

# Developing the Clang Static Analyzer

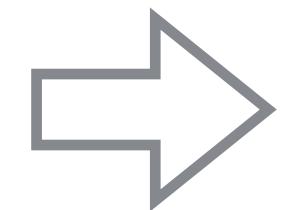
Artem Dergachev, Apple

# Clang Static Analyzer

- Finds bugs at compile time – by inspecting your source code
- Bugs it finds are more sophisticated than warnings or Clang-Tidy

# Clang Static Analyzer

```
1 int foo(int x) {  
2     int y = x;  
3     if (y == 0)  
4         return 24 / x;  
5     return 0;  
6 }
```



```
1 int foo(int x) {  
2     int y = x;  
3     if (y == 0)  
4  
5         1 Assuming 'y' is equal to 0 →  
6             2 ← Taking true branch →  
7                 return 24 / x;  
8  
9         3 ← Division by zero →  
10            return 0;  
11 }
```

# Clang Static Analyzer

- Natural!
  - Mimics normal program execution
  - Easy to understand why it “thinks” there is a bug
  - Takes all source code information into account
  - Explains bugs in terms of the source code

# Clang Static Analyzer

- Natural!
- Researchy!
  - Deals with a lot of open problems
  - People publish articles,  
defend BS/MS/Ph.D. theses on it
  - Fully Open Source – lives in Clang repo

# Clang Static Analyzer

- Natural!
- Researchy!
- Practical!
- Used in industry
- Shipped with IDEs
- Finds bugs in your code before your users do!

# Clang Static Analyzer

- Natural!
- Researchy!
- Practical!
- Extensible!
  - Finds over 50 kinds of bugs!
    - Memory leaks
    - Null dereferences
    - Use-after-free
    - Use-after-move
    - ...
  - “Building a Checker in 24 hours” – LLVM DevMtg 2012  
<https://youtu.be/kdxlsP5QVPw>

# Clang Static Analyzer

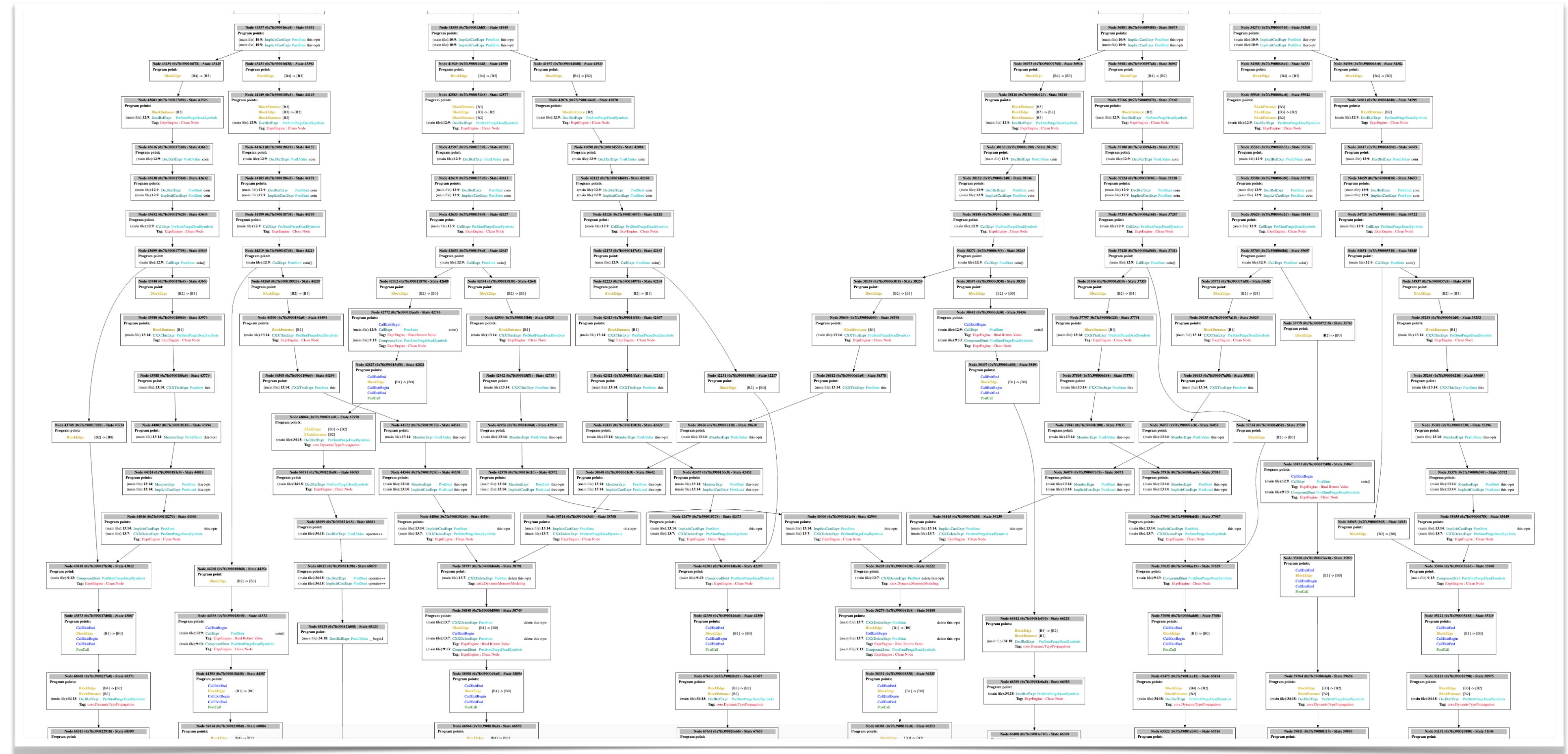
- Natural!  Symbolic Execution
  - Researchy!  Dead, Undead, Zombie and Schrödinger Symbols,  
The Reaper
  - Practical!
  - Extensible!
-  Spooky!

