

# SI231b: Matrix Computations

## Lecture 15: Rayleigh Quotient for Eigenvalues of Hermitian Matrices

Yue Qiu

[qiuyue@shanghaitech.edu.cn](mailto:qiuyue@shanghaitech.edu.cn)

School of Information Science and Technology  
ShanghaiTech University

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# Eigenvalues of Hermitian/real symmetric matrices

In this lecture, we focus on eigenvalues of Hermitian/real symmetric matrices whose eigenvalues are real, then we have simplified results.

- ▶ power iteration + deflation can be used to compute more eigenpairs, i.e., once  $(\lambda_1, v_1)$  is computed, applying power iteration to  $A = A - \lambda_1 v_1 v_1^H$  gives  $(\lambda_2, v_2)$ .
- ▶ QR iteration + Hessenberg reduction:
  - The Hessenberg reduction reduces  $A$  to a tridiagonal matrix
  - QR iteration (with shifts) forces to converge to a bi-diagonal matrix
- ▶ Subspace iteration: the orthogonal basis converge to the dominant invariant subspace (for general matrices)
  - the orthogonal basis converge to the associated eigenvectors of dominant eigenvalues

## Eigenvalue equation

$$f(\lambda, \mathbf{v}) = \mathbf{A}\mathbf{v} - \lambda\mathbf{v}.$$

To differentiate, we obtain

$$\delta f = (\mathbf{A} - \lambda \mathbf{I})\delta \mathbf{v} - (\delta \lambda)\mathbf{v}.$$

Newton's method gives

$$f(\lambda, \mathbf{v}) + \delta f = 0,$$

i.e., at the  $k$ -th Newton step,

$$\begin{aligned} 0 &= f(\lambda_k, \mathbf{v}_k) + \delta f(\lambda_k, \mathbf{v}_k) \\ &= (\mathbf{A} - \lambda_k \mathbf{I})(\mathbf{v}_k + \delta \mathbf{v}) - (\delta \lambda_k)\mathbf{v}_k \\ &= (\mathbf{A} - \lambda_k \mathbf{I})\mathbf{v}_{k+1} - (\delta \lambda_k)\mathbf{v}_k \end{aligned}$$

This gives  $\mathbf{v}_{k+1} = (\delta \lambda_k)(\mathbf{A} - \lambda_k \mathbf{I})^{-1}\mathbf{v}_k$ , where  $\delta \lambda_k$  is some normalizing constant.

How to update  $\lambda_k$ ?

## Least Square formulation of Eigenvalue Computation

Suppose  $\tilde{v}$  is an approximate eigenvector, we want to find the corresponding best approximate eigenvalue  $\tilde{\lambda}$ . This can be achieved by solving

$$\min_{\mu} \|A\tilde{v} - \mu\tilde{v}\|_2^2.$$

The best approximate  $\tilde{\lambda}$  is given by

$$\begin{aligned}\tilde{\lambda} &= \arg \min_{\mu} \|A\tilde{v} - \mu\tilde{v}\|_2^2 \\ &= \frac{\tilde{v}^H A \tilde{v}}{\tilde{v}^H \tilde{v}}\end{aligned}$$

## Rayleigh Quotient

For any  $x \in \mathbb{C}^n$  with  $x \neq 0$ , the Rayleigh Quotient is given by

$$r(x) = \frac{x^H A x}{x^H x}$$

The Rayleigh Quotient is a continuous function except at  $x = 0$ , and its gradient denoted by  $\nabla r(x)$  is given by

$$\nabla r(x) = \frac{2}{x^H x} (Ax - r(x)x)$$

- ▶ at an eigenvector of  $A$ , the gradient is a zero vector
- ▶ if  $\nabla r(x) = 0$ ,  $x$  is an eigenvector and  $r(x)$  is the corresponding eigenvalue.
- ▶ eigenvectors of  $A$  are the stationary points of the function  $r(x)$

Together with the Newton iteration for computing  $v_{k+1}$  and the Rayleigh Quotient, we obtain the [Rayleigh Quotient iteration](#) for computing the eigenpair.

