Advances in Loop Analysis Frameworks and Optimizations

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Loop Unrolling

```
for (x = 0; x < 6; x++) {
  foo(x);
}</pre>
```

Loop Unrolling

```
for (x = 0; x < 6; x += 2) {
  foo(x);
  foo(x + 1);
}</pre>
```

Loop Unrolling

```
foo(x);
foo(x + 1);
foo(x + 2);
foo(x + 3);
foo(x + 4);
foo(x + 5);
```

Unrolling: Pros and Cons

- + Removes loop overhead
- + Enables other optimizations
- Increases code size
- Increases compile time
- Might regress performance

New Heuristics

- Aim for bigger loops
- Analyze the loop body and predict potential optimization candidates for later passes

```
const int b[50] = \{1, 0, 0, ..., 0, 0\};
int foo(int *a) {
 int r = 0;
  for (int i = 0; i < 50; i++) {
   r += a[i] * b[i];
  return r;
```

```
const int b[50] = \{1, 0, 0, ..., 0, 0\};
int foo(int *a) {
 int r = 0;
 r += a[0] * b[0];
 r += a[1] * b[1];
 r += a[48] * b[48];
 r += a[49] * b[49];
  return r:
```

```
const int b[50] = \{1, 0, 0, ..., 0, 0\};
int foo(int *a) {
 int r = 0;
 r += a[0] * 1;
 r += a[1] * 0;
 r += a[48] * 0;
 r += a[49] * 0;
  return r:
```

```
const int b[50] = {1, 0, 0, ..., 0, 0};
int foo(int *a) {
  return a[0];
}
```

Analyzing Loop

- Simulate the loop execution instruction by instruction, iteration by iteration
- Try to predict possible simplifications of every instruction
- Compute accurate costs of the original loop and its unrolled version

Iteration 0

```
%r = 0
loop:
%y = b[i]
  %x = a[i]
  %t = %x * %y
  %r = %r + %t
  \%i = \%i + 1
  %cmp = %i < 50
  br %cmp, loop, exit
exit:
                                    Original loop
                                              Unrolled loop
                                      cost
   ret %r
```

cost

```
%r = 0
 loop:
- \text{%y} = \text{b[i]} = 1
   %x = a[i]
   %t = %x * %y
   %r = %r + %t
   \%i = \%i + 1
   %cmp = %i < 50
   br %cmp, loop, exit
 exit:
                                        Original loop
                                                    Unrolled loop
                                           cost
                                                       cost
   ret %r
```

```
%r = 0
  loop:
    %y = b[i] = 1
%x = a[i]
    %t = %x * %y
    %r = %r + %t
    \%i = \%i + 1
    %cmp = %i < 50
    br %cmp, loop, exit
  exit:
                                     Original loop
                                                Unrolled loop
                                        cost
                                                   cost
    ret %r
```

```
%r = 0
 loop:
   %y = b[i] = 1
   %x = a[i]
%t = %x * %y = %x
   %r = %r + %t
   \%i = \%i + 1
   %cmp = %i < 50
   br %cmp, loop, exit
 exit:
                                    Original loop
                                               Unrolled loop
                                       cost
                                                 cost
   ret %r
```

```
%r = 0
loop:
  %y = b[i] = 1
  %x = a[i]
  %t = %x * %y = %x
%r = %i + %t = %t
  \%i = \%i + 1
  %cmp = %i < 50
  br %cmp, loop, exit
exit:
                                    Original loop
                                              Unrolled loop
                                      cost
                                                 cost
   ret %r
```

Iteration 0

```
%r = 0
loop:
 %y = b[i] = 1
 %x = a[i]
 %t = %x * %y = %x
 %r = %i + %t = %t
 \%i = \%i + 1 = 1
  %cmp = %i < 50
  br %cmp, loop, exit
exit:
  ret %r
```



Original loop cost

Unrolled loop cost

Iteration 0

```
%r = 0
loop:
  %y = b[i] = 1
  %x = a[i]
  %t = %x * %y = %x
  %r = %i + %t = %t
  \%i = \%i + 1 = 1
 %cmp = \%i < 50 = true
  br %cmp, loop, exit
exit:
  ret %r
```



cost

Unrolled loop cost

Iteration 0

```
%r = 0
loop:
  %y = b[i] = 1
  %x = a[i]
  %t = %x * %y = %x
  %r = %i + %t = %t
  \%i = \%i + 1 = 1
  %cmp = \%i < 50 = true
br %cmp, loop, exit
exit:
   ret %r
```



Unrolled loop cost

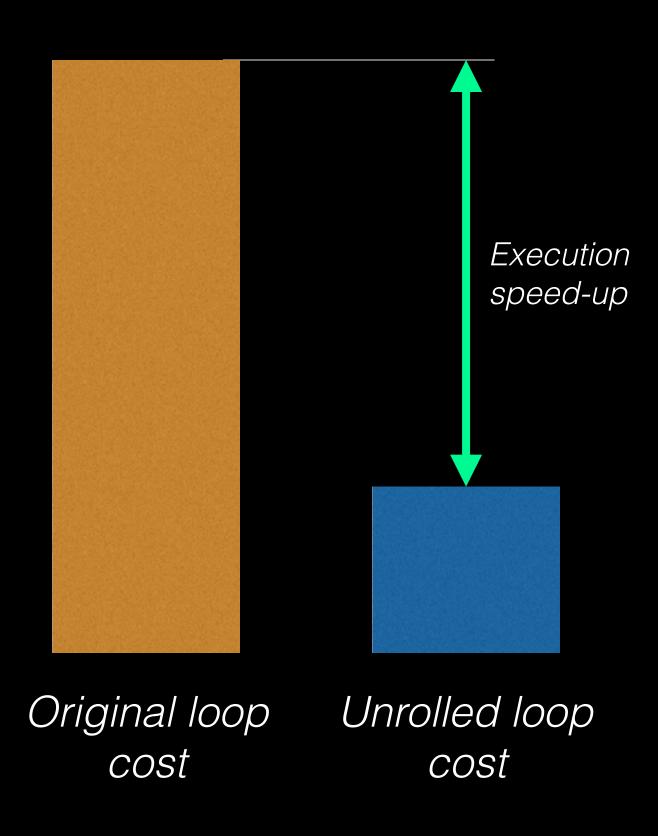
```
%r = 0
 loop:
%y = b[i]
   %x = a[i]
   %t = %x * %y
   %r = %r + %t
   \%i = \%i + 1
   %cmp = %i < 50
   br %cmp, loop, exit
exit:
                                    Original loop
                                               Unrolled loop
                                       cost
                                                  cost
   ret %r
```

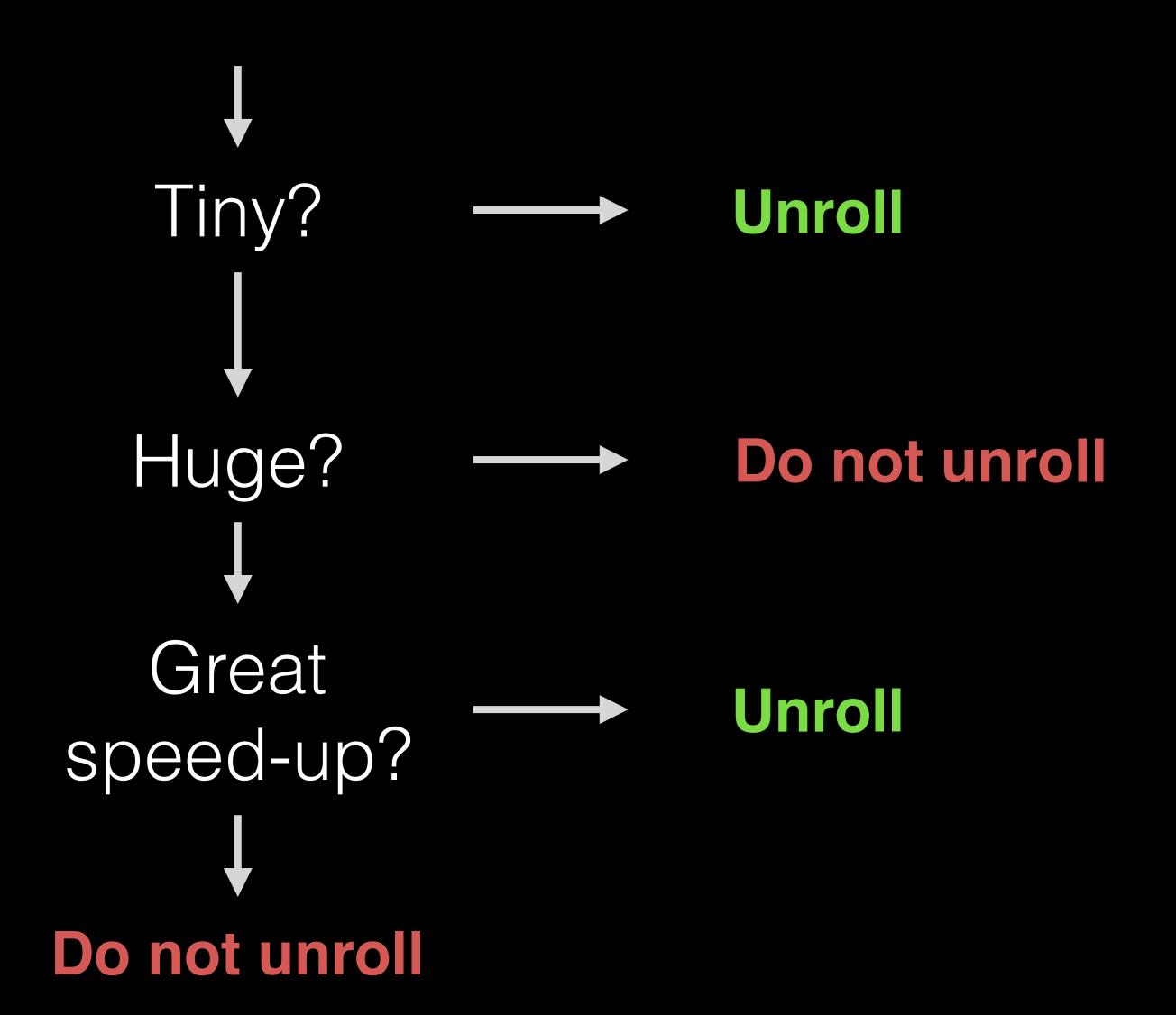
```
%r = 0
loop:
  %y = D[i] = 0
  %x _ :: [1]
  %t = %x * %y = 0
  %r = %r + %t = %r
  \%i = \%i + 1 = 2
  %cmp = \%i < 50 = true
br %cmp, loop, exit
exit:
  ret %r
```



Unrolled loop cost

```
%r = 0
loop:
  %y = b[i]
  %x = a[i]
  %t = %x * %y
  %r = %r + %t
  \%i = \%i + 1
  %cmp = %i < 50
  br %cmp, loop, exit
exit:
                                    Original loop
                                               Unrolled loop
                                       cost
                                                  cost
  ret %r
```





