

SI231 - Matrix Computations, Fall 2020-21

Homework Set #3

Prof. Yue Qiu and Prof. Ziping Zhao

Name: Tian Yun **Major:** Master in CS

Student No.: 2019232102 **E-mail:** tianyun@shanghaitech.edu.cn

Acknowledgements:

- 1) Deadline: **2020-11-01 23:59:00**
 - 2) Submit your homework at **Gradescope**. Entry Code: **MY3XBJ**. Homework #3 contains two parts, the theoretical part and the programming part.
 - 3) About the theoretical part:
 - (a) Submit your homework in **Homework 3** in gradescope. Make sure that you have correctly select pages for each problem. If not, you probably will get 0 point.
 - (b) Your homework should be uploaded in the **PDF** format, and the naming format of the file is not specified.
 - (c) You need to use \LaTeX in principle.
 - (d) Use the given template and give your solution in English. Solution in Chinese is not allowed.
 - 4) About the programming part:
 - (a) Submit your codes in **Homework 3 Programming part** in gradescope.
 - (b) Detailed requirements see in Problem 2 and Problem 3.
 - 5) **No late submission is allowed.**
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I. UNDERSTANDING PROJECTION

Problem 1. (5 points \times 3)

Suppose that $\mathbf{P} \in \mathbb{R}^{n \times n}$ is a projector onto a subspace \mathcal{U} along its orthogonal complement \mathcal{U}^\perp , then it is called the **orthogonal projector** onto \mathcal{U} .

- 1) Prove that an orthogonal projector must be singular if it is not an identity matrix.
- 2) What is the orthogonal projector onto \mathcal{U}^\perp along the subspace \mathcal{U} ?
- 3) Let \mathcal{U} and \mathcal{W} be two subspaces of a vector space \mathcal{V} , and denote $\mathbf{P}_\mathcal{U}$ and $\mathbf{P}_\mathcal{W}$ as the corresponding orthogonal projectors, respectively. Prove that $\mathbf{P}_\mathcal{U}\mathbf{P}_\mathcal{W} = 0$ if and only if $\mathcal{U} \perp \mathcal{W}$.

Solution.

- 1) **proof:** A projection on a vector space V is a linear operator $P : V \mapsto V$ such that $P^2 = P$. $P \in \mathbb{R}^{n \times m}$

$$\therefore P(I - P) = 0, \text{ and } \text{rank}(P(I - P)) = 0.$$

$$\therefore \text{rank}(P) + \text{rank}(I - P) - n \leq 0$$

$$\therefore \text{rank}(P) + \text{rank}(I - P) \leq n$$

$$\text{First case : } \text{rank}(P) = n, \quad \therefore \text{rank}(I - P) = 0, \therefore I = P$$

$$\text{Second case: } \text{rank}(P) \leq n, \quad \therefore P \text{ must be singular.}$$

- 2) **I - P.**

$\mathbf{P} \in \mathbb{R}^{n \times n}$ is a projector onto a subspace \mathcal{U} along its orthogonal complement \mathcal{U}^\perp

so for all $x \in \mathcal{U}$, $\mathbf{P}x = x$, for all $x \in \mathcal{U}^\perp$, $\mathbf{P}x = 0$

Now we want to find the orthogonal projector \mathbf{P}^* onto \mathcal{U}^\perp along the subspace \mathcal{U}

it means for all $y \in \mathcal{U}^\perp$, $\mathbf{P}^*y = y$, for all $y \in \mathcal{U}$, $\mathbf{P}^*y = 0$

$$\mathbf{P}^* = \mathbf{I} - \mathbf{P}.$$

because for all $x \in \mathcal{U}^\perp$, $(\mathbf{I} - \mathbf{P})x = x$ and for all $x \in \mathcal{U}$, $(\mathbf{I} - \mathbf{P})x = 0$, satisfied with the analysis above.

- 3) **proof:**

On the one hand, $\forall x_u \in \mathcal{U}, \exists x \in \mathcal{V}$ s.t. $x_u = \mathbf{P}_\mathcal{U}x \quad \forall x_w \in \mathcal{W}, \exists y \in \mathcal{V}$ s.t. $y_w = \mathbf{P}_\mathcal{W}y$.

if $\mathbf{P}_\mathcal{U}\mathbf{P}_\mathcal{W} = 0$, and $\mathbf{P}_\mathcal{U}x_u = 0$, $\mathbf{P}_\mathcal{W}y_w = 0$ so $x_u^T \mathbf{P}_\mathcal{U}^T \mathbf{P}_\mathcal{W} x_w = 0$