# Clang Static Analyzer

- Natural!
  - Researchy!
  - Practical!
  - Extensible!

🎃 Spooky!

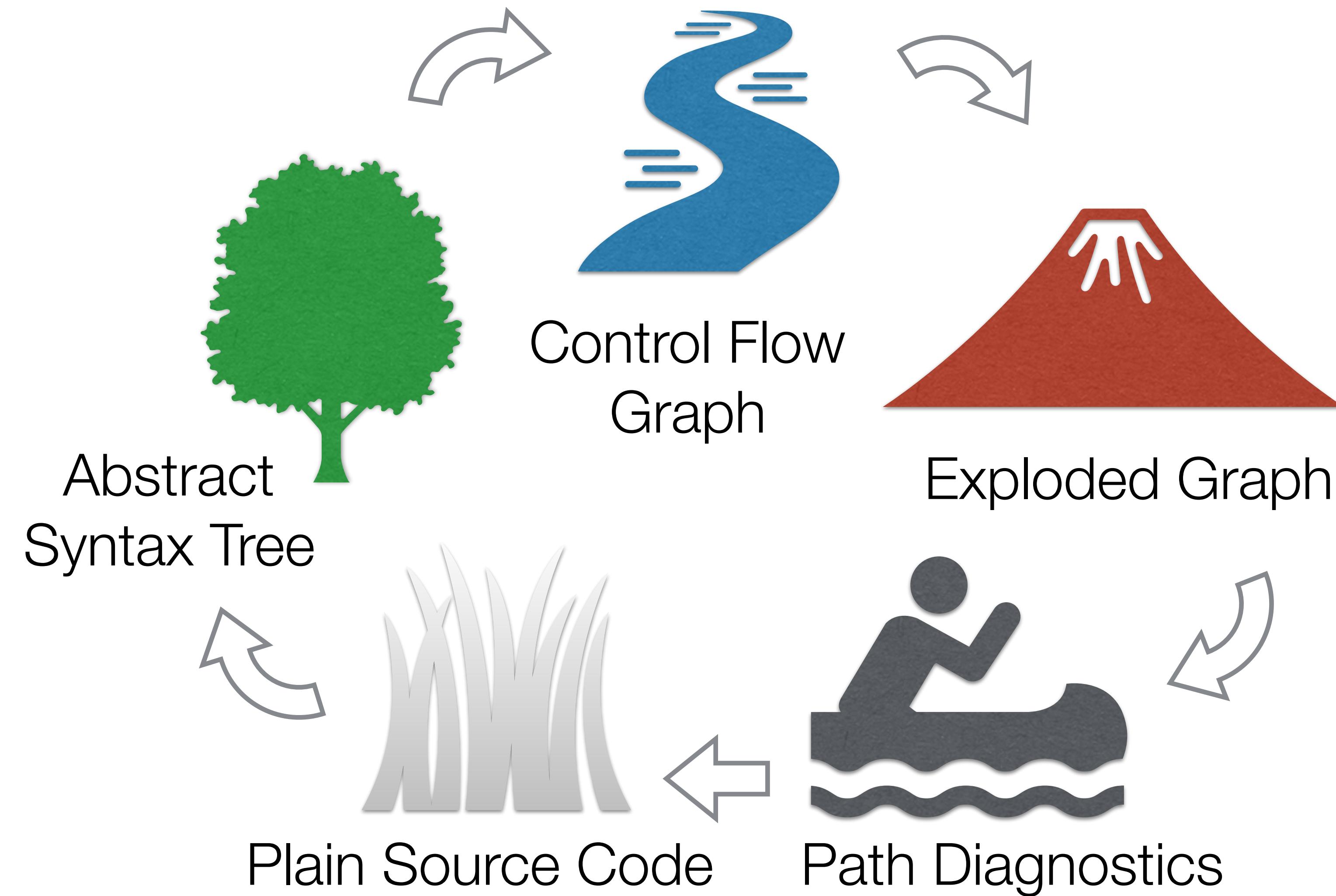
  - Exciting!