Unrolling: Results

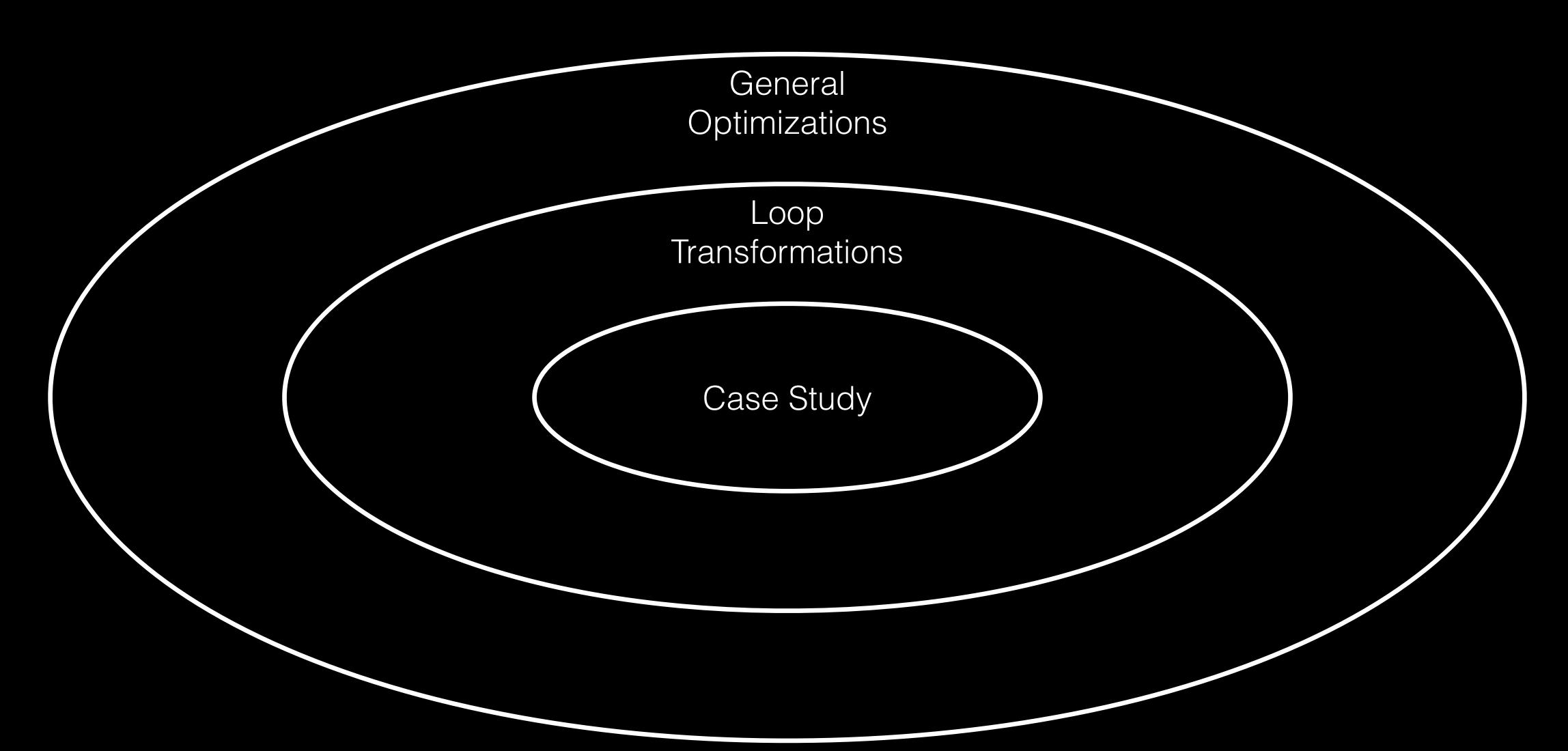
- Up to 70% performance gains on kernels
- Few performance gains across various testsuites
- No performance regressions
- Some compile time regressions

Unrolling: Future Work

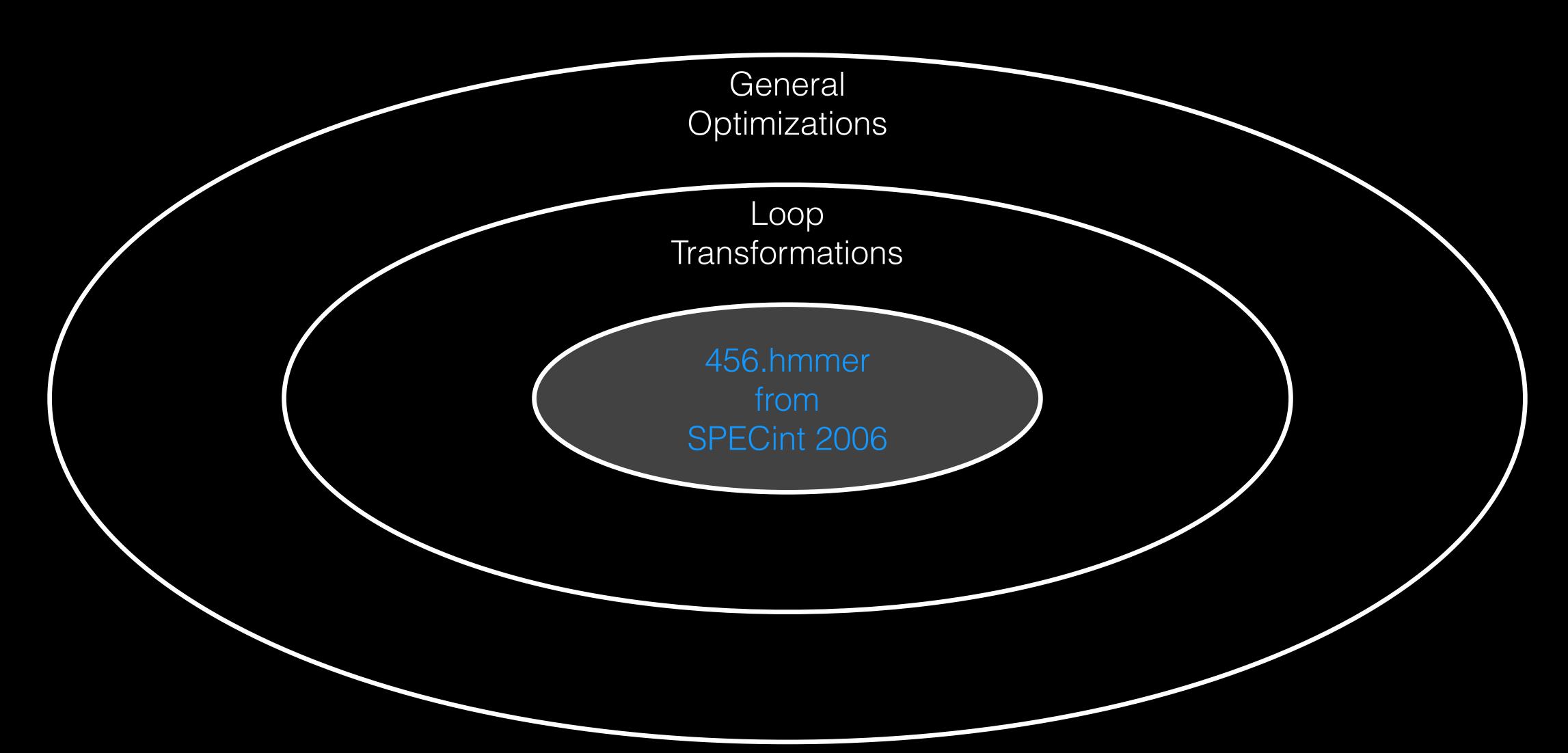
- Enable new heuristics by default after investigating compile time regressions
- Model other optimizations
- Find trip count

Next Up

Approach



Approach



```
for (k = 1; k \le M; k++) {
 mc[k] = mpp[k-1] + tpmm[k-1];
 if ((sc = ip[k-1] + tpim[k-1]) > mc[k]) mc[k] = sc;
 if ((sc = dpp[k-1] + tpdm[k-1]) > mc[k]) mc[k] = sc;
 if ((sc = xmb + bp[k]) > mc[k]) mc[k] = sc;
 mc[k] += ms[k];
 if (mc[k] < -INFTY) mc[k] = -INFTY;
 dc[k] = dc[k-1] + tpdd[k-1];
 if ((sc = mc[k-1] + tpmd[k-1]) > dc[k]) dc[k] = sc;
 if (dc[k] < -INFTY) dc[k] = -INFTY;
 if (k < M) {
   ic[k] = mpp[k] + tpmi[k];
   if ((sc = ip[k] + tpii[k]) > ic[k]) ic[k] = sc;
   ic[k] += is[k];
   if (ic[k] < -INFTY) ic[k] = -INFTY;
```

```
What does this loop do?
if ((sc = mc[k-1] + tpmd[k-1]) > dc[k]) dc[k] = sc;
```

```
mc[k] = mpp[k-1] + tpmm[k-1];
if ((sc = ip[k-1] + tpim[k-1]) > mc[k]) mc[k] = sc;
if ((sc = dpp[k-1] + tpdm[k-1]) > mc[k]) mc[k] = sc;
if ((sc = xmb + bp[k]) > mc[k]) mc[k] = sc;
mc[k] += ms[k];
if (mc[k] < -INFTY) mc[k] = -INFTY;</pre>
```

```
mc[k] = mpp[k-1] + tpmm[k-1];
if ((sc = ip[k-1] + tpim[k-1]) > mc[k]) mc[k] = sc;
if ((sc = dpp[k-1] + tpdm[k-1]) > mc[k]) mc[k] = sc;
if ((sc = xmb + bp[k]) > mc[k]) mc[k] = sc;
mc[k] += ms[k];
if (mc[k] < -INFTY) mc[k] = -INFTY;</pre>
```

```
dc[k] = dc[k-1] + tpdd[k-1];
if ((sc = mc[k-1] + tpmd[k-1]) > dc[k]) dc[k] = sc;
if (dc[k] < -INFTY) dc[k] = -INFTY;</pre>
```

```
dc[k] = dc[k-1] + tpdd[k-1];
if ((sc = mc[k-1] + tpmd[k-1]) > dc[k]) dc[k] = sc;
if (dc[k] < -INFTY) dc[k] = -INFTY;</pre>
```

```
if (k < M) {
   ic[k] = mpp[k] + tpmi[k];
   if ((sc = ip[k] + tpii[k]) > ic[k]) ic[k] = sc;
   ic[k] += is[k];
   if (ic[k] < -INFTY) ic[k] = -INFTY;
}</pre>
```

```
if (k < M) {
   ic[k] = mpp[k] + tpmi[k];
   if ((sc = ip[k] + tpii[k]) > ic[k]) ic[k] = sc;
   ic[k] += is[k];
   if (ic[k] < -INFTY) ic[k] = -INFTY;
}</pre>
```

```
How can we optimize this loop?
```

```
if (mc[k] < -INFTY) \c[k] - INFTY;

dc[k] = dc[k-1] + tpdd[k-1];</pre>
```

```
dc[k] = dc[k-1] + tpdd[k-1];
if ((sc = mc[k-1] + tpmd[k-1]) > dc[k]) dc[k] = sc;
if (dc[k] < -INFTY) dc[k] = -INFTY;
```

```
dc[k] = dc[k-1] + tpdd[k-1];
```

```
t = dc[k-1] + tpdd[k-1];
dc[k] = t;
```

```
t = dc[k-1] + tpdd[k-1];
dc[k] = t;
```

```
t = dc[k-1] + tpdd[k-1];
dc[k] = t;
```

```
Iteration K:
         = dc[k-1] + tpdd[k-1];
         dc[k]
    Iteration K+1:
                     = dc[k] + tpdd[k];
              t2
              dc[k+1] = t2;
```

```
dc[k] = dc[k-1] + tpdd[k-1];
if ((sc = mc[k-1] + tpmd[k-1]) > dc[k]) dc[k] = sc;
if (dc[k] < -INFIY) dc[k] = -INFTY;</pre>
```

```
for (k = 1; k <= M; k++) {
    mc[k] = mpp[k-1] + tpmm[k-1];
    if ((sc = ip[k-1] + tpim[k-1]) > mc[k]) mc[k] = sc;
    if ((sc = dpp[k-1] + tpdm[k-1]) > mc[k]) mc[k] = sc;
    if ((sc = xmb + bp[k]) > mc[k]) mc[k] = sc;
    mc[k] += ms[k];
    if (mc[k] < -INFTY) mc[k] = -INFTY;</pre>
```