So for any $z_u \in \mathcal{U}$ and because for any $z_w \in \mathcal{W}$ $z_u^T \perp z_w$

$$\therefore \mathcal{U} \perp \mathcal{W}$$

On the other hand, if $\mathcal{U} \perp \mathcal{W}$ So for any $z_u \in \mathcal{U}$ and for any $z_w \in \mathcal{W}$ $z_u^T z_w = 0$

so for any $z_u \in \mathcal{U}$, we can find a vector $x \in V$ s.t. $\mathbf{P}_\mathcal{U}x = z_u$

and for any $z_w \in \mathcal{W}$ we can find a vector $y \in V$ s.t. $\mathbf{P}_\mathcal{W}y = z_w$

$$\therefore z_u^T z_w = 0, \text{ therefore } (\mathbf{P}_\mathcal{U}x)^T \mathbf{P}_\mathcal{W}y = 0$$

$$\text{So } x^T \mathbf{P}_\mathcal{U}^T \mathbf{P}_\mathcal{W} y = 0$$

$$\therefore \mathbf{P}_\mathcal{U}^T \mathbf{P}_\mathcal{W} = 0$$

II. LEAST SQUARE (LS) PROGRAMMING.

Problem 2. (10 points + 10 points + 5 points)

Write programs to solve the least square problem with specified methods, any programming language is suitable.

$$\mathbf{x} = \arg \min_{\mathbf{x} \in \mathbb{R}^n} f(\mathbf{x}), \quad f(\mathbf{x}) = \|\mathbf{y} - \mathbf{Ax}\|_2^2$$

where $\mathbf{A} \in \mathbb{R}^{m \times n}$ is a matrix representing the predefined data set with m data samples of n dimensions ($m=1000$, $n=210$), and $\mathbf{y} \in \mathbb{R}^m$ represents the labels. The data samples are provided in the "data.txt" file, and the labels are provided in the "label.txt" file, you are supposed to load the data before solving the problem.

- 1) Solve the LS with gradient decent method.

The gradient descent method for solving problem updates \mathbf{x} as

$$\mathbf{x} = \mathbf{x} - \gamma \cdot \nabla_{\mathbf{x}} f(\mathbf{x}),$$

where γ is the step size of the gradient decent methods. We suggest that you can set $\gamma = 1e - 5$.

- 2) Solve the LS by the method of normal equation with Cholesky decomposition and forward/backward substitution.
- 3) Compare two methods above.
 - (a) Basing on the true running results from the program, count the number of "flops";
 - (b) Compare gradient norm and loss $f(\mathbf{x})$ for results $\mathbf{x} = \mathbf{x}_{LS}$ of above two algorithms.

Notation*: "flop": one flop means one floating point operation, i.e., one addition, subtraction, multiplication, or division of two floating-point numbers, in this problem each floating points operation $+$, $-$, \times , \div , $\sqrt{\cdot}$ counts as one "flop".

Hint for gradient decent programming:

- 1) **Step size selection:** to ensure the convergence of the method, γ is supposed to be selected properly (large step size may accelerate the convergence rate but also may lead to instability, A sufficiently small compensation always ensures that the algorithm converges).
- 2) **Terminal condition:** the gradient decent is an iteration algorithm that need a terminal condition. In this problem, the algorithm can stop when the gradient of the loss function $f(\mathbf{x})$ at current \mathbf{x} is small enough.

Remarks:

- The solution of the two methods should be printed in files named "sol1.txt" and "sol2.txt" and submitted in gradescope. The format should be same as the input file (210 rows plain text, each rows is a dimension of the final solution).
- Make sure that your codes are executable and are consistent with your solutions.

Solution.

- 1).The first part you can check in my uploaded file.
- 2). (1).Gradient descent method:

The stop condition used here is that the change value of cost is less than $1e-6$, and a total of 85 iterations are performed.

Each iteration needs to be calculated:

1. $X * \theta$
2. $X' * (predictions - y)$
3. $\theta = \theta - \alpha * updates$

Assuming that C iterations are required to complete, the time complexity of each iteration is $O(n * k + k)$, so the total time complexity is $O(C(n * k + k))$ flops

(2). Normal Equations:

1. $L = chol(X' * X);$
2. $temp = forwardsub(L', X' * y);$
3. $\theta = backwardsub(L, temp);$

chol: A total of $O(n^2)$ elements are calculated, and the calculation of each element contains the multiplication of $O(n)$, so it is $O(n^3)$

For forward substitution and backward substitution algorithms, since L or L' is a triangular matrix, the first one can be calculated in the order of y_1, y_2, \dots , the required calculation The complexity is $O(n^2)$, Similarly, the third one can also be calculated like this.

