Transactional Memory at BoostCon'10



BoostCon 2010

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TM Lineup







Vicente Botet Maurice Herlihy Justin Gottschlich Tatiana Shpeisman Michael Wong Mark Moir

Tatiana (Part II)







Library

o TBoost.STM Engine

Compiler

- Proposed C++ language extension
- o IBM
- o Sun / Oracle
- o Intel

Gottschlich



- Maurice Herlihy
- o Mark Moir
- Tatiana Shpeisman
- Michael Wong





Submit questions today



Why Should You Care?

- May affect you
 - C++ language extension
 - Library / Compiler
- Insight
 - Understand parallel algorithms
 - Understand how TM applies
- Contacts

The TBoost.STM Engine



BoostCon 2010

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Motivation

- Problem
 - o TM is not fast enough! (Cascaval et al., 2008)
- Reason
 - Conflict Detection and Opacity
 - Most TMs use Validation
- Our solution:
 - o Full Invalidation
 - o InvalSTM



TM Performance Bottleneck



Conflict Detection

- o Determine if transaction can commit
 - (Papadimitrou, "Theory of Database Concurrency Control," 1986)

Opacity

- Keep in-flight transactions consistent
 - (Guerraoui & Kapalka, PPoPP'08)

Conflict Detection

Conflict: $W_{T1} \cap (W_{T2} \cup R_{T2}) \neq \emptyset$



- Validation (T2)
 - Analyze the Past
 - Version # is same

- Invalidation (T1)
 - Analyze the Future
 - *T2.valid* = *false*

Opacity



- Validation
 - O Version # is same
- Invalidation
 - Oheck valid != false

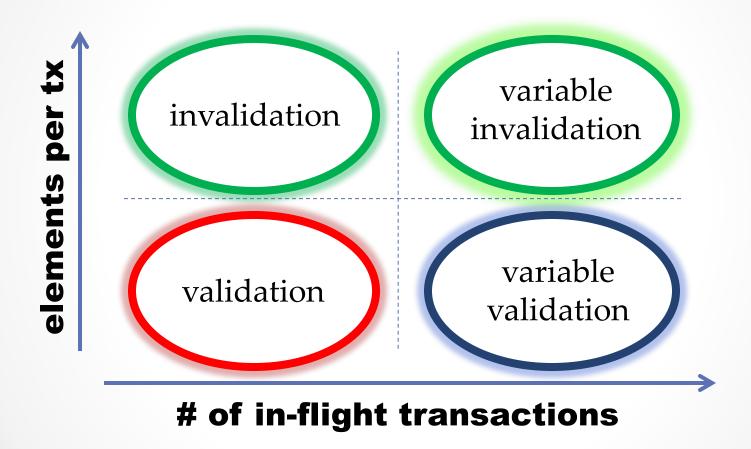
Opacity: Validation

```
atomic { return (x == 0 ? 0 : y / z); }
     read x \rightarrow store x's ver / ref
     read y \rightarrow store y's ver / ref,
               opacity: validate x
     read z \rightarrow store z's ver / ref,
               opacity: validate x, y
     commit \rightarrow conflict: validate x, y, z
```

Opacity: Invalidation

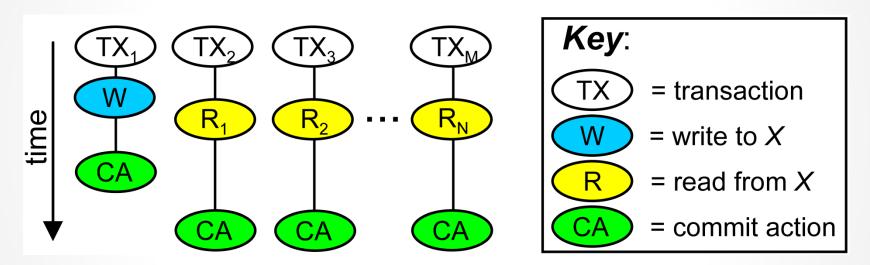
```
atomic { return (x == 0 ? 0 : y / z); }
     read x \rightarrow store ref to x
    read y \rightarrow store ref to y
               opacity: if (!valid) retry
     read z \rightarrow store ref to z
               opacity: if (!valid) retry
     commit → conflict: none
```

