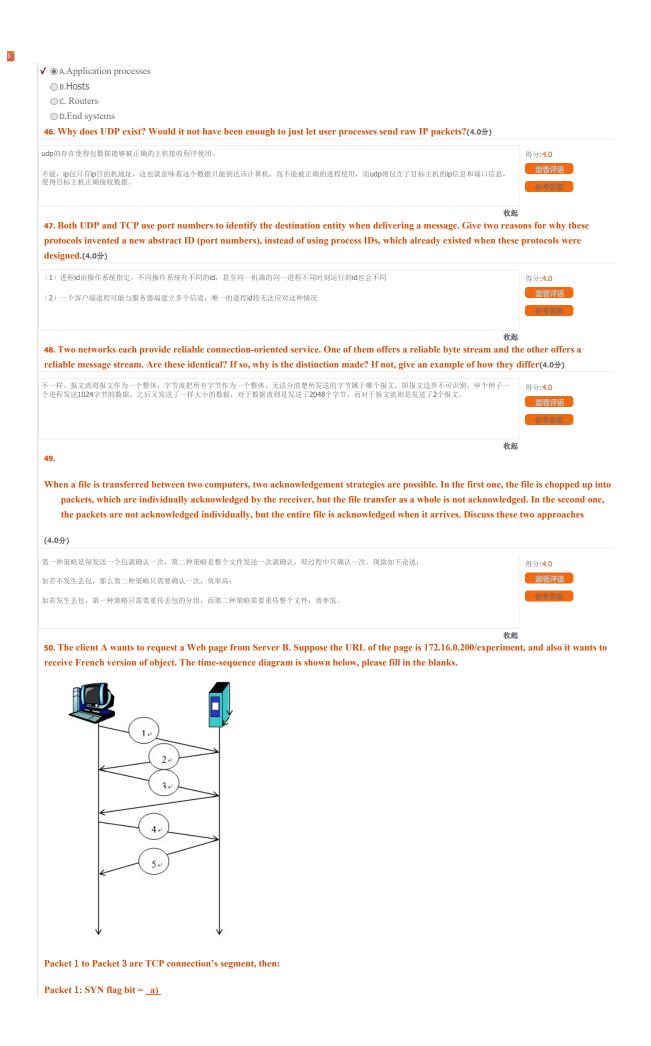
分组	
1.	
Th	
the transmission rate	otocol is highly inefficient when there is a large distance between source and destination and is high.
(1.0分)	
✓ A.True	
○ B.False	
2. Suppose that host A wa	nts to send data over TCP to host B, and host B wants to send data to host A over TCP. Two separate TCP
connections - one for each	direction - are needed.(1.0分)
○ A.True	
▼ B.False	
	uces the transmission rate at the sender when the receiver is overloaded(1.0分)
✓ ⑥ A.True	
	on timeout value is always set equal to the measured RTT.(1.0分)
O A.True	on timeout value is arrays see equal to the incasured KTT-(X.O.J.)
✓ B.False	
5.	
The Slow-Start algorit	thm increases a source's rate of transmission faster than "additive increase"
(1.0分)	
✓ A.True	
O B.False	
6. Flow control and conge	stion control are same at that they all limit the rate of the sender, but differ in that (1.0分)
✓ ● A.flow control limits its	rate by the size of RcvWindow, but congestion control by the traffic on the link
B.congestion control lin	nits its rate by the size of RcvWindow, but flow control by the traffic on the link
C.flow control mainly is	s accomplished by the sender, but congestion control by the receiver.
O.flow control mainly i	s accomplished by the receiver, but congestion control by the link.
	here is timeout event when the value of threshold is 32 and the size of congestion window is 16. According to the
	olicy, the new value of threshold and the new size of congestion window should be, respectively.(1.0分)
OA.16, 8	
OB.24, 8	
_ ′	
√ ⑥ c.8, 1	
✓	erval is a function of:(1.0分)
✓	erval is a function of:(1.0分)
√ ⊚ c.8, 1	sender
 ✓	sender d of a segment
✓ ● C.8, 1 ● D.16, 1 8. In TCP, the timeout int ✓ ● A. estimated RTT at the ■ B.MSS and the overhea	sender d of a segment ne receiver
✓ ● C.8, 1	sender d of a segment ne receiver indow TCP congestion control, which are(1.0分)
✓ ⑥ C.8, 1	sender d of a segment ne receiver indow TCP congestion control, which are(1.0分) stion avoidance
 ✓ ● C.8, 1 ● D.16, 1 8. In TCP, the timeout int ✓ ● A. estimated RTT at the ● B.MSS and the overhea ○ C.the size of buffer at th ○ D.the size of sending w 9. There are two states in ✓ ● A. slow start and congest ○ B.safe start and congest 	sender d of a segment he receiver indow TCP congestion control, which are(1.0分) stion avoidance ion avoidance
 ✓ ⑥ C.8, 1 ⑥ D.16, 1 8. In TCP, the timeout int ✓ ⑥ A. estimated RTT at the ⑥ B.MSS and the overhea ⑥ C.the size of buffer at th ⑥ D.the size of sending w 9. There are two states in ✓ ⑥ A. slow start and congest ⑥ B.safe start and congest ⑥ C.slow start and congest ⑥ C.slow start and congest 	sender d of a segment he receiver indow TCP congestion control, which are(1.0分) stion avoidance ion avoidance tion abandon
✓ ⑥ C.8, 1	sender d of a segment ne receiver indow TCP congestion control, which are
✓ ⑥ C.8, 1	sender d of a segment ne receiver indow TCP congestion control, which are
✓ ⑥ C.8, 1	sender d of a segment ne receiver indow TCP congestion control, which are
✓ ⑥ C.8, 1	sender d of a segment ne receiver indow TCP congestion control, which are
✓ ⑥ C.8, 1	sender d of a segment ne receiver indow TCP congestion control, which are
V ● C.8, 1	sender d of a segment ne receiver indow TCP congestion control, which are
✓ ⑥ C.8, 1	sender d of a segment he receiver indow TCP congestion control, which are(1.0分) stion avoidance ion avoidance tion abandon tion abandon perceive congestion?(1.0分) ent duplicate ACK-s event
V ● C.8, 1	sender d of a segment he receiver indow TCP congestion control, which are(1.0分) stion avoidance ion avoidance tion abandon tion abandon perceive congestion?(1.0分) ent duplicate ACK-s event
V ● C.8, 1	resender d of a segment the receiver indow TCP congestion control, which are(1.0分) stion avoidance tion avoidance tion abandon tion abandon perceive congestion?(1.0分) ent thuplicate ACK-s event descriptions about TCP connection management, which one is not correct?(1.0分) tessess participating in a TCP connection can end the connection
✓ ⑥ C.8, 1	sender d of a segment he receiver indow TCP congestion control, which are(1.0分) stion avoidance ion avoidance tion abandon tion abandon perceive congestion?(1.0分) ent duplicate ACK-s event descriptions about TCP connection management, which one is not correct?(1.0分) essesses participating in a TCP connection can end the connection o 1, it means that it wants to close the connection f the three-way handshake, the client and server randomly choose an initial sequence number of the three-way handshake, the SYN bit must be set to 1
V ● C.8, 1	resender d of a segment he receiver indow TCP congestion control, which are(1.0%) stion avoidance ion avoidance tion abandon tion abandon perceive congestion?(1.0%) ent duplicate ACK-s event Rescriptions about TCP connection management, which one is not correct?(1.0%) exesses participating in a TCP connection can end the connection o 1, it means that it wants to close the connection f the three-way handshake, the client and server randomly choose an initial sequence number
✓ ⑥ C.8, 1	sender d of a segment he receiver indow TCP congestion control, which are(1.0%) stion avoidance ion avoidance tion abandon tion abandon perceive congestion?(1.0%) ent duplicate ACK-s event descriptions about TCP connection management, which one is not correct?(1.0%) essesses participating in a TCP connection can end the connection of the three-way handshake, the client and server randomly choose an initial sequence number of the three-way handshake, the SYN bit must be set to 1
V ● C.8, 1	sender d of a segment he receiver indow TCP congestion control, which are(1.0%) stion avoidance ion avoidance tion abandon tion abandon perceive congestion?(1.0%) ent duplicate ACK-s event descriptions about TCP connection management, which one is not correct?(1.0%) essesses participating in a TCP connection can end the connection of the three-way handshake, the client and server randomly choose an initial sequence number of the three-way handshake, the SYN bit must be set to 1
✓ ⑥ C.8, 1	sender d of a segment he receiver indow TCP congestion control, which are(1.0%) stion avoidance ion avoidance tion abandon tion abandon perceive congestion?(1.0%) ent duplicate ACK-s event descriptions about TCP connection management, which one is not correct?(1.0%) essesses participating in a TCP connection can end the connection of the three-way handshake, the client and server randomly choose an initial sequence number of the three-way handshake, the SYN bit must be set to 1

