

### 108. Convert Sorted Array to Binary Search Tree

We will use recursion to construct this BST. TC is  $O(n)$ .

class Solution:

```
def sortedArrayToBST(self, nums: List[int]) -> TreeNode:
    if not nums:
        return None
    length = len(nums)
    if length == 1:
        return TreeNode(nums[0])
    left_nums = nums[:length // 2]
    right_nums = nums[length // 2 + 1:]
    root = TreeNode(nums[length // 2])
    root.left = self.sortedArrayToBST(left_nums)
    root.right = self.sortedArrayToBST(right_nums)
    return root
```

### 501. Find Mode in Binary Search Tree

Using inorder traversal. TC is  $O(n)$ , SC is  $O(1)$

class Solution:

```
def findMode(self, root: TreeNode) -> List[int]:
    self.count = 0
    self.result = []
    self.max = 0
    self.prev = None
    if not root:
        return self.result
    def inorderTraverse(node):
        if not node:
            return
        inorderTraverse(node.left)
        if node.val == self.prev or self.prev == None:
            self.count += 1
        else:
            if self.count > self.max:
                self.max = self.count
                self.result = [self.prev]
            elif self.count == self.max:
                self.max = self.count
                self.result.append(self.prev)
            self.count = 1
        self.prev = node.val
        inorderTraverse(node.right)
```

```

inorderTraverse(root)
if self.count > self.max:
    self.max = self.count
    self.result = [self.prev]
elif self.count == self.max:
    self.max = self.count
    self.result.append(self.prev)
return self.result

```

#### 239. Sliding Window Maximum

We will use a deque to manage our number, if  $\text{len}(\text{deque}) == k$ , we will `popleft()`, if the pop number is max, we will find a new `cur_max` from deque, then we will compare the new number with `cur_max`, if it's larger or equal to `cur_max`, we will clear deque, and let `cur_max = num`, and append num to memo. TC is  $O(kn)$

from collections import deque

class Solution:

```

def maxSlidingWindow(self, nums: List[int], k: int) -> List[int]:

```

```

    if not nums:
        return []
    if k == 1:
        return nums
    memo = deque(nums[:k])
    cur_max = max(memo)
    result = [cur_max]
    for num in nums[k:]:
        if num >= cur_max:
            cur_max = num
            memo.clear()
            memo.append(num)
    for num in nums[k:]:
        if len(memo) == k:
            temp = memo.popleft()
            if temp == cur_max:
                cur_max = max(memo)
        if num >= cur_max:
            cur_max = num
            memo.clear()
            memo.append(num)
        result.append(cur_max)
    return result

```

#### 4. Maximum Subarray

We will accumulate all numbers from left to right, so every sub sum from  $j - 1$  to  $i$  is  $\text{nums}[i] - \text{nums}[j]$ , so we only need to maintain our  $\text{cur\_min}$  and  $\text{cur\_max}$  in every iteration. In the end, we will return  $\text{cur\_max}$

class Solution:

```
def maxSubArray(self, nums: List[int]) -> int:
    cur_min = min(nums[0], 0)
    cur_max = nums[0]
    for i in range(1, len(nums)):
        nums[i] += nums[i - 1]
        cur_max = max(cur_max, nums[i] - cur_min)
        cur_min = min(cur_min, nums[i])
    return cur_max
```

### 5. Maximum Product Subarray

We will always remember  $\text{cur\_min}$  and  $\text{cur\_max}$ . And in each iteration, we will do  $\text{result} = (\text{result}, \text{cur\_max} * \text{num}, \text{cur\_min} * \text{num})$ . TC is  $O(n)$

class Solution:

```
def maxProduct(self, nums: List[int]) -> int:
    cur_min = min(1, nums[0])
    cur_max = max(1, nums[0])
    result = nums[0]
    for num in nums[1:]:
        num1, num2 = num * cur_min, num * cur_max
        result = max(result, num1, num2)
        cur_min = min(1, num1, num2)
        cur_max = max(1, num1, num2)
    return result
```