





COMPILER CONSTRUCTION

Lex





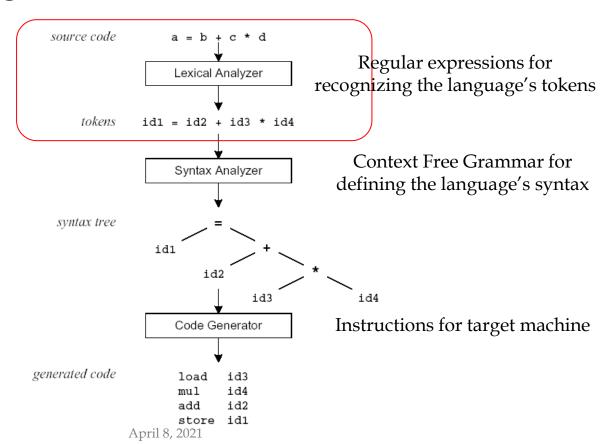






Where are we?

- Lex and Yacc are able to do the following
- Now, our target is Lex













Lex and Yacc

- Two compiler generation tools:
 - Lexical Tokens and their Order of Processing (Lex)
 - Context Free Grammar for LALR(1) (Yacc)
- Both Lex and Yacc have Long History in Computing
 - Lex and Yacc Earliest Days of Unix Minicomputers
 - Flex and Bison From GNU
 - JFlex Fast Scanner Generator for Java
 - BYacc/J Berkeley
 - PCLEX and PCYACC from Abacus
 - ANTLR, CUP, PCYACC



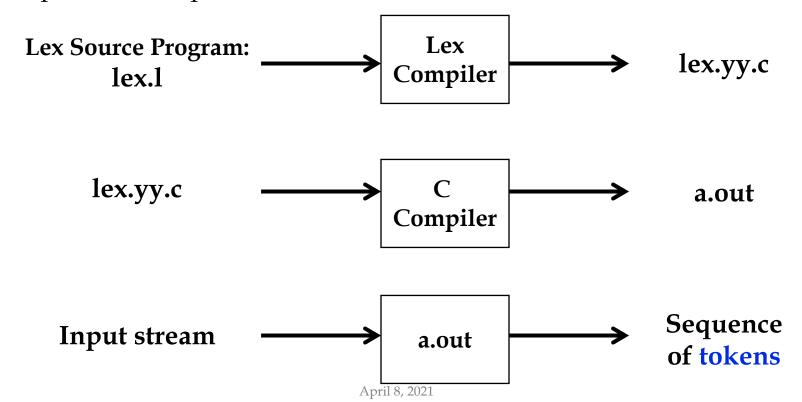






Lex - A Lexical Analyzer Generator

- A Unix Utility from early 1970s
- A Compiler that takes as source a specification for:
 - Tokens/Patterns of a Language
 - Generates a "C" Lexical Analyzer Program
- Inputs and outputs of Lex



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Lex



• The main job of a *lexical analyzer* (*scanner*) is to break up an input stream into more usable elements (*tokens*)

 $a = b + c * d; \leftarrow$ Input stream

ID ASSIGN ID PLUS ID MULT ID SEMI ← Tokens

• Lex is an utility to help you rapidly generate your scanners











Lex (Cont'd)

- Lexical analyzers tokenize input streams
- Tokens are the terminals of a language
 - English
 - words, punctuation marks, ...
 - Programming language
 - Identifiers, operators, keywords, ...
- Regular expressions define terminals/tokens













Lex Source to C Program

- The input (.l) is translated to a C program (lex.yy.c) which
 - reads an input stream,
 - partitions the input into strings which match the given expressions, and
 - copies it to an output stream if necessary









Lex vs. Yacc

- Lex
 - Generates C code for a lexical analyzer, or scanner
 - Uses patterns that match strings in the input and converts the strings to tokens

- Yacc
 - Generates C code for syntax analyzer, or parser
 - Uses grammar rules that allow it to analyze tokens from Lex and create a syntax tree

