# The Formal Semantics of Programming Languages / Chapter1-3

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## **Diagonal Argument**

#### Claim

A set X and  $\mathcal{P}ow(X)$  are never in 1-1 correspondence for any X.

#### Proof.

Proof by contradiction. Consider a set X and its powerset  $\mathcal{P}ow(x)$ . Let  $\theta: X \to \mathcal{P}ow(X)$  be a 1-1 correspondence between X and  $\mathcal{P}ow(X)$ . Suppose  $Y = \{x \in X \mid x \notin \theta(x)\}$ . Y is a subset of X and therefore in correspondence with a  $y \in X$ . So  $\theta(y) = Y$ . Thus either  $y \in Y$  or  $y \notin Y$ .

- If  $y \in Y$  then  $y \notin Y$ , because  $y \notin \theta(y)$ .
- If  $y \notin Y$  then  $y \in \theta(y)$ , so  $y \in Y$ .

In either case, we have contradiction.

## **Diagonal Argument**

#### Claim

A set X and  $\mathcal{P}ow(X)$  are never in 1-1 correspondence for any X.

#### Example

- $X = \{1, 2, 3\}$
- $\blacksquare \mathcal{P}ow(X) = \{ \phi, \{1\}, \{2\}, \{3\}, \{1,2\}, \{1,3\}, \{2,3\}, \{1,2,3\} \}$

but, if X is infinite

- $X = \{1, 2, 3, 4, \cdots\}$
- $Pow(X) = \{\phi, \{1\}, \{2\}, \{3\}, \cdots \}$

dose claim still hold?

answer is YES

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## Why it's Called "Diagonal" Arugument

consider following table where *i*th row and *j*th column is placed 1 if  $x_i \in \theta(x_j)$  and 0 otherwise

	$\theta(x_0)$	$\theta(x_1)$	$\theta(x_2)$	 $\theta(x_j)$	
$\overline{x_0}$	0	1	1	 1	
$x_1$	1	1	0	 0	
$x_2$	0	0	1	 1	
:	:	:	:	:	
$x_i$	0	1	0	 1	
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in above table, we can define Y by pick up  $\boldsymbol{x}_n$  which corresponding cell along diagonal is  $\mathbf{0}$ 

$$Y = \{x_0\}$$

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## **Direct and Inverse Image of Relation**

### Definition (direct and inverse image)

Let be  $R: X \times Y$  is a relation, A is a subset of X and B is a subset of Y, a set RA and  $R^{-1}B$  are defined as follows:

- $\blacksquare RA = \{ y \in Y \mid \exists x \in A((x, y) \in R) \},\$
- $\blacksquare \ R^{-1}B = \{x \in X \mid \exists y \in B((x,y) \in R)\}.$

The set RA is called direct image of A under R, and the set  $R^{-1}B$  is called inverse image of B under R.

#### **Example**

- $X = \{0, 1, 2, 3, 4\}, Y = \{5, 6, 7, 8, 9\}, R = \{(0, 9), (1, 5), (2, 6), (3, 6)\}, A = \{1, 3, 4\}, B = \{5, 7, 9\}$
- $\blacksquare RA = \{5, 6\}, R^{-1}B = \{0, 1\}$

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# **Equivalence Class**

## Definition (equivalence class)

If R is an equivalence relation on a set X then (R-) equivalence class  $\{x\}_R$  of an element  $x \in X$  is defined as follows;

$$\{x\}_R = \{y \in X \mid yRx\}.$$

### Example

- $\blacksquare R = \{(x, y) \in \mathbb{Z} \times \mathbb{Z} \mid x \equiv y \pmod{3}\}$
- $\blacksquare \{1\}_R = \{1, 4, 7, 10, \cdots\}$

## **Equivalence Relation**

## Definition (equivalence relation and equivalence class)

An equivalence relation is a relation  $R \subseteq X \times X$  on a set X which is

- reflexive:  $\forall x \in X(xRx)$ ,
- $\blacksquare$  symmetric:  $\forall x,y \in X(xRy \rightarrow yRx)$  and
- transitive:  $\forall x, y, z \in X((xRy \land yRz) \rightarrow xRz)$ .

#### Example

- $\blacksquare R = \{(x,y) \in B \times B \mid x \in A \land y \in A\} \text{ where } A \subseteq B$
- congruence of figures
- congruence of integers

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