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1. (A) Yes. $f(n) = n$ $\log_2 f(n) = \log_2 n$ $\log_2 g(n) = \log_2 (cn)$

$\log_2 c + \log_2 n$ always greater than $\log_2 n$ for any n , because c is always greater than 1 and $\log_2 c$ always greater than 0
 $\therefore \log_2 f(n)$ is $O(\log_2 g(n))$

(B) No. example: $f(n) = n$, $g(n) = 2n$, $2^{f(n)} = 2^n$ $O(g(n)^2) = O(4n^2)$
when $n=9$, $2^{f(n)} = 2^9 = 512$, $4n^2 = 4 \times 81 = 324$, $2^{f(n)} \geq g(n)^2$ when $n \geq 9$

(C) Yes, for example, $f(n) = n$, $f(n)^2 = n^2$, $g(n) = cn$ and $c > 1$
 $g(n)^2 = c^2 n^2$ $c^2 > c$, for any n , $c^2 n^2 > n^2$, so $f(n)^2$ is $O(g(n)^2)$

2. solution 1:

$$\lim_{n \rightarrow \infty} \frac{n \sqrt{g(n)}}{\sqrt{n}} = \lim_{n \rightarrow \infty} \sqrt{\frac{n^2 \cdot g(n)}{n}} = \lim_{n \rightarrow \infty} \sqrt{n \cdot g(n)} = \infty$$

$\therefore n \sqrt{g(n)}$ grows faster than \sqrt{n}

$$\lim_{n \rightarrow \infty} \frac{\sqrt{n}}{(g(n))^2} = \lim_{n \rightarrow \infty} \frac{\frac{1}{2} \cdot n^{-\frac{1}{2}}}{2 \cdot g(n) \cdot \frac{1}{n \cdot g(n)}} = \lim_{n \rightarrow \infty} \frac{n \cdot g(n) \cdot n^{-\frac{1}{2}}}{4 (g(n))} = \lim_{n \rightarrow \infty} \frac{\sqrt{n}}{g(n)} = \infty$$

$\therefore \sqrt{n}$ grows faster than $(g(n))^2$

$2^{\sqrt{g(n)}}$ is an exponential expression, it grows the fastest among the 4 expressions

\therefore the final order is $2^{\sqrt{g(n)}} > n \sqrt{g(n)} > \sqrt{n} > (g(n))^2$

Solution 2: let $\lg n = a \Rightarrow n = a^{10}$

$$2^{\sqrt{g(n)}} = 2^a$$

\therefore growth rate: $2^a > a^{10.5} > a^5 > a^2$

$$n \sqrt{g(n)} = a^{10} \cdot a^5 = a^{10.5}$$

$$\sqrt{n} = a^5$$

$$2^{\sqrt{g(n)}} > n \sqrt{g(n)} > \sqrt{n} > (g(n))^2$$

Time complexity
3. $f(n) = k$

sing-song(k):

index = 1

song = empty line

while index $\leq k$:

if index ≤ 4 :

concatenate a new line to song

print song

print a new line

else:

insert a new line to the third line in song

print song

print a new line

index++