Together they form the basis of a compiler

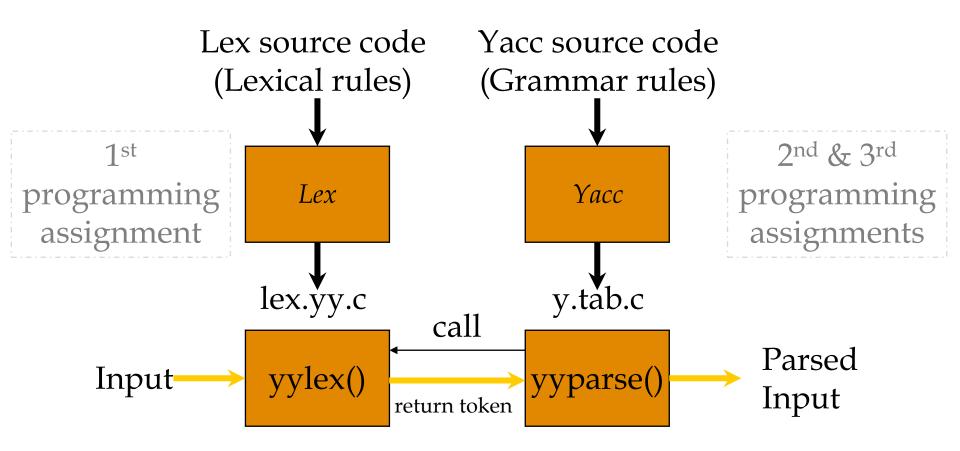








Lex & Yacc











Format of a Lexical Specification

• lex.l is divided into 3 parts:

1. Declarations:

- Defs, Constants, Types, #includes, etc. that can occur in a C Program
- Regular Definitions (expressions)

2. Translation rules:

- Pairs of (Regular Expression, Action)
- Informs lexical analyzer of action when pattern is recognized

3. Auxiliary procedures:

- Designer Defined C Code
- E.g., symbol table codes

lex.I file format:

DECLARATIONS

%%

TRANSLATION RULES

%%

AUXILIARY PROCEDURES

lex.l Example







%{ #define T IDENTIFIER 300 #define T INTEGER 301 #define T REAL 302 #define T STRING 303 #define T ASSIGN 304 #define T ELSE 305 #define T IF 306 #define T THEN 307 #define T EQ 308 #define T LT 309 #define T NE 310 #define T GE 311 #define T GT 312

%}

User defined values to each token (else lex will assign)

Regular expression rules for later token definitions

```
letter
                           [a-zA-Z]
diait
                            [0-9]
                            [ \t\n]+
WS
id
                           [A-Z\alpha-z][A-Z\alpha-zO-9]*
                            "(*"([^*]|\n|"*"+[^)])*"*"+")"
comment
                           [0-9]+/([^0-9]]"..")
integer
                           [0-9]+"."[0-9]*([0-9]|"E"[+-]?[0-9]+)
real
string
                           \'([^']|\'\)*\'
%%
```

Token definitions and action

```
":=" {printf(" %s ", yytext);return(T_ASSIGN);}

"else" {printf(" %s ", yytext);return(T_ELSE);}
```

lex.l Example (Cont'd)





Conditional compilation action



```
"then"
                          {#ifdef PRNTFLG
                           printf(" %s ", yytext);
                           #endif
                           return(T THEN);
"<="
                          {printf(" %s ", yytext);return(T EQ);}
"<"
                          {printf(" %s ", yytext);return(T LT);}
                          {printf(" %s ", yytext);return(T NE);}
                          {printf(" %s ", yytext);return(T GE);}
">="
">"
                          {printf(" %s ", yytext);return(T GT);}
{id}
                          {printf(" %s ", yytext);return(T IDENTIFIER);}
                          {printf(" %s ", yytext);return(T INTEGER);}
{integer}
                          {printf(" %s ", yytext);return(T REAL);}
{real}
                          {printf(" %s ", yytext);return(T STRING);}
{string}
                          {/* T COMMENT */}
{comment}
                          {/* spaces, tabs, newlines */}
{ws}
%%
                          {return 0;}
yywrap()
main()
 int i;
 do {
  i = yylex();
 } while (i!=0);
```

Token definitions and action

Discard

EOF for input

```
Three variables shared by
```

Lex&Yacc:

- 1. yytext = "currenttoken"
- 2. yylen = 12
- 3. yylval = 300









Internal Variables in Lex

- You may find the variables useful when writing translation rules
- char *yytext;
 - Pointer to current lexeme terminated by ' \setminus 0'
- int yylen;
 - Number of chacters in yytex but not '\0'
- yylval:
 - Global variable through which the token value can be returned to Yacc
 - Parser (Yacc) can access yylval, yylen, and yytext
- How are these used?