So the total computational complexity is $O(n^3 + n^2) = O(n^3)$ flops. The advantage of this algorithm is that when it is necessary to solve different x with relatively different b for the same A , the new part can be completed with $O(n^2)$ complexity.

III. UNDERSTANDING THE QR FACTORIZATION

Problem 3 [Understanding the Gram-Schmidt algorithm.]. (5 points + 7 points + 6 points + 7 points)

1) Consider the subspace \mathcal{S} spanned by $\{\mathbf{a}_1, \dots, \mathbf{a}_4\}$, where

$$\mathbf{a}_1 = \begin{bmatrix} 1 \\ 2 \\ 3 \\ 4 \end{bmatrix}, \quad \mathbf{a}_2 = \begin{bmatrix} 2 \\ 3 \\ 4 \\ 5 \end{bmatrix}, \quad \mathbf{a}_3 = \begin{bmatrix} 3 \\ 4 \\ 5 \\ 6 \end{bmatrix}, \quad \mathbf{a}_4 = \begin{bmatrix} 3 \\ 5 \\ 7 \\ 11 \end{bmatrix}.$$

Use the **classical** Gram-Schmidt algorithm (See Algorithm 1), find a set of orthonormal basis $\{\mathbf{q}_i\}$ for \mathcal{S} by hand (derivation is expected). Do not use decimals in your answers, fraction and n -th roots of numbers are accepted. Verify the orthonormality of the found basis.

Algorithm 1: Classical Gram-Schmidt algorithm

Input : A collection of linearly independent vectors $\mathbf{a}_1, \dots, \mathbf{a}_n$.

1 **Initilization:** $\tilde{\mathbf{q}}_1 = \mathbf{a}_1, \mathbf{q}_1 = \tilde{\mathbf{q}}_1 / \|\tilde{\mathbf{q}}_1\|_2$

2 **for** $i = 2, \dots, n$ **do**

3 $\tilde{\mathbf{q}}_i = \mathbf{a}_i - \sum_{j=1}^{i-1} (\mathbf{q}_j^T \mathbf{a}_i) \mathbf{q}_j$

4 $\mathbf{q}_i = \tilde{\mathbf{q}}_i / \|\tilde{\mathbf{q}}_i\|_2$

5 **end**

Output: $\mathbf{q}_1, \dots, \mathbf{q}_n$

2) Orthogonal projection of vector \mathbf{a} onto a nonzero vector \mathbf{b} is defined as

$$\text{proj}_{\mathbf{b}}(\mathbf{a}) = \frac{\langle \mathbf{a}, \mathbf{b} \rangle}{\langle \mathbf{b}, \mathbf{b} \rangle} \mathbf{b},$$

where $\langle \cdot, \cdot \rangle$ denotes the inner product of vectors. And for subspace \mathcal{M} with orthonormal basis $\{\mathbf{u}_1, \dots, \mathbf{u}_k\}$, the orthogonal projector onto subspace \mathcal{M} is given by

$$\mathbf{P} = \mathbf{U}\mathbf{U}^T, \quad \mathbf{U} = [\mathbf{u}_1 | \dots | \mathbf{u}_k].$$

In the context of **projection of vector** and **projection onto subspace** respectively, can you give another two understandings of the classical Gram-Schmidt algorithm?

3) Consider the subspace \mathcal{S} spanned by $\{\mathbf{a}_1, \mathbf{a}_2, \mathbf{a}_3\}$,

$$\mathbf{a}_1 = \begin{bmatrix} 1 \\ \epsilon \\ \epsilon \end{bmatrix}, \quad \mathbf{a}_2 = \begin{bmatrix} 1 \\ \epsilon \\ 0 \end{bmatrix}, \quad \mathbf{a}_3 = \begin{bmatrix} 1 \\ 0 \\ \epsilon \end{bmatrix},$$

where ϵ is a small real number such that $1 + k\epsilon^2 = 1$ ($k \in \mathbb{N}^+$). First complete the pseudo algorithm in Algorithm 2. Then use the **classical** Gram-Schmidt algorithm and the **modified** Gram-Schmidt algorithm respectively, find two sets of basis for \mathcal{S} by hand (derivation is expected). Are the two sets of basis the same? If not, which one is the desired orthonormal basis? Report what you have found.

Algorithm 2: Modified Gram-Schmidt algorithm

Input : A collection of linearly independent vectors $\mathbf{a}_1, \dots, \mathbf{a}_n$.