# Plan For Today!

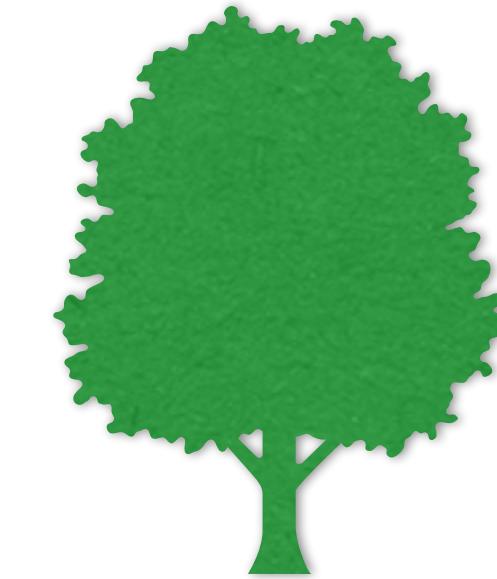
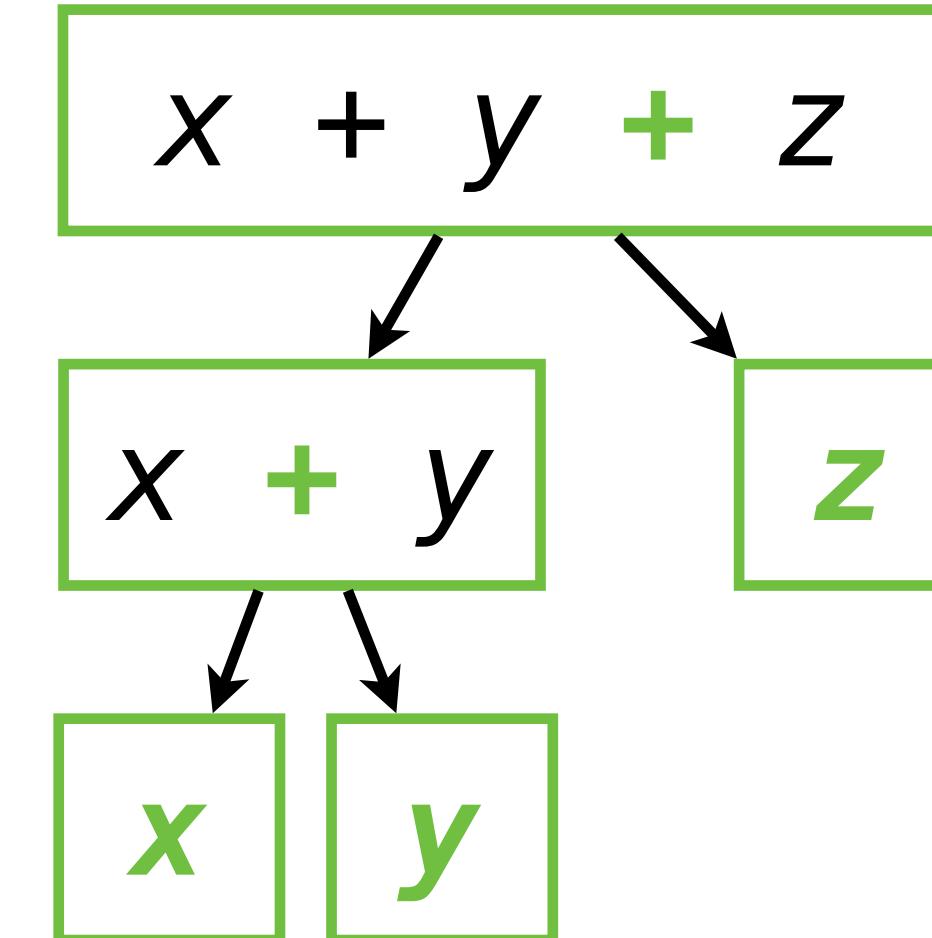
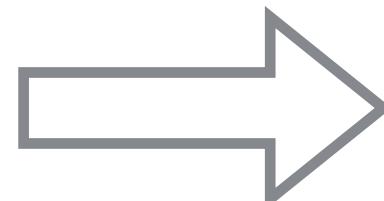
- Algorithms and Data Structures of the Static Analyzer
- How to Fix a Static Analyzer Bug in 24 minutes

# Algorithms and Data Structures of the Static Analyzer



# AST: How Compiler Sees Your Code

- *Nodes*: statements, declarations, types – annotated and cross-referenced
- *Edges*: “is-part-of” relation

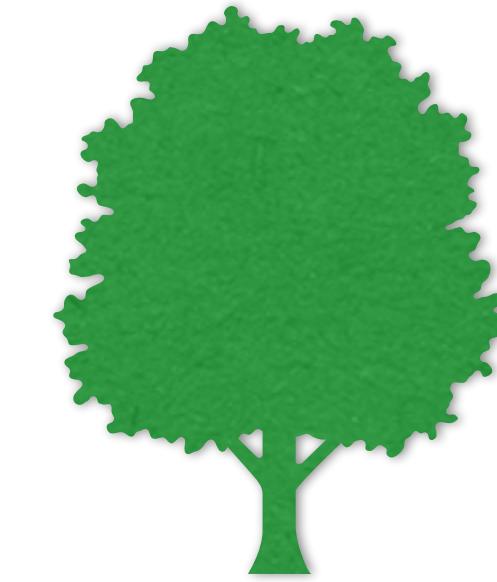
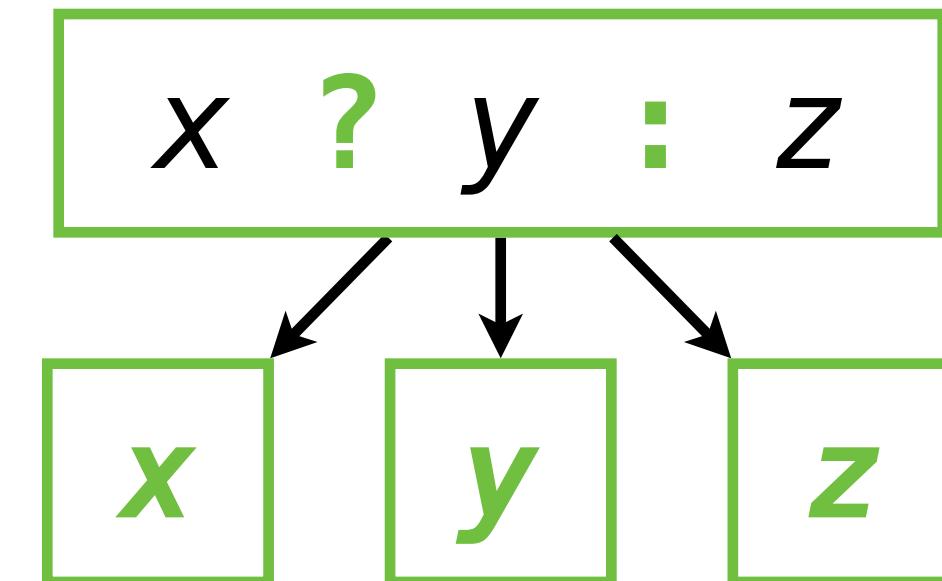
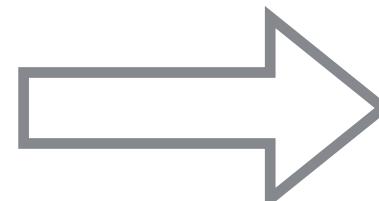
 $x + y + z$ 

