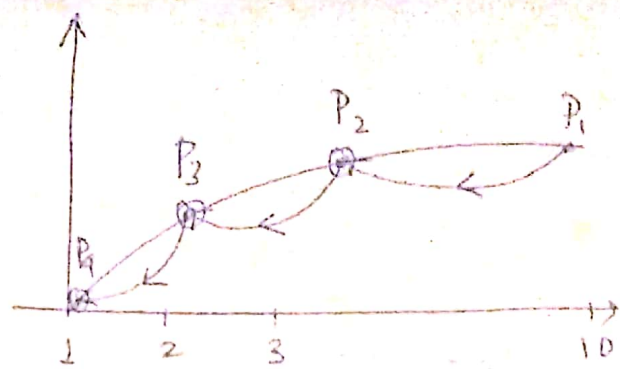


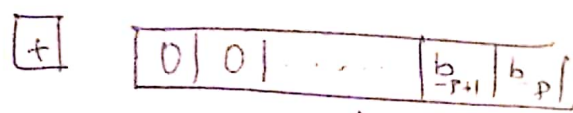
6. On starting off with a random number in $[1, 10]$ and carrying out successive square root operations; we find ourselves traversing the points $P_1 \rightarrow P_2 \rightarrow P_3 \rightarrow P_4 \rightarrow \dots \rightarrow P_n$.



If n is large; we find $P_n \rightarrow 1$.

So; carrying out square roots very large number of times leaves us with a number of the form $(1+d)$ where $d \ll 1$.

Now; in representing this number as a machine number it is fair to assume it would look like:-



some fixed bits in mantissa would be filled

$\times 2^0$

this is because we have an addition by 1 so; 2^0 is the exponent.

Now; each successive square root operation takes off one bit from the left in the mantissa. So; after 53 iterations (starting from a number in the range of 10^{-10} ;) we would be left with either a zero or 1 at the b_{-52} bit.

The fact that we have a 2-point degeneracy (all the numbers collapse to either of two values i.e; e^0 or e^2) is also a proof that only the last bit in the mantissa was different at the end of 53 square roots.

So, all the numbers we started of with will have a mantissa that looks like either of the following:-

$$A = \boxed{0 \mid 0 \mid 0 \mid \dots \mid 0 \mid 0} \quad \text{or} \quad B = \boxed{0 \mid 0 \mid 0 \mid \dots \mid 0 \mid 1}$$

raised to the exponent 2^0 .

Now; A can be represented as:-

$$(1 + 0) = 1 \quad \text{and squaring that 53 times will give 1.}$$

so, all numbers that have mantissa A collapse to 1.

For B; we can write it as:- $\left(1 + \frac{1}{2^{52}}\right)$ and squaring it 53 times :-

$$\left(1 + \frac{1}{2^{52}}\right)^{2^{53}} = \left\{ \left(1 + \frac{1}{2^{52}}\right)^{2^{52}} \right\}^2$$

Now; from the relation:- $\lim_{n \rightarrow \infty} \left(1 + \frac{1}{n}\right)^n = e$;

we can say that the number $\left\{ \left(1 + \frac{1}{2^{52}}\right)^{2^{52}} \right\}^2$ will be very close to e^2 .

Therefore; all the numbers that had mantissa B will collapse to e^2 .

Therefore; only e^0 and e^2 yield correct results after square-rooting and squaring 53 times.

Now; carrying out a similar treatment for $n=52$:-

However; this time the last 2 bits will be filled (because we start off with a range of 0-10 and the numbers closer to 10 sum up to 2 bits in mantissa after 52 square roots).

$$A_1 = \begin{array}{|c|c|c|c|c|c|c|c|} \hline 0 & 0 & 0 & \dots & 0 & 1 & 1 & \\ \hline \end{array}$$

(0's)

$$(0's) A_2 = \begin{array}{|c|c|c|c|c|c|c|c|} \hline 0 & 0 & \dots & 0 & 1 & 0 & \\ \hline \end{array}$$

(0's)

$$B_1 = \begin{array}{|c|c|c|c|c|c|c|c|} \hline 0 & 0 & \dots & 0 & 0 & 1 & \\ \hline \end{array}$$

$$(0's) B_2 = \begin{array}{|c|c|c|c|c|c|c|c|} \hline 0 & 0 & \dots & 0 & 0 & 0 & \\ \hline \end{array}$$

$$A_1 = \left(1 + \frac{1}{2^{51}} + \frac{1}{2^{52}} \right) = \left(1 + \frac{1}{2^{52}} \left(\frac{2}{3} \right) \right)$$

$$\text{and; } A_1^{52} = \left\{ \left(1 + \frac{1}{2^{52}} \left(\frac{2}{3} \right) \right)^{2^{52}} \left(\frac{2}{3} \right) \right\}^{3/2} \approx e^{3/2}$$