Validation Vs. Invalidation

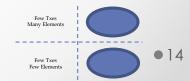


Contending + Concurrent Workload

1-Writer, N-Reader

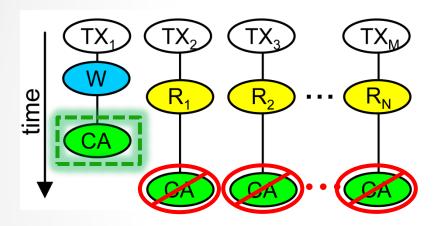


Commit to Executed Ratio: *Commits / Executed* Max = 1, Min = 0

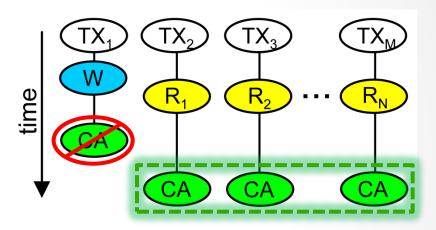


Side-By-Side Analysis

Validation



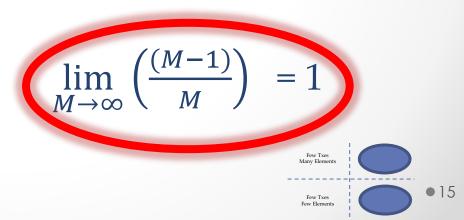
Invalidation



Commit / Executed: 1 / M

$$\lim_{M \to \infty} \left(\frac{1}{M} \right) = 0$$

Commit / Executed: (M-1) / M



Algorithmic Growth

$$Validation = \sum_{i=1}^{M} \sum_{j=1}^{r_i} j$$

Invalidation =
$$\sum_{i=1}^{M} \left(r_i + \sum_{j=1}^{F_i} w_i (s_{rj}(r_j) + s_{wj}(w_j)) \right)$$

Bloom Inval =
$$\sum_{i=1}^{M} (r_i + (2kw * Fi))$$



Efficient Read-Only Transactions

Validation Read-Only =
$$\sum_{i=1}^{n} \sum_{j=1}^{n} j$$

Invalidation =
$$\sum_{i=1}^{M} \left(r_i + \sum_{j=1}^{F_i} w_i(s_{j}(r_j) + s_{wj}(w_j)) \right)$$

Invalidation Read-Only =
$$\sum_{i=1}^{n} r_i$$

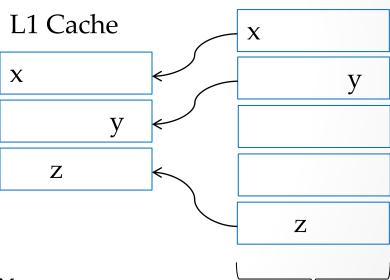


Validation + Memory

```
atomic { x = y / z; }
```

Addressable memory

Cache line size



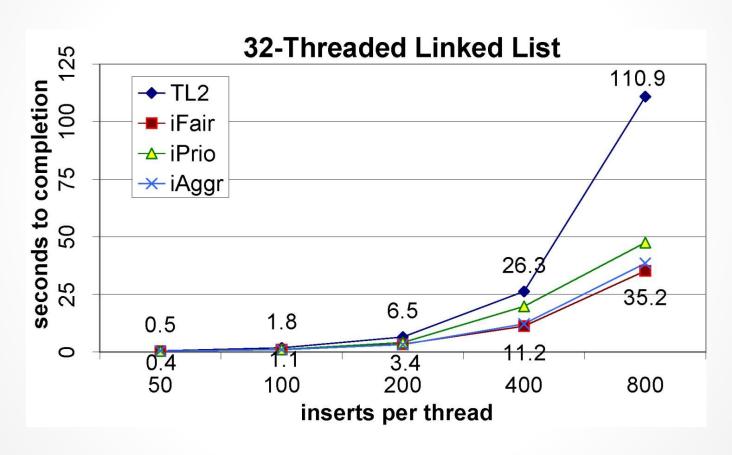
N = elements per tx $O(N^2)$ cache misses per tx

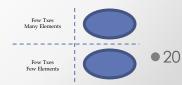
Invalidation + Memory

Addressable atomic $\{ x = y / z; \}$ memory L1 Cache Bloom Filter X X \mathbf{Z} X \mathbf{Z} y y Other Tx \mathbf{Z} M = # of in-flight txes N = # of read elements Cache line size O(N + M) cache miss per tx

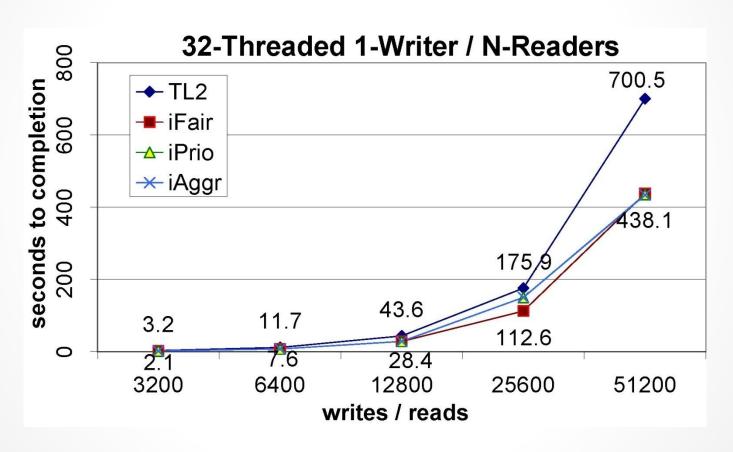
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Linked List



