

### 105. Construct Binary Tree from Preorder and Inorder Traversal

We will use recursion to build `TreeNode`, each time we will separate inorder and preorder array.

TC is  $O(n * n)$

class Solution:

```
def buildTree(self, preorder: List[int], inorder: List[int]) -> TreeNode:
    if not preorder:
        return None
    root = TreeNode(preorder[0])
    index = inorder.index(preorder[0])
    inorder_left = inorder[:index]
    inorder_right = inorder[index + 1:]
    preorder_left = preorder[1:len(inorder_left) + 1]
    preorder_right = preorder[len(inorder_left) + 1:]
    root.left = self.buildTree(preorder_left, inorder_left)
    root.right = self.buildTree(preorder_right, inorder_right)
    return root
```

### 2. Random Pick with Weight

We will accumulate weights from head to end. And then we will use binary search to get the random number from 1 to largest weight. Get largest number that less than or equal to the random number. In the end, we will return its index. TC is  $O(\log n)$

from bisect import \*

from random import \*

class Solution:

```
def __init__(self, w: List[int]):
    self.weights = []
    cur_sum = 0
    for i in w:
        cur_sum += i
        self.weights.append(cur_sum)
    self.total = cur_sum

def pickIndex(self) -> int:
    return bisect_left(self.weights, randint(1, self.total))
```

### 239. Sliding Window Maximum

We will use a deque to manage our number, if  $\text{len}(\text{deque}) == k$ , we will `popleft()`, if the pop number is max, we will find a new `cur_max` from deque, then we will compare the new number with `cur_max`, if it's larger or equal to `cur_max`, we will clear deque, and let `cur_max = num`, and append num to memo. TC is  $O(kn)$

from collections import deque

class Solution:

```

def maxSlidingWindow(self, nums: List[int], k: int) -> List[int]:
    if not nums:
        return []
    if k == 1:
        return nums
    memo = deque(nums[:k])
    cur_max = max(memo)
    result = [cur_max]
    for num in nums[k:]:
        if num >= cur_max:
            cur_max = num
            memo.clear()
            memo.append(num)
        for num in nums[k:]:
            if len(memo) == k:
                temp = memo.popleft()
                if temp == cur_max:
                    cur_max = max(memo)
            if num >= cur_max:
                cur_max = num
                memo.clear()
                memo.append(num)
            result.append(cur_max)
    return result

```

#### 4. Maximum Subarray

We will accumulate all numbers from left to right, so every sub sum from  $j - 1$  to  $i$  is  $nums[i] - nums[j]$ , so we only need to maintain our  $cur\_min$  and  $cur\_max$  in every iteration. In the end, we will return  $cur\_max$

class Solution:

```

def maxSubArray(self, nums: List[int]) -> int:
    cur_min = min(nums[0], 0)
    cur_max = nums[0]
    for i in range(1, len(nums)):
        nums[i] += nums[i - 1]
        cur_max = max(cur_max, nums[i] - cur_min)
        cur_min = min(cur_min, nums[i])
    return cur_max

```

#### 5. Maximum Product Subarray

We will always remember  $cur\_min$  and  $cur\_max$ . And in each iteration, we will do  $result = (result, cur\_max * num, cur\_min * num)$ . TC is  $O(n)$

class Solution:

```

def maxProduct(self, nums: List[int]) -> int:

```

```
cur_min = min(1, nums[0])
cur_max = max(1, nums[0])
result = nums[0]
for num in nums[1:]:
    num1, num2 = num * cur_min, num * cur_max
    result = max(result, num1, num2)
    cur_min = min(1, num1, num2)
    cur_max = max(1, num1, num2)
return result
```