1 **Initilization:**

2 *Complete your algorithm here...*

Output: $\mathbf{q}_1, \dots, \mathbf{q}_n$

4) **Programming part:** In this part, you are required to code both the **classical Gram-Schmidt** and the **modified Gram-Schmidt** algorithms. For $\epsilon = 1e-4$ and $\epsilon = 1e-9$ in sub-problem 2), give the outputs of two algorithms and calculate $\|\mathbf{Q}^T \mathbf{Q} - \mathbf{I}\|_F$, where $\mathbf{Q} = [\mathbf{q}_1, \mathbf{q}_2, \mathbf{q}_3]$.

Remarks:

- Coding languages are not restricted, but do not use built-in function such as [qr](#).
- When handing in your homework in gradescope, package all your codes into `your_student_id+hw3_code.zip` and upload. In the package, you also need to include a file named `README.txt/md` to clearly identify the function of each file.
- Make sure that your codes can run and are consistent with your solutions.

Solution.

$$1) \tilde{\mathbf{q}}_1 = \mathbf{a}_1 = \begin{bmatrix} 1 & 2 & 3 & 4 \end{bmatrix}^T.$$

$$\mathbf{q}_1 = \tilde{\mathbf{q}}_1 / \|\tilde{\mathbf{q}}_1\|_2 = \frac{1}{\sqrt{30}} \begin{bmatrix} 1 & 2 & 3 & 4 \end{bmatrix}^T$$

$$\tilde{\mathbf{q}}_2 = \mathbf{a}_2 - (\mathbf{q}_1^T \mathbf{a}_2) \mathbf{q}_1 = \begin{bmatrix} 2 & 3 & 4 & 5 \end{bmatrix}^T - \frac{1}{3} \begin{bmatrix} 4 & 8 & 12 & 16 \end{bmatrix}^T = \frac{1}{3} \begin{bmatrix} 2 & 1 & 0 & -1 \end{bmatrix}^T$$

$$\mathbf{q}_2 = \tilde{\mathbf{q}}_2 / \|\tilde{\mathbf{q}}_2\|_2 = \frac{1}{\sqrt{6}} \begin{bmatrix} 2 & 1 & 0 & -1 \end{bmatrix}^T$$

$\tilde{\mathbf{q}}_3 = \mathbf{a}_3 - (\mathbf{q}_1^T \mathbf{a}_3) \mathbf{q}_1 - (\mathbf{q}_2^T \mathbf{a}_3) \mathbf{q}_2 = 0$, So $(\mathbf{a}_1, \mathbf{a}_2, \mathbf{a}_3, \mathbf{a}_4)$ are not linearly independent. We can get the set of orthonormal basis from

$\mathbf{a}_1, \mathbf{a}_2, \mathbf{a}_4$

$$\tilde{\mathbf{q}}_3 = \mathbf{a}_4 - (\mathbf{q}_1^T \mathbf{a}_4) \mathbf{q}_1 - (\mathbf{q}_2^T \mathbf{a}_4) \mathbf{q}_2 = \frac{1}{5} \begin{bmatrix} 2 & -1 & -4 \end{bmatrix}^T$$

$$\mathbf{q}_3 = \tilde{\mathbf{q}}_3 / \|\tilde{\mathbf{q}}_3\|_2 = \frac{1}{\sqrt{30}} \begin{bmatrix} 2 & -1 & -4 & 3 \end{bmatrix}^T$$

$$\mathbf{q}_1^T * \mathbf{q}_2 = \frac{1}{6\sqrt{5}} * (2 + 2 + 0 - 4) = 0$$

$$\mathbf{q}_1^T * \mathbf{q}_3 = \frac{1}{30} * (2 - 2 - 12 + 12) = 0$$

$$\mathbf{q}_2^T * \mathbf{q}_3 = \frac{1}{6\sqrt{5}} * (4 - 1 + 0 - 3) = 0$$

2) In the context of projection of vector:

$$\tilde{\mathbf{q}}_i = \mathbf{a}_i - \sum_{j=1}^{i-1} \text{proj}_{\mathbf{q}}(\mathbf{a}_i)$$

In the context of projection onto subspace:

IV. SOLVING LS VIA QR FACTORIZATION AND NORMAL EQUATION

Problem 4 [Understanding the influence of the condition number to the solution.]. (4 points + 5 points + 4 points + 4 points + 3 points points)

