

# Metric Entropy and its Applications

Presented by Wanteng Ma

Department of Mathematics

Hong Kong University of Science and Technology

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# Metric Entropy

**Metric entropy** serves as a tool to measure the “size” of a set. It is also a measure of **statistical complexity**.

- First introduced by Kolmogorov.
- Broadly applied to theoretical statistics
  - nonparametric function estimation<sup>1</sup>
  - high-dimensional statistical inference<sup>2</sup>
  - statistical learning theory<sup>3</sup>

Generally, it reveals interesting connections between the complexity of the parameter space and the fundamental difficulty of the statistical problem.

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<sup>1</sup>Yihong Wu and Pengkun Yang. “Minimax rates of entropy estimation on large alphabets via best polynomial approximation”. In: *IEEE Transactions on Information Theory* 62.6 (2016), pp. 3702–3720.

<sup>2</sup>T Tony Cai, Zongming Ma, and Yihong Wu. “Sparse PCA: Optimal rates and adaptive estimation”. In: *The Annals of Statistics* 41.6 (2013), pp. 3074–3110.

<sup>3</sup>Vladimir Koltchinskii. “Local Rademacher complexities and oracle inequalities in risk minimization”. In: *The Annals of Statistics* 34.6 (2006), pp. 2593–2656.

# Definition

Consider a metric space  $(\mathbb{T}, \rho)$  with metric  $\rho(\theta, \theta')$  satisfying

- ① Non-negativity
- ② Symmetry
- ③ Triangular inequality

e.g.

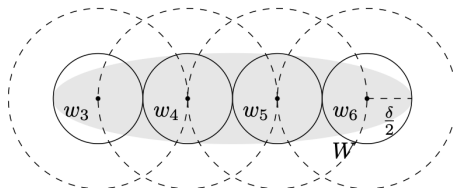
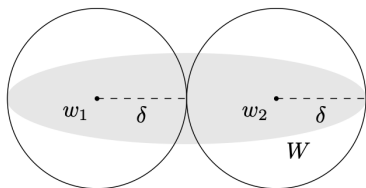
- (Euclidean metric)  $\rho(\theta, \theta') = \|\theta - \theta'\|_2$  on  $\mathbb{T} = \mathbb{R}^d$
- (Rescaled Hamming metric)  $\rho(\theta, \theta') = \frac{1}{d} \sum_{i=1}^d \mathbb{1}(\theta_i \neq \theta'_i)$  on  $\mathbb{T} = \{0, 1\}^d$
- Function space  $\mathcal{L}^2(\mu, [0, 1])$  with metric

$$\|f - g\|_2 = \left( \int_0^1 (f(x) - g(x))^2 d\mu(x) \right)^{1/2}$$

# Definition

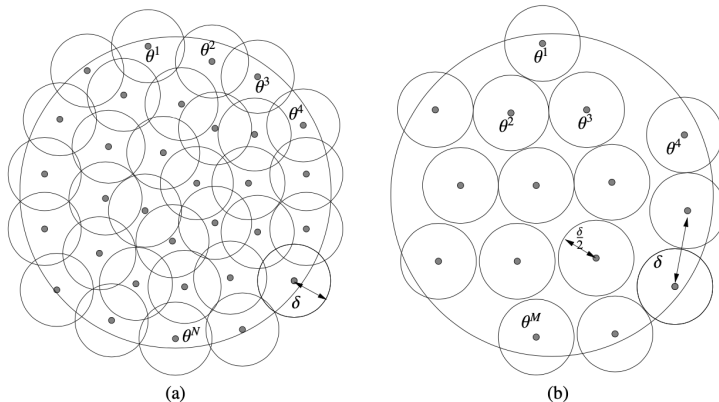
## Definition (Covering number)

A  $\delta$ -cover of a set  $\mathbb{T}$  with respect to a metric  $\rho$  is a set  $\{\theta_1, \dots, \theta_N\} \subseteq \mathbb{T}$  such that for each  $\theta \in \mathbb{T}$ , there exists some  $i \in \{1, \dots, N\}$  such that  $\rho(\theta, \theta_i) \leq \delta$ . The  $\delta$ -covering number  $N(\delta; \mathbb{T}, \rho)$  is the cardinality of the **smallest**  $\delta$ -cover.



