## Object Oriented Scientific Programming in C++ (WI4771TU)

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Numerical Analysis



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#### Overview

- Last lecture we started with polymorphism, that is, inheritance of one class from another class
  - Implement common functionality in base class (possibly realised as abstract class that cannot even be instantiated)
  - Derive specialised class(es) from the base class that
    - Implement the missing functionality (pure virtual functions)
    - Overwrite generic functionality by specialised variants (virtual functions)
    - Reuse all other functionality from the base class

### Overview, cont'd

- Today, a more careful view on polymorphism
  - Static polymorphism: static binding/method overloading
  - Dynamic polymorphism: dynamic binding/method overwriting
- C++11/14/17? Auto functionality
- Template meta programming
  - A powerful variant of static polymorphism

#### Task: Calculator

- Write a class (or for demonstrating purposes a hierarchy of classes) that provide(s) a member function to calculate the sum of two and three integer values, respectively
  - Use static polymorphism: method overloading
  - Use dynamic polymorphism: method overwriting

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### Static polymorphism

Method overloading (at compile time)

```
class Calc {
    public:
    int sum(int a, int b) { return a+b; }
    int sum(int a, int b, int c) {return sum(sum(a,b),c);}
};
```

 Class Calc has two member functions with identical names but different interface; it is decided at compile time which of the two functions should be called

```
std::cout << C.sum(1,2) << std::endl;
std::cout << C.sum(1,2,3) << std::endl;</pre>
```

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### Static polymorphism, cont'd

Method overloading (not working!)

```
class Calc {
    public:
    int sum(int a, int b) { return a+b; }
    void sum(int a, int b) {std::cout<< a+b <<std::endl;}
};</pre>
```

• Difference must be in the interface of arguments passed to the functions since compiler cannot distinguish between the two sum functions if they only differ in the return type

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### Static polymorphism, cont'd

 Method overloading: decision about which method to call is made at compile time; hence the compiler can decide to inline code to improve performance (no overhead due to function calls/copy of data to and from the stack!)

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### Dynamic polymorphism

 Method overwriting: reimplement a function inherited from base class with new function body (same interface!)

```
class BaseCalc {
public:
  int sum2(int a, int b) { return a+b; }
  int sum3(int a, int b, int c)
      { return sum2(sum2(a,b),c); }
class DerivedCalc: public BaseCalc {
public:
  int sum2(int a, int b) { return a+b; }
```

### Dynamic polymorphism, cont'd

Why?

```
DerivedCalc D;
std::cout << D.sum2(1,2) << std::endl;</pre>
-> DerivedCalc::sum2(a,b)
std::cout << D.sum3(1,2,3) << std::endl;</pre>
-> BaseCalc::sum3(a,b,c)
-> BaseCalc::sum2(a,b)
-> BaseCalc::sum2(a,b)
```

### Dynamic polymorphism, cont'd

- sum2 function is declared and implemented in the base class
   BaseCalc and the derived class DerivedCalc but not as virtual
  - D.sum2(1,2) calls the sum2 function from the derived class
    DerivedCalc::sum2(int a, int b)
    { return a+b; }
  - D.sum3(1,2,3) calls the sum3 function from the base class,
     which itself calls the sum2 function from the base class twice

```
BaseCalc::sum2(int a, int b)
     { return a+b; }
BaseCalc::sum3(int a, int b, int c)
     { return sum2(sum2(a,b),c);}
```

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### Dynamic polymorphism, cont'd

 Method overwriting: virtual specifier indicates that the sum2 function can be overwritten in a derived class

```
class BaseCalc {
public:
 virtual int sum2(int a, int b) { return a+b; }
  int sum3(int a, int b, int c)
      { return sum2(sum2(a,b),c); }
class DerivedCalc: public BaseCalc {
public:
  int sum2(int a, int b) { return a+b; }
```

