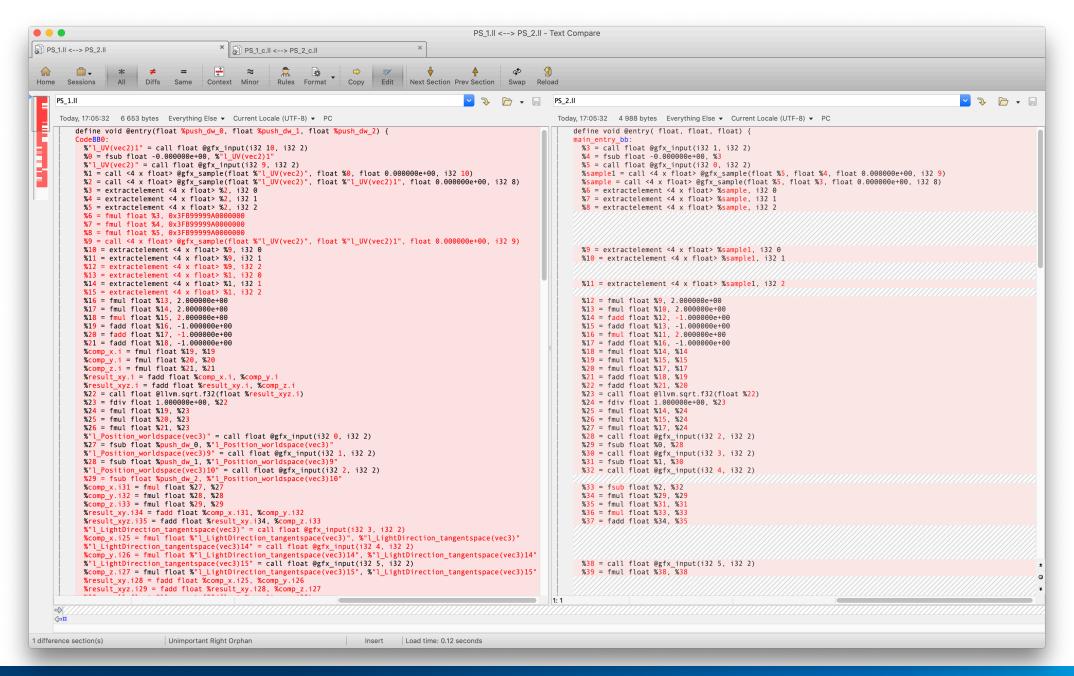


# LLVM-CANON: SHOOTING FOR CLEAR DIFFS

Presenting: Michal Paszkowski

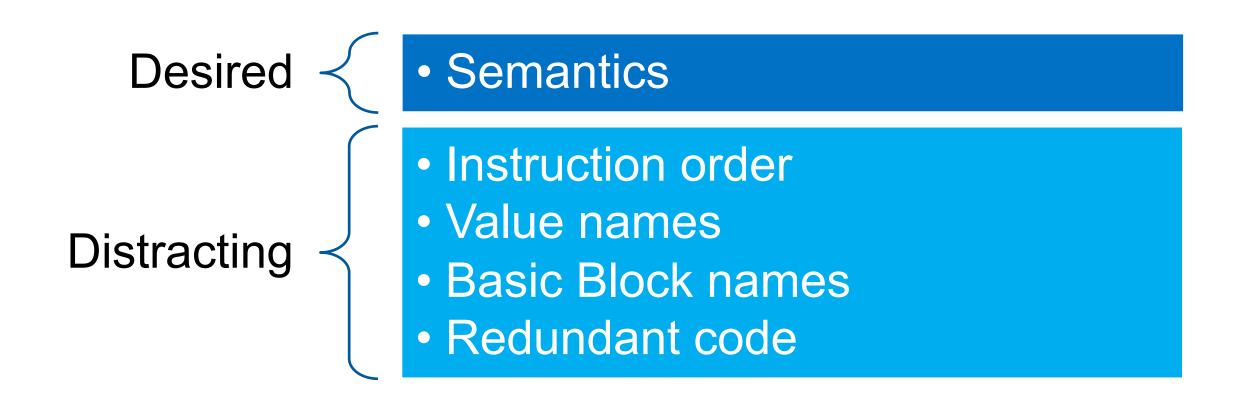
Research: Michal Paszkowski, Radoslaw Drabinski

Special thanks to: Julia Koval





#### It is all about the semantic differences!





#### How could we automate that?

We needed a <u>simple tool</u> that would makes comparing <u>semantic differences</u> between two modules easier.