# AST: How Compiler Sees Your Code

- *Nodes*: statements, declarations, types – annotated and cross-referenced
- *Edges*: “is-part-of” relation

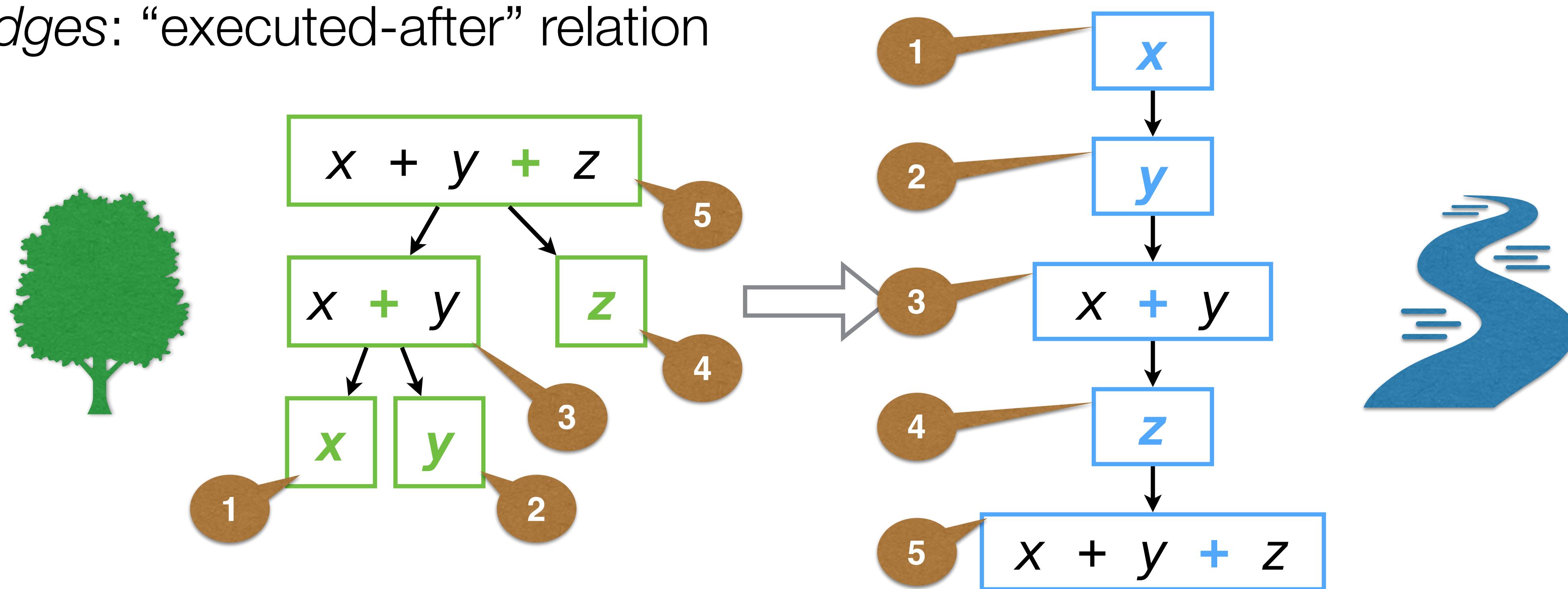


$x ? y : z$



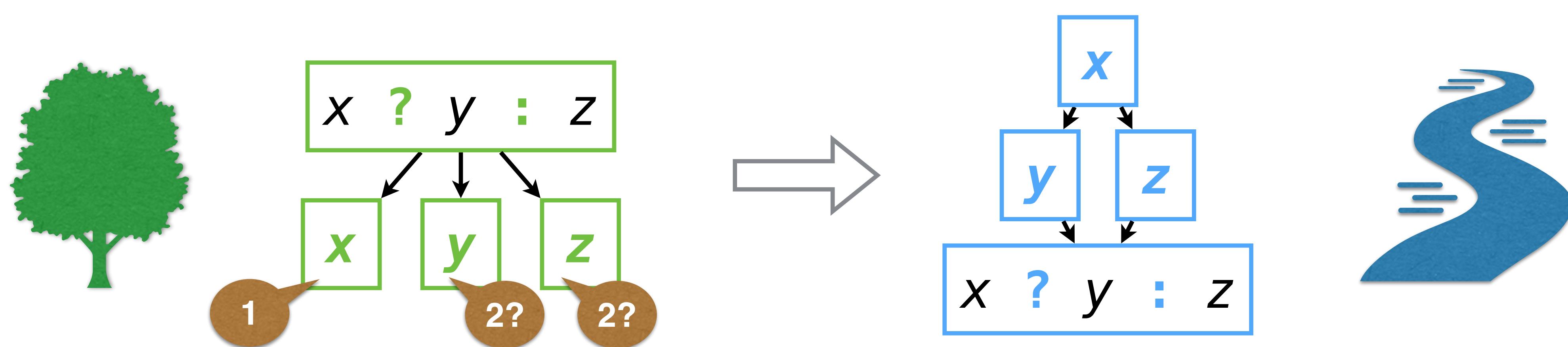
# CFG: Order in which Statements are Executed

