## 

X	Y	Z
m	20	T
m	10	F
0	30	T
$\overline{n}$	30	T
0	20	T

Table 1: Legal instance of schema S for question 2.1

(a)	) Which of the following dependencies are <i>violated</i> by the instances of $S$ in Table 1?					
	i. [2 points] $\nabla$ Yes	$\square$ No : $X \to Y$ is violated.				
		$\square$ No : $Z \to X$ is violated.				
		$\nabla$ No : $Y \to Z$ is violated.				
	iv. [2 points] □ Yes	$\bigvee$ No : $XY \to Z$ is violated.				
	v. [2 points]  Yes	□ No : $YZ → X$ is violated. $□$ No : $XZ → Y$ is violated.				
	vi. [2 points]	$\square$ No : $XZ \rightarrow Y$ is violated.				
(b)	(b) [3 points] By only observing the instance of S in Table 1, can you identify the functional dependencies that hold on schema S? Why?					
	□ Yes No No	系的一个空侧不能推断出关系模式的函数该较				

Question 2: Functional Dependencies II  For the next set of questions consider the relational the set of functional dependencies FD:	
$Q \;  o \; U$	(1)
$U \rightarrow V$	(2)
$PQ \rightarrow W$	VST (3)
SU  o T	R (4)
$VT \rightarrow R$	W (5)
$R \rightarrow W$	(6)
(a) <b>[8 points]</b> Which of the following is a minim if none, mark accordingly, and give your <i>own</i> .  i. The given FDs (Eq 1-6), is a minimum co ii. $\{Q \to U, U \to V, PQ \to S, SU \to T, SU \in \{Q \to U, U \to V, PQ \to S, SU \to T, PQ \in \{Q \to U, U \to V, PQ \to S, SU \to T, V\}$	answer. ver already. $V \to R, VT \to R, VT \to W, R \to W$ $Q \to W, VT \to R, PQ \to T, R \to W$
v. $\{Q \to U, U \to V, PQ \to S, SU \to T, SU$ vi. none of the above - the cover is	$Y \to R, VT \to R, PQ \to T, R \to W$
<ul> <li>(b) Yes/No: Which of the following functional dependencies (Eq. (1)-(6))?</li> <li>i. [3 points] □ Yes □ No: Q→V</li> <li>ii. [3 points] □ Yes □ No: QU → R</li> <li>iii. [3 points] □ Yes □ No: SQ → T</li> <li>iv. [3 points] □ Yes □ No: SQ → W</li> <li>v. [3 points] □ Yes □ No: PQ → R</li> <li>vi. [3 points] □ Yes □ No: VT → Q</li> <li>(c) [3 points] □ True or False: The attribute closure</li> <li>□ True □ False</li> </ul>	e $\{Q\}^+$ is $\{Q,U,V\}$ .
(d) [3 points] True or False: The attribute closure  □ True □ False	$\{PQ\}^+ \text{ is } \{P,Q,W,S,T\}.$

Question 3: Decompositions.....[20 points]

For this set of questions, consider the relation with attributes,  $\mathcal{X} = \{A, B, C, D, E, F\}$ , Let the following functional dependencies FD be defined over the relation  $\mathcal{X}$ :

$$A \to B$$
$$B \to CD$$
$$E \to F$$

- (a) [2 points] Provide the attribute closure of  $\{AB\}$ .  $\{AB\}^{\dagger} = \{A, B, C, D\}$ (b) Consider the decomposition AB, BCD, EF. Mark 'True' or 'False':
- - i. [3 points] □ True □ False: It is lossless
  - ii. [3 points] \to True □ False: It is dependency-preserving
- (c) Consider the decomposition AB, BCDF, EF. Mark 'True' or 'False':
  - False: It is lossless i. [3 points] □ True
- (d) Consider the decomposition ABCEF, EBD. Mark 'True' or 'False':
  - i. [3 points] True  $\Box$  False: It is lossless
  - ii. [3 points] ✓ True □ False : It is dependency-preserving

Cons	ion 4: Normal Forms
	$PQ \rightarrow R$ (7)
	$PQ \rightarrow S$ (8)
	$R \rightarrow P$ (9)
	$S \rightarrow Q$ (10)
(a)	[6 points] List all the candidate key(s) for $\mathcal{E}$ . $\{P, Q\} \} \{R, S\} \}$
	[2 points] Is the relation $\mathcal{E}$ in BCNF? $\square$ Yes $\bigvee$ No
(c)	From the list below, select all applicable choices to justify whether $\mathcal E$ is (or is not) in BCNF.
	<b>Note</b> : when we refer to the <i>main requirement</i> for BCNF, we mean: <i>every determinant is a super key</i> .
	i. [1 point]   True False: All FD's satisfy the main requirement.
	ii. [1 point]   True False: FD (7) violates the main requirement.
	iii. [1 point] $\Box$ True $\checkmark$ False : FD (8) violates the main requirement.
	iv. [1 point] True
	v. [1 point] True
d)	[2 points] Is the relation $\mathcal{E}$ in 3NF? $\bigvee$ Yes $\square$ No
e)	From the list below, select all applicable choices to justify whether $\mathcal{E}$ is (or is not) in 3NF.
	Note: when we refer to the secondary requirement for 3NF, we mean: for every FD
	$X \to A$ , A is part of a candidate key. i. [1 point] $\ \Box$ True $\ \Box$ False: All FD's satisfy the secondary requirement.
	ii. [1 point]  True  False : All FD's satisfy the secondary requirement.
	iii. [1 point] $\Box$ True $\bigvee$ False : FD (8) violates the secondary requirement.
	iv. [1 point] $\Box$ True $\nabla$ False : FD (9) violates the secondary requirement.
	v. [1 point] $\square$ True $\square$ False : FD (10) violates the secondary requirement.
(f)	[5 points] Give a 3NF decomposition of $\mathcal{E}$ that is lossless, dependency preserving, and has as few tables as possible.
(g)	[8 points] Give a BCNF decomposition of E that is iossiess, and has as few tables as $\{R, P, S, P, P, S, P, P,$
_	possible.
<b>①</b>	R-> P 在 E 上违反BCNF 治 E 为例为 A = {R.P}
	A, noththan FDs为 R-> P
	A2的缺为(RS)(RQ) FDs为5->R
3	5-1R在A上违反BCNF 格品分词为[5. R] 和[5. R]
	此时为各几个要求
r.)	最终分射结果为 [R.P] 15.R], 15.Q)