#### The tool should...

- "Reduce" non-semantic differences
- Process modules independently
- Leverage existing diff tools



#### How could we automate that?

Therefore, the tool should transform a module into a canonical form.

How will that "canonical form" help us?

• Two semantically identical <u>canonicalized</u> modules should show no differences when diffed.

And more importantly...

When the modules are not identical the <u>semantic differences</u> should stand out.



#### How do we arrive at this canonical form?

Let's start with instruction ordering...

Assumptions for comparing an identical module after two different transformations:

- Side-Effects should be roughly the same
- Control-Flow Graphs should be similar



• *def-use* distance reduction

```
%1 = ...
%2 = ...
%T1 = call float @inputV.f32(i32 15, i32 2)
%A = fsub float %T1, %T1
%B = fmul float %A, %T1
%C = fsub float %T1, %B
%X = fadd float %C, %A
store float %C, float addrspace(479623)* %1
store float %X, float addrspace(283111)* %2
ret void
```



- def-use distance reduction
  - 1. Collect instructions with side-effects and "ret" instructions

```
%1 = ...
%2 = ...
%T1 = call float @inputV.f32(i32 15, i32 2)
%A = fsub float %T1, %T1
%B = fmul float %A, %T1
%C = fsub float %T1, %B
%X = fadd float %C, %A
store float %C, float addrspace(479623)* %1
store float %X, float addrspace(283111)* %2
ret void
```



- def-use distance reduction
  - 1. Collect instructions with side-effects and "ret" instructions
  - 2. Walk the instructions with sideeffects (top-down) and on each instruction their operands (left-right)

```
%1 = ...
%2 = ...
%T1 = call float @inputV.f32(i32 15, i32 2)
%A = fsub float %T1, %T1
%B = fmul float %A, %T1
%C = fsub float %T1, %B
%X = fadd float %C, %A
store float %C, float addrspace(479623)* %1
store float %X, float addrspace(283111)* %2
ret void
```



- def-use distance reduction
  - 1. Collect instructions with side-effects and "ret" instructions
  - 2. Walk the instructions with sideeffects (top-down) and on each instruction their operands (left-right)
  - 3. For each operand, bring their definition as close as possible to the using instruction

```
%1 = ...
%2 = ...
%T1 = call float @inputV.f32(i32 15, i32 2)
%A = fsub float %T1, %T1
%B = fmul float %A, %T1
■%C = fsub float %T1, %B
%X = fadd float %C, %A
store float %C, float addrspace(479623)* %1
store float %X, float addrspace(283111)* %2
ret void
```



```
%T1 = call float @inputV.f32(i32 15, i32 2)
%A = fsub float %T1, %T1
%B = fmul float %A, %T1
%C = fsub float %T1, %B
%X = fadd float %C, %A
store float %C, float addrspace(479623)* %1
store float %X, float addrspace(283111)* %2
ret void
```



```
%T1 = call float @inputV.f32(i32 15, i32 2)
%A = fsub float %T1, %T1

%B = fmul float %A, %T1

%C = fsub float %T1, %B
%X = fadd float %C, %A
```

Select the 1<sup>st</sup> side-effecting instruction.