Consider such two LS problems:

$$\begin{aligned} \min_{\mathbf{x} \in \mathbb{R}^n} \|\mathbf{Ax} - \mathbf{b}\|_2^2 \\ \min_{\mathbf{x} \in \mathbb{R}^n} \|\mathbf{Ax} - (\mathbf{b} + \delta\mathbf{b})\|_2^2 \end{aligned} \quad (1)$$

with $\mathbf{A} \in \mathbb{R}^{m \times n}$. For $\mathbf{b} = \begin{bmatrix} 1 & 3/2 & 3 & 6 \end{bmatrix}^T$ and $\delta\mathbf{b} = \begin{bmatrix} 1/10 & 0 & 0 & 0 \end{bmatrix}^T$,

1) Computing solution to the problem (1) via QR decomposition when

$$\mathbf{A} = \begin{bmatrix} 1 & 2 & 3 \\ 2 & 3 & 5 \\ 3 & 4 & 7 \\ 4 & 5 & 11 \end{bmatrix}.$$

2) For a full-rank matrix \mathbf{A} , consider the equation $\mathbf{Ax} = \mathbf{b}$, after adding some noise $\delta\mathbf{b}$ to \mathbf{b} , we have $\mathbf{A}(\mathbf{x} + \delta\mathbf{x}) = \mathbf{b} + \delta\mathbf{b}$, and then proof

$$\frac{1}{\|\mathbf{A}\| \|\mathbf{A}^\dagger\|} \frac{\|\delta\mathbf{b}\|}{\|\mathbf{b}\|} \leq \frac{\|\delta\mathbf{x}\|}{\|\mathbf{x}\|} \leq \|\mathbf{A}\| \|\mathbf{A}^\dagger\| \frac{\|\delta\mathbf{b}\|}{\|\mathbf{b}\|},$$

and give it a plain interpretation.

3) Computing the solutions to the two LS problems via the normal equation $\mathbf{A}^T \mathbf{Ax}_{LS} = \mathbf{A}^T \mathbf{b}$ when

$$\mathbf{A} = \begin{bmatrix} 1 & 2 & 2 \\ 2 & 2 & 2 \\ 3 & 3 & 3 \\ 1 & 1 & 0 \end{bmatrix}.$$

4) Computing the solutions to the two LS problems via the normal equation $\mathbf{A}^T \mathbf{Ax}_{LS} = \mathbf{A}^T \mathbf{b}$ when

$$\mathbf{A} = \begin{bmatrix} 1 & 1 & 1 \\ 1 & 2 & 4 \\ 1 & 3 & 9 \\ 1 & 4 & 16 \end{bmatrix}.$$

5) Compare the 2-norm condition number $\|\mathbf{A}\| \|\mathbf{A}^\dagger\|$ for \mathbf{A} in 3) and 4) and the influence on the solution to problem (1) resulted by the additional noise $\delta\mathbf{b}$.

Hint: Show the influence on the solution by $\frac{\|\delta\mathbf{x}\|}{\|\mathbf{x}\|}$.

Remarks: You can use MATLAB for some matrix computations (deviation is expected) in 3), 4), 5). Do not use decimals in your answers, fraction and n -th roots of numbers are accepted.

Solution.

1) 1) $\mathbf{A} = \begin{bmatrix} a_1 & a_2 & a_3 \end{bmatrix}$, we have

$$\begin{aligned}
q_1 &= \frac{a_1}{\|a_1\|} = \frac{\sqrt{30}}{30} \begin{bmatrix} 1 & 2 & 3 & 4 \end{bmatrix}^T \\
\bar{q}_2 &= a_2 - (q_1^T a_2) q_1 = \frac{1}{3} \begin{bmatrix} 2 & 1 & 0 & -1 \end{bmatrix}^T \quad q_2 = \frac{\bar{q}_2}{\|\bar{q}_2\|} = \frac{\sqrt{6}}{6} \begin{bmatrix} 2 & 1 & 0 & -1 \end{bmatrix}^T \\
\bar{q}_3 &= a_3 - (q_1^T a_3) q_1 - (q_2^T a_3) q_2 = \frac{1}{5} \begin{bmatrix} 2 & -1 & -4 & 3 \end{bmatrix}^T \quad q_3 = \frac{\bar{q}_3}{\|\bar{q}_3\|} = \frac{\sqrt{30}}{30} \begin{bmatrix} 2 & -1 & -4 & 3 \end{bmatrix}^T \\
r_{11} &= \|\bar{a}_1\| = \frac{\sqrt{30}}{30} \\
r_{22} &= \|\bar{q}_2\| = \frac{\sqrt{6}}{3} \\
r_{33} &= \|\bar{q}_3\| = \frac{\sqrt{30}}{5} \\
r_{12} &= q_1^T a_2 = \sqrt{30} \\
r_{13} &= q_1^T a_3 = \frac{13\sqrt{30}}{5} \\
r_{23} &= q_2^T a_3 = 0
\end{aligned}$$