## Definition (packing number)

A  $\delta$ -packing of a set  $\mathbb{T}$  with respect to a metric  $\rho$  is a set  $\{\theta_1, \dots, \theta_M\} \subseteq \mathbb{T}$  such that for all distinct  $i, j \in \{1, \dots, M\}$ ,  $\rho(\theta_i, \theta_j) > \delta$ . The  $\delta$ -packing number  $M(\delta; \mathbb{T}, \rho)$  is the cardinality of the **largest**  $\delta$ -packing.



From the definition of covering and packing set, we have

- 1 If  $M$  is the cardinality of a  $\delta$ -packing set, then we need at least  $M$  balls with radius  $\frac{\delta}{2}$  to cover the space.
- 2 A  $\delta$ -packing set with the largest cardinality  $M(\delta; \mathbb{T}, \rho)$  is also a  $\delta$ -cover.

For all  $\delta > 0$ , the packing and covering numbers are related as follows:

$$M(2\delta; \mathbb{T}, \rho) \leq N(\delta; \mathbb{T}, \rho) \leq M(\delta; \mathbb{T}, \rho),$$

which indicates

- ④ The packing and covering numbers exhibit the same scaling behavior as  $\delta \rightarrow 0$
- ② We can find upper/lower bound of covering/packing numbers by constructing cover set or packing set.
- ③ We can investigate the growth rate of covering/packing numbers by studying *metric entropy*  $\log(N(\delta; \mathbb{T}, \rho))$ .

## Example 1 (Covering and packing of unit cubes)

Consider the interval  $[-1, 1]$  in  $\mathbb{R}$ , with the metric  $\rho(\theta, \theta') = |\theta - \theta'|$ .

- Divide the interval into  $L = \lfloor \frac{1}{\delta} \rfloor + 1$  sub-intervals, with each one centered at  $\theta = -1 + (2i - 1)\delta$
- By construction, for any point  $\theta \in [-1, 1]$ , there is some  $j \in [L]$  such that  $|\theta_j - \theta| \leq \delta$ , i.e.,  $N(\delta, [-1, 1], |\cdot|) \leq \frac{1}{\delta} + 1$ . Generalize this analysis to higher dimensions we can show  $N(\delta, [-1, 1]^d, \|\cdot\|_\infty) \leq (\frac{1}{\delta} + 1)^d$ .
- Consider  $\{\theta_i : i \in [L - 1]\}$ . Any two elements  $\theta_i, \theta_j$  in this set have the separation  $|\theta_i - \theta_j| \geq 2\delta$ , which implies  $M(2\delta, [-1, 1], |\cdot|) \geq \lfloor \frac{1}{\delta} \rfloor$ . This also means  $M(2\delta, [-1, 1]^d, \|\cdot\|_\infty) \geq \lfloor \frac{1}{\delta} \rfloor^d$
- Combined with the relation of packing number and covering number, we have

$$\log N(\delta; [0, 1]^d, \|\cdot\|_\infty) \asymp d \log(1/\delta)$$

## Lemma (Volume ratios and metric entropy)

Consider a pair of norms  $\|\cdot\|$  and  $\|\cdot\|'$  on  $\mathbb{R}^d$ , and let  $\mathbb{B}$  and  $\mathbb{B}'$  be their corresponding unit balls (i.e.,  $\mathbb{B} = \{\theta \mid \|\theta\| \leq 1\}$ , with  $\mathbb{B}'$  similarly defined). Then we have

$$\left(\frac{1}{\delta}\right)^d \frac{\text{vol}(\mathbb{B})}{\text{vol}(\mathbb{B}')} \stackrel{(a)}{\leq} N(\delta; \mathbb{B}, \|\cdot\|') \stackrel{(b)}{\leq} \frac{\text{vol}(\frac{2}{\delta}\mathbb{B} + \mathbb{B}')}{\text{vol}(\mathbb{B}')}.$$