Don't move it, otherwise semantics may not be preserved!

```
store float %C, float addrspace(479623)* %1
store float %X, float addrspace(283111)* %2
ret void
```



```
%T1 = call float @inputV.f32(i32 15, i32 2)
%A = fsub float %T1, %T1

%B = fmul float %A, %T1

%C = fsub float %T1, %B
%X = fadd float %C, %A
```

Take the 1<sup>st</sup> operand of the side-effecting instruction.

```
store float %C, float addrspace(479623)* %1
store float %X, float addrspace(283111)* %2
ret void
```



```
%T1 = call float @inputV.f32(i32 15, i32 2)
%A = fsub float %T1, %T1
%B = fmul float %A, %T1
%C = fsub float %T1, %B
%X = fadd float %C, %A
%C = fsub float %T1,
store float %C, float addrspace(479623)* %1
store float %X, float addrspace(283111)* %2
ret void
```

Move it closer to the user.

Def-Use sequence may be temporarily broken.



```
%T1 = call float @inputV.f32(i32 15, i32 2)
%A = fsub float %T1, %T1
%B = fmul float %A, %T1
%X = fadd float %C, %A
%C = fsub float %T1, %B
store float %C, float addrspace(479623)* %1
store float %X, float addrspace(283111)* %2
ret void
```

Select the 1<sup>st</sup> operand of the moved instruction.



```
%T1 = call float @inputV.f32(i32 15, i32 2)
%B = fmul float %A, <a href="#">%T1</a>
%X = fadd float %C, %A
%T1 = call float @inputV.f32(i32 15, i32 2)
%C = fsub float %T1, %B
store float %C, float addrspace(479623)* %1
store float %X, float addrspace(283111)* %2
ret void
```

Move it closer to the user.

Def-Use sequence may be temporarily broken.



```
%A = fsub float %T1, %T1
%B = fmul float %A, %T1
%X = fadd float %C, %A
%T1 = call float @inputV.f32(i32 15, i32 2)
%C = fsub float %T1, %B
store float %C, float addrspace(479623)* %1
store float %X, float addrspace(283111)* %2
ret void
```

Select the 2<sup>nd</sup> operand of the previously moved instruction.



```
%A = fsub float %T1, %T1
%B = fmul float %A, %T1
%X = fadd float %C, %A
%T1 = call float @inputV.f32(is2 15, i32 2)
%B = fmul float %A, %T1
%C = fsub float %T1, %B
store float %C, float addrspace(479623)* %1
store float %X, float addrspace(283111)* %2
ret void
```

Move it closer to the user. Def-Use sequence is being repaired.



```
%X = fadd float %C, %A
%T1 = call float @inputV.f32(i32 15, i32 2)
%B = fmul float %A, %T1
%C = fsub float %T1, %B
store float %C, float addrspace(479623)* %1
store float %X, float addrspace(283111)* %2
ret void
```

%A = fsub float %T1, %T1

Select the 1<sup>st</sup> operand of the moved instruction.



```
%A = fsub float %T1, %T1
%X = fadd float %C, %A
%T1 = call float @inputV.f32
%A = fsub float %T1, %T1
%B = fmul float %A, %T1
%C = fsub float %T1, %B
store float %C, float addrspace(479623)* %1
store float %X, float addrspace(283111)* %2
ret void
```

Move it closer to the user.

Def-Use sequence is being repaired.



```
%X = fadd float %C, %A
%T1 = call float @inputV.f32(i32 15, i32 2)
%A = fsub float %T1, %T1
%B = fmul float %A, %T1
%C = fsub float %T1, %B
store float %C, float addrspace(479623)* %1
store float %X, float addrspace(283111)* %2
ret void
```

Select the 1st operand of the moved instruction.

Don't move it! %T1 has been already moved!

Select the 2nd operand of the moved instruction.

Don't move it! %T1 has been already moved!



```
%X = fadd float %C, %A
%T1 = call float @inputV.f32(i32 15, i32 2)
%A = fsub float %T1, %T1
%B = fmul float %A, %T1
%C = fsub float %T1, %B
store float %C, float addrspace(479623)* %1
store float %X, float addrspace(283111)* %2
ret void
```

Select the 2<sup>nd</sup> operand of the previously moved instruction. Don't move it, %T1 has been moved before.

