.html
  - Adds invocation points to each line of code
- lcov
  - Set of Linux tools to derive human-readable coverage results
  - genhtml to create HTML report

#### gcov

- Hint: Clean all .gcno results from previous runs
- 2 Add -coverage switch to gnatmake
   (or alternatively -f -cargs -fprofile-arcs
   -ftest-coverage -largs -fprofile-arcs)
- Run your program once
- 4 Run gcov to generate coverage results (\*.gcno) (there was a reason why we used lcov -gcov-tool gcov instead)

#### lcov

- Run lcov with the -directory switch
- 2 Run genhtml
- Hint: Remove uninteresting packages with lcov -r before running genhtml

# Hoare Logic

#### Hoare Logic

- Grammar for expressions, assignments, conditions, etc.
- From 1969 but very close to today programming
- Set of rules to transform expressions
- Task: Derive post-conditions from pre-conditions and transformation rules

```
function Exponentiate(N: Natural; B: Natural) is
  {Pre-Condition := N >= 0, B >= 0}
  X: Natural;
  Y: Natural;
begin
    X := 0;
  Y := 1;
  while X /= N loop
    X := X + 1;
    Y := Y * B;
  end loop;
  {Post-Condition := Y = B^N}
  return Y;
end Exponentiate;
```

#### **Program Correctness**

- A program is P correct ⇔ for all possible inputs that fulfill the pre-conditions, P always fulfills all post-conditions.
- Partial correctness:
  - P is correct if it terminates.
- Total correctness:
  - P is correct and always terminates.

- How can one prove that a loop terminates?
- $\Rightarrow$  Loop variant  $\vee$ :
  - Decreases with every iteration
  - Loop is exited when V <=</p>

```
{N > 0, B > 3,
X := 0;
Y := 1;
{X = 0, Y = 1}
{Y = B^X} (Loop Invariant)
while X /= N loop
                                                                                   \{X /= N\} (Condition)
                                                                                   {X = X'Old + 1} (Assignment)
Loop is exited when V <=
0
Trivial for for loops:
V := Range'Last - I

Y := Y * B;
{Y = Y'Old * B} (Assignment)
{Y = B^X} (Loop Invariant)
end loop;
{X = N} (Inverse Condition)
{Y = B^X} (Loop Invariant)
{Y = B^N} (Implication)</pre>
```

#### Loops

- How to prove that loops lead to correct post-conditions?
- $\Rightarrow$  Loop invariant IV :=  $Y = B^X$ :
  - Must be valid before the loop
  - Must be valid at each loop iteration

```
X := 0;
Y := 1;
{X = 0, Y = 1}
{Y = B^X} (Loop Invariant)
  \{X /= N\} (Condition)
   {X = X'Old + 1} (Assignment)
  Y := Y * B;

{Y = Y'Old * B} (Assignment)

{Y = B^X} (Loop Invariant)
{I = B A; (LOOP III)
end loop;
{X = N} (Inverse Condition)
{Y = B^X} (Loop Invariant)
{Y = B^N} (Implication)
```