### Dynamic polymorphism, cont'd

Now:

```
DerivedCalc D;
std::cout << D.sum2(1,2) << std::endl;</pre>
-> DerivedCalc::sum2(a,b)
std::cout << D.sum3(1,2,3) << std::endl;</pre>
-> BaseCalc::sum3(a,b,c)
-> DerivedCalc::sum2(a,b)
-> DerivedCalc::sum2(a,b)
```

### Static polymorphism, cont'd

- Method overwriting: decision about which virtual method to call is made at run time; hence no inlining possible
- Common design pattern
  - Specify expected minimal functionality of a group of classes in abstract base class via pure virtual member functions
  - Implement generic common functionality of a group of classes abstract base class via virtual member functions
  - Implement expected functionality of a particular class by overwriting the pure virtual member function

### Example: inner product space

 In linear algebra, an inner product space is a vector space V that is equipped with a special mapping (inner product)

$$\langle \cdot, \cdot \rangle : V \times V \longrightarrow \mathbb{R} \text{ or } \mathbb{C}$$

Inner product spaces have a naturally induced norm

$$\|x\| = \sqrt{\langle x, x \rangle}$$

### Example: inner product space, cont'd

 Class InnerProductSpaceBase declares inner product as pure virtual and implements the naturally induced norm

```
class InnerProductBase {
public:
   pure virtual double inner_product(... x,... y) = 0;
   double norm(x) { return inner_product(x,x); }
};
```

Derived InnerProductSpace class implements inner product

```
class InnerProductSpace : public InnerProductSpaceBase {
  public:
    double inner_product(... x, ... y) = { return x*y; }
};
```

#### Task: Calculator2

- Extend the calculator class so that it can handle numbers of integer, float and double type at the same time
  - Prevent manual code duplication
  - Prevent explicit type casting
  - Make use of auto-functionality (C++11/14/17?)
  - Make use of template meta programming

### Vanilla implementation in C++

```
Class Calc2 {
public:
   int sum(int a, int b)
       { return a+b; }
   int sum(int a, int b, int c)
       { return sum(sum(a,b),c); }
int main() {
   Calc2 C;
   std::cout << C.sum(1,2) << std::endl;</pre>
   std::cout << C.sum(1,2,3) << std::endl;</pre>
```

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### Automatic return type deduction (C++11)

Explicit definition of the function return type

```
int sum(int a, int b)
{ return a+b; }
```

Automatic function return type (since C++11)

```
auto sum(int a, int b) -> decltype(a+b)
     { return a+b; }
```

- Using decltype, the return type of the sum function is determined automatically as the type of operator+(a,b)
- Feature is termed cxx\_trailing\_return\_types in Cmake (see target\_compile\_features line in CMakeLists.txt file)

### Automatic return type deduction, cont'd

Implementation in C++11

```
class Calc2 {
public:
  virtual auto sum(int a, int b) -> decltype(a+b)
      { return a+b; }
  auto sum(int a, int b, int c) -> decltype(a+b+c)
      { return sum(sum(a,b),c); }
int main() {
  Calc2 C;
  std::cout << C.sum(1,2) << std::endl;</pre>
  std::cout << C.sum(1,2,3) << std::endl;</pre>
```

### Automatic return type deduction, cont'd

decltype specifier (C++11) queries the type of an expression

```
class Calc2 {
public:
  virtual auto sum(int a, int b) -> decltype(a+b)
      { return a+b; }
  auto sum(int a, int b, int c) -> decltype(sum(sum(a,b),c))
      { return sum(sum(a,b),c); }
int main() {
  Calc2 C;
  std::cout << C.sum(1,2) << std::endl;</pre>
  std::cout << C.sum(1,2,3) << std::endl;</pre>
```

### Automatic type deduction (C++14)

• C++14 (not fully implemented by all compilers) allows you to even deduce the type of parameters automatically

```
auto sum(int a, int b)  // no -> decltype(a+b)
    { return a+b; }
auto sum(int a, int b, int c) // no -> decltype(a+b+c)
    { return sum(sum(a,b),c); }
```

 Remark: Feature helps to improve readability and prevents deduction errors (due to forgotten/inconsistent deduction rule by the programmer) but it does not solve the problem to pass arguments of different type to the same function

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### Generic functions (C++17?)