13. In the following four descriptions about Rcv-Window, which one is correct? (1.0分)	
A.The size of the TCP RcvWindow never changes throughout the duration of the connection	
⊕ B.The size of the TCP RcvWindow will change with the size of the TCP RcvBuffer	
▼ (a) C.The size of the TCP RcvWindow must be less than or equal to the size of the TCP RcvBuffer	
On D. Suppose host A sends a file to host B over a TCP connection, the number of unacknowledged bytes that A sends can be a send of the connection of the number of unacknowledged bytes that A sends can be a send of the connection of the number of unacknowledged bytes that A sends can be a send of the connection of the connection of the number of unacknowledged bytes that A sends can be a send of the connection of the conne	annot exceed the size of the
size of the RcvWindow.	
14. TCP provides flow control by having the sender maintain a variable called the (1.0分)	
✓ A.Receive window	
B.Congestion window	
© C.Sliding window	
O D.buffer	
15 is a speeding-matching servicematching the rate which the sender is sending against the rate at which	the receiving application
is reading.(1.0分)	
O.A.congestion control	
✓ ⑥ B.flow control	
C.sliding-window control D.variable control	
16. Fast retransmit means in the case that duplicate ACK-s are received, the TCP sender resend segment between the transmit means in the case that	fore timer expires (1.04)
✓ ® A.3	fore timer expires.(1.03/)
0 B.4	
OC. 5	
© D.6	
17. What is the main difference between stop-and-wait and pipelined reliable data transfer protocol?(1.0分)	
A.The pipelined protocol uses the NAK packets, whereas in the stop-and-wait protocol senders always wait for ACK	packets.
✓ ⑥ B.With the pipelined protocol, the sender can send several packets in row, whereas in the stop-and-wait protocol the	sender cannot send the pack
in row.	
© C.With the pipelined protocol, the receiver must send one ACK for several packets (cumulative ACK), whereas in the	e stop-and-wait protocol the
receiver can not send the cumulative ACK.	
\bigcirc D.The pipelined protocol uses timeouts, whereas the stop-and-wait protocol does not use the timeout.	
18. Which of the following is not a pipelining protocol.(1.0分)	
√ (a) A.Rdt1.0	
○B.Go-Back-N	
© C.Selective repeat	
○ D.TCP	
19. If we define N to be the window size, base to be the sequence number of the oldest unacknowledged packet, an	•
smallest unused sequence number, then the interval [nextseqnum,base+N-1] corresponds to packet that(1.05	分)
✓ ⑥A. can be sent immediately	
B.have already been transmitted and acknowledged	
© C.cannot be used	
© D.have been sent but not yet acknowledged 20. Receive TCP only cally only cally only to the first missing byte in the street. TCP is said to provide the provider of the first missing byte in the street.	0/3
20. Because TCP only acknowledges bytes up to the first missing byte in the stream, TCP is said to provide(1 ✓ ⑥ A.Cumulative acknowledgements	1.0分)
B.Selective acknowledgements	
© C.3 duplicate ACKs	
© D.positive ACKs	
21. Suppose host A sends host B one TCP segment with sequence number 418, acknowledgement number 571, and	d 4 bytes of data. Then the
sequence number in the acknowledgement to this segment is (1.0%)	
OA.422	
● B.418	
√ ○ C.571	
OD.575	
22. There are two 16-bit integers: 1110 0110 0110 0110, 1101 0101 0101 01	
√ ⊚ A.0100010001000011	
© B.1011101110111100	
OC.1111111111111	
O D.1000000000000000000000000000000000000	
23. Provided α =0.125, current value of Estimated-RTT is 0.4s, Sample-RTT is 0.8s, then the new value of Estimated Part of the contract of t	ed-RTT iss.(1.0分)
√ ⊚ A.0.45	
© B.0.6	
© c.0.7	
O D.0.8	