### Proof:

- For (a), if  $\{\theta_1, \dots, \theta_N\}$  a  $\delta$ -covering of  $\mathbb{B}$ , then  $\mathbb{B} \subseteq \bigcup_{j=1}^N \{\theta_j + \delta\mathbb{B}'\}$ , which implies that  $\text{vol}(\mathbb{B}) \leq N \text{vol}(\delta\mathbb{B}') = N\delta^d \text{vol}(\mathbb{B}')$
- For (b), let  $\{\theta_1, \dots, \theta_M\}$  be a maximal  $\delta$ -packing of  $\mathbb{B}$  in the  $\|\cdot\|'$ -norm; by maximality, this set must also be a  $\delta$ -covering. The balls  $\{\theta_j + \frac{\delta}{2}\mathbb{B}' \mid j \in [M]\}$  are all disjoint and contained within  $\mathbb{B} + \frac{\delta}{2}\mathbb{B}'$ .



- This indicates that  $M \operatorname{vol}(\frac{\delta}{2}\mathbb{B}') \leq \operatorname{vol}(\mathbb{B} + \frac{\delta}{2}\mathbb{B}')$ . Taking  $\frac{\delta}{2}$  out we have  $\operatorname{vol}(\frac{\delta}{2}\mathbb{B}') = (\frac{\delta}{2})^d \operatorname{vol}(\mathbb{B}')$ , which justifies (b).
- A special case: when  $\mathbb{B}' \subseteq \mathbb{B}$ , we have:

$$\operatorname{vol}\left(\frac{2}{\delta}\mathbb{B} + \mathbb{B}'\right) \leq \operatorname{vol}\left(\left(\frac{2}{\delta} + 1\right)\mathbb{B}\right) = \left(\frac{2}{\delta} + 1\right)^d \operatorname{vol}(\mathbb{B})$$

- If further we take  $\mathbb{B}' = \mathbb{B}$ , then we have the following bounds for metric entropy of a unit ball.

$$d \log(1/\delta) \leq \log N(\delta; \mathbb{B}, \|\cdot\|) \leq d \log\left(1 + \frac{2}{\delta}\right)$$

## Example 2 (Lipschitz functions on the unit interval)

Now consider the class of Lipschitz functions  $\mathcal{F}_L :=$

$$\{g : [0, 1] \rightarrow \mathbb{R} \mid g(0) = 0, \text{ and } |g(x) - g(x')| \leq L|x - x'|, \forall x, x' \in [0, 1]\}.$$

### Lipschitz functions on the unit interval

The metric entropy of  $\mathcal{F}_L$  with respect to the sup-norm scales as

$$\log N_\infty(\delta, \mathcal{F}_L) \asymp (L/\delta)$$

### Proof:

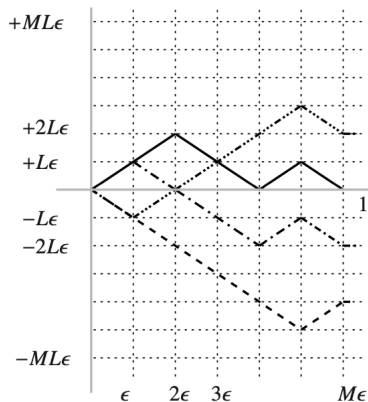
- Prove the lower bound by constructing a sufficiently large packing of the set  $\mathcal{F}_L$ . For any  $\epsilon$ , define  $M = \lfloor \frac{1}{\epsilon} \rfloor$ , consider the points in  $[0, 1]$  given by

$$x_i = (i - 1)\epsilon, \text{ for } i \in [M], \text{ and } x_{M+1} = M\epsilon.$$

- We construct a set of segmented functions that each function is piecewise linear over the intervals  $[x_{i-1}, x_i]$  with slope either  $L$  or  $-L$ .

## Example 2 (Lipschitz functions on the unit interval)

- Define the unit function:  $\phi(u) := \begin{cases} 0 & \text{for } u < 0 \\ u & \text{for } u \in [0, 1] \\ 1 & \text{otherwise} \end{cases}$
- For each binary sequence  $\beta \in \{-1, 1\}^M$ , define  $f_\beta(y) = \sum_{i=1}^M \beta_i L \epsilon \phi\left(\frac{y-x_i}{\epsilon}\right)$ .



