### Optimistic Assumptions in Polyhedral Compilation

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## Be Optimistic!

*Programs* might be nasty but *programmers* are not.

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Optimistic Assumptions & Speculative Versioning

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### **Optimistic Assumptions**



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1 Make optimistic assumptions to (better) optimize loops



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- Make optimistic assumptions to (better) optimize loops
- Derive runtime conditions that imply these assumptions



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- 3 Version the code based on the assumptions made and conditions derived.



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- Make optimistic assumptions to (better) optimize loops
- Derive runtime conditions that imply these assumptions
- 3 Version the code based on the assumptions made and conditions derived.

### Speculative Versioning

```
if (/* Runtime Conditions */)
   /* Optimized Loop Nest */
else
   /* Original Loop Nest */
```

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**Runtime Conditions** 

#### **Runtime Conditions**

- Fast to derive (compile time)
- Fast to verify (runtime)
- High probability to be true
- ToDo: A feedback/profile driven approach



- (A) Applicability/Correctness
  - No Alias Assumption<sup>1</sup>
  - No Wrapping Assumption<sup>2</sup>
  - 3 Finite Loops Assumption<sup>2</sup>
  - 4 Array In-bounds Assumption<sup>2</sup>
  - 5 Valid Multidimensional View Assumption (Delinearization)<sup>3</sup>
- (B) Optimizations
  - Array In-bounds Check Hoisting<sup>2</sup>
  - Parametric Dependence Distances<sup>4</sup>
  - 3 Possibly Invariant Loads

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<sup>&</sup>lt;sup>1</sup>Joint work Fabrice Rastello (INRIA Grenoble) & others. [OOPSLA'15]

<sup>&</sup>lt;sup>2</sup>Joint work with Tobias Grosser (ETH)

<sup>&</sup>lt;sup>3</sup>Tobias Grosser & Sebastian Pop (Samsung) [ICS'15]

<sup>&</sup>lt;sup>4</sup>Joint work with Zino Benaissa (Qualcomm)



```
void mem_copy(int N, float *A, float *B) {
  if (
    #pramga vectorize
    for (i = 0; i < N; i++)
        A[i] = B[i+5];
} else {
    /* original code */
}
</pre>
```



```
void mem_copy(int N, float *A, float *B) {
   if (&A[0] >= &B[N+5] ||

    #pramga vectorize
   for (i = 0; i < N; i++)
        A[i] = B[i+5];
} else {
    /* original code */
}
</pre>
```



No Alias Assumptions

```
void mem_copy(int N, float *A, float *B) {
   if (&A[0] >= &B[N+5] || &A[N] <= &B[5]) {
    #pramga vectorize
   for (i = 0; i < N; i++)
        A[i] = B[i+5];
} else {
    /* original code */
}
</pre>
```

Compare minimal/maximal accesses to possible aliasing arrays



No Alias Assumptions

```
void evn_odd(int N, int *Evn, int *Odd, int *A, int *B) {
  if
    for (int i = 0: i < N: i += 2)
      if (N % 2)
        Odd[i/2] = A[i+1] - B[i+1];
      else
        Evn[i/2] = A[i] + B[i];
  } else {
    /* original code */
```

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- Compare minimal/maximal accesses to possible aliasing arrays
- Do not compare accesses to read-only arrays



No Alias Assumptions

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void evn_odd(int N, int *Evn, int *Odd, int *A, int *B) {
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  } else {
    /* original code */
```

- Compare minimal/maximal accesses to possible aliasing arrays
- Do not compare accesses to read-only arrays
- Use the iteration domain of the accesses



#### No Alias Assumptions

```
void evn odd(int N. int *Evn. int *Odd. int *A. int *B) {
  if (N\%2 ? ((\&B[N+1] \le \&Odd[0] | | \&Odd[(N+1)/2] \le \&B[1]) \&\&
                 (&A[N+1] \le &Odd[0] | &Odd[(N+1)/2] \le &A[1])
          : ((\&B[N] \le \&Evn[0] | \&Evn[(N+1)/2] \le \&B[0]) \&\&
                 (&A[N] \le &Evn[O] \mid &Evn[(N+1)/2] \le &A[O]))) {
    for (int i = 0: i < N: i += 2)
      if (N % 2)
        Odd[i/2] = A[i+1] - B[i+1];
      else
        Evn[i/2] = A[i] + B[i]:
  } else {
    /* original code */
```

- Compare minimal/maximal accesses to possible aliasing arrays
- Do not compare accesses to read-only arrays
- Use the iteration domain of the accesses



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```
void mem_shift(unsigned char N, float *A) {
  if (N <= 128) {

    #pramga vectorize
    for (unsigned char i = 0; i < N; i++)
        A[i] = A[N + i];

} else {
    /* original code */
}
}</pre>
```

- Finite bit width can cause integer expressions to "wrap around"
- Wrapping causes multiple addresses for one memory location