dc[k] = dc[k-1] + tpdd[k-1]; if | Ope=V+Ctor|Zap|+ dc[k] = sc; if (dc[k] < -INFTY) dc[k] = -INFTY;</pre>

```
if (k < M) {
   ic[k] = mpp[k] + tpmi[k];
   if ((sc = ip[k] + tpii[k]) > ic[k]) ic[k] = sc;
   ic[k] += is[k];
   if (ic[k] < -INFTY) ic[k] = -INFTY;
}</pre>
```

```
for (k = 1; k \le M; k++) 
 \frac{dc[k] = dc[k-1] + tpdd[k-1];}{10m + V + Ctorrizable dc[k] = sc;}
 if (k < M) {
   ic[k] = mpp[k] + tpmi[k];
   if ((sc = ip[k] + tpii[k]) > ic[k]) ic[k] = sc;
   ic[k] += is[k];
   if (ic[k] < -INFTY) ic[k] = -INFTY;
```

```
if (k < M) {
   ic[k] = mpp[k] + tpmi[k];
   if ((sc = ip[k] + tpii[k]) > ic[k]) ic[k] = sc;
   ic[k] += is[k];
   if (ic[k] < -INFTY) ic[k] = -INFTY;
}</pre>
```

```
if (k < M) {
  ic[k] = mpp[k] + tpmi[k];
  if ((sc = ip[k] + tpii[k]) > ic[k]) ic[k] = sc;
  ic[k] += is[k];
  if (ic[k] < -INFTY) ic[k] = -INFTY;
}</pre>
```

```
for (k = 1; k \le M; k++) 
  if ((sc = ip[k-1] + tpim[k-1]) > mc[k]) mc[k] = sc;
if ((sc = constant)) e^{mc[k]} mc[k] = sc;
if ((sc = constant)) e^{mc[k]} e^{mc[k]} = sc;
  Mon-vectorizable
```

```
for (k = 1; k \le M; k++) {
for (k = 1; k \le M; k++) {
  Non-vectorizable
```

Plan

- Distribute loop
- Let LoopVectorizer vectorize top loop

-> Partial Loop Vectorization

Loop Distribution

Pros and Cons

- + Partial loop vectorization
- + Improve memory access pattern:
 - Cache associativity
 - Number of HW prefetcher streams
- + Reduce spilling
- Loop overhead
- Instructions duplicated across new loops
- Instruction-level parallelism

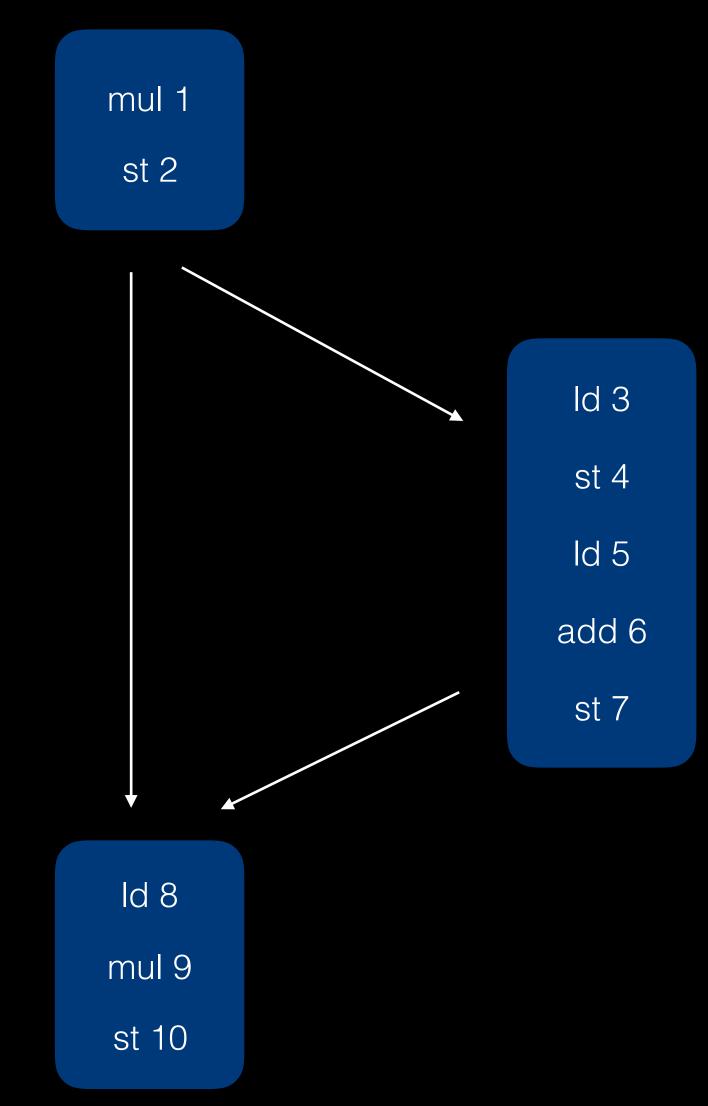
Legality

```
for (k = 1; k \le M; k++) {
 if ((sc = xmb + bp[k]) > mc[k]) mc[k] = sc;
for (k = 1; k <= M; k++) {
 if ((sc = mc[k-1] + tpmd[k-1]) > d
 if (dc[k] < -INFTY) dc[k] = -INF
                                           Run-time
                            Loop
   ic[k] = mpp[k] + tpmi
   if ((sc = ip[k] + tp
                        Dependence
                                             Alias
                          Analysis
                                           Checks
```

Loop Access Analysis

- Born from the Loop Vectorizer
- Generalized as new analysis pass
- Computed on-demand and cached
- New Loop Versioning utility

- Light-weight
 - Uses only LoopAccessAnalysis
 - No Program Dependence Graph
 - No Control Dependence
- Inner loops only
- Different from textbook algorithm
 - No reordering of memory operations