~~B_1 and B_2 are e^2 and e^0 etc~~

$$A_2^{52} = \left(1 + \frac{1}{2^{51}} \right)^{2^{52}} \approx e^2$$

$$B_1 = \left(1 + \frac{1}{2^{52}} \right)^{2^{52}} \approx e \quad \text{and} \quad B_2 = (1+0)^{2^{52}} = e^0$$

so; we have 4-point degeneracy = e^0, e, e^2, e^3

However; e^3 falls outside our domain and so we cannot see it.

so; we have 3-point degeneracy = e^0, e^1, e^2 for 52 square roots followed by 52 squares.

For 51 square roots and 51 squares,

The range of mantissa configurations-

$$A_1: \begin{array}{|c|c|c|c|c|c|c|c|} \hline 0 & 0 & 0 & \dots & 0 & 1 & 1 & 1 \\ \hline \end{array}$$

(05)

$$B_1: \begin{array}{|c|c|c|c|c|c|c|c|} \hline 0 & 0 & \dots & 0 & 1 & 1 & 0 \\ \hline \end{array}$$

$$A_2: \begin{array}{|c|c|c|c|c|c|c|c|} \hline 0 & 0 & 0 & \dots & 0 & 1 & 0 & 1 \\ \hline \end{array}$$

(05)

$$B_2: \begin{array}{|c|c|c|c|c|c|c|c|} \hline 0 & 0 & \dots & 0 & 0 & 1 & 1 \\ \hline \end{array}$$

$$A_3: \begin{array}{|c|c|c|c|c|c|c|c|} \hline 0 & 0 & \dots & 0 & 0 & 1 & 0 \\ \hline \end{array}$$

(05)

$$B_3: \begin{array}{|c|c|c|c|c|c|c|c|} \hline 0 & 0 & \dots & 0 & 0 & 0 & 1 \\ \hline \end{array}$$

$$A_4: \begin{array}{|c|c|c|c|c|c|c|c|} \hline 0 & 0 & \dots & 0 & 0 & 0 & 0 \\ \hline \end{array}$$

$$B_4: \begin{array}{|c|c|c|c|c|c|c|c|} \hline 0 & 0 & \dots & 0 & 1 & 0 & 0 \\ \hline \end{array}$$

$$A_1^{2^{51}} = \left(1 + \frac{1}{2^{52}} (4 + 2 + 1) \right)^{2^{51}} = \left\{ \left(1 + \frac{1}{2^{52} \times (1/7)} \right)^{2^{52} \times \frac{1}{7}} \right\}^{7 \times \frac{1}{2}} \approx e^{3.5}$$

$$A_2^{2^{51}} = \left(1 + \frac{1}{2^{52}} \left(\frac{1}{5} \right) \right)^{2^{51}} = \left\{ \left(1 + \frac{1}{2^{52} \times \frac{1}{5}} \right)^{2^{52} \times \frac{1}{5}} \right\}^{5/2} \approx e^{2.5}$$

$$A_3^{2^{51}} = \left(1 + \frac{1}{2^{51}} \right)^{2^{51}} \approx e^1; \quad A_4^{2^{51}} = e^0 = 1$$

$$B_1^{2^{51}} = \left\{ \left(1 + \frac{1}{2^{52} \times \frac{1}{3}} \right)^{2^{52} \times \frac{1}{3}} \right\}^3 \approx e^3; \quad B_2^{2^{51}} = \left\{ \left(1 + \frac{1}{2^{52} \times \frac{1}{3}} \right)^{2^{52} \times \frac{1}{3}} \right\}^{3/2} \approx e^{1.5}$$

$$B_3^{2^{51}} = \left\{ \left(1 + \frac{1}{2^{52}} \right)^{2^{52}} \right\}^{1/2} \approx e^{0.5}$$

$$B_4^{2^{51}} = \left\{ \left(1 + \frac{1}{2^{50}} \right)^{2^{50}} \right\}^2 \approx e^2$$

However, we can go till e^2 in our domain. So, the values are:-

$$e^0, e^{0.5}, e^1, e^{1.5}, e^2$$

which remain the same and other values collapse to these!