Future C++ standards a likely to support the following code

```
auto sum(auto a, auto b)
    { return a+b; }
auto sum(auto a, auto b, auto c)
    { return sum(sum(a,b),c); }
```

Remark: above code already compiles with GNU g++ 5.x

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### **Function templates**

- Until C++17?, template meta programming is the standard technique to deal with arbitrary function parameters
- Function templates: allow you to implement so-called parameterised functions for generic parameter types

```
template<typename R, typename A, typename B>
R sum(A a, B b)
{
   return a+b;
}
```

### Function templates

Types must be specified explicitly when function is called

```
int s1 = sum<int, int, int>(1, 2);
double s2 = sum<double, double, int>(1.2, 2);
double s3 = sum<double, float, double>(1.4, 2.2);
```

This can be slightly simplified using the auto specifier

```
auto s1 = sum<int, int, int>(1, 2);
auto s2 = sum<double, double, int>(1.2, 2);
auto s3 = sum<double, float, double>(1.4, 2.2);
```

### Function templates, cont'd

C++11: automatic return type deduction

C++14: automatic type deduction

Usage

```
auto s1 = sum < int, int > (1, 2);
```

### Function templates, cont'd

How to convert this function into a templated function

```
int sum(int a, int b, int c)
{
    return sum(sum(a,b), c);
}
```

### Function templates, cont'd

Use explicit return type parameter (ugly!)

```
template<typename R, typename A, typename B, typename C>
auto sum(A a, B b, C c)
{
    return sum<R,C>(sum<A,B>(a,b), c);
}
```

Guess what this function call will return

```
auto s1 = sum<int,double,double,double>(1.1,2.2,3.3)
```

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### Function templates, cont'd

Here it is, the holy grail

```
template<typename A, typename B>
auto sum(A a, B b) -> decltype(a+b) // omit in C++14
    return a+b;
template<typename A, typename B, typename C>
auto sum(A a, B b, C c)
    return sum<decltype(sum<A,B>(a,b)),C>
              (sum<A,B>(a,b), c);
```

### Function templates, cont'd

Now, we can call sum functions as follows

```
auto s1 = sum<int, int>(1, 2);
auto s2 = sum<double, int>(1.2, 2);
auto s3 = sum<float, double>(1.4, 2.2);
```

- Since the compiler needs to duplicate code and substitute A,B,C for each combination of templated types both the compile time and the size of the executable will increase
- Template meta programming is simplest of the code resides in header files only; later we will see how to use TMP in combination with pre-compiled libraries

### Task: generic Vector class

- Write Vector class that can store real values (float/double) and complex values and supports the following operations:
  - Addition of two vectors of same length (possibly different type)
  - Multiplication of a vector with a scalar (possible different type)
  - Dot product of a vector with another one (possibly different type)

### Vector class prototype

Implementation of Vector-of-double class

```
class Vector {
private:
   double* data;
   int n;
public:
   Vector() : n(0), data(nullptr) {}
   Vector(int n) : n(n), data(new double[n]) {}
   ~Vector() { n=0; delete[] data; }
```

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### Vector class prototype, cont'd

```
Vector& operator+=(const Vector& other) {
    for (auto i=0; i<n; i++)
        data[i] += other.data[i];
    return *this; }
Vector& operator*=(double scalar) {
    for (auto i=0; i<n; i++)
        data[i] *= scalar;
    return *this; }
double dot(const Vector& other) const {
    double d=0;
    for (auto i=0; i<n; i++)
        d += data[i]*other.data[i];
    return d; }
```

### **Brainstorming**

- Function templates alone will not help since the type of a class attribute needs to be templated -> class templates
- Some member functions can be implemented generically, e.g., addition of two vectors and multiplication of a vector with a scalar value since they are the same for all types
- Some member functions must be implemented in different manners for real and complex values -> specialisation

$$x \cdot y = \sum_{i=1}^{n} x_i y_i, \quad x, y \in \mathbb{R}, \qquad x \cdot y = \sum_{i=1}^{n} x_i \overline{y}_i \quad x, y \in \mathbb{C}$$

