

COSC2473

Digital Logic

Logic Gates & Bit Masking
Boolean Algebra



Digital Logic

- Binary can be considered as truth values:
0 = False, 1 = True
- Boolean Algebra: the laws of algebra for truth values
 - consists of operations on truth values, the result of which is another truth value
 - fundamental for the design of logic circuits and for tests in computer programming
- CPU hardware consists of circuits built from logic gates
 - logic gates perform Boolean Algebra

Boolean Operators in Python

- Python does not have a Boolean data type directly.
 - There is variable that you can declare as Boolean, but there is a Boolean value type, and a function `bool()` that returns it.
- A Boolean value type is True or False (not true or TRUE)
 - Any value can be converted to Boolean
 - `Bool("hello")=True`, `bool(3)=True`, `bool(0)=False`, `bool("")=False` \overline{x}
- In Java and Python, Boolean operations come in:
 - Real Boolean operators with True/False values
 - These are constructions within the language (e.g. Objects)
 - Bitwise Boolean operators
 - these operate on the bits of (usually) integers
- The bitwise operators are generally more useful.

We will show only bitwise operators from now on.

Bitwise operators

- The bitwise operators work the same in Java and Python

<u>Notation</u>	<u>Operation</u>
$\sim a$	1's complement of a
$a \& b$	AND
$a b$	OR
$a \wedge b$	XOR
$a \ll n$	Left shift by n bits
$a \gg n$	Right shift by n bits

- aa

NOT Operator

- Boolean Operator: NOT
 - NOT returns the opposite value
 - i.e. True becomes False, False becomes True

NOT	-	(written as /x or -x or sometimes !x or \overline{x})
0		1
1		0

Java:

```
int x, a=1;           // a = 0x01, 0000 00012
x = ~a;               // x = 0xFE, 1111 11102 = 254
```

Python

```
a = 0
~a = -1               # since '11111...11' = -1
```

AND Operator

- AND is true when ALL inputs are true
 - written as: $x.y$ or $x \wedge y$ or xy

AND	.
0 . 0	0
0 . 1	0
1 . 0	0
1 . 1	1

Java:

```
int x, a = 0x18, b = 0x11;  
x = a & b;           // a = 0001 10002  
                     // b = 0001 00012  
                     // x = 0001 00002 = 0x10 = 16
```

OR Operator

- OR is true when ANY input is true
 - written as: $x+y$ or $x \vee y$

OR	+
0 + 0	0
0 + 1	1
1 + 0	1
1 + 1	1

Java:

```
int x, a = 0x18, b = 0x11;
```

```
x = a | b;
```

```
// a = 0001 10002
```

```
// b = 0001 00012
```

```
// x = 0001 10012 = 0x19 = 25
```

XOR Operator

- XOR (exclusive OR) is similar to OR, except it is true when only one, not both, input is true
 - written as $x \oplus y$

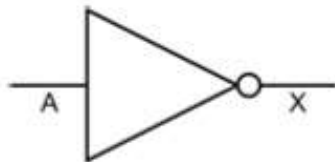
XOR	\oplus
0 . 0	0
0 . 1	1
1 . 0	1
1 . 1	0

Java:

```
int x, a = 0x18, b = 0x11;  
x = a ^ b;           // a = 0001 10002  
                     // b = 0001 00012  
                     // x = 0000 10012 = 0x09 = 9
```


Logic NOT gate

- *Computer hardware is made from logic gates*
 - *Basic logic gates: AND, OR, NOT*
 - *Derived negated gates: NAND, NOR, XOR, XNOR*
- *Any logic circuits can be build from combinations of just AND, OR, NOT*
 - *alternatively, use just NAND, NOR*
- *NOT gate*