24. In the following four options, which one is correct?(1.0分)
✓ ● A.The variations in the SampleRTT are smoothed out in the computation of the EstimatedRTT
○B. The timeout should be less than the connection's RTT
© C.Suppose that the last SampleRTT in a TCP connection is equal to 1 sec. Then the current value of TimeoutInterval will necessarily be≥1 sec.
D.Suppose that the last SampleRTT in a TCP connection is equal to 1 sec. Then the current value of TimeoutInterval will necessarily be≤1 sec.
25 is the byte sequence numbers of next byte expected from other side. (1.0分)
A.Sequence number
√ ® B.ACK number
○ C.Checksum
O.port number
26 is the byte stream number of first byte in the segment.(1.0%)
✓ A.Sequence number
OB.ACK number
○ c.Checksum
O D.port number
27.
Which of the following field is not used for connection setup and teardown?
(1.0分)
O.A.Sequence number
✓ ⑥B.TST
OC.SYN
O D.FIN
28. In the following four descriptions about MSS and MTU, which one is not correct?(1.0分)
A.The MSS is the maximum amount of application-layer data in the segment
✓ ○B.The MSS is the maximum size of the TCP segment including headers
O.T. The MSS is typically set by MTU
On The MTU means the largest link-layer frame
29. Which of the following about TCP connection is not correct?(1.0%)
✓ ● A.It is a broadcast connection
B.It is a point-to-point connection C. It is a mindlined connection.
Oc. It is a pipelined connection D.It is a full duplex connection
30. The field of Length in UDP segment specifies the length of . (1.0分)
A.the UDP segment, not including the header
✓
C.the UDP segment's header
On the Length field
31. The UDP header has only four fields, they are .(1.0分)
✓
B.Source port number, destination port number, source IP and destination IP
© C.source IP, destination IP, source MAC address and destination MAC address
D.source IP, destination IP, sequence number and ACK sequence number
32.
The port numbers ranging fromto are called well-known port number and are restricted.
(1.0分)
√ (a A.0, 1023
©B.0, 65535
©c.0, 127
©D,0,255
33. Port number's scope is to (1.0分)
©A.0, 1023
√ ® B.0, 65535
© C.0, 127
© 0.0,127
34. The following four descriptions about multiplexing and de-multiplexing, which one is correct?(1.0%)
A. A UDP socket is identified by a two-tuples consisting of a source port number and a destination port number.
✓ ○ B.If two UDP segment have different source port number, they may be directed to the same destination process.