## Rayleigh Quotient Iteration:

```
random selection of  $v_0 \in \mathbb{C}^n$   
 $\lambda_0 = r(v_0) = \frac{v_0^H A v_0}{v_0^H v_0}$   
for  $k = 1, 2, \dots$   
     $v_k = (A - \lambda_{k-1} I)^{-1} v_{k-1}$     solve  $(A - \lambda_{k-1} I) v_k = v_{k-1}$   
     $v_k = \frac{v_k}{\|v_k\|_2}$   
     $\lambda_k = (v_k)^H A v_k$   
end
```

- ▶ inverse iteration with shift, and shift varies per iteration
- ▶ **cubic convergence** for Hermitian/real symmetric matrices

**Theorem**[Rayleigh-Ritz]. Let  $A$  be a Hermitian matrix. It holds that

$$\lambda_{\min} \|x\|_2^2 \leq x^H A x \leq \lambda_{\max} \|x\|_2^2$$
$$\lambda_{\min} = \min_{x \in \mathbb{C}^n, \|x\|_2=1} x^H A x, \quad \lambda_{\max} = \max_{x \in \mathbb{C}^n, \|x\|_2=1} x^H A x$$

► provides information about  $\lambda_1$  and  $\lambda_n$  for  $A$

► Proof:

- by a change of variable  $y = V^H x$ , we have

$$x^H A x = y^H \Lambda y = \sum_{i=1}^n \lambda_i |y_i|^2 \leq \lambda_1 \sum_{i=1}^n |y_i|^2 = \lambda_1 \|V^H x\|_2^2 = \lambda_1 \|x\|_2^2$$

- we thus have  $\max_{\|x\|_2=1} x^H A x \leq \lambda_1$
- since  $v_1^H A v_1 = \lambda_1$ , the above equality is attained
- the results  $x^H A x \geq \lambda_n \|x\|_2^2$  and  $\min_{\|x\|_2=1} x^H A x = \lambda_n$  are proven in the same way

**Question:** how about  $\lambda_k$  for any  $k \in \{1, \dots, n\}$ ? Do we have a similar variational characterization as that in the Rayleigh-Ritz theorem?

**Theorem**[Courant-Fischer]. Let  $A$  be a Hermitian matrix, and let  $\mathcal{S}_k$  denote any subspace of  $\mathbb{C}^n$  and of dimension  $k$ . For any  $k \in \{1, \dots, n\}$ , it holds that

$$\begin{aligned}\lambda_k &= \min_{\mathcal{S}_{n-k+1} \subseteq \mathbb{C}^n} \max_{x \in \mathcal{S}_{n-k+1}, \|x\|_2=1} x^H A x \\ &= \max_{\mathcal{S}_k \subseteq \mathbb{C}^n} \min_{x \in \mathcal{S}_k, \|x\|_2=1} x^H A x\end{aligned}$$

(requires a proof)

- ▶ optima is achieved when  $\mathcal{S}_k = \text{span}\{v_1, v_2, \dots, v_k\}$  and  $\mathcal{S}_{n-k+1} = \text{span}\{v_k, v_{k+1}, \dots, v_n\}$
- ▶ The Rayleigh-Ritz Theorem is a special case of the Courant-Fischer minimax theorem when  $k = 1$  and  $k = n$



From the Courant-Fischer minimax theorem, we know that

- ▶ given any  $k$ -dimensional subspace  $\mathcal{S}$ , the smallest value of the Rayleigh quotient over that subspace is a lower bound on  $\lambda_k$  and an upper bound on  $\lambda_{n-k+1}$ .
- ▶ practically used to bound eigenvalues (Hermitian/real symmetric matrices)
- ▶ numerical range is often used to bound eigenvalues of non-Hermitian matrices (self-study for your own interest)

One step further, we have the [Cauchy interlace theorem](#), which relates the eigenvalues of a block Rayleigh quotient to the eigenvalues of the corresponding matrix.