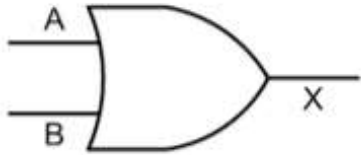


- *The above triangle is actually the symbol for an “Op Amp”, an operational amplifier whose function doubles to clean up the digital signal as it traverses the circuit. The circle to the left is the thing that denotes the actual NOT operation*

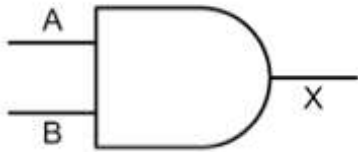


Logic OR / AND gates

- OR gate



- AND gate



Logic XOR / NAND gates

- XOR gate

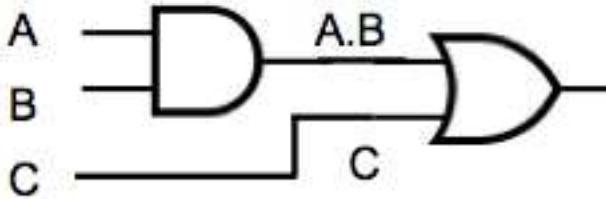


- NAND gate (a combination of NOT AND)
 - Note the little circle to the right.

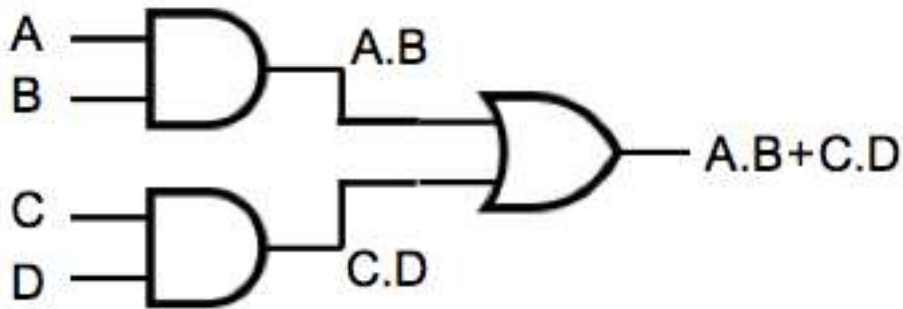


Logic Circuits

- $A.B+C$



- $A.B+C.D$



Gating

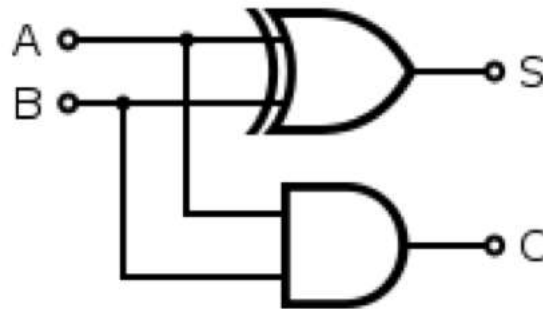
- Normally, AND, OR, XOR have pure signal inputs, but what if one of the signals is a constant?
- AND
 - $0 \text{ AND } A = 0$, always
 - $1 \text{ AND } A = A$, always
- OR
 - $0 \text{ OR } A = A$, always
 - $1 \text{ OR } A = 1$, always
- XOR
 - $0 \text{ XOR } A = A$
 - $1 \text{ XOR } A = \sim A$

Half Adder

- Say we wanted to construct an adder, which adds two bits together and outputs the sum (as a single bit) and a carry bit
 - called a *half adder*
 - truth table:

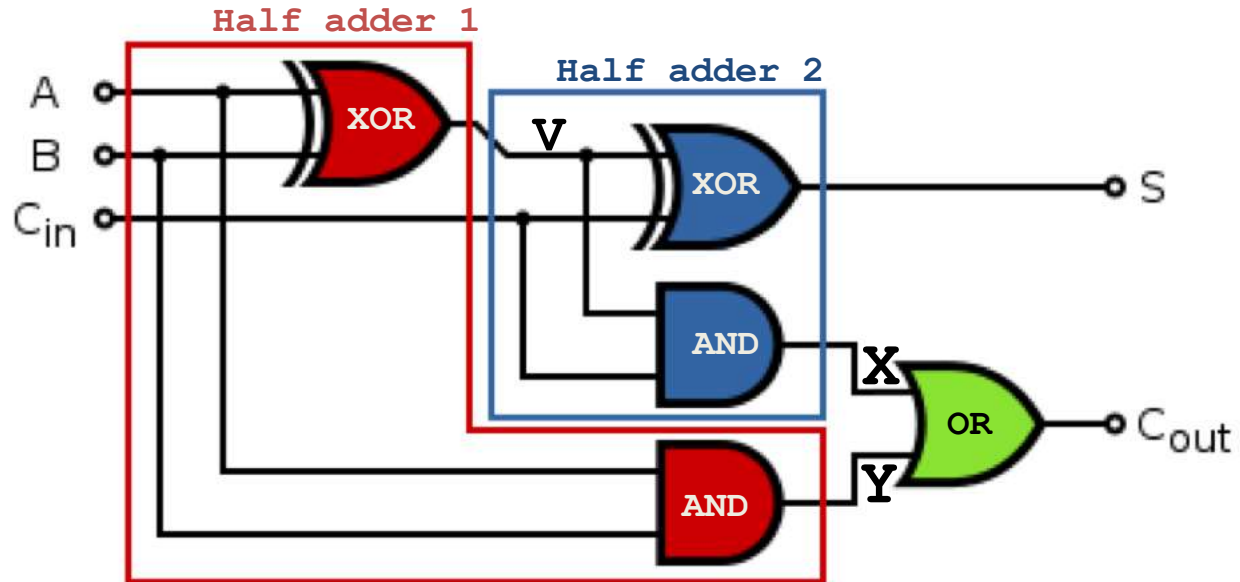
A	B	Carry	Sum
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	0

$$\text{Sum} = A \oplus B$$
$$\text{Carry} = A.B$$



Two Half Adders = A Full Adder

- We could combine two half adders to make a full adder (which accepts 3 inputs: A, B and a carry-in)



- Full adders can be cascaded together to make a parallel adder that can add multi-bit binary numbers

Full Adder Truth Table

A	B	C _{in}	V A ⊕ B	sum S V ⊕ C _{in}	X V • C _{in}	.Y A • B	carry C _{out} X + Y
0	0	0					
0	0	1					
0	1	0					
0	1	1					
1	0	0					
1	0	1					
1	1	0					
1	1	1					

$$V = A \oplus B$$

$$X = V \cdot C_{in}$$

$$Y = A \cdot B$$

$$S = V \oplus C_{in}$$

$$C_{out} = X + Y$$

Full Adder Truth Table (Solution)

A	B	C _{in}	V A \oplus B	sum S V \oplus C _{in}	X V \cdot C _{in}	Y A \cdot B	carry C _{out} X + Y
0	0	0	0	0	0	0	0
0	0	1	0	1	0	0	0
0	1	0	1	1	0	0	0
0	1	1	1	0	1	0	1
1	0	0	1	1	0	0	0
1	0	1	1	0	1	0	1
1	1	0	0	0	0	1	1
1	1	1	0	1	0	1	1