- We can check that  $\{f_\beta \mid \beta \in \{-1, 1\}^M\} \subseteq \mathcal{F}_L$

## Example 2 (Lipschitz functions on the unit interval)

- From the construction of  $f_\beta$  we know that for any distinct  $i, j \in \{-1, 1\}^M$ ,  $\|f_i - f_j\|_\infty \geq 2L\epsilon$ , showing that  $\{f_\beta | \beta \in \{-1, 1\}^M\}$  is a  $2L\epsilon$ -packing with cardinality  $2^M$ . Thus

$$\log N(\delta; \mathcal{F}_L, \|\cdot\|_\infty) \gtrsim L/\delta$$

- For the upper bound, we show that  $\{f_\beta \mid \beta \in \{-1, 1\}^M\}$  is also a  $\delta$ -covering of set  $\mathcal{F}_L$ , which can be justified by induction.
- An extension of this argument shows that

$$\log N_\infty(\delta, \mathcal{F}_L([0, 1]^d)) \asymp (L/\delta)^d$$

# Gaussian and Rademacher complexity

Metric entropy plays a fundamental role in [understanding the behavior of stochastic processes](#). Define the Gaussian and Rademacher complexity as following:

## Definition (Gaussian complexity)

Given a set  $\mathbb{T} \in \mathbb{R}^d$ , the random variable  $w$  follows  $d$ -dimensional standard Gaussian distribution, then the Gaussian complexity of  $\mathbb{T}$  is defined as

$$\mathcal{G}(\mathbb{T}) := \mathbb{E} \left[ \sup_{\theta \in \mathbb{T}} \langle \theta, w \rangle \right].$$

## Definition (Rademacher complexity)

Given a set  $\mathbb{T} \in \mathbb{R}^d$ ,  $\varepsilon$  is a  $d$ -dimensional Rademacher variable, then the Rademacher complexity of  $\mathbb{T}$  is defined as

$$\mathcal{R}(\mathbb{T}) := \mathbb{E} \left[ \sup_{\theta \in \mathbb{T}} \langle \theta, \varepsilon \rangle \right].$$

# Gaussian and Rademacher complexity

Trivial relation between Rademacher/Gaussian complexity:

$$\mathcal{R}(\mathbb{T}) \leq \sqrt{\frac{\pi}{2}} \mathcal{G}(\mathbb{T})$$

**Proof:**

- Consider a fixed  $\varepsilon \in \{-1, 1\}^d$ . For  $\forall \theta \in \mathbb{T}$ , integrate  $\langle \theta, w \rangle$  over the corresponding  $\frac{1}{2^d}$ -half space:

$$\int_{R(\varepsilon)} \langle \theta, w \rangle f(w) dw = \frac{1}{2^{d-1} \sqrt{2\pi}} \langle \theta, \varepsilon \rangle$$

- Taking supremum and Summing all the  $\varepsilon$

$$\begin{aligned} \mathbb{E} \left[ \sup_{\theta \in \mathbb{T}} \langle \theta, \varepsilon \rangle \right] &= \frac{1}{2^d} \sum_{\varepsilon} \sup_{\theta \in \mathbb{T}} \langle \theta, \varepsilon \rangle \leq \sqrt{\frac{\pi}{2}} \sum_{\varepsilon} \sup_{\theta \in \mathbb{T}} \int_{R(\varepsilon)} \langle \theta, w \rangle f(w) dw \\ &\leq \sqrt{\frac{\pi}{2}} \int \sup_{\theta \in \mathbb{T}} \langle \theta, w \rangle f(w) dw = \sqrt{\frac{\pi}{2}} \mathcal{G}(\mathbb{T}). \end{aligned}$$

# Gaussian and Rademacher complexity

Some prerequisite knowledge on Gaussian variables sequences: Suppose  $X \in \mathbb{R}^n$  and  $\{X_i\}_{i=1}^n$  is an i.i.d. sequence of  $N(0, \sigma^2)$  variables, then we have

- (Concentration bounds on  $\chi^2$ ) The growth of  $\|X\|_2$  follows

$$\frac{\mathbb{E}\|X\|_2}{\sigma\sqrt{n}} = 1 - o(1).$$

with upper bound  $\mathbb{E}\|X\|_2 \leq \sigma\sqrt{n}$ .

