

(2.8.) ($\&$: address of)

$$x30 = \&A[1]$$

$$x31 = \&A$$

$$A[1] = \&A$$

$$x30 = A[1] = \&A$$

$$f = \&A + \&A = 2 * (\&A)$$

(2.9.)

	opcode	rs1	rd	type
① addi x30, x10, 8 :	0010011	01010	11110	I-type
② addi x31, x10, 0 :	0010011	01010	11111	I-type
③ sd x31, 0(x30) :	0100011	11110	Stype rd	S-type
④ ld x30, 0(x30) :	0000011	11110	11110	I-type
⑤ add x5, x30, x31 :	0110011	11110	00101	R-type

Immediate	rs2	funct3	funct7
① 0000000001000		000	
② 00000000000000		000	
③	11111	011	(imm 11=5) 0000111
④ 00000000000000		011	
⑤	11111	000	00000000

(2.16)

(1) R-type

register field $5 \rightarrow 7$, opcode field $7 \rightarrow 9$

funct	rs2	rs1	funct	rd	opcode
1	7	7	1	7	9

(2) I-type

immediate	rs1	funct	rd	opcode
8	7	1	7	9

(3)

因為指令集變大，可以引入更多有效率的指令，來取代舊的指令，比如 swap 指令可取代原本需要 3.4 行的指令，因此 program size 會變小。

但是像是 I-type 指令，原本在 12 bits 以下可表示的 immediate 都可以一個指令結束，現在超過 8 bits 即需要更多指令來完成任務，因此 program size 也可能變大。

Report on Matrix Multiplication

1. it takes $128*128*128*8 = 16777216$ cycles.
2. it needs roughly $128*128*128*2 = 4194304$ loads and $128*128 = 16384$ stores.
3. Using blocking technique can ensure that in a single basic block, the registers can be used as much as possible.
4. it needs roughly $128*128*128 = 2097152$ loop controls.