Sanity check

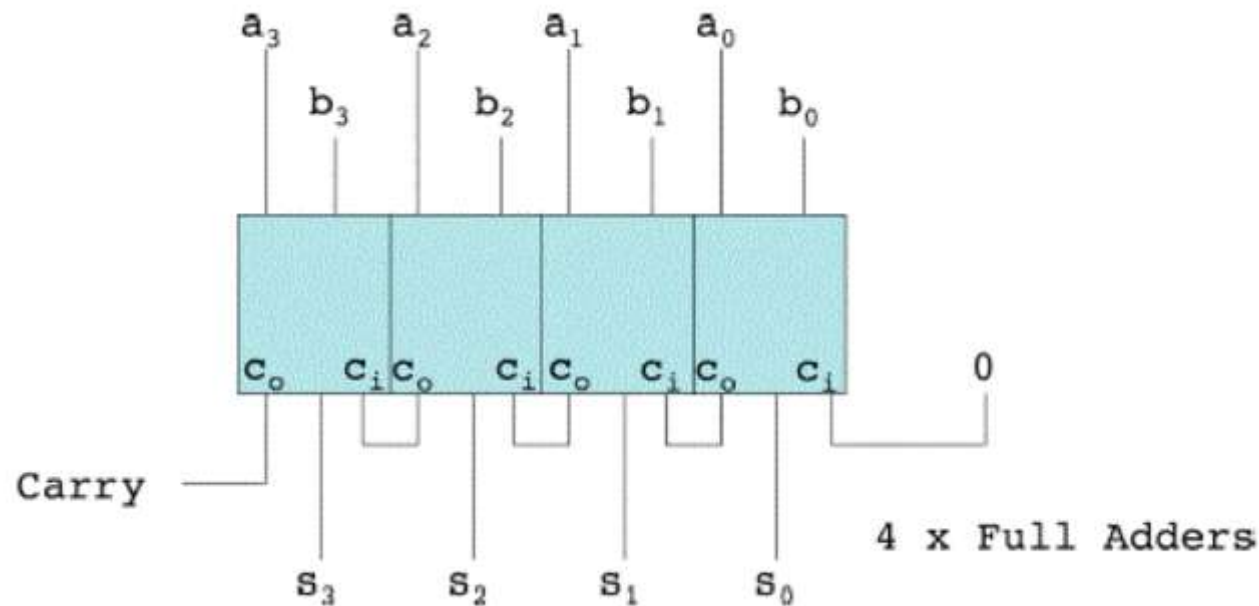
S = 1, whenever **A+B+C_{in} = 1 or 3, else 0**