© C.If two TCP segments with different source port number, they may be directed to the same destination process.	
© D.If two TCP segments with same destination IP address and destination port number, they must be the same TCP connection.	
35. This job of delivering the data in a transport-layer segment to the correct socket is called(1.0分)	
○ A.multiplexing✓ ® B.de-multiplexing	
© C.forwarding	
© D.routing	
36. The job of gathering data chunks, encapsulating each data chunk with header information to create segments and passing	the seaments
to the network is called(1.0分)	the segments
✓ ● A. multiplexing	
© B.de-multiplexing	
© C.forwarding	
© D.routing	
37. Which of the following applications normally uses UDP services?(1.0分)	
© A.SMTP	
▼ ® B.Streaming multimedia	
© C.FTP	
○ D.HTTP	
38. UDP offers which of the following benefits relative to TCP? (1.0分)	
▼	
B.UDP supports a self-regulating "throttle" feature that prevents network saturation	
© C.UDP guarantees that Individual packets of a transmission will arrive "in order"	
D.None of the above	
39. These two minimal transport-layer services andare the only two services that UDP provides(1.0分)	
▼	
B.congestion control, reliable data transfer	
C.flow control, congestion control	
D.In-order data transfer, error checking	
40. Which of the following services is not provided by TCP?(1.0分)	
√	
B.Reliable data transfers and flow controls	
© C. Congestion controls	
© D.In-order data transfers	
41. Services provided by transport layer include(1.0分)	
A.HTTP and FTP B.TCP and IP	
✓ © C.TCP and UDP	
© D. SMTP	
42. The receive side of transport layer reassembles segments into messages, passes tolayer.(1.0\(\frac{1}{2}\))	
✓ ②A.Application	
B.Networking	
© C.Physical	
© D.MAC	
43. In transport layer, the send side breaks application messages into , passes to network layer.(1.0分)	
© A.Frames	
✓ ⑥ B. Segments	
© C.Data-grams	
D. bit streams	
44.	
Transport-layer protocols run in	
(1.0%) A.Servers	
B.Clients	
© C.Routers	
✓ ● D.End systems	
45.	
A transport-layer protocol provides for logical communication between	
(1.0分)	



```
ACK flag bit = _b)
Sequence number = 92

Packet 2: SYN flag bit = 1

ACK flag bit = _c)
Sequence number = 100

Packet 3: SYN flag bit = _d)

ACK flag bit = 1

Sequence number = _e)

(4.0分)

1010193

(4.0分)

51.

Host A wants to send a 12KB file F over a TCP connection. Several assumptions:
```

