CS 61C Fall 2015 Guerrilla Section 2: MIPS

Question 1 (SP14 Section 4): MIPS Warm-Up.

- 1. How should \$sp be used? When do we add or subtract from \$sp?
- 2. Which registers need to be saved or restored before using jr to return from a function? 3. Which registers need to be saved before using jal?
- 4. How do we pass arguments into functions?
- 5. What do we do if there are more than four arguments to a function?
- 6. How are values returned by functions?

Question 2: It's a bloody MIPStery. . .

a) Write the MAL MIPS function reverse_str(char *string, int string_length), that can reverse strings (with an even length) in-place. The MIPS should be non-delayed branch, and you will probably not use all the lines. In your solution, register \$a0 should signify the parameter char *string and register \$a1 should signify the parameter int string length.

| reverse_str:b | peq \$a1 \$0 done |
|---------------|--|
| | |
| | |
| done:jr \$ra | j reverse_str |
| | de below, using at most two TAL MIPS instructions, so that the function contains an R-type instruction and true otherwise. |
| NotRType: | |
| | |

Question 3: "free at last, thank gosh we are free at last..."

We wish to free a linked list of strings (example below) whose nodes are made up of this struct. Complete the code below; we have started you off with some filled in. You may use fewer lines, but do not add any.

```
// Assume compiler packs tightly
struct node {
        char *string;
        struct node *next;
};

void FreeLL(struct node *ptr) {
        if (ptr == NULL) return;
        else {
            FreeLL(ptr->next);
            free(ptr->string);
            free(ptr);
        }
}
```

| FreeLL:beq | ,, NULL_CASE |
|------------|-----------------|
| | |
| | |
| | |
| | jal FreeLL |
| | lw \$a0 0(\$sp) |
| | |
| | |
| | jal free |
| | |
| | |

NULL CASE:jr \$ra

head

I love CS61C midterms!

Question 4 (su13 q2): MIPStifying (9 points, 20 minutes)

Answer the questions below about the following MIPS function. Answer each part separately, assuming each time that mystery() has not been called yet.

mystery: 1 andi \$a0, \$a0, 3 2 ori \$t0, \$0, 1 3 sll \$t0, \$t0, 6 Lb11: beq \$a0, \$0, Lb12 4 5 sll \$t0, \$t0, 5 addi \$a0, \$a0, -1 6 7 j Lb11 Lb12: la \$s0, Lb13 8 8 lw \$s1, 0(\$s0) 9 add \$s1, \$s1, \$t0 \$s1, 0(\$s0) 10 SW Lb13: add \$v0, \$0, \$0 11 12 jr \$ra

- A. Which instruction (number) gets modified in the above function?
- B. Write an equivalent arithmetic (not logical) C expression to instruction 1. a0 = _____
- C. Which instruction field gets modified when mystery is called with \$a0 = 3?
- D. How many times can mystery(2) be called before the behavior of mystery() changes?
- E. How many times can mystery(0) be called before the behavior of mystery() changes?
- F. A program calls mystery with the following sequence of arguments: 0, 1, 2, 3, 4, 5. What MIPS instruction gets stored in memory?

Question 5 (SP07 Final M1): Do You Remember?

Decode the binary numbers into MIPS instructions with proper register names (\$s0, \$t0, etc.).If there are any memory addresses, represent them in hex.

| Address | 32-bit Binary Instruction | Type (R, I, J) | MIPS Instruction w/args |
|------------|--|-------------------|-------------------------|
| 0xAFFFFF8 | 0000 0001 0000 1000 0100 0000 0010 0110 | | |
| 0xAFFFFFC | 0001 0100 0000 1000 1111 1111 1111 1110 | | |
| 0xB0000000 | 0000 1000 0000 0000 0000 0000 0000 0001 | | |
| 0xB0000004 | whatever | whatever | ori \$v0, \$0, 0x61C |
| 0xB0000008 | whatever | whatever | jr \$ra |

You can replace the first instruction with a *new* instruction and save 2 clock cycles on a single-cycle non-delayed branch MIPS machine. What is it (in MIPS)? *Careful!*