- Nodes: usually AST statements
- Edges: “executed-after” relation

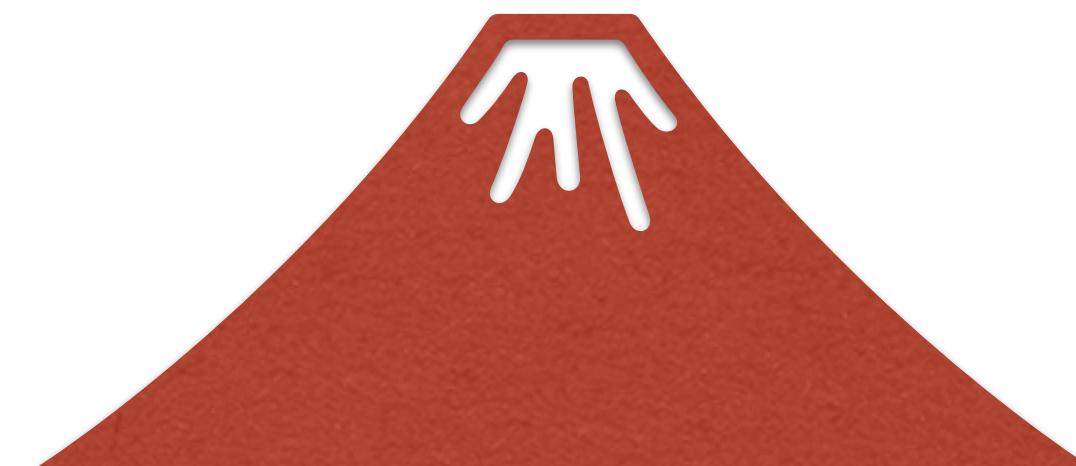
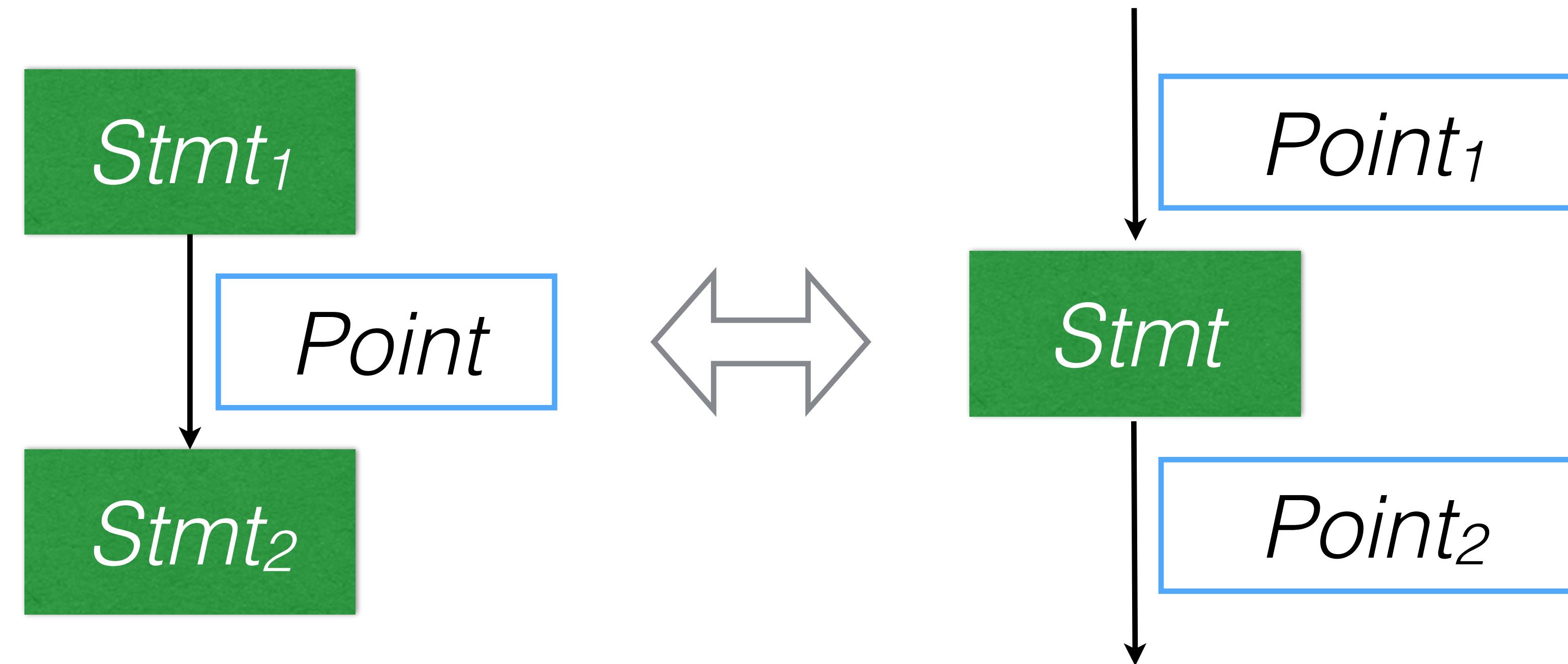