mul 1

st 2

ld 3

st 4

ld 5

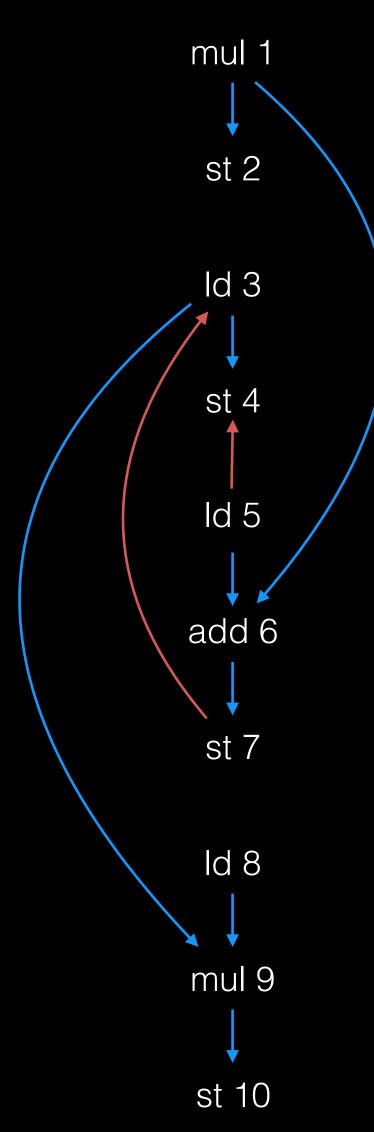
add 6

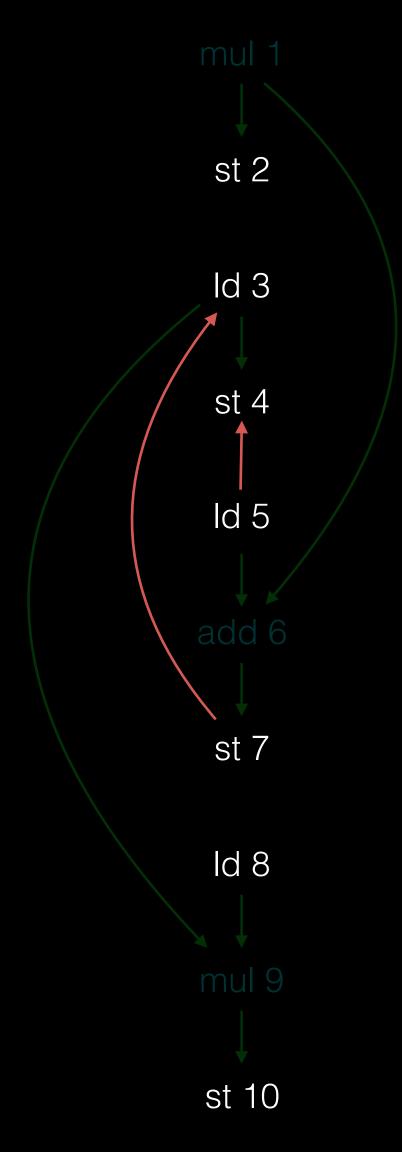
st 7

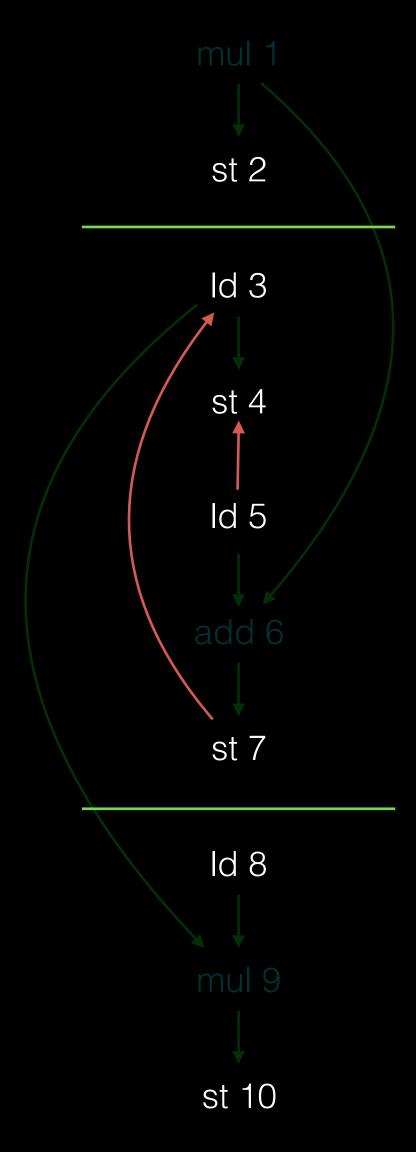
ld 8

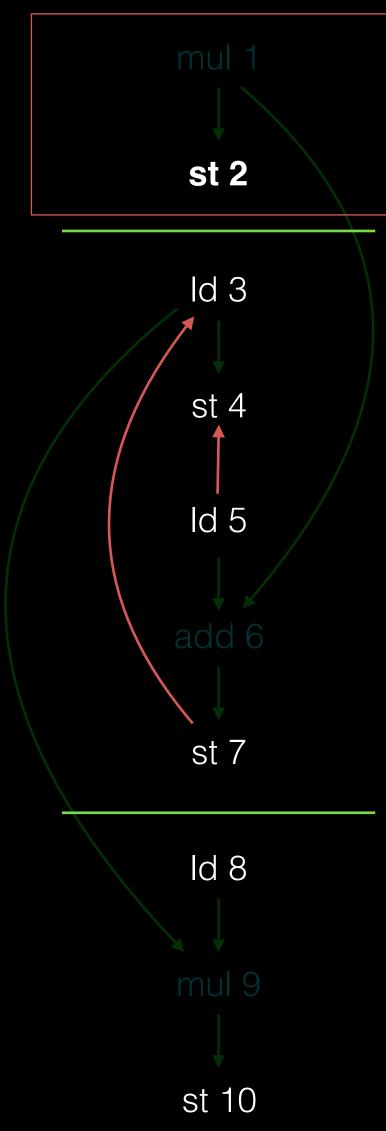
mul 9

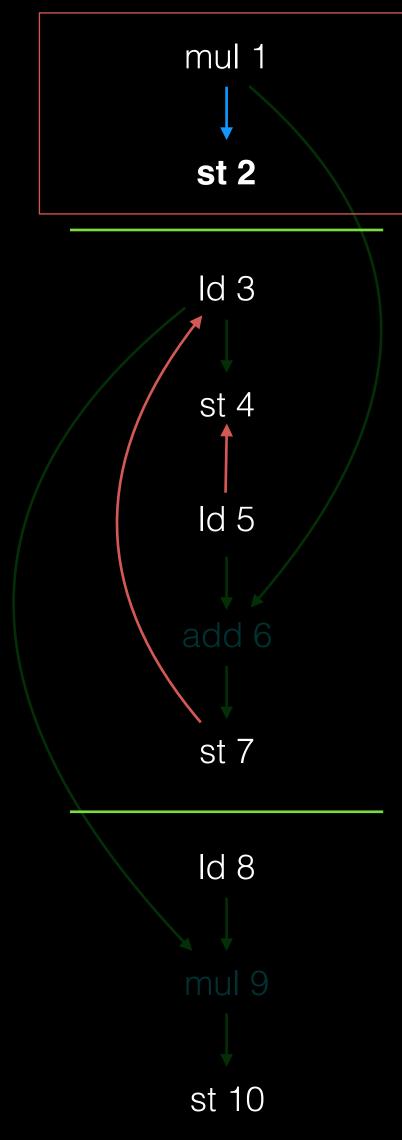
st 10

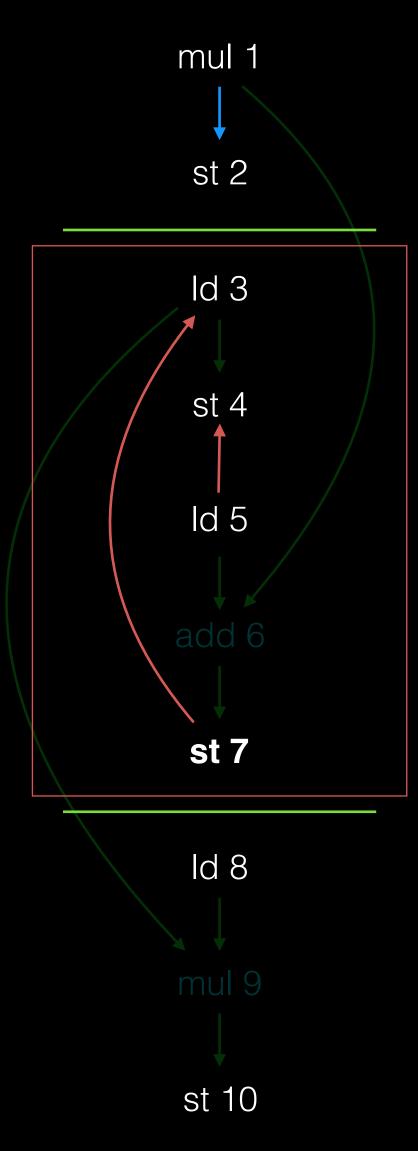


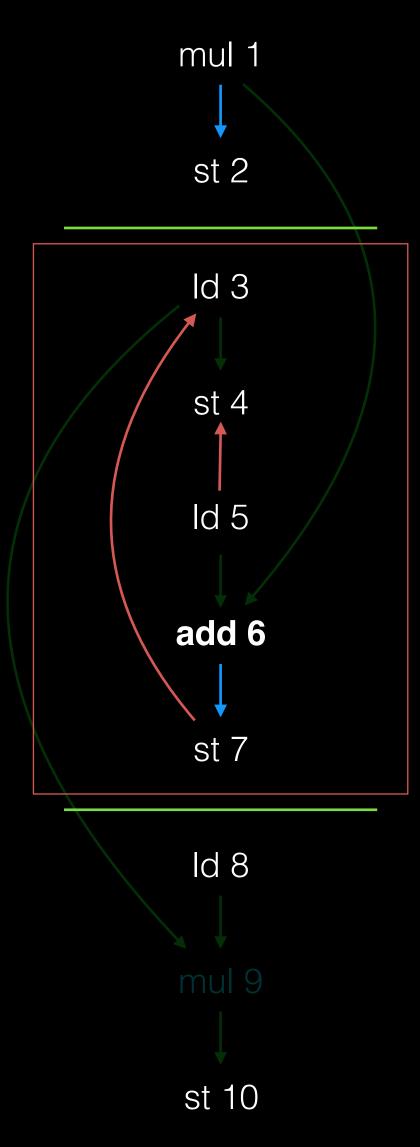