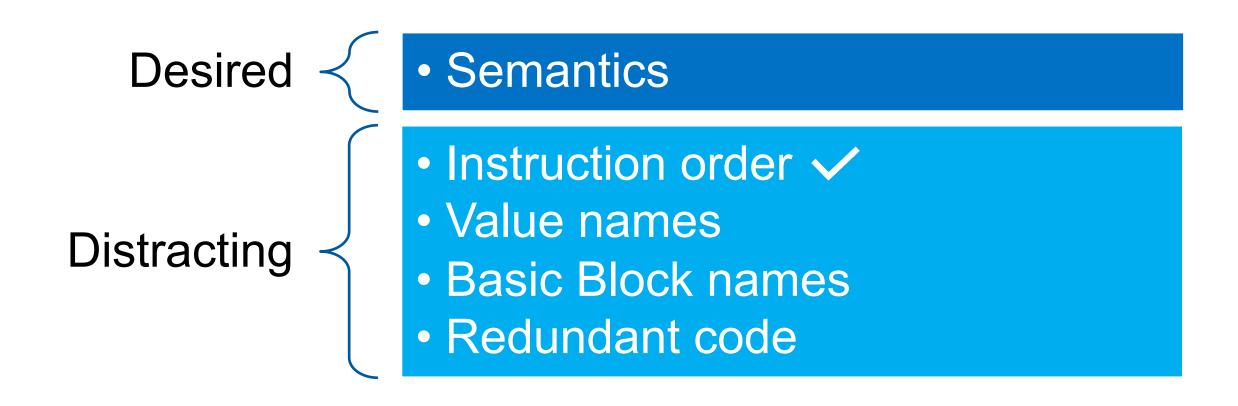


```
%X = fadd float %C, %A
%T1 = call float @inputV.f32(i32 15, i32 2)
%A = fsub float %T1, %T1
%B = fmul float %A, %T1
%C = fsub float %T1, %B
store float %C, float addrspace(479623)* %1
store float %X, float addrspace(283111)* %2
ret void
```

We repeat the process for all operands in all side-effecting instructions.



#### How do we arrive at this canonical form?





#### Naming instructions: Linear

- Numbers all instructions top-down <u>after reordering</u> vn
- We were hoping that maybe the reordering mechanism could be used as a 'seed' for instruction naming

```
%T1 = @inputV.f32(i32 15, i32 2)

%A = fsub float %T1, %T1

%B = fmul float %A, %T1

%C = fsub float %T1, %T1

%X = fadd float %C, %A

store float %C, float addrspace(479623)* %1

store float %X, float addrspace(283111)* %2

ret void
```

```
%v0 = @inputV.f32(i32 15, i32 2)

%v1 = fsub float %v0, %v0

%v2 = fmul float %v1, %v0

%v3 = fsub float %v0, %v0

%v4 = fadd float %v3, %v1

store float %v3, float addrspace(479623)* %a3
store float %v4, float addrspace(283111)* %a4
ret void
```



## Naming instructions: Linear

- Numbers all instructions top-down after reordering vn
- We were hoping that maybe the reordering mechanism could be use
   'seed' for instruction naming

```
%T1 = @inputV.f32(i32 15, i32 2)
%A = fsub float %T1, %T1
%B = fmul float %A, %T1
%C = fsub float %T1, %T1
%X = fadd float %C, %A
store float %C, float addrspace(479623)* %1
store float %X, float addrspace(283111)* %2
ret void
```

```
%v0 = @inputV.f32(i32 15, i32 2)

%v1 = fsub float %v0, %v0

%v2 = fmul float %v1, %v0

%v3 = fsub float %v0, %v0

%v4 = fadd float %v3, %v1

store float %v3, float addrspace(479623)* %a3
store float %v4, float addrspace(283111)* %a4
ret void
```



# Naming instructions: "Graph naming"

#### Two types of instructions:

#### 1. Initial instructions

- Instructions with only immediate operands
- Numbered according to positions of outputs using that instruction after sorting

#### 2. Regular instructions

 "Graph naming": Differences in defs are reflected in uses

```
%v0| = call float @gfx_input(i32 9, i32 2)
%v1| = call float @gfx_input(i32 10, i32 2)
%"op(v0, v1)" = fmul float %v0, %v1
%"op(op(v0, v1),v0)" = fadd float %"op(v0, v1)", %v0
```