Consider **integer tokens**:

yylval = ascii_to_integer (yytext);

→ Conversion from *string* to actual *integer value*













Internal Variables in Lex (Cont'd)

- FILE *yyin
 - The input of the lex, pointing to the current file position
 - Default is set to stdin
- FILE *yyout
 - The output of the lex program
 - Default is set to stdout
- yylineno
 - The current line number of yyin













Lex Library Routines

- yylex()
 - The default main() contains a call of yylex()
- yymore()
 - return the next token
- yyless(n)
 - retain the first n characters in yytext
- yywarp()
 - is called whenever Lex reaches an end-of-file
 - The default yywarp() always returns 1













Lex Regular Expressions (Extended Regular Expressions)

- A regular expression matches a set of strings
- Regular expression
 - Operators
 - Character classes
 - Arbitrary character
 - Optional expressions
 - Alternation and grouping
 - Context sensitivity
 - Repetitions and definitions









Operators

- Considered as the meta-character of the regular expressions used in Lex
- If they are to be used as **text characters**, an escape should be used

 Every character is always a text character, except blank, tab (\t), newline (\n) and the list above









Character Classes []

- [abc]
 - matches a single character, which may be a, b, or c
- Every operator meaning is ignored
 - except \ and ^
- Examples:

[ab] => a or b

[a-z] => a or b or c or ... or z

[-+0-9] => all the digits and the two signs

 $[^a-zA-z]$ => any character which is not a letter











Arbitrary Character .

- The operator character.
 - is the class of all characters, except newline

An escape character example:

- [\40-\176]
 - matches all printable characters in the ASCII character set, from $40_{\rm octal}$ (space) to $176_{\rm octal}$ (tilde~)











Optional & Repeated Expressions

- a? => zero or one instance of a
- a* => zero or more instances of a
- a+ => one or more instances of a

Examples:

ab?c => ac or abc

[a-z]+ => all strings of lower case letters

[a-zA-Z][a-zA-Z0-9]*

=> all alphanumeric strings with a leading alphabetic character











Precedence of Operators

- Level of precedence
 - 1. Kleene closure (*), ?, +
 - 2. concatenation
 - 3. alternation (|)
- All operators are left associative
- Ex: a*b|cd* = ((a*)b)|(c(d*))











Pattern Matching Primitives

Metacharacter	Matches
	any character except newline
\n	newline
*	zero or more copies of the preceding expression
+	one or more copies of the preceding expression
?	zero or one copy of the preceding expression
^	beginning of line / complement
\$	end of line
a b	a or b
(ab)+	one or more copies of ab (grouping)
[ab]	a or b
a{3}	3 instances of a
"a+b"	literal "a+b" (C escapes still work)









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Regular Expression and its Action

```
// Input stream
                           // Input stream
a = b + c;
                           a = b + c;
                           // Output
                           a operator: ASSIGNMENT b + c;
응응
<regexp> <action>
<regexp> <action>
응응
응응
[abcd]
               {printf("%s ", yytext);}
"="
                {printf("operator: ASSIGNMENT");}
```









Transition Rules

- regexp <one or more blanks> action (C code);
- regexp <one or more blanks> { actions (C code) }
- A null statement; will ignore the input (no actions)
 - Example:

```
[\t\n];
```

Causes the three spacing characters to be ignored

```
// Input stream
a = b + c;
d = b * c;

// Corresponding output for the rule
a=b+c;d=b*c;
```











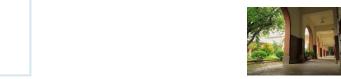
Transition Rules (cont'd)

- Four special options for actions:
 |, ECHO;, BEGIN, and REJECT;
- | indicates that the action for this rule is from the action for the next rule

The unmatched token is using a default action:
 ECHO from the input to the output













Transition Rules (cont'd)

• REJECT

- Go do the next alternative rule
- It causes whatever rule was second choice after the current rule to be executed

- The matching rules depend on the tool you use:
 - Matching rules for Flex:
 - Match the longest possible token
 - 2. Of the tokens with the same length, prefer the pattern earlier in the source file













Transition Rules (cont'd)

- BEGIN
 - For conditional rules

BEGIN name1;

- Executing the action statement enters a start condition, which changes the start condition to **name1**

BEGIN 0;

• Executing the action statement resets the initial condition of the Lex automaton interpreter

```
Example: %START AA BB CC
%%

^a {ECHO; BEGIN AA;}
^b {ECHO; BEGIN BB;}
^c {ECHO; BEGIN CC;}
\n {ECHO; BEGIN O;}
<AA>magic printf("first");
<BB>magic printf("second");
```