# CFG: Order in which Statements are Executed

- *Nodes*: usually AST statements
- *Edges*: “executed-after” relation

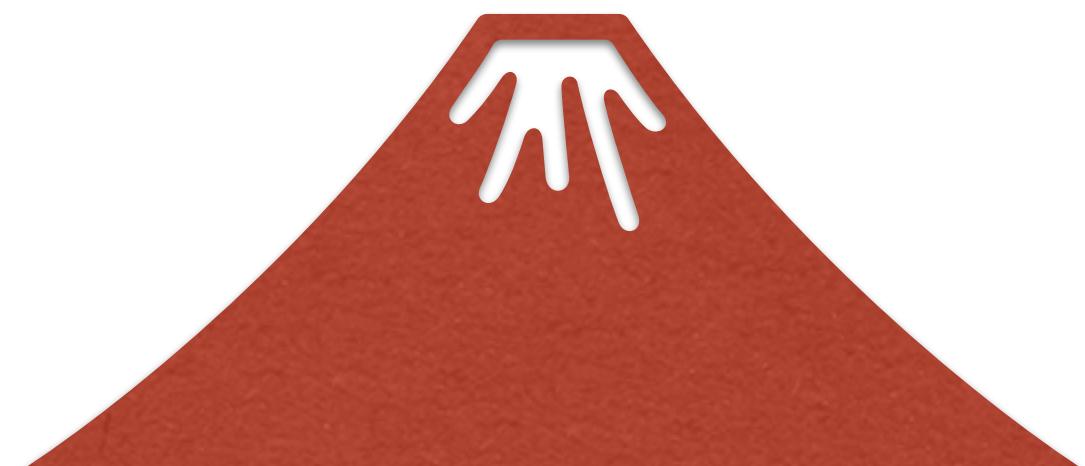


# Program Points

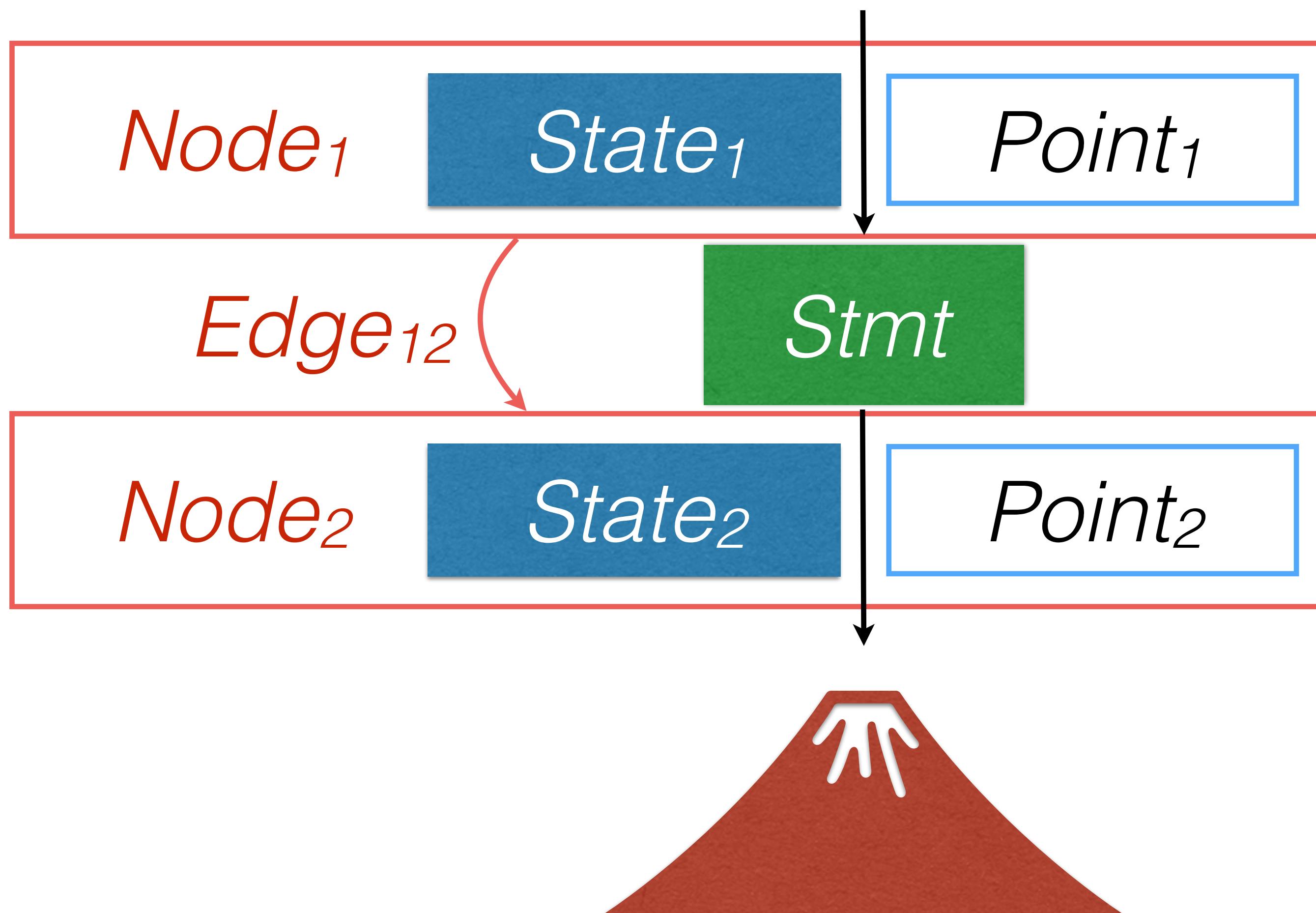


# Exploded Graph: Paths Through CFG

- Nodes:  $(Point, State)$  pairs
  - **Program Point:** A point between statements (usually)
  - **Program State:** A record of effects of statements evaluated so far
- Edges: An edge from  $(Point_1, State_1)$  to  $(Point_2, State_2)$  means that the statement between  $Point_1$  and  $Point_2$  updates  $State_1$  to  $State_2$