- (Gaussian maxima) The Gaussian maxima  $Z_n := \max_{i=1, \dots, n} |X_i| = \|X\|_\infty$  follows

$$\frac{\mathbb{E}[Z_n]}{\sqrt{2\sigma^2 \log n}} \rightarrow 1 \text{ as } n \rightarrow +\infty.$$

with upper bound  $\mathbb{E}Z_n \leq 2\sigma\sqrt{\log n}$ .

## Example 3 (Rademacher/Gaussian complexity of unit balls)

Compute the Rademacher and Gaussian complexities of the Euclidean ball of unit norm  $\mathbb{B}_2^d$ .

- Rademacher complexity:

$$\mathcal{R}(\mathbb{B}_2^d) = \mathbb{E} \left[ \sup_{\|\theta\|_2 \leq 1} \langle \theta, \varepsilon \rangle \right] = \mathbb{E} \left[ \left( \sum_{i=1}^d \varepsilon_i^2 \right)^{1/2} \right] = \sqrt{d}$$

- Gaussian complexity: we have  $\mathcal{G}(\mathbb{B}_2^d) = \mathbb{E}[\|w\|_2]$ , and preliminary knowledge of Gaussian distribution shows

$$\frac{\mathcal{G}(\mathbb{B}_2^d)}{\sqrt{d}} = 1 - o(1).$$



## Example 3 (Rademacher/Gaussian complexity of unit balls)

Compute the Rademacher and Gaussian complexities of the ball of unit  $l_1$  norm  $\mathbb{B}_1^d$ .

- Rademacher complexity:

$$\mathcal{R}(\mathbb{B}_1^d) = \mathbb{E} \left[ \sup_{\|\theta\|_1 \leq 1} \langle \theta, \varepsilon \rangle \right] = \mathbb{E} \|\varepsilon\|_\infty = 1.$$

- Gaussian complexity: we have  $\mathcal{G}(\mathbb{B}_1^d) = \mathbb{E} [\|w\|_\infty]$ , and preliminary knowledge of Gaussian distribution shows

$$\frac{\mathcal{G}(\mathbb{B}_1^d)}{\sqrt{2 \log d}} = 1 \pm o(1).$$

# Metric entropy and sub-Gaussian processes

## Definition

A collection of zero-mean random variables  $\{X_\theta | \theta \in \mathbb{T}\}$  is a sub-Gaussian process with respect to a metric  $\rho_X$  if

$$\mathbb{E} \left[ e^{\lambda(X_\theta - X_{\tilde{\theta}})} \right] \leq e^{\frac{\lambda^2 \rho_X^2(\theta, \tilde{\theta})}{2}}, \text{ for all } \theta, \tilde{\theta} \in \mathbb{T}, \text{ and } \lambda \in \mathbb{R}$$

For the sub-Gaussian process, we have the famous Massart's finite class lemma:

## Lemma (Massart's)

Let  $A$  be some finite set with  $\text{Diam}(A) = \sup_{\theta, \theta' \in \mathbb{T}} \rho_X(\theta, \theta') \leq R$ , then

$$\sup_{\theta \in \mathbb{T}} X_\theta \leq R \sqrt{2 \log |A|}.$$

# Massart's finite class lemma

## Proof:

- Denote  $\mu = \mathbb{E} \sup_{\theta \in \mathbb{T}} (X_\theta - X_{\theta'})$ , we have

$$e^{\lambda \mu} \leq \mathbb{E} \sup_{\theta \in \mathbb{T}} e^{\lambda (X_\theta - X_{\theta'})} \leq \mathbb{E} \sum_{\theta \in A} e^{\frac{\lambda^2}{2} \rho_X^2(\theta, \theta')} \leq |A| e^{\frac{\lambda^2}{2} R^2}, \text{ for } \forall \lambda$$

- Taking logarithm we have

$$\mu \leq \frac{1}{\lambda} \log |A| + \frac{\lambda}{2} R^2$$

- Taking minimum value of RHS, we have

$$\sup_{\theta \in \mathbb{T}} X_\theta \leq R \sqrt{2 \log |A|}.$$

From this argument we can also derive the Gaussian maxima, together with bound of  $\sup_{\theta \in \mathbb{T}} |X_\theta|$  or  $\sup_{\theta, \theta' \in \mathbb{T}} (X_\theta - X_{\theta'})$ .