# Naming instructions: "Graph nam

#### Two types of instructions:

#### 1. Initial instructions

- Instructions with only immediate operands
- Numbered according to positions of outputs using that instruction after sorting

#### 2. Regular instructions

 "Graph naming": Differences in defs are reflected in uses

```
""" op(v0, v1),v0)" = fadd float %"op(v0, v1)", %v0
""" op(v0, v1),v0)" = fadd float %"op(v0, v1)", %v0
```

# Naming instructions: "Graph nam

# ALMOS

#### Two types of instructions:

#### 1. Initial instructions

- Instructions with only immediate operands
- Numbered according to positions of outputs using that instruction after sorting

#### 2. Regular instructions

 "Graph naming": Differences in defs are reflected in uses Instructions with different opcodes but same users got the same names.

Extremely long names!

Differences in defs should be reflected only in outputs



1. Initial instructions (those with only immediate operands)

$$%"vli2345Foo'(2, 5)" = ...$$
hash callee operands

- Hash calculated considering instruction's opcode and the "output footprint"
- Called function name only included in case of a CallInst
- Immediate operands list (sorted in case of commutative instructions)



#### 2. Regular instructions

- Hash calculated considering instruction's and its operands' opcodes
- Called function name only included in case of a CallInst
- Short operand names



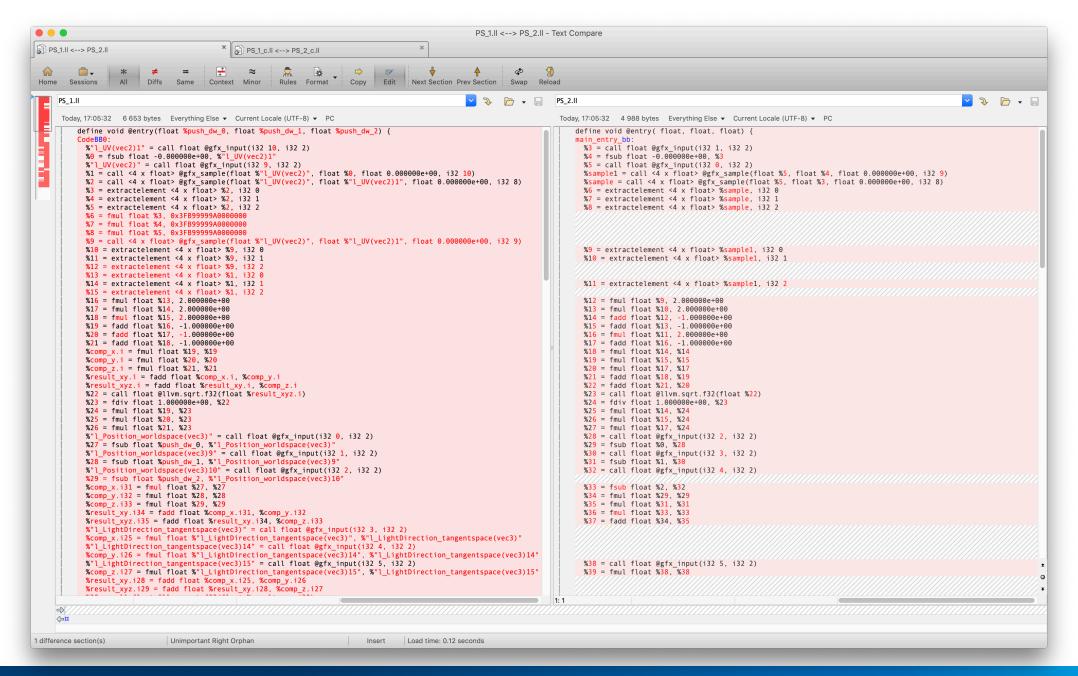
3. Output instructions (instructions with side-effects and their relative operands)

- Same as regular instructions, but...
- recursively generated long operand list is kept, so...
- by just looking at an output we see what impacts its semantics in the diff.