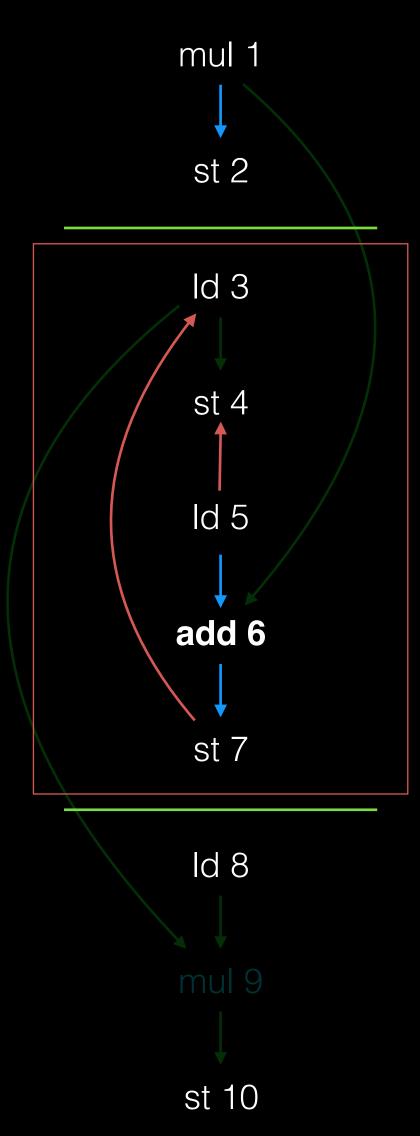


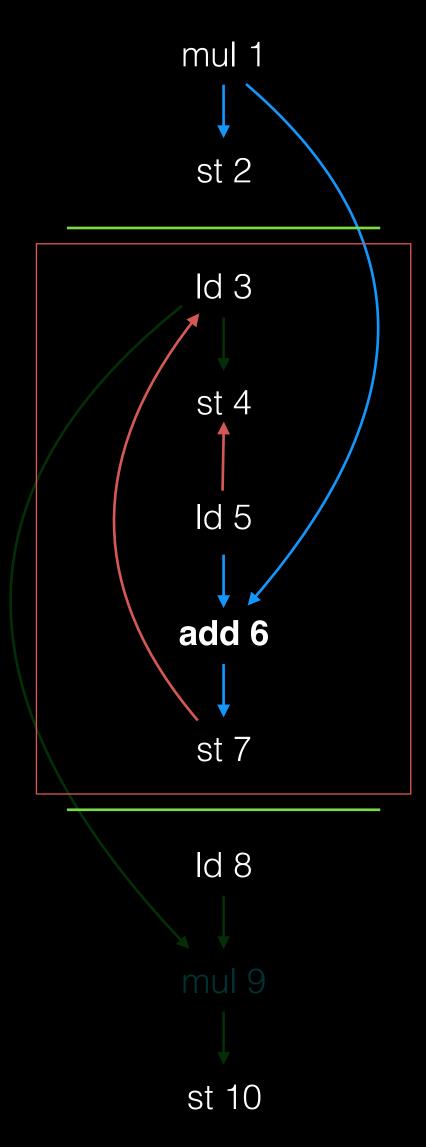


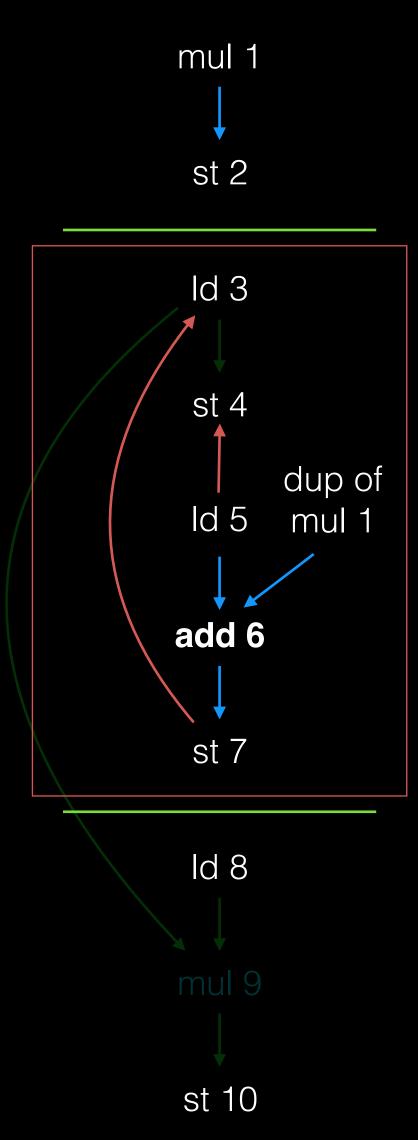


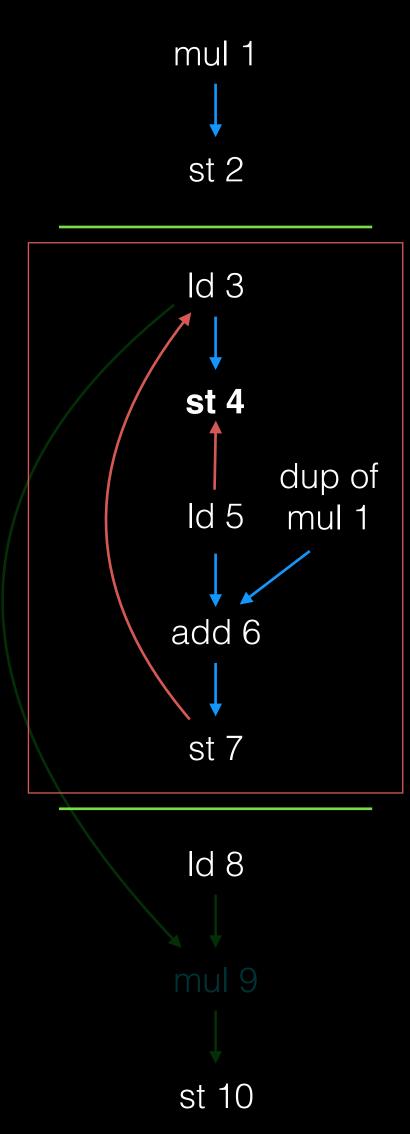


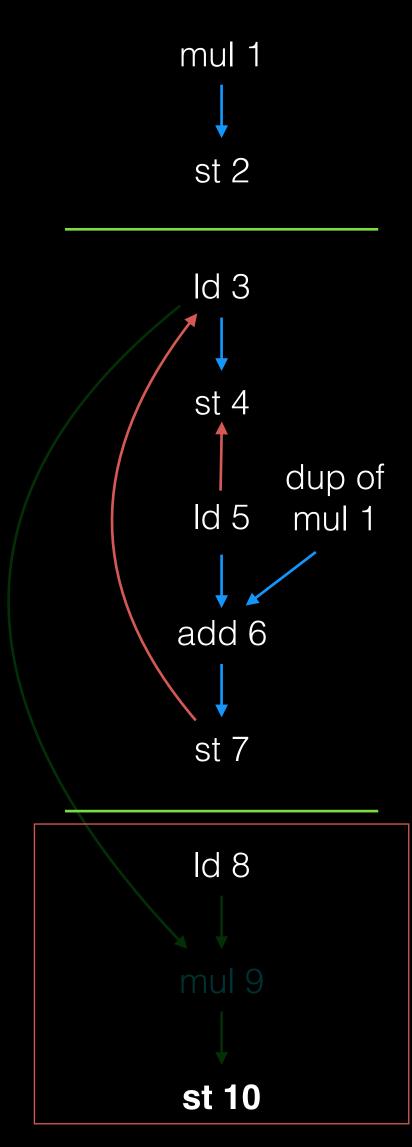


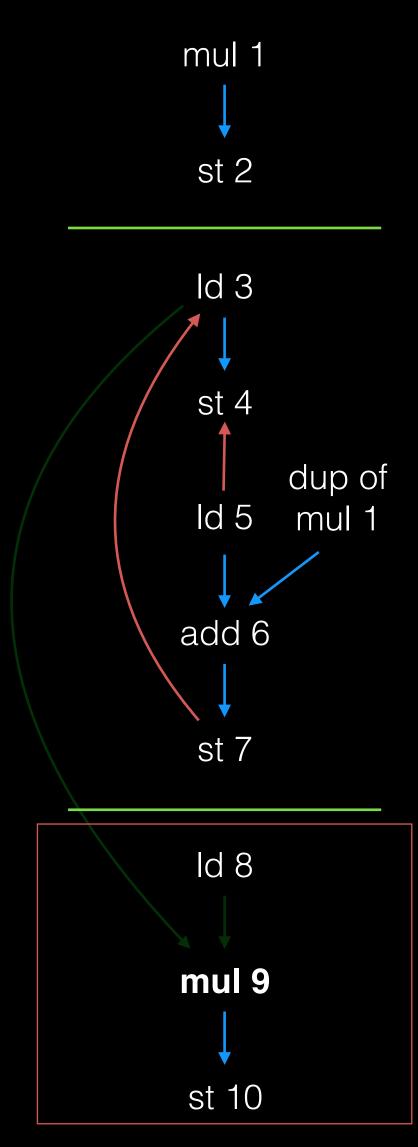


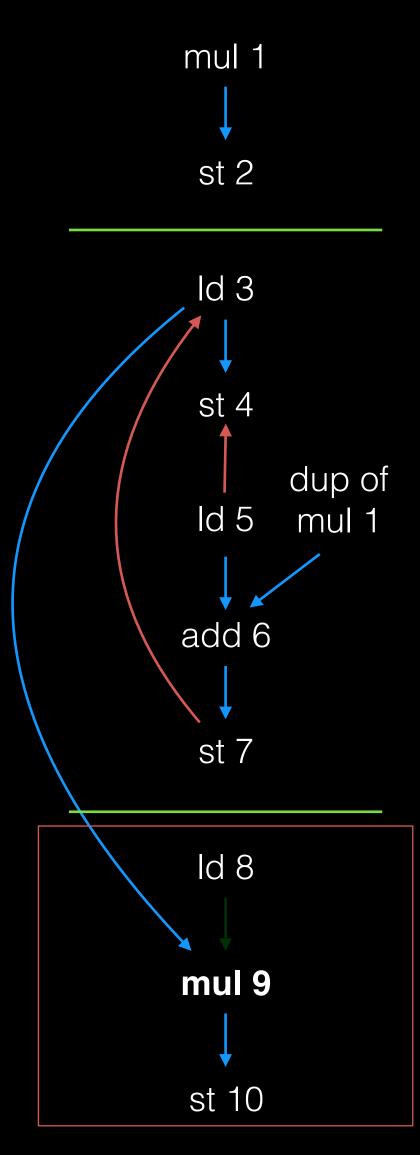


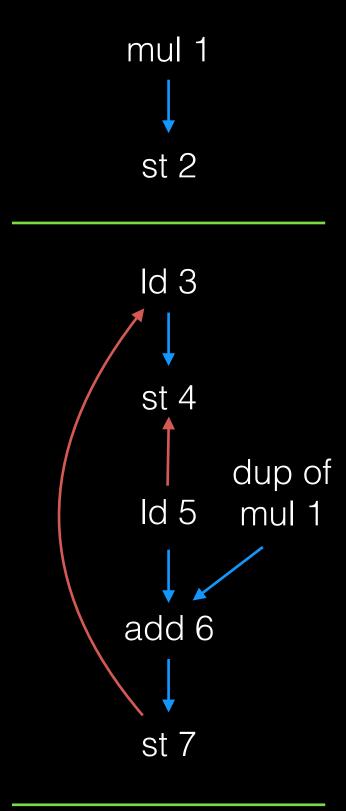