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### Class template

Implementation of Vector-of-anything class

```
template<typename T>
class Vector {
private:
   T* data;
   int n;
public:
   Vector() : n(0), data(nullptr) {}
   Vector(int n) : n(n), data(new T[n]) {}
   ~Vector() { n=0; delete[] data; }
```

### Class template, cont'd

Template parameter must be explicitly specified

```
Vector<int> x(10); // Vector-of-int with length 10
Vector<double> y; // Empty Vector-of-double
Vector<float> z(5); // Vector-of-float with length 5
```

 Remark: if you want to pass a Vector-of-anything to a function in the templated class Vector you have to write

```
Vector<T> v
Vector<T>& v
instead of
Vector v
Vector& v
```

```
template <typename T>
class Vector {
public:
    Vector<T>& operator+=(const Vector<T>& other)
        (for auto i=0; i<n; i++)
            data[i] += other.data[i];
        return *this;
```

```
template <typename T>
class Vector {
public:
    Vector<T>& operator*=(T scalar)
         (for auto i=0; i<n; i++)
            data[i] *= scalar;
        return *this;
```



```
template <typename T>
class Vector {
public:
    T dot(const Vector<T>& other) const
        T d=0;
        for (auto i=0; i<n; i++)
            d += data[i]*other.data[i];
        return d;
```

With the single class template parameter T we can do

```
Vector<int> x1(5), x2(5);
x1 += x2;
x1 *= (int)2;
int x1 = x1.dot(x2);
```

How about?

```
Vector<int> x2(5);
x2 *= (double)1.2;
```



#### Intermezzo

Class templates and function templates can be combined

```
template<typename T>
class Vector {
    template<typename S>
    Vector<T>& multiply<S scalar>
        (auto i=0; i<n; i++)
            data[i] *= S;
        return *this;
```

### Intermezzo, cont'd

At first glance, this seems to be more flexible

```
Vector<double> x1(5);
x1.multiply<int>(5);
```

But be really careful since strange things can happen

```
Vector<int> x1(5);
x1.multiply<double>(5.5);
```

• Rule of thumb: before using extensive templating like this think about all(!) implications; it can help to think if the planned functionality has a meaningful mathematical counterpart, e.g. dot product of  $x \in \mathbb{R}, y \in \mathbb{C}$ 

### Specialisation

The dot product needs special treatment since

```
T dot(const Vector<T>& other) const
{
    T d=0;
    for (auto i=0; i<n; i++)
        d += data[i]*other.data[i];
    return d;
}</pre>
```

- Lacks the complex conjugate of other and yields the wrong return type in case of complex-valued vectors
- Remedy: implement a specialised variant for this case

### Specialisation, cont'd

Generic dot product implemented in Vector class

```
#include <complex>
template<typename T>
class Vector { ...
   T dot(const Vector<T>& other) const {...}
};
```

 This function is used whenever no specialised implementation for a concrete type is available

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### Specialisation, cont'd

Specialised dot product for Vectors-of-complex-float

```
template<>
std::complex<float> Vector< std::complex<float> >::
    dot(const Vector< std::complex<float> >& other) const
        std::complex<float> d=0;
        // special treatment of dot product
        for (auto i=0; i<n; i++)
            d += data[i]*std::conj(other.data[i]);
        return d;
```

### Specialisation, cont'd

Current implementation yields

```
Vector<float> x1(5), x2(5);
auto x1.dot(x2); // calls generic implementation
Vector<std::complex<float> > y1(5), y2(5);
auto y1.dot(y2); // calls specialised implementation
Vector<std::complex<double> > z1(5), z2(5);
auto z1.dot(z2); // calls generic implementation
auto x1.dot(y1); // does not compile(!)
```

#### Outlook on next session

C++ allows you to partially specialise class templates
 Note that this code will not compile (we will see why!)

• Welcome to where the magic begins!