- <!--[if !supportLists]--> \bullet <!--[endif]-->The MSS is 1KB.
- <!--[if!supportLists]-->● <!--[endif]-->The receiver's advertised window is initially 4 KB.
- <!--[if !supportLists]--> <!--[endif]-->Unless indicated otherwise, all segments were received properly and received in the same order as they were sent
- <!--[if !supportLists]--> <!--[endif]-->The receiver will send ACK immediately, once receiving one data segment.
- <!--[if !supportLists]--> <!--[endif]-->It takes the sender 10 ms to "push" the segment onto the network. This means that if the first data segment is pushed onto the network starting at time 0, then the second segment can started to be pushed onto the network at 10 ms.
- <!--[if !supportLists]--> <!--[endif]-->Unless indicated otherwise, each successfully transmitted segment has a round trip time of exactly 60ms (30 ms each way). This time includes transmission time.
- $<!--[if ! support Lists]--> \\ \bullet <!--[endif]--> \\ The timer on host A of this TCP connection is always set as 100 ms.$
- <!--[if !supportLists]--> <!--[endif]-->The seq number and ack number of the first data segment are 0 and 100, respectively.
- <!--[if !supportLists]-->• <!--[endif]-->Under the set of assumptions above, Find the seq number, ack number and the time to send of each data segments if:
- <!--[if !supportLists]-->
 -->

 -
- <!--[if!supportLists]-->• <!--[endif]-->The first transmission of the 4th data segment is not received by the receiver.
- <!--[if !supportLists]-->• <!--[endif]-->In the ACKs for the 8th data segment and subsequently, the receiver's advertised window is reset to 2 KB

Please fill in the value of time, seq number, and ack number for each segments A will send.

Sending time	seq	ack
0	0	100

(5.0分)

10 1k 100 得分:**4.0** 20 2k 100 30 3k 100 40 4k 100 140 4k 100 200 5k 100 210 6k 100 270 7k 100 280 8k 100 290 9k 100 300 10k 100 360 11k 100 370 12k 100 收起 52.

<!--[if !supportLists]-->3.1 <!--[endif]-->Fill in the value of the congestion window size (number of segments) for each transmission round. Assume the threshold starts at 30 segments and the following events occur:

<!--[if!supportLists]-->• <!--[endif]-->triple duplicate ACK during round 10

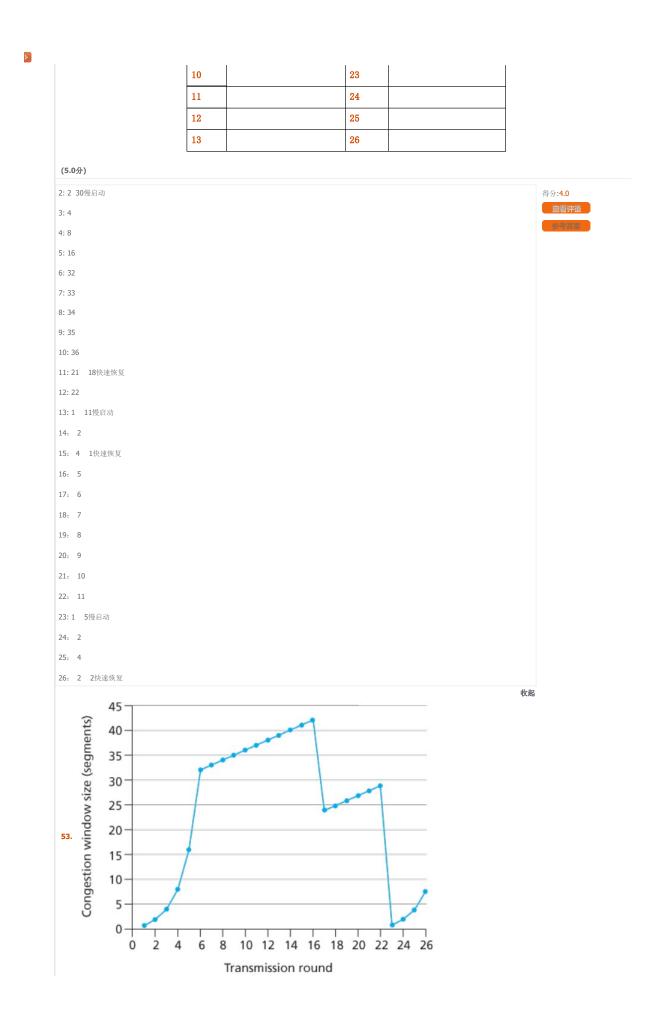
<!--[if !supportLists]--> <!--[endif]-->timeout during round 12

<!--[if!supportLists]-->
 <!--[endif]-->triple duplicate ACK during round 14

<!--[if!supportLists]-->
• <!--[endif]-->timeout during round 22

• timeout during round 25

Round	Congestion Window Size	Round	Congestion Window Size
1	1	14	
2		15	
3		16	
4		17	
5		18	
6		19	
7		20	
8		21	
9		22	



- <!--[if !supportLists]-->3.54 <!--[endif]-->You are sending a message with 10 segments, each having a sequence number from 1 to 10. The flow control algorithm used between the sender and receiver is Go-Back-N using cumulative ACKs where the sender window size is 3. Under the set of assumptions below, how long does it take before the sender deems the transmission complete if:
 - <!--[if!supportLists]-->• <!--[endif]-->The first transmission of segment 5 is not received by the receiver.
 - <!--[if !supportLists]-->• <!--[endif]-->The second transmission of segment 5 is received by the receiver but the ACK is slow, taking 60 ms (instead of 20ms as mentioned below). It still takes 20 ms for the segment to be sent from the sender to the receiver.
 - <!--[if!supportLists]-->• <!--[endif]--->The first ACK sent by the receiver after receiving segment 8 is not received.

Several assumptions:

Sequence numbers are assigned to segments rather than bytes. Unless indicated otherwise, all segments were receiv

Show your work by filling in the table on the following page. For each time interval, write the sequence number of any segment that is sent (i.e., starting to be pushed onto the network) in the SEND column. Similarly, write the sequence number included in any ACK received in the RECV column. To help you get started, segments 1 is done for you. Not all of the rows will be necessary. (10 points)

Time	SEND	RECV	Time	SEND	RECV
0	1		220		
10			230		
20			240		
30			250		
40		1	260		
50			270		
60			280		
70			290		
80			300		
90			310		

100			320		
110			330		
120			340		
130			350		
140			360		
150			370		
160 170			380		
180			400		
190			410		
200			420		
210			430		
(10.0分)					
) 1 null					得分:5.0
LO 2 null					宣看评语
20 3 null					参考答案
30 null null					
10 4 1					
50 5 2					
50 6 3					
70 null null					
30 7 4					
00 null null					
LOO null null					
110 null null					
.20 null null					
130 null null					
L40 5 null		第三次作业(周三班) (总分:100)			
150 6 null					
160 7 null					
170 null null					
180 null null					
190 8 6					
200 9 5,7					
210 10 null					
220 null null					
230 null null					
240 null 9					
250 null 10					
55 Refering to	problem 3.54, assuming S	Selective Reneated is	used instead of CRN place	ase fill the table (10 poi	收起 nts)
Time	SEND	RECV	Time	SEND	RECV
0	1		220		
10			230		
			240		
20		1	250		
20 30					
20 30 40		1	260		
20 30 40 50		1	260 270		
20 30 40		1	260		

90	310	
100	320	
110	330	
120	340	
130	350	
140	360	
150	370	
160	380	
170	390	
180	400	
190	410	
200	420	
210	430	

(10.0分)

