

# COSC2473

## Digital Logic

Logic Gates & Bit Masking  
Boolean Algebra



# Boolean Algebra

- Boolean algebra has many of the same rules as normal high school level algebra, and indeed the symbols used (+ for OR, \* for AND, etc) reflect this.
  - And just like algebra, you can use various methods to simplify.
  - We will show here two methods:
- Algebraic Simplification
  - Can work for any number of variables
- Graphical Simplification using Karnaugh maps
  - This works for up to 4 variables, but beyond that, it gets too complicated.

# Boolean Logic Rules

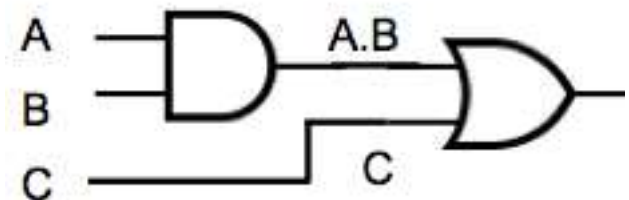
- Reminder of the Boolean Logic rules

		NOT	AND	OR	XOR
A	B	$\sim A$	$\bullet$	$+$	$\oplus$
0	0	1	0	0	0
0	1	1	0	1	1
1	0	0	0	1	1
1	1	0	1	1	0

- Typical Boolean expressions are:

- $C + AB = \text{true}$

- if C is true or both A and B are true and is equivalent to the circuit at right



# Boolean Algebra Rules

Commutative law	$A + B = B + A$
Associative law	$A + (B + C) = A + B + C = (A + B) + C$ $A(BC) = ABC = (AB)C$
Distributive law	$A(B + C) = AB + AC$
de Morgan's theorems	$\overline{A + B} = \overline{A} \cdot \overline{B}$ $\overline{A \cdot B} = \overline{A} + \overline{B}$
Identities	$1 + \text{any} = 1, \quad 0 + A = A$ $1 \cdot A = A, \quad 0 \cdot \text{any} = 0$ $A \cdot A = A, \quad A + A = A$ $A + \overline{A} = 1, \quad A \cdot \overline{A} = 0$ $AB + \overline{A}B = B$ $A + BC = (A + B)(A + C)$

# Algebraic Simplification and Python

Example of simplification and how it is used in programming

$$\begin{aligned} A + \neg A C &= (A + \neg A)(A + C), \\ &\text{using the pattern } A + BC = (A+B)(A+C) \\ &= 1 (A + C = A + C, \text{ since } (A + \neg A) = 1 \end{aligned}$$

## Python Code

Let's test the above. Suppose we have two sensors, S1, S2 = +/- 5 volts, and an activation decision Z

```
>>> S1=-5; S2=-5; A=S1>0; C=S2>0; Z = A or (not A and C); print(A,C,Z)
False False False
>>> S1=-5; S2=+5; A=S1>0; C=S2>0; Z = A or (not A and C); print(A,C, Z)
False True True
>>> S1=+5; S2=-5; A=S1>0; C=S2>0; Z = A or (not A and C); print(A,C, Z)
True False True
>>> S1=+5; S2=+5; A=S1>0; C=S2>0; Z = A or (not A and C); print(A,C, Z)
True True True
>>>
```

## Proof

The above values of A,B,Z look like a truth table for  $Z=A + C$

# Algebraic Simplification

- Example of simplification

1. Eliminating common variables

$$A + AB = A(\underline{1 + B}) = A, \quad \text{since } 1 + \text{anything} = 1$$

$$AB + \neg AB = (\underline{A + \neg A})B = B, \quad \text{where } \neg A = \text{NOT } A$$

2. More Complex Factoring

$$\begin{aligned} A + BC &= (A + B)(A + C) \\ &= \underline{A.A} + A.C + \underline{B.A} + B.C, \\ &= A + A.C + A.B + B.C, \\ &= \underline{A.1} + A.C + AB + B.C, \\ &= A.(\underline{1 + C + B}) + B.C, \\ &= A.(1) + B.C, \\ &= A + B.C \end{aligned}$$