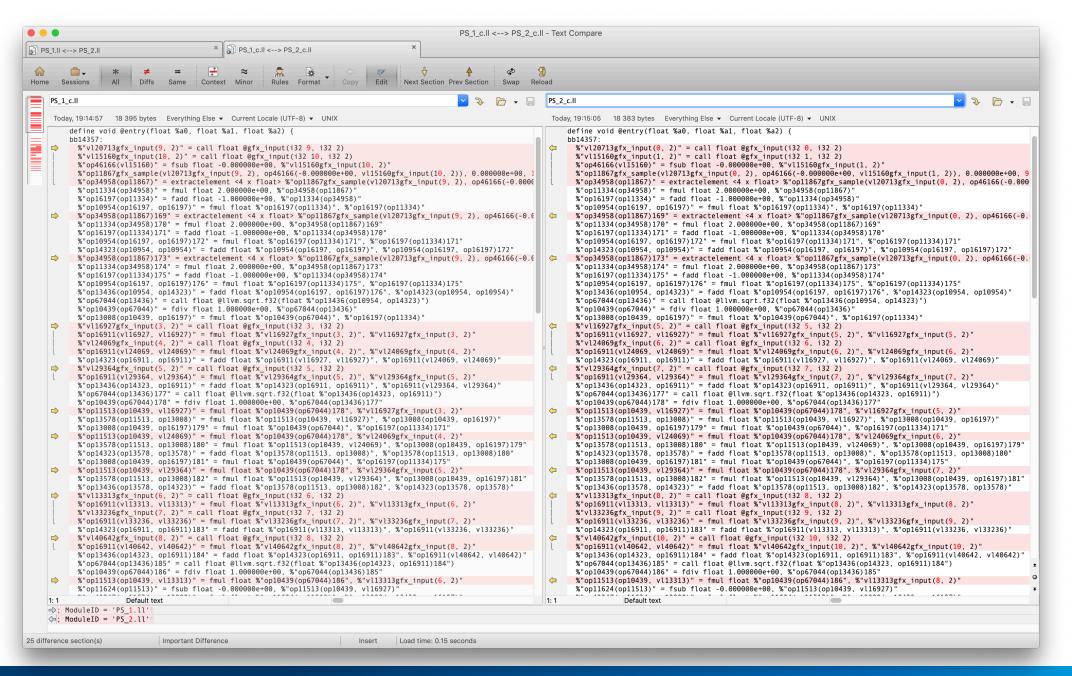




**Output instruction** 









# Other little things...

- Naming basic blocks
- Numbering function arguments
- Sorting values in PHI nodes



#### Ilvm-canon vs Ilvm-diff

Why not integrate with IIvm-diff?

- We wanted to use this tool to also spot differences in just one file.
- We wanted to leverage existing diff tools.

However, Ilvm-canon could be used as a prepass before using Ilvm-diff:

```
bin — -bash

[Michas-MBP:bin michalpaszkowski$ ./llvm-diff test1.ll test2.ll 2>&1 >/dev/null | wc -l
9

Michas-MBP:bin michalpaszkowski$ ./llvm-diff test1_c.ll test2_c.ll 2>&1 >/dev/null | wc -l
0

Michas-MBP:bin michalpaszkowski$ |
```



#### Ilvm-canon vs Ilvm-diff

```
bin — -bash

[Michas-MBP:bin michalpaszkowski$ ./llvm-diff test1.ll test2.ll 2>&1 >/dev/null | wc -l 9

Michas-MBP:bin michalpaszkowski$ ./llvm-diff test1_c.ll test2_c.ll 2>&1 >/dev/null | wc -l 0

Michas-MBP:bin michalpaszkowski$ |
```



## Special thanks

Thanks to Puyan Lotfi for his talk on MIR-Canon during 2018 EuroLLVM and a head start on canonicalization techniques.

Special thanks to Radoslaw Drabinski and Julia Koval.

I would like to also thank the LLVM community for excellent code review and coming to this talk.



# Q&A