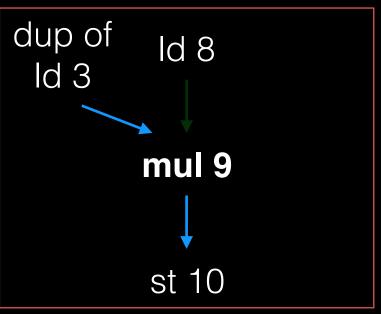


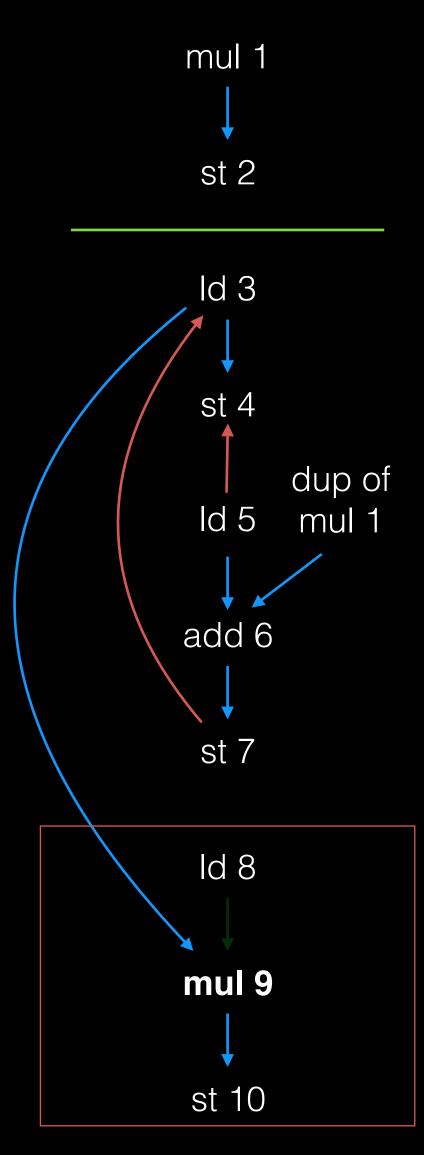


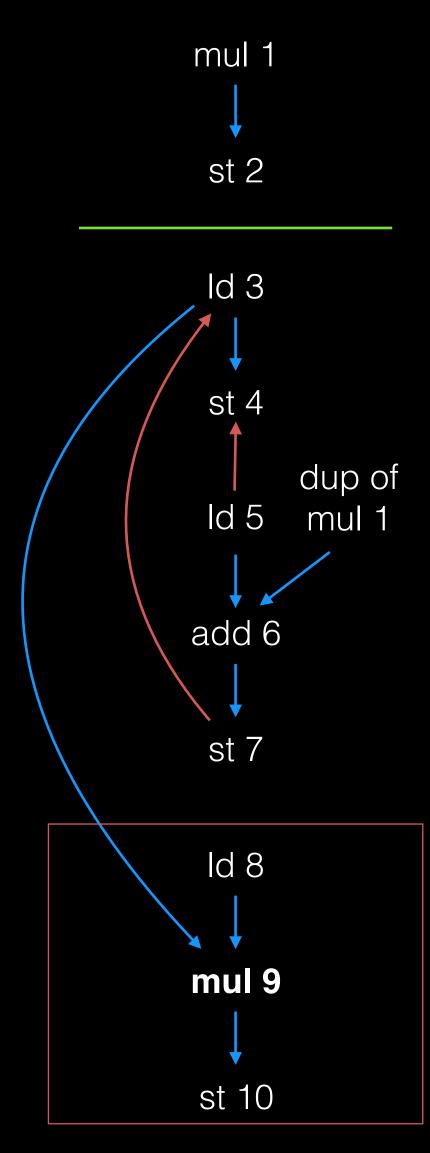


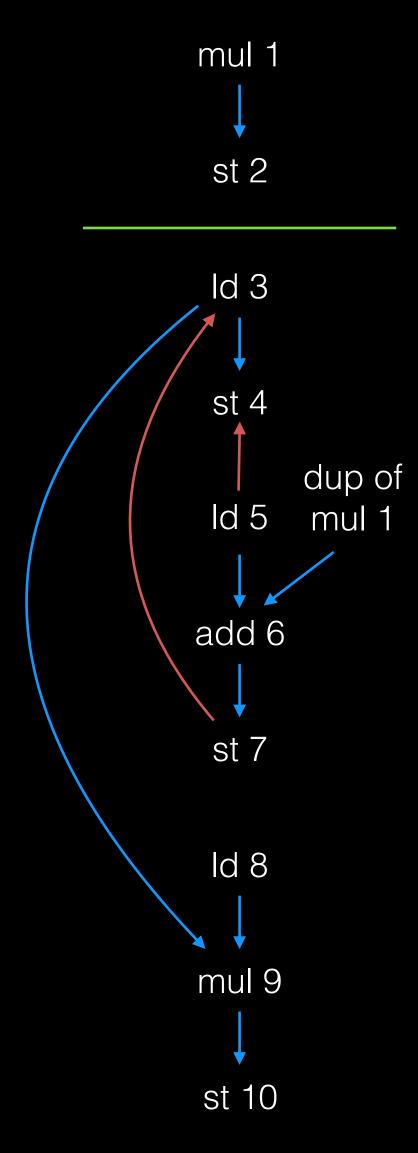












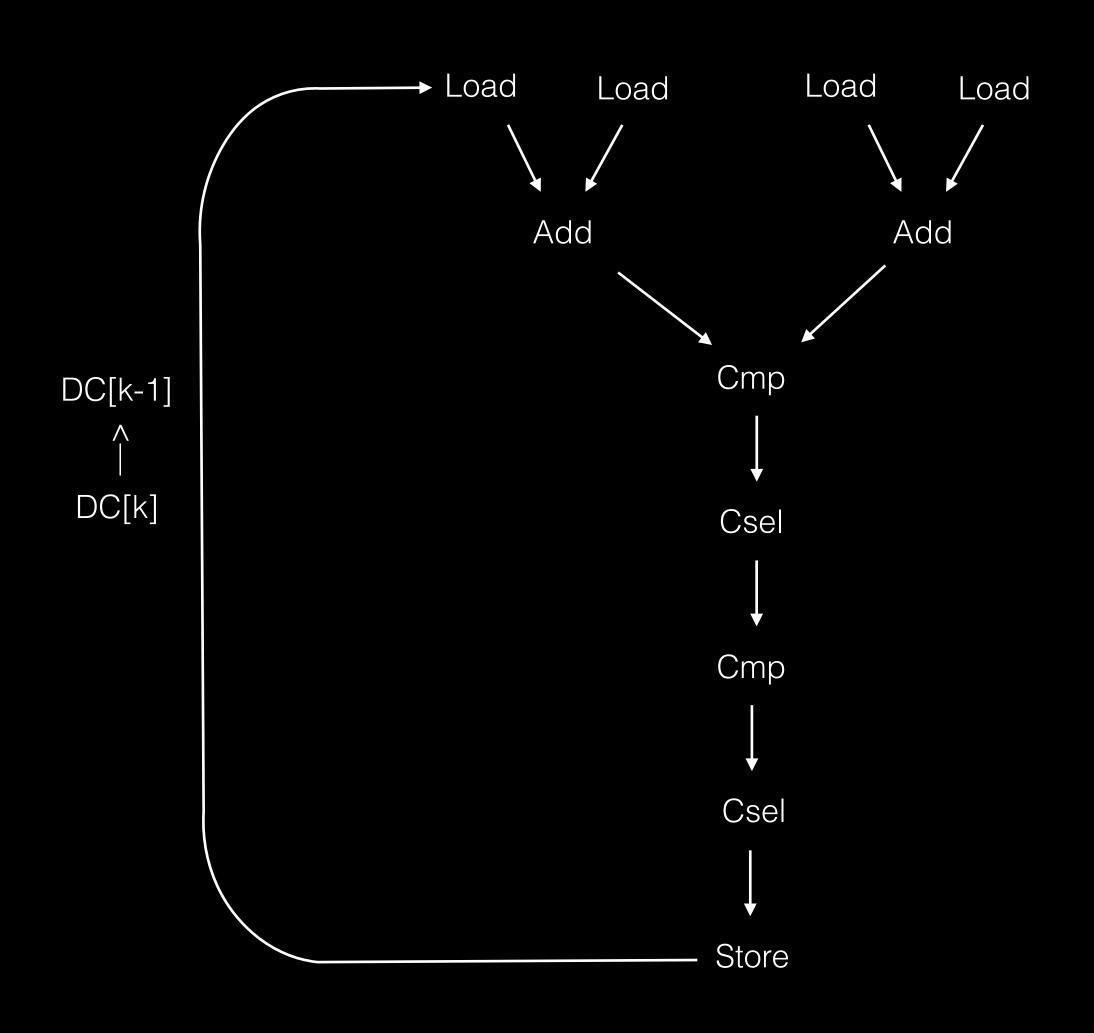
Recap

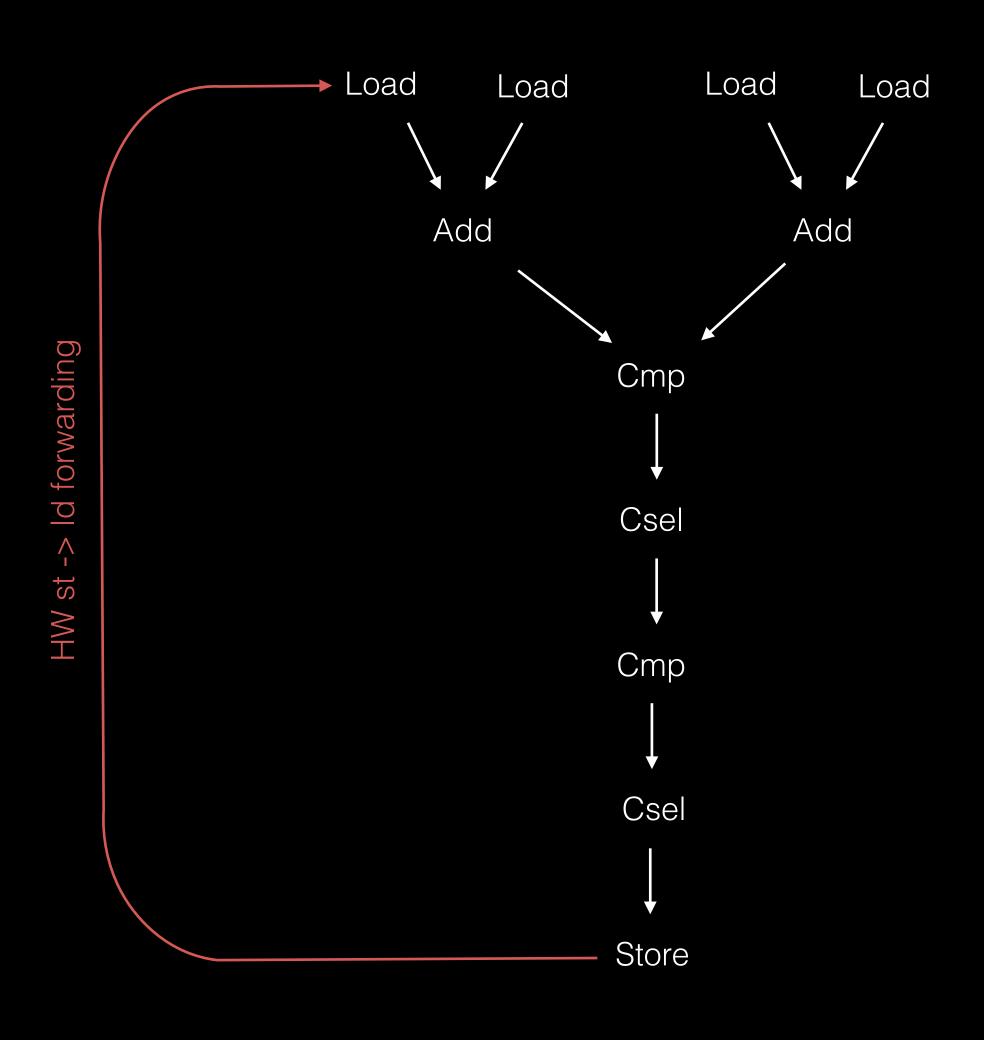
- Distributed loop
 - Versioned with run-time alias checks
- Top loop vectorized