$$\begin{aligned}
\mathbf{Q} &= \begin{bmatrix} \frac{\sqrt{30}}{30} & \frac{2\sqrt{6}}{6} & \frac{2\sqrt{30}}{300} \\ \frac{2\sqrt{30}}{30} & \frac{\sqrt{6}}{6} & \frac{-\sqrt{30}}{30} \\ \frac{3\sqrt{30}}{30} & 0 & \frac{-4\sqrt{30}}{30} \\ \frac{4\sqrt{30}}{30} & \frac{-\sqrt{6}}{6} & \frac{3\sqrt{30}}{30} \end{bmatrix} \\
\mathbf{R} &= \begin{bmatrix} \frac{\sqrt{30}}{30} & \frac{4\sqrt{30}}{3} & \frac{1.3\sqrt{30}}{5} \\ 0 & \frac{\sqrt{6}}{3} & 0 \\ 0 & 0 & \frac{\sqrt{30}}{5} \end{bmatrix}
\end{aligned}$$

2) since $x + \delta x = \mathbf{A}^\dagger(b + \delta b)$ and $x = \mathbf{A}^\dagger b$, we have $\delta x = \mathbf{A}^\dagger \delta b$,

$$\text{so } \|\delta x\| = \|\mathbf{A}^\dagger \delta b\| \leq \|\mathbf{A}^\dagger\| \|\delta b\|.$$

$$\text{since } \mathbf{A}x = b, \text{ we have } \frac{1}{\|x\|} \leq \frac{\|b\|}{\|x\|},$$

$$\text{so } \frac{\|\delta x\|}{\|x\|} \leq \|\mathbf{A}\| \|\mathbf{A}^\dagger\| \frac{\|\delta b\|}{\|b\|}.$$

Similarly, $\mathbf{A}(x + \delta x) = b + \delta b$, we have $\delta b = \mathbf{A} \delta x$,

$$\text{so } \|\delta b\| \leq \|\mathbf{A}\| \|\delta x\|, x = \mathbf{A}^\dagger b,$$

$$\text{we have } \|x\| \leq \|\mathbf{A}^\dagger\| \|b\|, \text{ so } \frac{1}{\|\mathbf{A}\| \|\mathbf{A}^\dagger\|} \frac{\|\delta b\|}{\|b\|} \leq \frac{\|\delta x\|}{\|x\|}$$

Above all

$$\frac{1}{\|\mathbf{A}\| \|\mathbf{A}^\dagger\|} \frac{\|\delta b\|}{\|b\|} \leq \frac{\|\delta x\|}{\|x\|} \leq \|\mathbf{A}\| \|\mathbf{A}^\dagger\| \frac{\|\delta b\|}{\|b\|},$$

3)

$$\mathbf{A}^T \mathbf{A} = \begin{bmatrix} 15 & 16 & 15 \\ 16 & 18 & 17 \\ 15 & 17 & 17 \end{bmatrix} \quad \mathbf{A}^T \mathbf{b} = \begin{bmatrix} 19 \\ 20 \\ 14 \end{bmatrix}$$

Solve the equation $\mathbf{A}^T \mathbf{A} x_{LSi} = \mathbf{A}^T b$, we have

$$\mathbf{x}_{LS} = \begin{bmatrix} \frac{11}{13} \\ \frac{67}{13} \\ -\frac{66}{13} \end{bmatrix}$$

4)

$$\mathbf{A}^T \mathbf{A} = \begin{bmatrix} 4 & 10 & 30 \\ 10 & 30 & 100 \\ 30 & 100 & 357 \end{bmatrix} \quad \mathbf{A}^T \mathbf{b} = \begin{bmatrix} \frac{23}{2} \\ 37 \\ 130 \end{bmatrix}$$

Solve the equation $\mathbf{A}^T \mathbf{A} \mathbf{x}_{LS} = \mathbf{A}^T \mathbf{b}$, we have

$$\mathbf{x}_{LS} = \begin{bmatrix} \frac{15}{8} \\ -\frac{59}{40} \\ \frac{5}{8} \end{bmatrix}$$