**Theorem**[[Cauchy interlace](#)]. Suppose  $A$  is Hermitian/real symmetric, and let  $V$  be a matrix with  $m$  orthonormal columns. Then the eigenvalues of  $V^H A V$  interlace the eigenvalues of  $A$ . That is, if  $A$  has eigenvalues  $\alpha_1, \alpha_2, \dots, \alpha_n$  and  $V^H A V$  has eigenvalues  $\beta_j$ , then

$$\beta_j \in [\alpha_{n-m+j}, \alpha_j]$$

## Special Case of Cauchy Interlace Theorem

If  $B$  is a principle submatrix of  $m \times m$  for an  $n \times n$  Hermitian/real symmetric matrix  $A$ . Suppose  $A$  has eigenvalues  $\lambda_1 \geq \lambda_2 \geq \cdots \geq \lambda_n$ , and  $B$  has eigenvalues  $\beta_1 \geq \beta_2 \geq \cdots \geq \beta_m$ , then

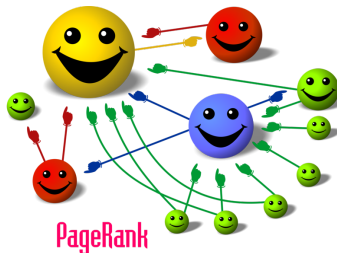
$$\lambda_{n-m+j} \leq \beta_j \leq \lambda_j.$$

Specifically, when  $m = n - 1$ ,

$$\lambda_n \leq \beta_{n-1} \leq \lambda_{n-1} \leq \cdots \leq \lambda_2 \leq \beta_1 \leq \lambda_1$$

# Case Study: PageRank Problem

- ▶ PageRank is an algorithm used by Google to rank the pages of a search result.
- ▶ the idea is to use counts of links of various pages to determine pages' importance.



Source: Wiki.

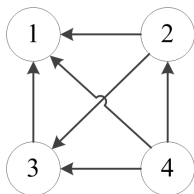
- ▶ further reading: [\[Bryan-Tanya2006\]](#)
  - K. Bryan and L. Tanya. "The 25, 000, 000, 000 eigenvector: The linear algebra behind Google," *SIAM Review*, vol. 48, no. 3, pp. 569–581, 2006.

► Model:

$$\sum_{j \in \mathcal{L}_i} \frac{v_j}{c_j} = v_i, \quad i = 1, \dots, n,$$

where  $c_j$  is the number of outgoing links from page  $j$ ;  $\mathcal{L}_i$  is the set of pages with a link to page  $i$ ;  $v_i$  is the importance score of page  $i$ .

► example:



$$\overbrace{\begin{bmatrix} 0 & \frac{1}{2} & 1 & \frac{1}{3} \\ 0 & 0 & 0 & \frac{1}{3} \\ 0 & \frac{1}{2} & 0 & \frac{1}{3} \\ 0 & 0 & 0 & 0 \end{bmatrix}}^A \overbrace{\begin{bmatrix} v_1 \\ v_2 \\ v_3 \\ v_4 \end{bmatrix}}^v = \overbrace{\begin{bmatrix} v_1 \\ v_2 \\ v_3 \\ v_4 \end{bmatrix}}^v.$$