C_{out} = 1, whenever **A+B+C_{in} = 2 or 3, else 0**

Parallel Adders

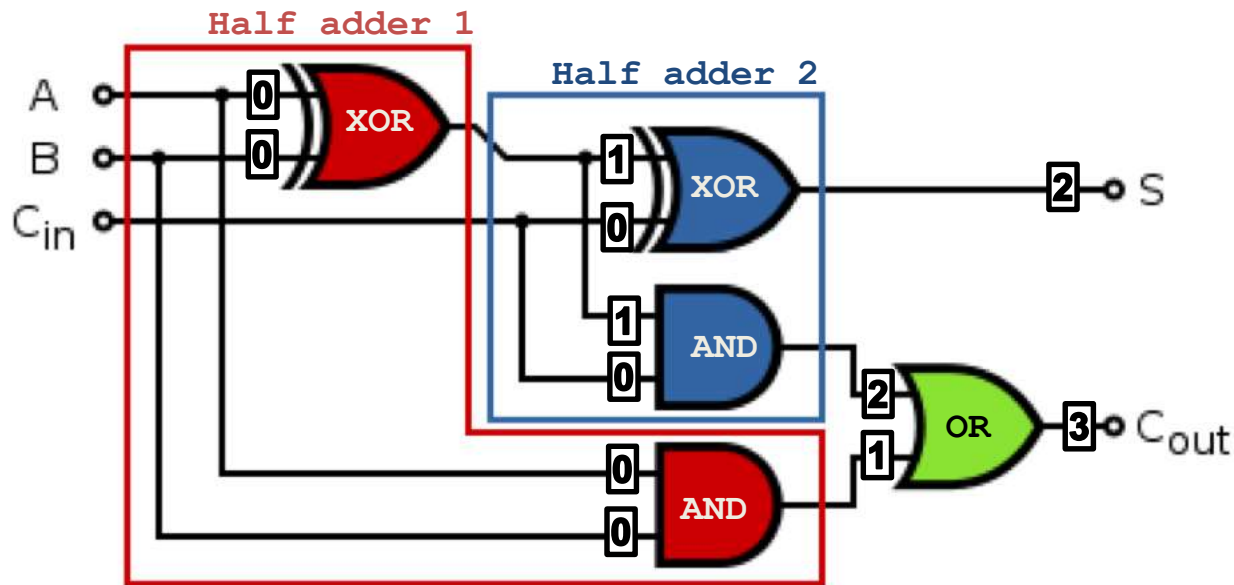
Parallel adder example

- 4 cascaded full-adders to add two 4-bit binary numbers
- the carry-out of each full-adder feeds into carry-in of the next full-adder
- the first carry-in is set to 0



Propagation Delay

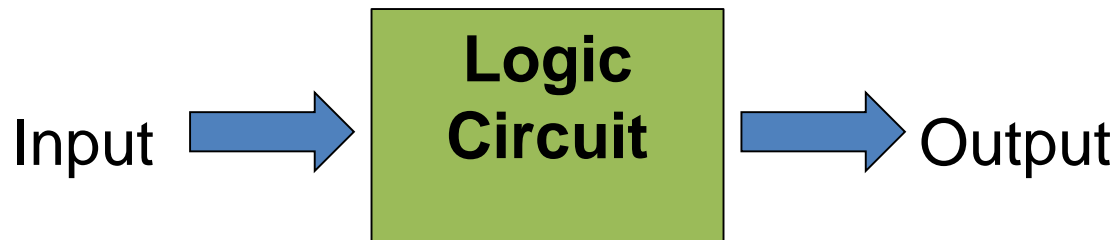
- Notice how in the full adder below, a signal can pass through 1, 2 or 3 gates before reaching C_{out} .
 - Thus, if A, B, C_{in} change, it may take up to 3 cycles before S and C_{out} take their correct values.
 - This is in general called propagation delay and has to be taken into account when building circuits.



- So it can take up to 3 clock ticks for the output to correctly reflect the input

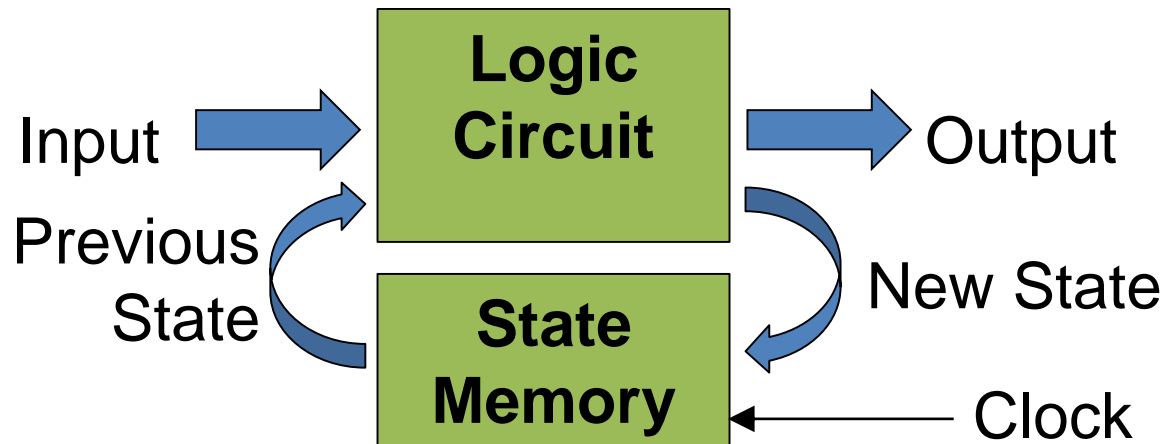
Combinational Circuits

- A **combinational circuit** is one where the output state should instantly change whenever the input changes.
 - The full adder circuit is a combinational circuit.
 - Mathematically it is equivalent to the statement
$$1 + 3 = 2 + 2, \text{ or}$$
$$A + 5 = B + 3, \text{ when } A = 3 \text{ and } B = 5$$
Which is always true. The symbol A is **substituted** the number 3, and so for B.
- Clearly the output is a direct function of the input