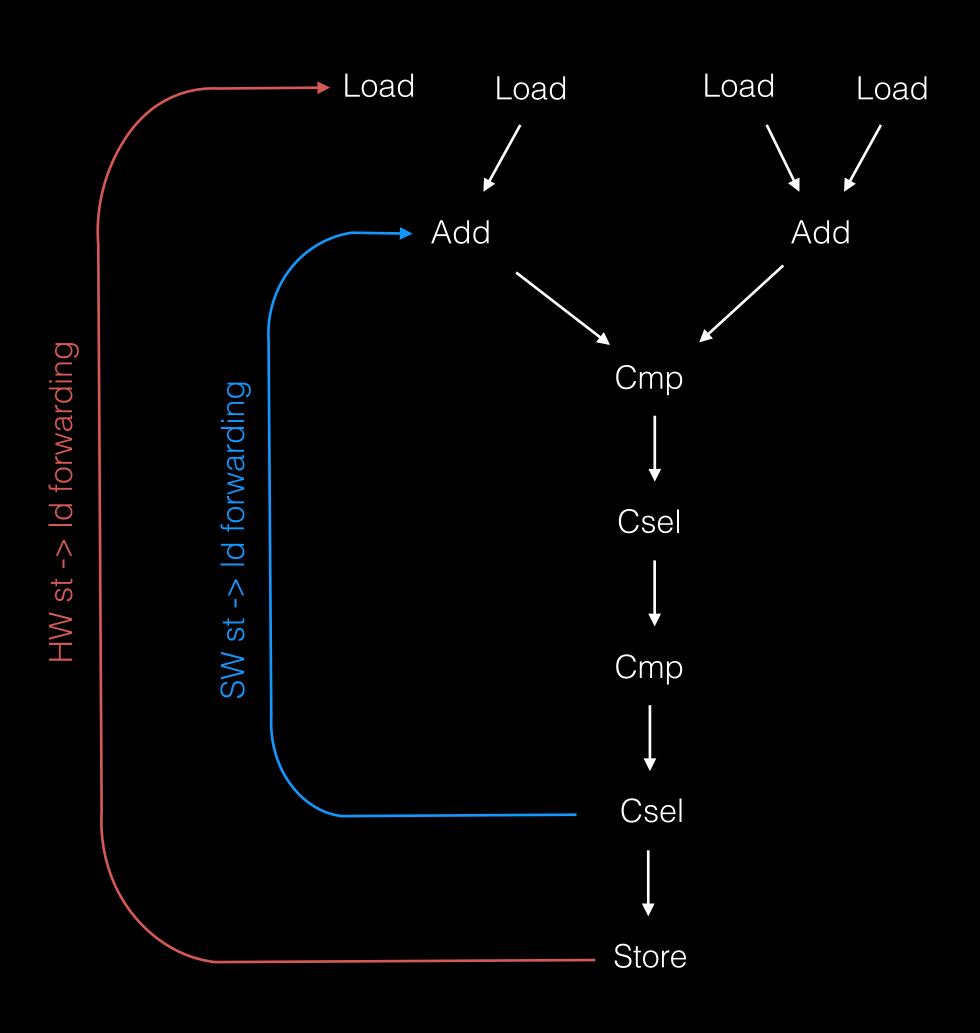
```
for (k = 1; k \le M; k++) {
  if ((sc = ip[k-1] + tpim[k-1]) > mc[k]) mc[k] = sc;
if ((sc = xmb+ cp[k]) / mc[k] > mc[k]) mc[k] = sc;
if ((sc = xmb+ cp[k]) / mc[k] = sc;
for (k = 1; k \le M; k++) {
```

```
for (k = 1; k \le M; k++) {
 dc[k] = dc[k-1] + tpdd[k-1];
 if ((sc = mc[k-1] + tpmd[k-1]) > dc[k]) dc[k] = sc;
 if (dc[k] < -INFTY) dc[k] = -INFTY;
 if (k < M) {
   ic[k] = mpp[k] + tpmi[k];
   if ((sc = ip[k] + tpii[k]) > ic[k]) ic[k] = sc;
   ic[k] += is[k];
   if (ic[k] < -INFTY) ic[k] = -INFTY;
```

```
dc[k] = dc[k-1] + tpdd[k-1];
if ((sc = mc[k-1] + tpmd[k-1]) > dc[k]) dc[k] = sc;
if (dc[k] < -INFTY) dc[k] = -INFTY;</pre>
```







```
dc[k] = dc[k-1] + tpdd[k-1];
if ((sc = mc[k-1] + tpmd[k-1]) > dc[k]) dc[k] = sc;
if (dc[k] < -INFIY) dc[k] = -INFTY;</pre>
```

Loop Load Elimination

- 1. Find loop-carried dependences with iteration distance of one
- 2. Between store -> load?
- 3. No (may-)intervening store
- 4. Propagate value stored to uses of load

```
for (k = 1; k \le M; k++) {
 dc[k] = dc[k-1] + tpdd[k-1];
 if ((sc = mc[k-1] + tpmd[k-1]) > dc[k]) dc[k] = sc;
 if (dc[k] < -INFTY) dc[k] = -INFTY;
```

```
for (k = 1; k \le M; k++) {
 dc[k] = T = dc[k-1] + tpdd[k-1];
 if ((sc = mc[k-1] + tpmd[k-1]) > dc[k]) dc[k] = T = sc;
 if (dc[k] < -INFTY) dc[k] = T = -INFTY;
```

```
for (k = 1; k \le M; k++) {
 dc[k] = T = T + tpdd[k-1];
 if ((sc = mc[k-1] + tpmd[k-1]) > dc[k]) dc[k] = T = sc;
 if (dc[k] < -INFTY) dc[k] = T = -INFTY;
```

```
T = dc[0];
for (k = 1; k \le M; k++) {
  dc[k] = T = T + tpdd[k-1];
  if ((sc = mc[k-1] + tpmd[k-1]) > dc[k]) dc[k] = T = sc;
  if (dc[k] < -INFTY) dc[k] = T = -INFTY;
```

```
T = dc[0];
for (k = 1; k \le M; k++) {
    ic[k] = mpp[k] + tpmi[k];
    if ((sc = ip[k] + tpii[k]) > ic[k]) ic[k] = sc;
    ic[k] += is[k];
   if (ic[k] < -INFTY) ic[k] = -INFTY;
```

Loop Load Elimination

- Simple and cheap using Loop Access Analysis
 - With Loop Versioning can optimize more loops
- GVN Load-PRE can be simplified to not worry about loop cases

Recap

- Distributed loop into two loops
 - Versioned with run-time alias checks
- Vectorized top loop
- Store-to-load forwarding in bottom loop
 - Versioned with run-time alias checks

Results

- 20-30% gain on 456.hmmer on ARM64 and x86
- Loop Access Analysis pass
- Loop Versioning utility
- Loop Distribution pass
- Loop Load Elimination pass

Future Work

- Commit Loop Load Elimination
- Tune Loop Distribution and turn it on by default
- Loop Distribution with Program Dependence Graph