# Upper bound by one-step discretization

## One-step discretization

Consider the zero-mean sub-Gaussian process  $X_\theta$  with  $\text{Diam}(\mathbb{T}) \leq D$ . For any appropriate  $\delta$ , we have

$$\mathbb{E} \left[ \sup_{\theta, \tilde{\theta} \in \mathbb{T}} (X_\theta - X_{\tilde{\theta}}) \right] \leq 2 \mathbb{E} \left[ \sup_{\substack{\gamma, \gamma' \in \mathbb{T} \\ \rho_X(\gamma, \gamma') \leq \delta}} (X_\gamma - X_{\gamma'}) \right] + 4 \sqrt{D^2 \log N_X(\delta; \mathbb{T})}$$

**Proof:**

- For a given  $\delta$  and associated  $\delta$ -covering  $\{\theta_1, \dots, \theta_N\}$  with  $N = N_X(\theta, \mathbb{T})$ .  
For any  $\theta$  we can find some  $\theta_j$  such that  $\rho_X(\theta, \theta_j) \leq \delta$ . Then

$$\begin{aligned} X_\theta - X_{\theta_1} &= X_\theta - X_{\theta_j} + X_{\theta_j} - X_{\theta_1} \\ &\leq \sup_{\substack{\gamma, \gamma' \in \mathbb{T} \\ \rho_X(\gamma, \gamma') \leq \delta}} (X_\gamma - X_{\gamma'}) + \max_{i=1,2,\dots,N} |X_{\theta_i} - X_{\theta_1}| \end{aligned}$$

- the same upper bound also holds for  $X_{\theta_1} - X_{\tilde{\theta}}$ .

# One-step discretization

- By Massart's lemma we have

$$\mathbb{E} \left[ \max_{i=1, \dots, N} |X_{\theta_i} - X_{\theta_1}| \right] \leq 2\sqrt{D^2 \log N},$$

which yields the claim.

## Rmk

Since the zero-mean condition means that

$$\mathbb{E} \left[ \sup_{\theta \in \mathbb{T}} X_{\theta} \right] = \mathbb{E} \left[ \sup_{\theta \in \mathbb{T}} (X_{\theta} - X_{\theta_0}) \right] \leq \mathbb{E} \left[ \sup_{\theta, \tilde{\theta} \in \mathbb{T}} (X_{\theta} - X_{\tilde{\theta}}) \right]$$

We can also use this one-step discretization to bound  $\mathbb{E} [\sup_{\theta \in \mathbb{T}} X_{\theta}]$ , e.g., Rademacher/Gaussian complexity. In this case, the factor 2 is not needed because we can fix the  $X_{\tilde{\theta}}$ .

# One-step discretization

In the  $l_2$  norm, if we define  $\tilde{\mathbb{T}}(\delta) := \{\gamma - \gamma' \mid \gamma, \gamma' \in \mathbb{T}, \|\gamma - \gamma'\|_2 \leq \delta\}$ , then the Gaussian complexity follows

$$\mathcal{G}(\mathbb{T}) \leq \min_{\delta \in [0, D]} \left\{ \mathcal{G}(\tilde{\mathbb{T}}(\delta)) + 2\sqrt{D^2 \log N_2(\delta; \mathbb{T})} \right\}.$$

Here  $\mathcal{G}(\tilde{\mathbb{T}}(\delta))$  is referred to as a **localized Gaussian complexity**, and probably can be controlled by  $\delta$ .