*multiply out, so  
but  $AA=A$ ,  $BA=AB$   
but  $A = A.1$ , so  
factor out  $A$ , so  
but  $(1+\text{any})=1$ , so  
but  $A.1 = A$ , so  
...QED*

# Using Truth Table as Proof

- We can also prove the previous equation by showing all combinations in a truth table. We show that Left hand side (LHS) = right hand side (RHS)

A	B	C	X	LHS	Y	Z	RHS
			B.C	A + X	A + B	A + C	Y . Z
0	0	0	0	0	0	0	0
0	0	1	0	0	0	1	0
0	1	0	0	0	1	0	0
0	1	1	1	1	1	1	1
1	0	0	0	1	1	1	1
1	0	1	0	1	1	1	1
1	1	0	0	1	1	1	1
1	1	1	1	1	1	1	1

# De Morgan's Theorem

- One way in which Boolean algebra is different to other forms is the following
  - Most commonly used to change a term involving ORs into one involving ANDs, and vice versa
  - It works due to the symmetry of the OR and AND truthtables

$$\overline{A + B} = \overline{A} \cdot \overline{B}$$

$$\overline{A \cdot B} = \overline{A} + \overline{B}$$

A	B	$\overline{A}$	$\overline{B}$	A.B	$\overline{A+B}$
0	0	1	1	0	1
0	1	1	0	0	1
1	0	0	1	0	1
1	1	0	0	1	0

- One is basically the upside-down inverse of the other
  - You can get upside-down by complementing the operands A and B
  - and then simply map AND to OR
- de Morgan's Laws are also applied to Set Theory, where
  - OR  $\Leftrightarrow \cup$  (set union)
  - AND  $\Leftrightarrow \cap$  (set intersection)
  - NOT  $\Leftrightarrow '$  (set complement)



# A Practical Example

- Consider a platform elevator on Flinders Street Station
  - Would you know how to program a 2 level elevator?
  - Where do you start?
  - Consider the cases
    - 2 levels: Concourse, Platform
    - Up and Down buttons on each level outside of the car
    - Up and Down buttons inside the car
    - Activate motor to travel up or down
    - Door open sensor to decide whether to move the car
    - Door Open/Close on arrival or based on inside/outside buttons
- Not so simple, right?

# Flinders Street Platform Elevator 1

- Consider a platform elevator on Flinders Street Station
  - 2 levels: Concourse (C), Platform (/C) C
  - Up button on lower level outside car  $B_{OU}$
  - Down Button on upper level outside car  $B_{OD}$
  - Up button inside car  $B_{IU}$
  - Down button inside car  $B_{ID}$
  - Door open sensor (1 = closed) S
  - Travel Up for the car  $M_U$
  - Travel Down for the car  $M_D$
  - Door open / close Activator A

# Flinders Street Platform Elevator 2

- We can simplify further
  - Position of car (1=Concourse, 0= Platform) C
  - Up button on lower level outside car U
  - Down Button on upper level outside car D
  - Lever inside car L = 1 (Up) or 0 (Down) L
  - Motor = 1 (travel Up) or -1 (down) or 0 (do nothing) M
  - Assume doors open automatically on arrival, or button push when already there.

# Flinders Street Platform Elevator 3

- We can simplify further
  - Position of car (1=Concourse, 0= Platform) C
  - Outside Button calls the car B
  - Lever inside car L = 1 (Up) or 0 (Down) L
  - Motor = 1 (travel Up) or -1 (down) or 0 (do nothing) M

- What is needed for the motor to start?