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$$\underline{i} * \underline{c_0} + \underline{p} * \underline{c_1} \equiv_{p} (\underline{i} * \underline{c_0} + \underline{p} * \underline{c_1}) \mod 2^k$$



$$\underline{i} * \underline{c_0} + \underline{p} * \underline{c_1} \equiv_{p} (\underline{i} * \underline{c_0} + \underline{p} * \underline{c_1}) \mod 2^{k}$$

$$i \in [0, N-1] \land N \in [0, 2^{8}]$$

$$(N + i) \equiv_{p} (N + i) \mod 2^{8}$$

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$$i \in [0, N-1] \land N \in [0, 2^{8}]$$

$$(N + i) \equiv_{p} (N + i) \mod 2^{8}$$

$$\Longrightarrow (N + i) \leq_{p} 255$$



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$$i \in [0, N-1] \land N \in [0, 2^8]$$

$$(N + i) \equiv_{p} (N + i) \mod 2^8$$

$$\Longrightarrow (N + i) \leq_{p} 255$$

$$\Longrightarrow N \leq 128$$



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Finite Loops Assumption

```
void mem_shift(unsigned N, float *A) {
  if (N % 2 == 0) {
    #pramga vectorize
    for (unsigned i = 0; i != N; i+=2)
        A[i+4] = A[i];
} else {
    /* original code */
}
}
```

- Allows to provide other LLVM passes real loop bounds
- Infinite loops create unbounded optimization problems





Array In-bounds Assumptions



```
void stencil(int N, int M, float A[128][128]) {
    if (N <= 128) {

        #pragma loop interchange
        for (int i = 0; i < N; i++)
            for (int j = 0; j < M; j++)
                 A[2*j][i] += A[2*j+1][i];

    } else {
        /* original code */
    }
}</pre>
```

Out-of-bound accesses introduce multiple addresses for one memory location (e.g., &A[1][0] == &A[0][128])







Valid Multidimensional View Assumption

- Define multi-dimensional view of a linearized (one-dimensional) array
- Derive conditions that accesses are in-bounds for each dimension



```
struct SafeArray { int Size, int *Array };
inline void set(SafeArray A, int idx, int val) {
  if (idx < 0 || A.Size <= idx)
      throw OutOfBounds;
  A.Array[idx] = val;
}

void set_safe_array(int N, SafeArray A) {
  for (int i = 0; i < N; i++)
      for (int j = 0; j < i/2; j++)
      set(A, i+j, 1); /* Throws out-of-bounds */
}</pre>
```



```
struct SafeArray { int Size, int *Array };
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Array In-bounds Check Hoisting

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struct SafeArray { int Size, int *Array };
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Array In-bounds Check Hoisting

```
struct SafeArray { int Size, int *Array };
inline void set(SafeArray A, int idx, int val) {
  if (idx < 0 \mid | A.Size \le idx)
    throw OutOfBounds;
  A.Array[idx] = val;
void set_safe_array(int N, SafeArray A) {
    if (
      for (int i = 0; i < N; i++)
        for (int j = 0; j < i/2; j++)
          A[i+j] = 1;
    } else {
      /* original code */
```





Array In-bounds Check Hoisting

```
struct SafeArray { int Size, int *Array };
inline void set(SafeArray A, int idx, int val) {
  if (idx < 0 \mid | A.Size \le idx)
    throw OutOfBounds;
  A.Array[idx] = val;
void set_safe_array(int N, SafeArray A) {
    if ((3*N)/2 \le A.Size) {
      for (int i = 0; i < N; i++)
        for (int j = 0; j < i/2; j++)
          A[i+j] = 1;
    } else {
      /* original code */
```

■ Hoist in-bounds access conditions out of the loop nest





#### **Check Hoisting**

```
void copy(int N, double A[N][N], double B[N][N]) {
    if (DebugLevel <= 5) {</pre>
      #pragma parallel
      for (int i = 0; i < N; i++)
        #pragma simd
        for (int j = 0; j < N; j++)
          A[i][i] = B[i][i];
    } else {
      for (int i = 0; i < N; i++) {
        for (int j = 0; j < N; j++)
          A[i][i] = B[i][i];
        if (DebugLevel > 5)
          printf("Column_{\sqcup}\%d_{\sqcup}copied\n", i)
      }
```





Parametric Dependence Distances

```
void vectorize(int N, double *A) {
   if (c >= 4) {

    #pragma vectorize width(4)
   for (int i = c; i < N+c; i++)
        A[i-c] += A[i];

} else {
    /* original code */
}
</pre>
```

■ Assume *large enough* dependence distance, e.g., for vectorization



```
void may_load(int *size0, int *size1) {
  for (int i = 0; i < *size0; i++)
    for (int j = 0; j < *size1; j++)
    ...
}</pre>
```

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Possibly Invariant Loads

```
void may_load(int *size0, int *size1) {
  int size0val = *size0;
  int size1val = 1;

  if
    size1val = *size1;

  for (int i = 0; i < size0val; i++)
    for (int j = 0; j < size1val; j++)
    ...
}</pre>
```

Hoist invariant loads

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Possibly Invariant Loads

```
void may_load(int *size0, int *size1) {
  int size0val = *size0;
  int size1val = 1;

  if (size0val > 0)
      size1val = *size1;

  for (int i = 0; i < size0val; i++)
      for (int j = 0; j < size1val; j++)
      ...
}</pre>
```

- Hoist invariant loads
- Keep conditions for conditionally executed loads

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Possibly Invariant Loads

```
void may_load(int *size0, int *size1) {
  int size0val = *size0;
  int size1val = 1;

if (size0val > 0)
    size1val = *size1;

for (int i = 0; i < size0val; i++)
    for (int j = 0; j < size1val; j++)
    ...
}</pre>
```

- Hoist invariant loads
- Keep conditions for conditionally executed loads
- Powerful in combination with alias checks

### **Ongoing Work**



- Use profiling to minimize non-beneficial code duplication
- Derive more powerful checks (e.g., generate inspector loops)
- Find more opportunities to speculatively optimize using runtime conditions



## Thank You.