

1. (a) Augmenting path Residual capacity

s -0/10-> a -0/5-> t 5

(Indicating path in *original* network N with current flow and capacity. And remember: "augmenting f along path P" means to add the residual capacity to each forward edge, and to subtract it from each backward edge. In this case, set $f(s,a) = 5$, $f(a,t) = 5$.)

s -0/5-> d -0/10-> t 5

s -0/8-> b -0/10-> c -0/8-> t 8

s -5/10-> a -0/3-> c <-8/10- b -0/3-> d -5/10-> t 3

No more augmenting paths.

Final flows:

$f(s,a) = 8$	$f(b,a) = 0$	$f(a,t) = 5$
$f(s,b) = 8$	$f(b,c) = 5$	$f(d,c) = 0$
$f(s,d) = 5$	$f(b,d) = 3$	$f(d,t) = 8$
	$f(a,c) = 3$	$f(c,t) = 8$

Flow value $|f| = f(s,a) + f(s,b) + f(s,d) = 8 + 8 + 5 = 21$.

- (b) Forward edges across X_0 : (s,a), (b,a), (c,t), (d,t).
Backward edges across X_0 : (a,c).

- (c) Capacity of cut X_0 :

$$\begin{aligned} c(X_0) &= c(s,a) + c(b,a) + c(c,t) + c(d,t) \\ &= 10 + 3 + 8 + 10 \\ &= 31. \end{aligned}$$

Flow across cut X_0 :

$$\begin{aligned} f(X_0) &= f(s,a) + f(b,a) + f(c,t) + f(d,t) - f(a,c) \\ &= 8 + 0 + 8 + 8 - 3 \\ &= 21. \end{aligned}$$

- (d) Start with $X_1 = (\{s\}, \{a,b,c,d,t\})$ and flow from part (a).
Edge (s,a) crosses cut forward with residual capacity 2, so set $X_1 = (\{s,a\}, \{b,c,d,t\})$.
All edges forward across cut have $f = c$: (s,b), (s,d), (a,c), (a,t).
All edges backward across cut have $f = 0$: (b,a).
Cut $X_1 = (\{s,a\}, \{b,c,d,t\})$ has capacity

$$\begin{aligned} c(X_1) &= c(s,d) + c(s,b) + c(a,c) + c(a,t) \\ &= 5 + 8 + 3 + 5 \\ &= 21 \\ &= |f|. \end{aligned}$$

2. Add "super-source" s with edges (s,s₁), ..., (s,s_k) each of capacity oo; add "super-sink" t with edges (t₁,t), ..., (t_l,t) each of capacity oo. (Instead of using oo, can set capacity to sum of outgoing/incoming capacities).

Max flow in resulting network N' = max flow in original network N:

- any flow in original network can be extended to a flow in resulting network (for new edges from super-source to source, set flow equal

- to total flow out of source; for new edges from sink to super-sink, set flow equal to total flow into sink) -- hence, $\max \text{ flow in } N' \geq \max \text{ flow in } N$;
- any flow in N' induces flow in N (flow out of every source and into every sink limited only by edges in original network because of "infinite" capacities on new edges) -- hence, $\max \text{ flow in } N \geq \max \text{ flow in } N'$.