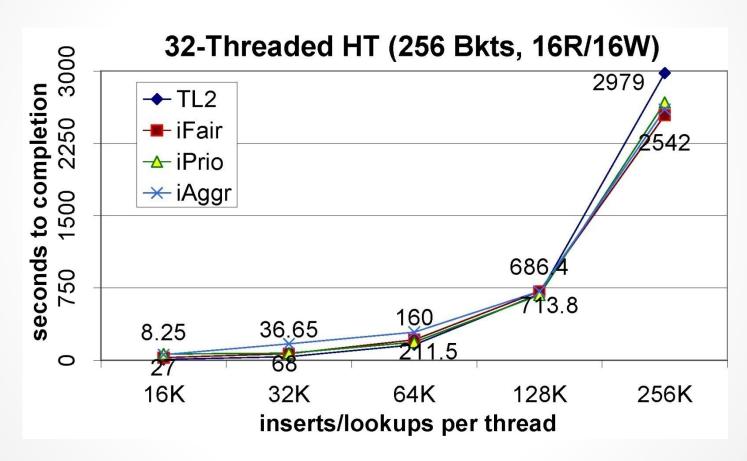


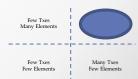
1-Writer / N-Readers





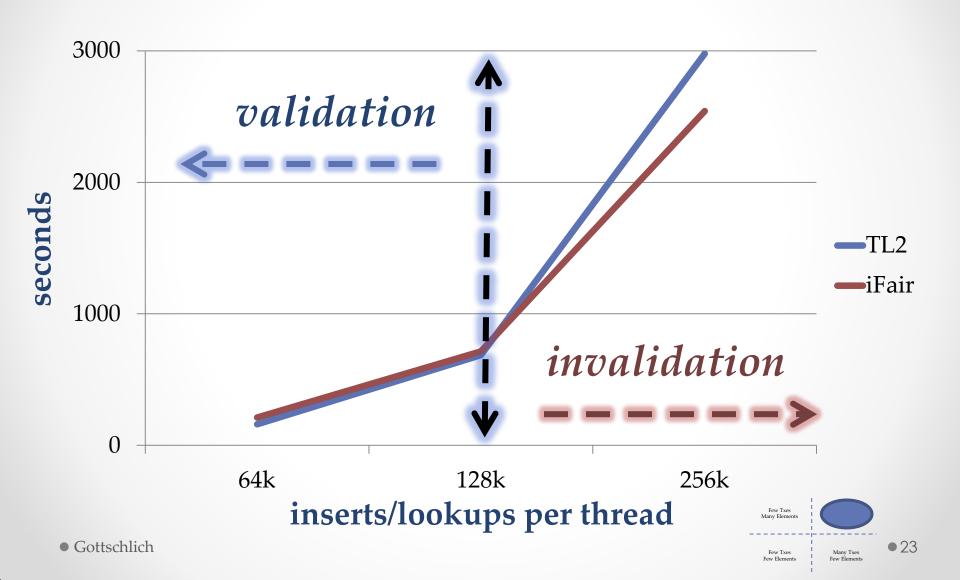
Hash Table





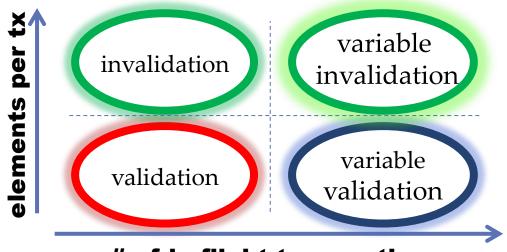
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Zoomed Hash Table



Conclusion

- InvalSTM (TBoost.STM) competitive with TL2
- Invalidation (InvalSTM) can be efficient



of in-flight transactions

- Next up
 - Proof of correctness for Full Invalidation
 - InvalSTM + STAMP
- Special thanks to Spear and Herlihy

Gottschlich

Questions?



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Gottschlich

What is the most important TM problem to solve in the next 1-2 years?

What is the most important TM problem to solve in the next 5-10 years?

Is TM on track to becoming a real parallel programming solution?

If so, when will it be realized?

If not, what are we doing wrong and what do we need to do to fix it?

In what types of environments (systems or programs) do you see TM being most useful?

Why?

Where is TM least useful?

Why?

What will parallel programs look like in the future?

What should a transaction do when an exception occurs within it?

Abort? Commit?