



Vidyavardhini's College of Engineering and Technology

Department of Artificial Intelligence & Data Science

Experiment No.2
Selection Sort
Date of Performance:
Date of Submission:



Experiment No. 2

Title: Selection Sort

Aim: To implement Selection Comparative analysis for large values of 'n'

Objective: To introduce the methods of designing and analyzing algorithms

Theory:

Selection sort is a sorting algorithm, specifically an in-place comparison sort. Selection sort is noted for its simplicity, and it has performance advantages over more complicated algorithms in certain situations, particularly where auxiliary memory is limited.

The algorithm divides the input list into two parts: the sub list of items already sorted, which is built up from left to right at the front (left) of the list, and the sub list of items remaining to be sorted that occupy the rest of the list. Initially, the sorted sub list is empty and the unsorted sub list is the entire input list. The algorithm proceeds by finding the smallest (or largest, depending on sorting order) element in the unsorted sub list, exchanging it with the leftmost unsorted element (putting it in sorted order), and moving the sublist boundaries one element to the right.

Example:

`arr[] = 64 25 12 22 11`

`// Find the minimum element in arr[0...4] // and place it at beginning`

`11 25 12 22 64`

`// Find the minimum element in arr[1...4] // and place it at beginning of arr[1...4]`

`11 12 25 22 64`

`// Find the minimum element in arr[2...4] // and place it at beginning of arr[2...4]`



11 12 22 25 64

// Find the minimum element in arr[3...4] // and place it at beginning of arr[3...4]

11 12 22 25 64

Algorithm and Complexity:

Alg.: SELECTION-SORT(A)		
	cost	Times
$n \leftarrow \text{length}[A]$	c_1	1
for $j \leftarrow 1$ to $n - 1$	c_2	$n-1$
do $\text{smallest} \leftarrow j$	c_3	$n-1$
for $i \leftarrow j + 1$ to n	c_4	$\sum_{j=1}^{n-1} (n-j+1)$
$\approx n^2/2$ comparisons, do if $A[i] < A[\text{smallest}]$	c_5	$\sum_{j=1}^{n-1} (n-j)$
then $\text{smallest} \leftarrow i$	c_6	$\sum_{j=1}^{n-1} (n-j)$
$\approx n$ exchanges, exchange $A[j] \leftrightarrow A[\text{smallest}]$	c_7	$n-1$

Implementation:

Code:

```
#include <stdio.h>
```

```
int main()
```

```
{
```

```
int array[100], n, c, d, position, t;
```

```
printf("Enter number of elements\n");
```

```
scanf("%d", &n);
```

```
printf("Enter %d integers\n", n);
```



```
for (c = 0; c < n; c++)
```

```
    scanf("%d", &array[c]);
```

```
for (c = 0; c < (n - 1); c++) // finding minimum element (n-1) times
```

```
{
```

```
    position = c;
```

```
    for (d = c + 1; d < n; d++)
```

```
    {
```

```
        if (array[position] > array[d])
```

```
            position = d;
```

```
    }
```

```
    if (position != c)
```

```
    {
```

```
        t = array[c];
```

```
        array[c] = array[position];
```

```
        array[position] = t;
```

```
    }
```

```
}
```

```
printf("Sorted list in ascending order:\n");
```

```
for (c = 0; c < n; c++)
```



```
printf("%d\n", array[c]);
```

```
return 0;
```

```
}
```

Output:

```
Enter number of elements
5
Enter 5 integers
12 5 1 78 9
Sorted Element In Selection Sort:
1
5
9
12
78
```

Conclusion:

In conclusion, implementing Selection Comparative analysis for large values of 'n' provides valuable insights into the efficiency and performance of selection algorithms under varying data sizes. Through this experiment, we can ascertain the scalability, resource utilization, and time complexity of these algorithms, aiding in informed decision-making for real-world applications.



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