# Exploded Graph Edges



# Effects of Assignments: Store

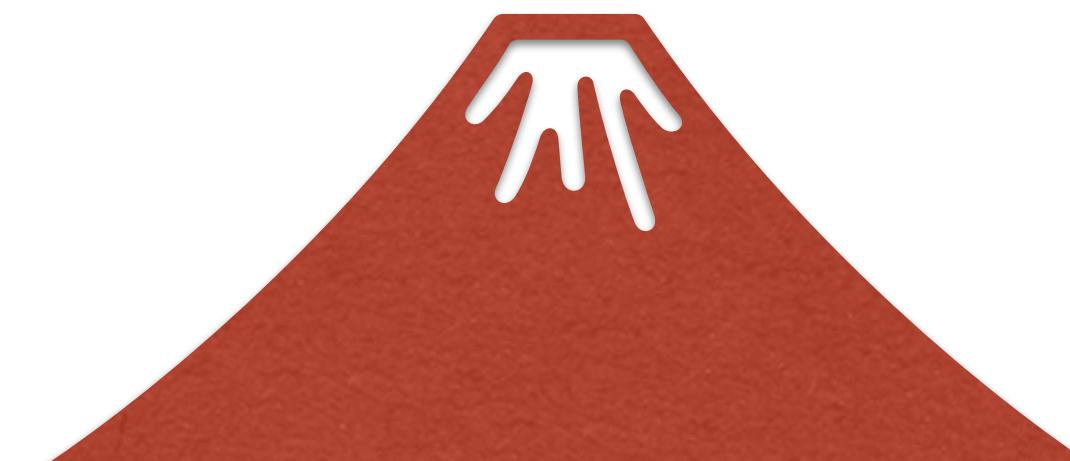
**Program State:**  
Nothing Yet!

