- How to organize carpools such that nobody drives "much more" than he/she would driving alone from his/her home?
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- Can we do both?

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- The company has $N \approx 500$ employees, all living in a 50 mile radius of the workplace.
- We know the location of each employee's home, and can compute the driving distance from any one to any other one.
- We assume that in a carpool some of the employees drive to another one's home, park there, and then travel together in one car; maybe they pick up additional employees on their way to the workplace. For sake of simplicity, we assume that all cars can accommodate 4 people the driver and 3 passengers.
- For "fairness" we require that in a carpool nobody travels "much more" than his/her own travel distance to the workplace. (Say, no more than M extra miles.)

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SETS:

• Employees E – each with its home location.

- Workplace w with its location.
- Driving distances $d[i, j], i \in E, j \in E \cup \{w\}$ (in miles).
- Extra mileage allowed to be traveled: M (in miles).
- Car capacities: s[i] is the number of seats in the car of employee $i \in E$. We assume that those who arrive to the home of i will travel together in the car of employee i; for simplicity, we assume s[i] = 4 for all $i \in E$.

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Ride Sharing: Decision Variables

VARIABLES:

- f[i, j] denotes the number of people traveling together from i to $j, i \in E, j \in E \cup \{w\}$. (Sounds like a **FLOW** looks like a **FLOW**, smells like a **FLOW** it is a **FLOW**!)
- Looks like a transshipment problem: employee homes $i \in E$ have a supply = 1, and workplace w has a demand = N.

$$\begin{split} 1 + \sum_{i \in E} f[i,j] &= \sum_{k \in E} f[j,k] \quad \forall j \in E, \\ \sum_{i \in E} f[i,w] &= N, \\ f[i,j] &\leq s[i] \qquad \forall i \in E, \ j \in E \cup \{w\}. \end{split}$$

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Ride Sharing: How to Formulate the Objective?

OBJECTIVE FUNCTION:

• Total mileage driven:

$$\min \sum_{\substack{i \in E \\ j \in E \cup \{w\}}} d[i, j] * f[i, j]$$

• Is this correct? What does it measure?

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VARIABLES:

• $y \in \{0,1\}^{E \times E \cup \{w\}}$

$$y[i,j] = \begin{cases} 1 & \text{if } f[i,j] > 0, \\ 0 & \text{otherwise.} \end{cases}$$

Enforced by

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- Introduce $\ell[i]$ to denote the distance traveled by employee $i \in E$; $\ell[w] = 0$.
- Then we have the following recursive update for this quantity:

$$\ell[i] = \min_{i \in E} y[i, j] * (d[i, j] + \ell[j]) \quad \forall i \in E,$$

- Does NOT work if $\sum_{i \in E} y[i,j] > 1$ for some $i \in E!$
- Assume (restricting solutions)

$$\sum_{j \in E \cup \{w\}} y[i,j] = 1 \quad \forall \ i \in E.$$

$$\ell[i] = \sum_{i \in E \cap I_m \setminus 1} y[i,j] * (d[i,j] + \ell[j]) \quad \forall \ i \in E.$$

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Linearizing the product of a binary variable and a linear function²

- Assume that L is a linear expression of our variables and X is a binary variable.
- We would like to enforce Z = X * L for all feasible solutions.
- Assume that $0 \le L \le U$ for all feasible solutions, where U is a constant.
- We have Z = X * L if and only if

$$Z \leq L,$$

 $Z \leq U * X,$
 $Z \geq 0,$
 $Z \geq L - U * (1 - X).$

²Glover (1975), Adams and Forester (2005)



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Linearizing the product of a binary variable and a linear $function^2$

- Assume that L is a linear expression of our variables and X is a binary variable.
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- Assume that $0 \le L \le U$ for all feasible solutions, where U is a constant.
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$$\begin{array}{ll} Z & \leq & L, \\ Z & \leq & U * X, \\ Z & \geq & 0, \\ Z & > & L - U * (1 - X). \end{array}$$



VARIABLES:

- Introduce $\ell[i]$ to denote the distance traveled by employee $i \in E$. $\ell[w] = 0$
- Introduce $z[i,j] = y[i,j] * (d[i,j] + \ell[j])$, enforced by

$$z[i,j] \leq d[i,j] + \ell[j] \qquad \forall i \in E, \ j \in E \cup \{u \\ z[i,j] \leq 100 * y[i,j] \qquad \forall i \in E, \ j \in E \cup \{u \\ z[i,j] \geq 0 \qquad \forall i \in E, \ z[i,j] \geq 0 \qquad \forall z[i,j]$$

Then we can write

$$\ell[i] = \sum_{i \in F: \{i,j\}} z[i,j] \quad \forall \ i \in E.$$

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