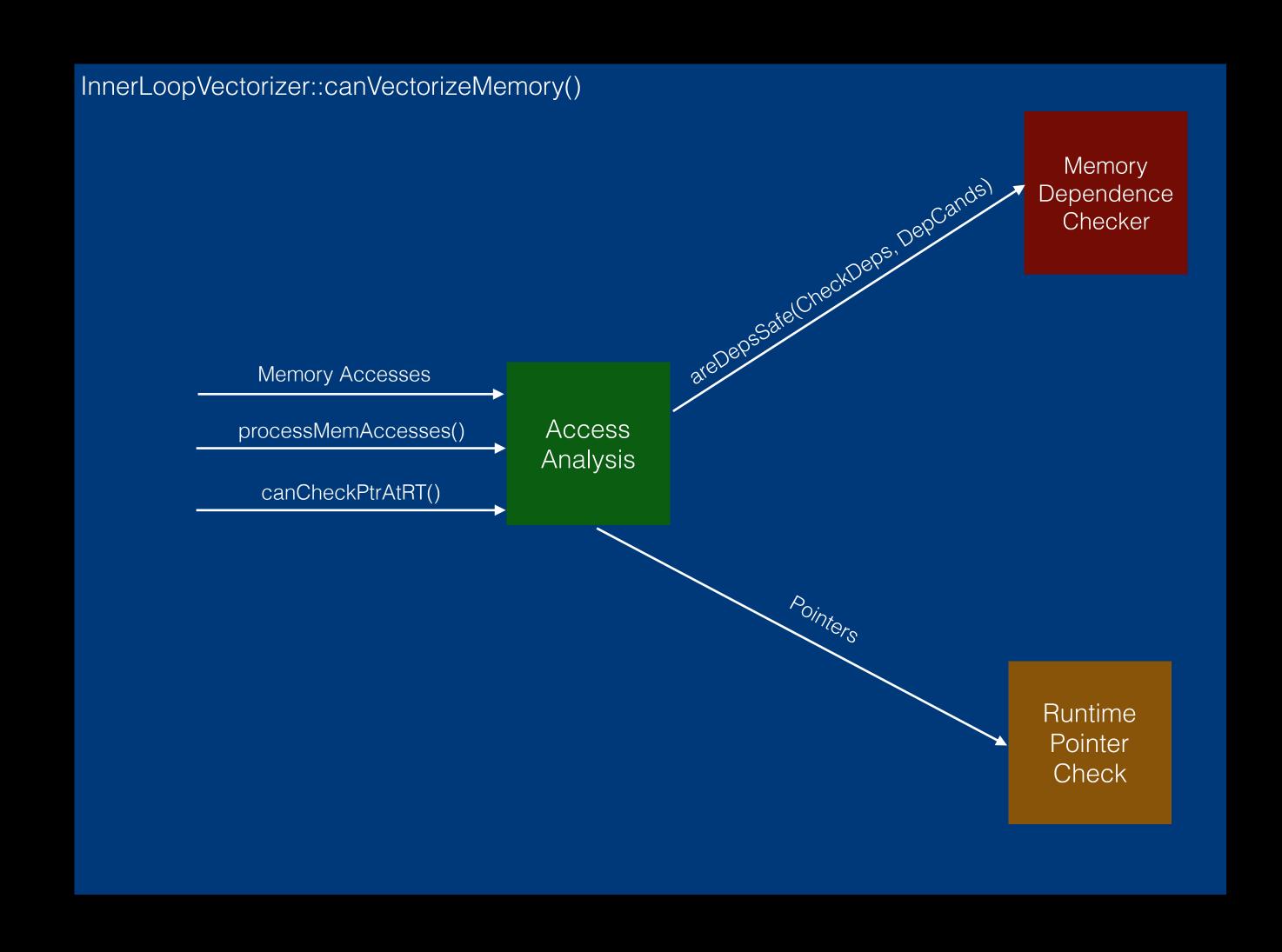
Acknowledgements

- Chandler Carruth
- Hal Finkel
- Arnold Schwaighofer
- Daniel Berlin

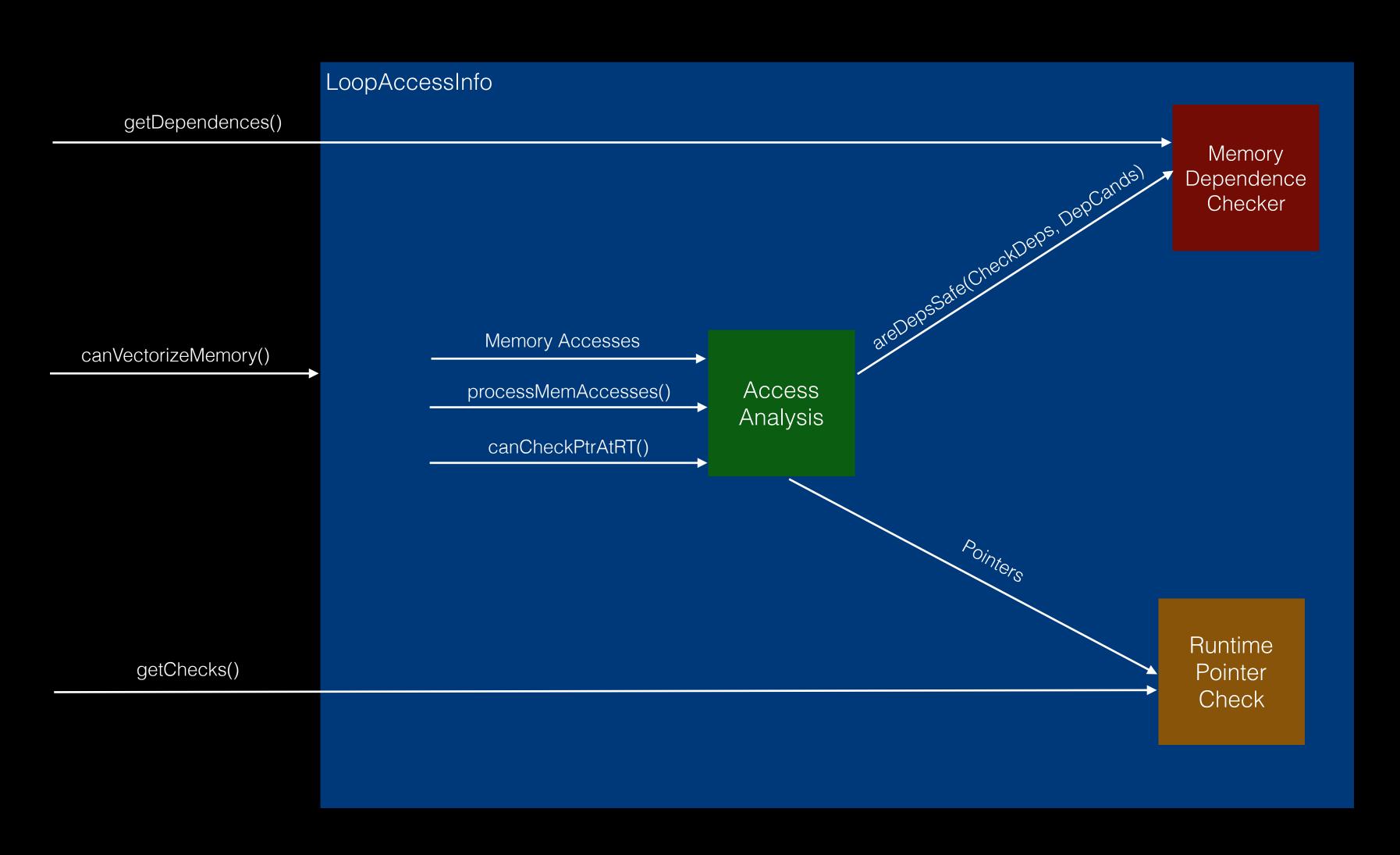
Q&A

Back-up

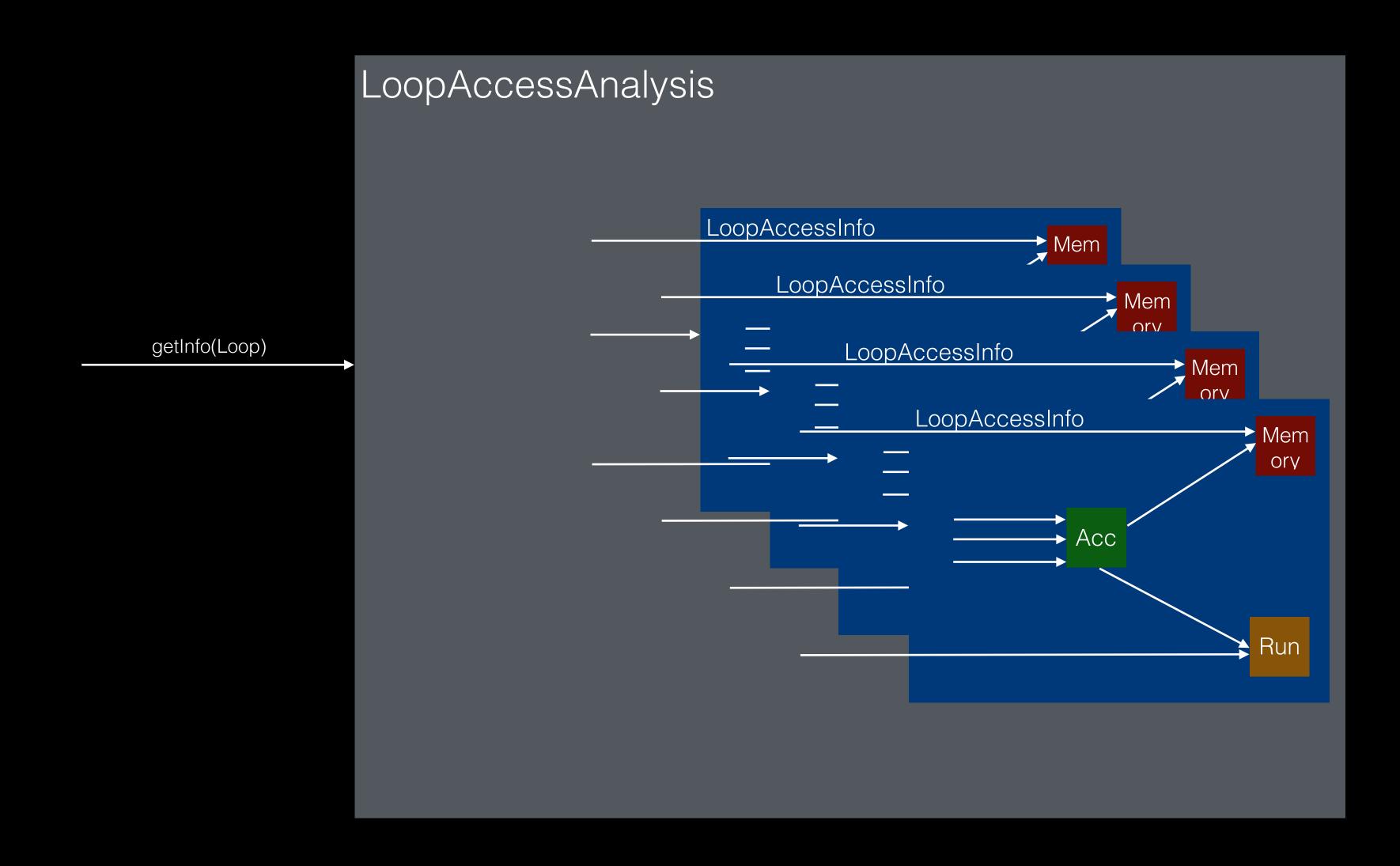
Loop Vectorizer

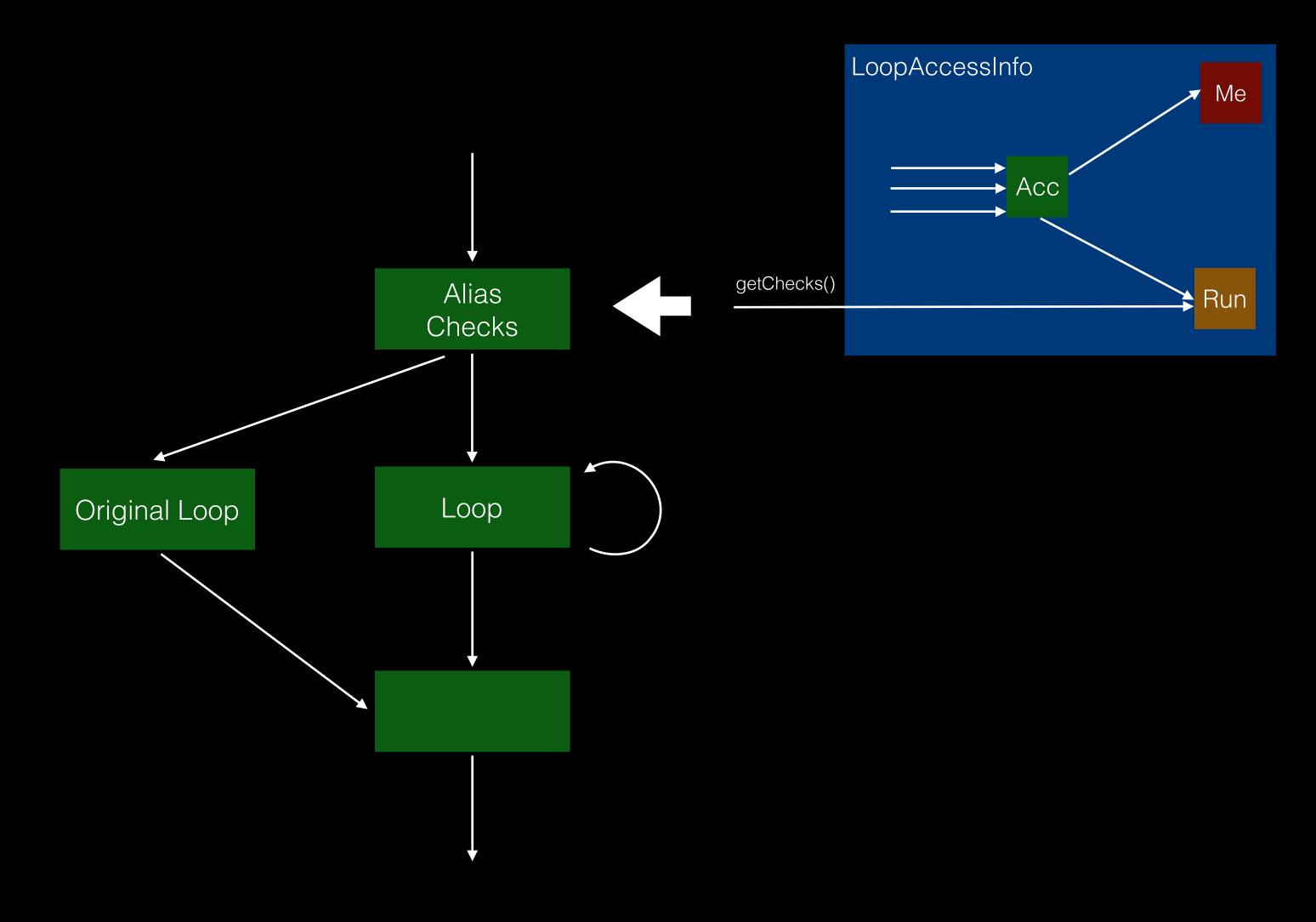


Loop Access Analysis



Loop Access Analysis





- Users:
 - Loop Distribution
 - Loop Load Elimination
 - WIP LICM-based Loop Versioning
- Future work:
 - Run-time trip count check
 - Merge versions into a slow path and a fast path