**Statement:**

$x = 7$



**Program State:**  
**Store:**  
 $x \rightarrow 7$



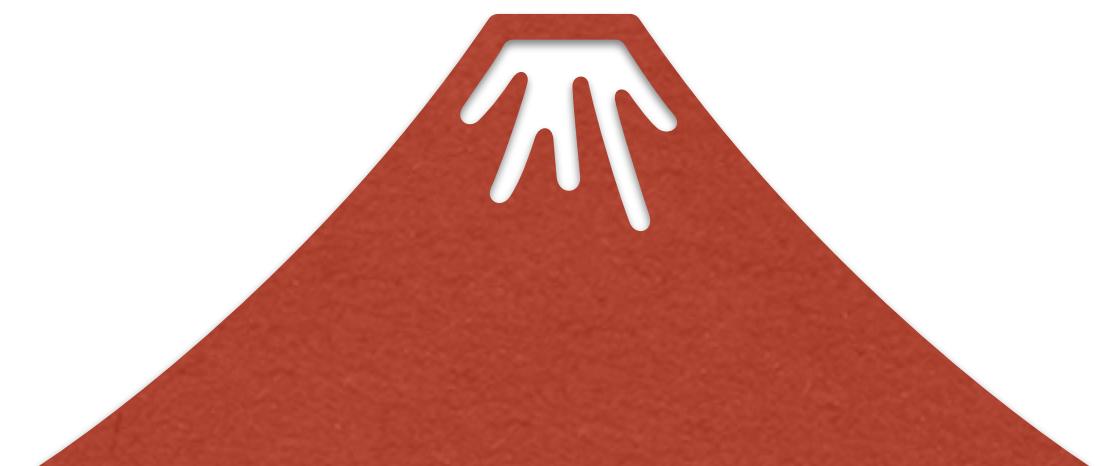
# Values of Expressions: Environment

**Program State:**  
**Store:**  
 $x \rightarrow 7$

**Statement:**  
 $x + 5;$



**Program State:**  
**Store:**  
 $x \rightarrow 7$   
**Exprs:**  
 $x + 5 \rightarrow 12$



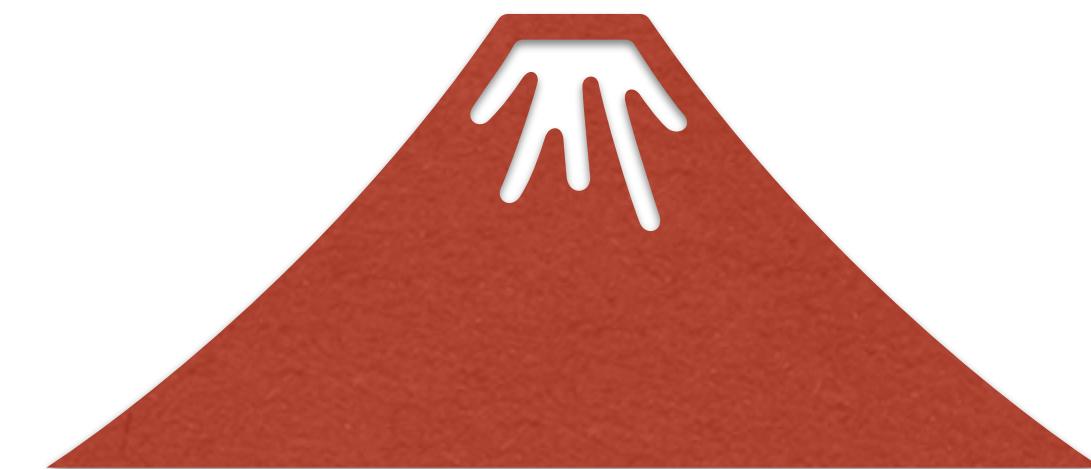
# Focus on One Operation at a Time

**Program State:**  
**Exprs:**  
 $x + 5 \rightarrow 12$

**Statement:**  
 $(x + 5) / 2;$



**Program State:**  
**Exprs:**  
 $(x + 5) / 2 \rightarrow 6$



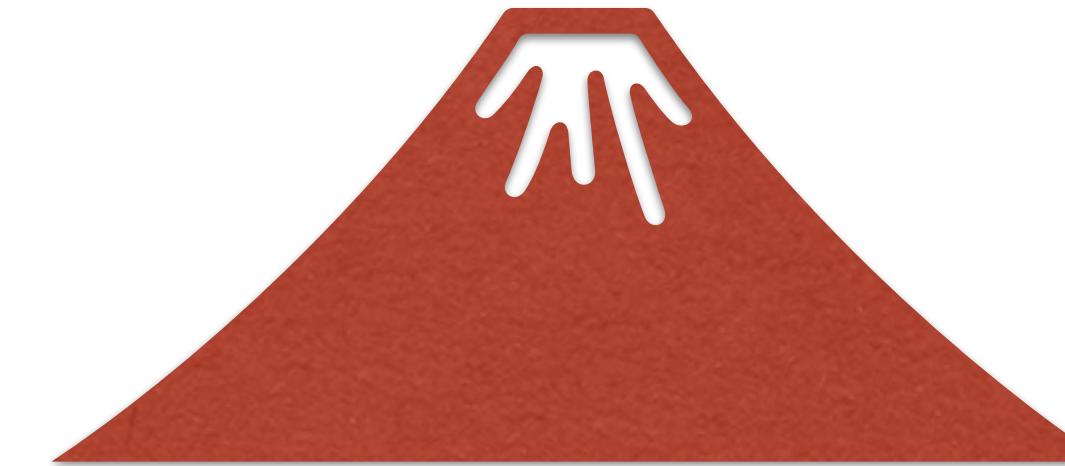
# What If It's Not In The Store?

**Program State:**  
Nothing Yet!

**Statement:**  
 $x;$

**Program State:**  
**Exprs:**  
 $x \rightarrow \text{reg\_\$0<int } x>$

```
// example:  
int foo(int x) {  
    return x;  
}
```

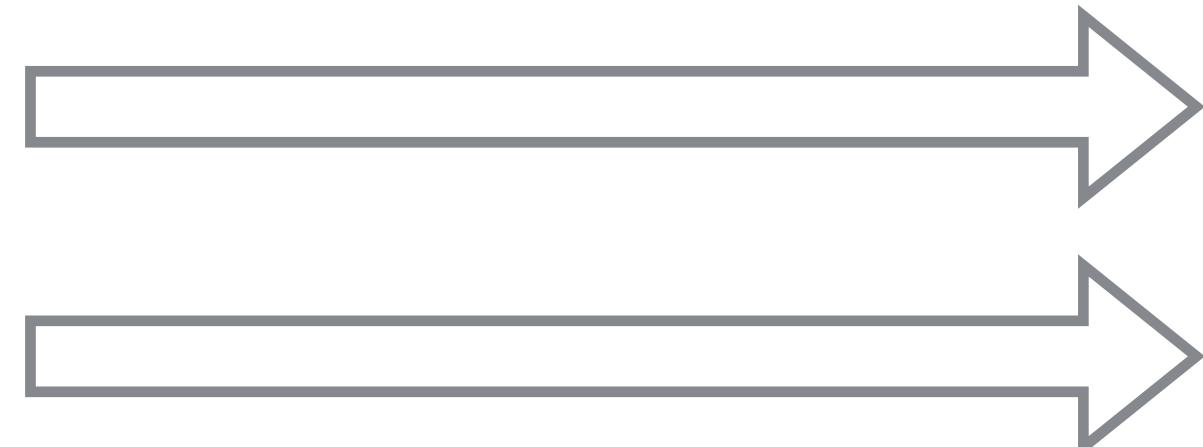


# Effects of Branches: Constraints

Program State:  
Exprs:

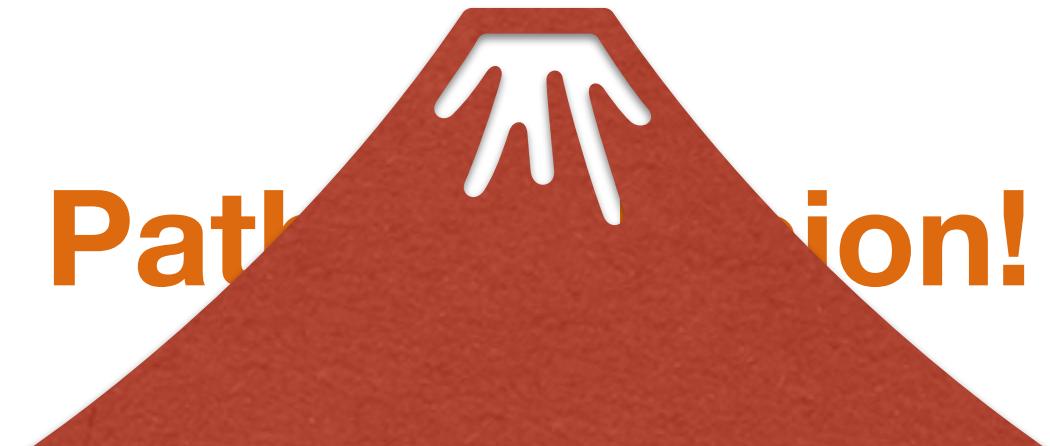
$x \rightarrow \text{reg\_\$0<int } x>$

**Statement:**  
**if** ( $x > 5$ ) ...



**Program State:**  
**Ranges:**  
 $\text{reg\_\$0<int } x> > 5$

**Program State:**  
**Ranges:**  
 $\text{reg\_\$0<int } x> \leq 5$



# Symbolic Execution Recipe

- Just execute the program as you normally would
- Don't know the value? – Denote it with a symbol
- Branch depends on a symbol? – Split up, record constraints
- Don't explore paths on which constraints contradict each other



Demo:  
How to Fix a Static Analyzer Bug  
in 24 minutes!



# Summary!

- Static Analyzer finds bugs by exploring sequences of events that may occur during the execution of the program.
- You can understand and study the internal logic of the static analyzer by looking at exploded graph dumps and setting conditional breaks on individual nodes.
- Sometimes these graphs are huge, so you should use `utils/analyzer/exploded-graph-rewriter.py` with various flags to extract useful information from the dump.
- See [clang-analyzer.llvm.org](http://clang-analyzer.llvm.org) for more information!