<CC>magic

- (AA) Copy the input to the output, except change the word magic to the word first on every line that begins with the letter a.
- **(BB)** Change magic to second on every line that begins with b.
- (CC) Change magic to third on every line that begins with c. Here is how the problem might be handled with a flag.

printf("third");











Usage

- To run Lex on a source file, type lex scanner.l
 - It produces a file named lex.yy.c which is a C program for the lexical analyzer
- To compile lex.yy.c, type cc lex.yy.c -ll
- To run the lexical analyzer program, type

./a.out < inputfile









Versions of Lex

- Versions:
 - AT&T Lex
 - http://www.combo.org/lex_yacc_page/lex.html
 - GNU Flex
 - https://github.com/westes/flex
 - Win32 version of Flex:
 - http://gnuwin32.sourceforge.net/packages/flex.htm
 - Lex on Cygwin: https://www.cygwin.com/
- Each Lex implementation may has its own character
 - Please refer to the online manual for Lex on The Lex & Yacc Page









QUESTIONS?











A lex.l example for tokens in **Pascal**

A Pascal lex.l

```
%{
#include "y.tab.h"
%}
letter
                      [a-zA-Z]
digit
                      [0-9]
                      [ \t\n]+
WS
id
                      [A-Za-z][A-Za-z0-9]*
                      "(*"([^*]|\n|"*"+[^)])*"*"+")"
comment
                      [0-9]+/([^0-9]|"..")
integer
                      [0-9]+"."[0-9]*([0-9]|"E"[+-]?[0-9]+)
real
                      \'([^']|\'\')*\'
string
%%
":="
                      {return(T_ASSIGN);}
...
                      {return(T_COLON);}
                      {return(T_ARRAY);}
"array"
"begin"
                      {return(T_BEGIN);}
"case"
                      {return(T_CASE);}
"const"
                      {return(T_CONST);}
"downto" {return(T_DOWNTO);}
"do"
                      {return(T_DO);}
"else"
                      {return(T_ELSE);}
"end"
                      {return(T_END);}
"file"
                      {return(T_FILE);}
"for"
                      {return(T_FOR);}
```







A Pascal lex.l (Cont'd)

```
"function" {return(T_FUNCTION);}
/* "goto"
                      {return(T GOTO);} */
"if"
                      {return(T_IF);}
"label"
                      {return(T LABEL);}
"nil"
                      {return(T_NIL);}
"not"
                      {return(T_NOT);}
"of"
                      {return(T OF);}
                      {return(T_PACKED);} */
/* "packed"
"procedure"
                      {return(T PROCEDURE);}
"end"
                      {return(T_END);}
"program" {return(T_PROGRAM);}
          {return(T RECORD);}
"record"
"repeat"
          {return(T REPEAT);}
"set"
                      {return(T SET);}
                      {return(T_THEN);}
"then"
"to"
                      {return(T TO);}
"type"
                      {return(T TYPE);}
"until"
                      {return(T_UNTIL);}
"var"
                      {return(T VAR);}
"while"
                      {return(T_WHILE);}
/* "with"
           {return(T WITH);} */
"+"
                      {return(T_PLUS);}
"_"
                      {return(T_MINUS);}
"or"
                      {return(T_OR);}
"and"
                      {return(T_AND);}
                      {return(T_DIV);}
"div"
"mod"
                      {return(T_MOD);}
"/"
                      {return(T_RDIV);}
```







A Pascal lex.l (Cont'd)

```
11 🛠 11
                     {return(T MULT);}
                     {return(T_LPAREN);}
")"
                     {return(T_RPAREN);}
                     {return(T_EQ);}
                     {return(T COMMA);}
                     {return(T_RANGE);}
                     {return(T_PERIOD);}
"F"
                     {return(T_LBRACK);}
                     {return(T_RBRACK);}
"<="
                     {return(T_EQ);}
"<"
                     {return(T_LT);}
"<>"
                     {return(T_NE);}
">="
                     {return(T_GE);}
">"
                     {return(T_GT);}
"in"
                     {return(T_IN);}
|| V ||
                      {return(T UPARROW);}
...
                     {return(T_SEMI);}
{id}
                     {return(T_IDENTIFIER);}
{integer}
          {return(T INTEGER);}
                     {return(T_REAL);}
{real}
          {return(T_STRING);}
{string}
                     {/* T COMMENT */}
{comment}
{ws}
                     {/* spaces, tabs, newlines */}
```