If	(B /C L)	M = +1
else if	(B C /L)	M = -1
else		M = 0

M = 0

If (B)

if	(C and not L)	M = -1
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if	(L and not C)	M = +1
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# Graphical Simplification

- A Karnaugh Map is a pictorial representation of the bit patterns arising from a truth table

- Suppose we have 2 variables A,B

$$X = A + \neg AB$$

$$= T_1 + T_2$$

1. We populate the grid with our terms

	A	$\neg A$
B	1	2
$\neg B$	1	

- More: <https://www.youtube.com/watch?v=3vkMgTmieZI>

# Karnaugh Maps – 2 Variables

- A Karnaugh Map is also “***a special form of truth table that enables easier pattern recognition***” by humans.
- Suppose we have 2 variables A,B
$$X = A + /AB$$
$$= T_1 + T_2$$
  1. We populate the grid with our terms
  2. We then draw the largest possible overlapping rectangles that can cover all terms

	A	/A
B	1	2
/B	1	

# Karnaugh Maps – 2 Variables

- Suppose we have 2 variables A,B

$$\begin{aligned} X &= A + /AB \\ &= T_1 + T_2 \end{aligned}$$

- We populate the grid with our terms
- We then draw the largest possible overlapping rectangles that can cover all terms
- We then describe these rectangles as:

$$\begin{aligned} X &= \text{vert-rect} + \text{horiz-rect} \\ &= A + B \end{aligned}$$

	A	/A
B	1	2
/B	1	

# Karnaugh Maps – 2 Variables

- Equivalent Boolean Logic simplification

$$X = A + \overline{A}B$$

$$= A(\underline{B + \overline{B}}) + \overline{A}B$$

$$= AB + A\overline{B} + \overline{A}B + \underline{AB}$$

$$= A(\underline{B + \overline{B}}) + B(\underline{\overline{A} + A})$$

$$= A + B$$

$$= \overline{\overline{A}\overline{B}}$$

...de Morgan's

Extra Term = Overlap

	A	$\overline{A}$
B	1	1
$\overline{B}$	1	

- See: Karnaugh Maps Introduction

<https://www.youtube.com/watch?v=A0XupfXiKlo>



# Karnaugh Maps – Term 1

- Now we have 3 variables A, B, C on a 4 variable map
  - $X = AB + AC + /AB/C + /B/C$
  - $X = T_1 + T_2 + T_3 + T_4$  'minterms'

– AB

	AB	A/B	/A/B	/AB
CD	1			
C/D	1			
/C/D	1			
/CD	1			

# Karnaugh Maps – Term 2

- Now we have 3 variables A, B, C on a 4 variable map
  - $X = AB + AC + /AB/C + /B/C$
  - $X = T_1 + T_2 + T_3 + T_4$  'minterms'

– AC

	AB	A/B	/A/B	/AB
CD	2	2		
C/D	2	2		
/C/D				
/CD				

# Karnaugh Maps – Term 3

- Now we have 3 variables A, B, C on a 4 variable map
  - $X = AB + AC + /AB/C + /B/C$
  - $X = T_1 + T_2 + T_3 + T_4$  'minterms'

–  $/AB/C$

	AB	A/B	/A/B	/AB
CD				
C/D				
/C/D				3
/CD				3

# Karnaugh Maps – Term 4

- Now we have 3 variables A, B, C on a 4 variable map
  - $X = AB + AC + /AB/C + /B/C$
  - $X = T_1 + T_2 + T_3 + T_4$  'minterms'

–  $/B/C$

	AB	A/B	/A/B	/AB
CD				
C/D				
/C/D		4	4	
/CD		4	4	

# Karnaugh Maps – A

- Karnaugh map rects must be as large as possible
- Now we have 3 variables A, B, C on a 4 variable map
  - $X = AB + AC + /AB/C + /B/C$
  - $X = T_1 + T_2 + T_3 + T_4$  'minterms'

- Make largest overlapping rect

- A

	AB	A/B	/A/B	/AB
CD	12	2		
C/D	12	2		
/C/D	1	4	4	3
/CD	1	4	4	3

# Karnaugh Maps – /C

- Karnaugh map rects must contain  $2^n$  elements only
- Now we have 3 variables A, B, C on a 4 variable map
  - $X = AB + AC + /AB/C + /B/C$
  - $X = T_1 + T_2 + T_3 + T_4$  'minterms'