- ▶ let  $A \in \mathbb{R}^{n \times n}$  be a matrix such that  $a_{ij} = 1/c_j$  if  $j \in \mathcal{L}_i$  and  $a_{ij} = 0$  if  $j \notin \mathcal{L}_i$
- ▶ **Problem:** find a **non-negative**  $v$  such that  $Av = v$ 
  - $A$  is extremely large and sparse, and we want to use the power method
- ▶ **Questions:**
  - does a solution to  $Av = v$  exist? Or, is  $\lambda = 1$  an eigenvalue of  $A$ ?
  - does  $Av = v$  have a non-negative solution? Or, does a non-negative eigenvector associated with  $\lambda = 1$  exist?
  - is the solution to  $Av = v$  unique? Or, would there exist more than one eigenvector associated with  $\lambda = 1$ ?
    - ▶ a unique solution is desired for this problem
  - is  $\lambda = 1$  the only eigenvalue that is the largest in modulus?
    - ▶ this is required for the power method

► notation:

- $x \geq y$  means that  $x_i \geq y_i$  for all  $i$
- $x > y$  means that  $x_i > y_i$  for all  $i$
- $x \not\geq y$  means that  $x \geq y$  does not hold
- the same notations apply to matrices

► conventions:

- $x$  is said to be non-negative if  $x \geq 0$ , and non-positive if  $-x \geq 0$
- $x$  is said to be positive if  $x > 0$ , and negative if  $-x > 0$
- the same conventions apply to matrices
- a square  $A$  is said to be **column-stochastic** if  $A \geq 0$  and  $A^T \mathbf{1} = \mathbf{1}$ 
  - a column-stochastic  $A$  has every column  $a_i$  satisfying  $a_i^T \mathbf{1} = \sum_{j=1}^n a_{ji} = 1$

- ▶ in PageRank,  $A$  is column-stochastic if all pages have outgoing links
  - see the literature to see how to deal with cases where some pages do not have outgoing links (dangling nodes)

**Property:** Let  $A$  be column-stochastic. Then,

1.  $\lambda = 1$  is an eigenvalue of  $A$
2.  $|\lambda| \leq 1$  for any eigenvalue  $\lambda$  of  $A$

- ▶ Implications:
  - a solution to  $Av = v$  does exist, though it doesn't say if  $v \geq 0$  or not
  - $\lambda = 1$  is an eigenvalue that has the largest modulus, but we don't know if it is the *only* eigenvalue that has the largest modulus
- ▶ we resort to non-negative matrix theory to answer the rest of the questions

**Theorem**[Perron-Frobenius] Let  $A$  be square positive. There exists an eigenvalue  $\rho$  of  $A$  such that

1.  $\rho$  is real and  $\rho > 0$
2.  $|\lambda| < \rho$  for any eigenvalue  $\lambda$  of  $A$  with  $\lambda \neq \rho$
3. there exists a positive eigenvector associated with  $\rho$
4. the algebraic multiplicity of  $\rho$  is 1 (so the geometric multiplicity of  $\rho$  is also 1)

A weaker result for general non-negative matrices:

**Theorem.** Let  $A$  be square non-negative. There exists an eigenvalue  $\rho$  of  $A$  such that

1.  $\rho$  is real and  $\rho \geq 0$
2.  $|\lambda| \leq \rho$  for any eigenvalue  $\lambda$  of  $A$
3. there exists a non-negative eigenvector associated with  $\rho$



- ▶ further implication by the above theorem:
  - a non-negative solution to  $Av = v$  exists, though it doesn't say if there exists another solution
  - even worse, it is not known if there exists another solution  $v$  such that  $v \not\geq 0$
- ▶ PageRank actually considers a modified version of  $A$

$$\tilde{A} = (1 - \beta)A + \beta \begin{bmatrix} 1/n & \dots & 1/n \\ \vdots & & \vdots \\ 1/n & \dots & 1/n \end{bmatrix}$$

where  $0 < \beta < 1$  (typical value is  $\beta = 0.15$ )

- ▶  $\tilde{A}$  is positive
- ▶ further implications
  - $\lambda = 1$  is the *only* eigenvalue that has the largest modulus
  - there exists *only* one eigenvector associated with  $\lambda = 1$ ; that eigenvector is either positive or negative