Sequential Circuits

- A **sequential circuit** is one that has some form of memory. It remembers its previous state, and this can affect its new state.
 - In practice, due to propagation delay, which can be seen as a form of memory, most circuits are treated as if they are sequential, even when they are not.
- **We will return to sequential logic in a future lecture.**



Python Code

- Consider the following Python code

```
def avg(a,b,c):  
    return (a+b+c) / 3
```

- Clearly, the function avg() need not remember what it did the last time. It just adds the three values. This is the **combinational circuit**.

- add
- Now consider the following

```
>>> avgbuf = [0, 0]  
>>> def runavg(x):  
    global avgbuf  
    s = x + avgbuf[0] + avgbuf[1]  
    avgbuf[0] = avgbuf[1]  
    avgbuf[1] = x  
    return s / 3
```

```
>>> runavg(3)  
1  
>>> runavg(3)  
2  
>>> runavg(3)  
3  
>>> avgbuf  
[3, 3]
```

The program remembered the values it was given and so the answer depends on what happened before.

This is equivalent to a **sequential circuit**.

Bit Masking

- Bit masks uses Boolean operations to access individual bits from binary data
- Used in 'low level' programming
 - device drivers and other hardware configuration/communication
 - data packet encoding/decoding
 - low level graphics

Bit Masking (Set)

- To set (i.e. make 1) a bit in a byte, OR it with a mask of all 0's except the bit to set

- e.g. set b4 of 10001101

10000101

00001000 ← turn ON bits where there is a '1'

----- Boolean OR

10001101

- In Java:

```
short data = 0x85, maskON = 0x08;
```

```
short result = (short) (data | maskON);    // = 0x8D
```

- In Python:

```
data = 0b10000101
```

```
maskON = 0b00001000
```

```
result = data | maskON                                # = 0b10001101
```


Bit Masking (Reset)

- To reset (i.e. make 0) a bit in a byte, AND it with a mask of all 1's except the bit to reset

- e.g. reset the LSB of 10001101

10001101

11111110 ← turn OFF bits where there is a '0'

----- Boolean *AND*

10001100

- In Java:

```
short data = 0x8D,    maskOFF = 0x01;
```

```
short result = (short) (data & ! maskOFF);    // = 0x8C
```

- In Python:

```
data = 0b10000101
```

```
maskOFF = 0b00000001
```

```
result = data & ~maskOFF
```

```
# = 0b10001100
```

Bit Masking (Flip)

- To flip (i.e. $1 \rightarrow 0, 0 \rightarrow 1$) a bit in a byte, XOR it with a mask of all 0's except the bit(s) to flip

- e.g. flip b4 and the LSB of 10001101

10000101

00001001 \leftarrow complement bits where there is a '1'

----- Boolean XOR

10001100

- In Java:

```
short data = 0x85, maskFLIP = 0x08 | 0x01;
```

```
short result = (short) (data ^ maskFLIP);           // = 0x8d
```

- In Python:

```
data = 0b10000101
```

```
maskFLIP = 0x08 ^ 0x01
```

```
result = data ^ maskON                                # = 0b10001100
```

Quiz

- Given the equations for a full adder

$$S = (A \oplus B) \oplus C_{in}$$

$$C_{out} = (A \oplus B) C_{in} + AB$$

Show that when $A = B$, then

$$S = C_{in} \text{ always, and}$$

$$C_{out} = C_{in} \text{ always}$$

using truth tables.

- If you were to use such a circuit in a parallel adder or 8-bit integers:
 - What is it useful for?
 - What does it do?
 - When $C_{in} = 0$,