

Relational algebra operations

Operation	Notation	Definition	Schema Description	Degree	Cardinality (c)
Selection	$\sigma_\theta(r)$ Returns tuples of r satisfying the predicate θ .		R	n	$0 \leq c \leq N$
Projection	$\pi_{A_1, A_2 \dots}(r)$ Returns tuples of r consisting only of attributes $A_1, A_2 \dots$ (eliminates duplicates).		$\{A_1, A_2 \dots\}$	$\ \{A_1, A_2 \dots\}\ $	$1 \leq c \leq N$
Union ¹	$r \cup s$ Returns all tuples of r and all of s (eliminates duplicates).		R	n, m	$\max\{N, M\} \leq c \leq N + M$
Difference ¹	$r - s, r \setminus s$ Returns all tuples of r that are not in s .		R	n, m	$\max\{0, N - M\} \leq c \leq N$
Cartesian product	$r \times s$ Concatenates each tuple of r with each tuple of s .		R, S	$n + m$	NM
Intersection ¹	$r \cap s$ $r - (r - s)$ Returns all tuples of r that are also in s .		R	n, m	$0 \leq c \leq \min\{N, M\}$
Theta join	$r \bowtie_\theta s$ $\sigma_\theta(r \times s)$ Returns tuples of cartesian product of r and s satisfying the predicate θ .		R, S	$n + m$	$0 \leq c \leq NM$
Equijoin	$r \bowtie_{\theta_=} s$ $\sigma_{\theta_=}(r \times s)$ Returns tuples of cartesian product of r and s satisfying the predicate $\theta_=$.		R, S	$n + m$	$0 \leq c \leq NM$
Natural join ²	$r \bowtie s$ $\pi_{R \cup S}(\sigma_{r, A_1=s, A_2 \dots}(r \times s))$ Returns equijoin of r and s over all common attributes where one occurrence of each common attribute is eliminated.		$R \cup S$	$n + m - cnm$	$0 \leq c \leq NM$
Left outer join	$r \bowtie s$ /		$R \cup S$	$n + m - cnm$	$N \leq c \leq NM$
Right outer join	$r \bowtie s$ /		$R \cup S$	$n + m - cnm$	$M \leq c \leq NM$
Full outer join	$r \times s$ /		$R \cup S$	$n + m - cnm$	$N + M \leq c \leq NM$
Semijoin	$r \triangleright_\theta s$ $\pi_R(\sigma_\theta(r \times s))$ Returns tuples of r that participate in theta join of r and s (due to predicate θ).		R	n	$0 \leq c \leq N$
Division ⁴	$r/s, r \div s$ $\pi_{R-S}(r) - \pi_{R-S}((\pi_{R-S}(r) \times s) - r)$ Returns tuples of r consisting only of attributes in $R - S$ (denoted x), thus that for every tuple in s (denoted y) there exists a tuple in r equal to concatenation of x and y .		$R - S$	$n - m$	$0 \leq c \leq N/M$
Aggregate	$\tau_{AL}(r)$ /	Applies aggregate function list AL to the relation r .		$ AL $	1
Grouping	$GA \tau_{AL}(r)$ /	Groups tuples of r by grouping attributes GA and then applies aggregate function list AL to these groups.		$ GA + AL $	$1 \leq c \leq \prod_{A_i \in GA} \text{dom}(A_i) $
Rename	$\rho_{r(A_1, A_2 \dots)}(s)$ /	Renames the relation s and all of its attributes (new names are r and $A_1, A_2 \dots$).	$\{A_1, A_2 \dots\}$	$\ \{A_1, A_2 \dots\}\ $	M
Assignment	$r(A_1, A_2 \dots) \leftarrow s$ /	Assigns name r to relation s and $A_1, A_2 \dots$ to its attributes.	$\{A_1, A_2 \dots\}$	$\ \{A_1, A_2 \dots\}\ $	M

$r, s \rightarrow$ relations, with relation schemas R, S ($Sh(r) = R, Sh(s) = S$)

$n, m \rightarrow$ degrees of the relations ($\deg(r) = n, \deg(s) = m$)

$cnm \rightarrow$ number of common attributes in the relations r and s

$N, M \rightarrow$ cardinalities of the relations ($\text{card}(r) = N, \text{card}(s) = M$)

$A_i \rightarrow$ attribute

$\theta \rightarrow$ logical predicate

$\theta_= \rightarrow$ logical predicate which contains only equality comparisons and logical conjunction \wedge

$AL \rightarrow$ list of aggregate functions⁵ and corresponding attributes (e.g. "COUNT $A_1, SUM A_2 \dots$ ")

$GA \rightarrow$ grouping attributes (e.g. " $A_1, A_2 \dots$ "). They must differ from those in AL

⁰Author takes no responsibility for the errors in the text.

¹ r and s must be union-compatible.

²When r and s have no common attributes, natural join equals cartesian product.

³Unknown values are set to $NULL$.

⁴Attributes of s must be a subset of attributes of $r, S \subseteq R$.

⁵Can be one of COUNT, SUM, AVG, MAX or MIN.