5)

$$\mathbf{A}^\dagger = (\mathbf{A}^T \mathbf{A})^{-1} \mathbf{A}^T$$

$$\begin{aligned} \|\mathbf{A}_1\|_2 = 16.75 & \quad \left\| \mathbf{A}_1^\dagger \right\|_2 = 5.297 \rightarrow \|\mathbf{A}_1\|_2 \left\| \mathbf{A}_1^\dagger \right\|_2 = 88.895 \\ \|\mathbf{A}_2\|_2 = 19.621 & \quad \left\| \mathbf{A}_2^\dagger \right\|_2 = 3.756 \rightarrow \|\mathbf{A}_2\|_2 \left\| \mathbf{A}_2^\dagger \right\|_2 = 73.695 \end{aligned}$$

V. UNDERDETERMINED SYSTEM

Problem 5 [Solving Underdetermined System by QR]. (10 points + 5 points)

Consider the following underdetermined system $\mathbf{Ax} = \mathbf{b}$ with $\mathbf{A} \in \mathbb{R}^{m \times n}$ and $m < n$. Let

$$\mathbf{A} = \begin{bmatrix} 1 & 2 & 2 & 0 \\ 0 & -2 & 2 & 1 \\ 2 & 5 & 6 & 1 \end{bmatrix}, \mathbf{b} = \begin{bmatrix} b_1 \\ b_2 \\ b_3 \end{bmatrix},$$

- 1) Use Householder reflection to give the full QR decomposition of tall \mathbf{A}^T , i.e., $\mathbf{A}^T = \mathbf{QR}$ with \mathbf{Q} being a square matrix with orthonormal columns.
- 2) Give one possible solution via QR decomposition of \mathbf{A}^T , write down your solution using \mathbf{b} .

Solution.

1)

$$\mathbf{A}^T = \begin{bmatrix} 1 & 0 & 2 \\ 2 & -2 & 5 \\ 2 & 2 & 6 \\ 0 & 1 & 1 \end{bmatrix} \mathbf{e}_1 = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix} \mathbf{u}_1 = \begin{bmatrix} 1 \\ 2 \\ 2 \\ 0 \end{bmatrix} - 3 \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix} = \begin{bmatrix} -2 \\ 2 \\ 2 \\ 0 \end{bmatrix} \quad P_1 = \mathbf{I} - 2 \frac{\mathbf{u}_1 \mathbf{u}_1^T}{\mathbf{u}_1^T \mathbf{u}_1} = \frac{1}{3} \begin{bmatrix} 1 & 2 & 2 & 0 \\ 2 & 1 & -2 & 0 \\ 2 & -2 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

$$P_1 \mathbf{A}^T = \begin{bmatrix} 1 & 0 & 2 \\ 2 & -2 & 5 \\ 2 & 2 & 6 \\ 0 & 1 & 1 \end{bmatrix} \mathbf{e}_2 = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix} \mathbf{u}_2 = \begin{bmatrix} -2 \\ 2 \\ 2 \\ 1 \end{bmatrix} - 3 \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix} = \begin{bmatrix} -5 \\ 2 \\ 2 \\ 1 \end{bmatrix} \quad P_2^* = \mathbf{I} - 2 \frac{\mathbf{u}_2 \mathbf{u}_2^T}{\mathbf{u}_2^T \mathbf{u}_2}, \quad P_2 = \begin{bmatrix} 1 & 0 \\ 0 & P_2^* \end{bmatrix}$$

$$P_2 P_1 \mathbf{A}^T = \begin{bmatrix} 3 & 0 & 8 \\ 0 & 3 & 1 \\ 0 & 0 & \frac{4}{5} \\ 0 & 0 & \frac{3}{5} \end{bmatrix} \mathbf{e}_2 = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix} \mathbf{u}_3 = \begin{bmatrix} -\frac{4}{5} \\ \frac{3}{5} \\ \frac{3}{5} \end{bmatrix} - \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} = \begin{bmatrix} -\frac{9}{5} \\ \frac{3}{5} \\ \frac{3}{5} \end{bmatrix} \quad P_3^* = \mathbf{I} - 2 \frac{\mathbf{u}_3 \mathbf{u}_3^T}{\mathbf{u}_3^T \mathbf{u}_3}, \quad P_3 = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & P_3^* \end{bmatrix}$$

$$P_3 P_2 P_1 \mathbf{A}^T = \begin{bmatrix} 3 & 0 & 8 \\ 0 & 3 & 1 \\ 0 & 0 & 1 \\ 0 & 0 & 0 \end{bmatrix}, \mathbf{A}^T = P_1 P_2 P_3 \mathbf{R} = \mathbf{QR} = \begin{bmatrix} \frac{1}{3} & \frac{2}{3} & \frac{2}{3} & 0 \\ 0 & -\frac{2}{3} & \frac{2}{3} & \frac{1}{3} \\ -\frac{2}{3} & \frac{1}{3} & 0 & \frac{2}{3} \\ \frac{2}{3} & 0 & -\frac{1}{3} & \frac{2}{3} \end{bmatrix} \times \begin{bmatrix} 3 & 0 & 8 \\ 0 & 3 & 1 \\ 0 & 0 & 1 \\ 0 & 0 & 0 \end{bmatrix}$$