- Make largest overlapping rect

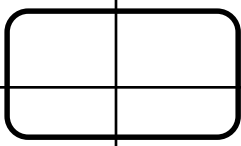
- $A + /C$

	AB	A/B	/A/B	/AB
CD	12	2		
C/D	12	2		
/C/D	1	4	4	3
/CD	1	4	4	3

# K-Map – Confirm using deMorgan's

- Confirm using deMorgan's Theorem
- Now we have 3 variables A, B, C on a 4 variable map
  - $X = AB + AC + /AB/C + /B/C$
  - $X = T_1 + T_2 + T_3 + T_4$  'minterms'

- Make largest overlapping rect but for the unoccupied parts
- Then invert this.

	AB	A/B	/A/B	/AB
CD	12	2		
C/D	12	2		
/C/D	1	4	4	3
/CD	1	4	4	3

- $/ ( /AC )$

# Karnaugh Maps – 4 Variables

- Karnaugh map rects must contain no gaps
- Now change a term to add a fourth variable, D

–  $X = AB + \underline{ACD} + /AB/C + /B/C$

–  $X = T_1 + T_2 + T_3 + T_4$  'minterms'

– So  $T_2$  now is  $ACD$

– The 2's are changed

– Gap in  $A/BC/D$

– So cannot use A rect

– Now

$X = AB + \underline{ACD} + /C$

	AB	A/B	/A/B	/AB
CD	12	2		
C/D	1			
/C/D	1	4	4	3
/CD	1	4	4	3



# Karnaugh Maps – Wrap-around Rects

- Karnaugh map rects can also wrap around.
- Now change a term to show wrap on 3 variable K-map
  - $X = AB + \underline{BC} + /AB/C + /B/C$
  - $X = T_1 + T_2 + T_3 + T_4$  'minterms'
  - So  $T_2$  now is  $BC$
  - The rect for  $T_2$  now wraps around
  - The biggest rect for  $T_2$  is  $1 \times 2$  that wraps around and corresponds to  $B$
  - This entirely covers  $T_1$ , so we can remove it.
  - Finally  

$$X = \underline{B} + /C$$

	AB	A/B	/A/B	/AB
C	12			2
/C	1	4	4	3

# K-Maps – General Rules

- General comments about K-maps
  - Expressions must be organised into “minterms”, which are “sums of products” of the form
$$X = T_1 + T_2 + \dots$$
where  $T$  is a group of ANDed variables.
  - Area of rect  $R = 2^n / m$ , where  $n$  is the total number of variables, and  $m$  is the number of variables in this particular term.
    - In previous slide,  $n=3$  variables,  $m$  for  $T_2$  (BC) = 2, so  $R=2^3/2 = 4$
- For example:
$$Y = A(B + C)$$
is in the wrong form, multiply out
$$Y = AB + AC = 2 \text{ rects of } 4 \text{ elements each}$$
before populating the map.

# Summary

- So we have three ways of simplifying Boolean expressions
  - First Principles
    - Good as a first choice for when you know the situation, such as the elevator example where we simplified the algorithm first.
  - Boolean Algebra
    - Good when you have a few Boolean condition variables, culminating into a small number of decision variables
  - Karnaugh Maps
    - As a graphical way of simplification, but it does not scale well beyond about 6 variables (if you can manage the 3D).
  - Truth Tables
    - The lowest level, but also the most robust, and the easiest to scale and automate.

# Quiz

- Using K-maps, simplify the expression

$$X = A.B + A./B .C + D( /C.+ /BC)$$

using the following steps.

1. Rewrite expression in terms of minterms
2. Assign minterm numbers
3. Draw each minterm on the K-Map
4. Draw as few maximal overlapping rectangles as possible to cover all the non-blank entries
5. List the terms corresponding to each rectangle.

	AB	A/B	/A/B	/AB
CD				
C/D				
/C/D				
/CD				