If  $\mathbb{T}$  is a subset of  $\mathbb{R}^d$ , we have

$$\mathcal{G}(\tilde{\mathbb{T}}(\delta)) = \mathbb{E} \left[ \sup_{\theta \in \tilde{\mathbb{T}}(\delta)} \langle \theta, w \rangle \right] \leq \delta \mathbb{E} [\|w\|_2] \leq \delta \sqrt{d}$$

Thus we have naive discretization bound

$$\mathcal{G}(\mathbb{T}) \leq \min_{\delta \in [0, D]} \left\{ \delta \sqrt{d} + 2\sqrt{D^2 \log N_2(\delta; \mathbb{T})} \right\}.$$

## Example 3 ( $l_2$ norm of sub-Gaussian random matrix)

Let  $W \in \mathbb{R}^{x \times d}$  be a random matrix with zero-mean i.i.d. entries  $W_{ij}$ , each sub-Gaussian with parameter  $\sigma = 1$ . Then we have

$$\mathbb{E} [\|W\|_2 / \sqrt{n}] \lesssim 1 + \sqrt{\frac{d}{n}}$$

**Proof:**

- Define matrix class  $\mathbb{M}^{n,d}(1) := \{\Theta \in \mathbb{R}^{n \times d} \mid \text{rank}(\Theta) = 1, \|\Theta\|_F = 1\}$ .

First prove the variational representation:

$$\|W\|_2 = \sup_{\Theta \in \mathbb{M}(1)} X_\Theta, \text{ where } X_\Theta = \langle W, \Theta \rangle = \sum_{i,j} W_{ij} \Theta_{ij}$$

- Apply the discretization bound

$$\mathbb{E} [\|W\|_2] \leq 2\mathbb{E} \left[ \sup_{\substack{\text{mank}(\Gamma)=\text{rank}(\Gamma')=1 \\ \|\Gamma-\Gamma'\|_F \leq \delta}} \langle \Gamma - \Gamma', W \rangle \right] + 6\sqrt{\log N_F(\delta; M^{n,d}(1))}$$

# One-step discretization

- Use SVD of rank 2 matrix to prove the local complexity bound

$$\mathbb{E} \left[ \sup_{\substack{\text{rank}(\mathbf{\Gamma})=\text{rank}(\mathbf{\Gamma}')=1 \\ \|\mathbf{\Gamma}-\mathbf{\Gamma}'\|_F \leq \delta}} \langle \mathbf{\Gamma} - \mathbf{\Gamma}', \mathbf{W} \rangle \right] \leq \sqrt{2} \delta \mathbb{E} [\|\mathbf{W}\|_2].$$

- Use upper bound of metric entropy on unit ball to prove

$$\log N_F \left( \delta; \mathbb{M}^{n,d}(1) \right) \leq (n + d) \log \left( 1 + \frac{2}{\delta} \right)$$

Recall the metric entropy on Lipschitz function class  $\mathcal{F}_L$ . Given sample  $x_1^n$ , we have  $\log N_\infty(\delta, \mathcal{F}_L) \leq \frac{cL}{\delta}$ , thus

$$\mathcal{G}(\mathcal{F}_L(x_1^n)/n) \leq \frac{1}{\sqrt{n}} \inf_{\delta \in (0, \delta_0)} \left\{ \delta \sqrt{n} + 3 \sqrt{\frac{cL}{\delta}} \right\}$$

Choose  $\delta = \frac{1}{n^3}$ , we have

$$\mathcal{G}(\mathcal{F}_L(x_1^n)/n) \lesssim n^{-1/3}$$



# Chaining and Dudley's entropy integral

Define the Dudley's entropy integral  $\mathcal{J}(\delta; D) := \int_\delta^D \sqrt{\log N_X(u; \mathbb{T})} du$ . We have

Dudley's

$$\mathbb{E} \left[ \sup_{\theta, \tilde{\theta} \in \mathbb{T}} (X_\theta - X_{\tilde{\theta}}) \right] \leq 2 \mathbb{E} \left[ \sup_{\substack{\gamma, \gamma' \in \mathbb{T} \\ \rho_X(\gamma, \gamma') \leq \delta}} (X_\gamma - X_{\gamma'}) \right] + 32 \mathcal{J}(\delta/4; D)$$

- we pursue a more refined chaining argument. Denote  $\mathbb{U} = \{\theta^1, \dots, \theta^N\}$