- 2) Solving Underdetermined Systems by QR For $\mathbf{A} \in \mathbb{R}^{m \times n}$ with $m < n$ and $\text{rank}(\mathbf{A}) = m$, we have

$$\mathbf{A}^T = \mathbf{QR} = \begin{bmatrix} \mathbf{Q}_1 & \mathbf{Q}_2 \end{bmatrix} \begin{bmatrix} \mathbf{R}_1 \\ \mathbf{0} \end{bmatrix} = \mathbf{Q}_1 \mathbf{R}_1 + \mathbf{Q}_2 \mathbf{0}$$

$$\mathbf{Ax} = \mathbf{R}_1^T \mathbf{Q}_1^T \mathbf{x} + \mathbf{0}^T \mathbf{Q}_2^T \mathbf{x} = \mathbf{b}$$

which indicates

$$\mathbf{Q}_1^T \mathbf{x} = \mathbf{R}_1^{-T} \mathbf{b}$$

and $Q_2^T x$ can be anything I set it to be 0

$$\begin{bmatrix} \mathbf{Q}_1^\top \mathbf{x} \\ \mathbf{Q}_2^\top \mathbf{x} \end{bmatrix} = \mathbf{Q}^T \mathbf{x} = \begin{bmatrix} \mathbf{R}_1^{-T} \mathbf{b} \\ \mathbf{0} \end{bmatrix}$$

hence, one possible solution is

$$R_1^{-T} b = \begin{bmatrix} 3 & 0 & 8 \\ 0 & 3 & 1 \\ 0 & 0 & 1 \end{bmatrix}^{-1} b = \begin{bmatrix} \frac{1}{3}b_1 - \frac{8}{3}b_3 \\ \frac{1}{3}b_2 - \frac{1}{3}b_3 \\ b_3 \\ 0 \end{bmatrix}$$

$$x = Q \begin{bmatrix} R_1^{-T} b \\ 0 \end{bmatrix} = \begin{bmatrix} \frac{1}{3} & \frac{2}{3} & \frac{2}{3} & 0 \\ 0 & -\frac{2}{3} & \frac{2}{3} & \frac{1}{3} \\ -\frac{2}{3} & \frac{1}{3} & 0 & \frac{2}{3} \\ \frac{2}{3} & 0 & -\frac{1}{3} & \frac{2}{3} \end{bmatrix} \times \begin{bmatrix} \frac{1}{3} & 0 & -\frac{8}{3} & 0 \\ 0 & \frac{1}{3} & \frac{1}{3} & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix} \begin{bmatrix} b_1 \\ b_2 \\ b_3 \\ b_4 \end{bmatrix} = \begin{bmatrix} \frac{1}{9} & \frac{2}{9} & -\frac{2}{3} & \frac{2}{3} \\ 0 & -\frac{2}{9} & -\frac{2}{9} & \frac{2}{3} \\ -\frac{2}{9} & \frac{1}{9} & \frac{17}{9} & \frac{2}{3} \\ \frac{2}{9} & 0 & -\frac{16}{9} & -\frac{1}{3} \end{bmatrix} \begin{bmatrix} b_1 \\ b_2 \\ b_3 \\ b_4 \end{bmatrix}$$

VI. SOLVING LS VIA PROJECTION

Problem 6. (Bonus question, 6 points + 4 points)

Consider the Least Square (LS) problem:

$$\min_{\mathbf{x} \in \mathbb{R}^n} \|\mathbf{A}\mathbf{x} - \mathbf{y}\|_2^2 \quad (2)$$

where $\mathbf{A} \in \mathbb{R}^{m \times n}$ ($m > n$) may not be full rank. Denote

$$X_{\text{LS}} = \{\mathbf{x} \in \mathbb{R}^n | \mathbf{A}^T \mathbf{A} \mathbf{x} = \mathbf{A}^T \mathbf{y}\}$$

as the set of all solutions to (2), and

$$\mathbf{x}_{\text{LS}} = \mathbf{A}^\dagger \mathbf{y}$$

where $\mathbf{A}^\dagger \in \mathbb{R}^{n \times m}$ is the *pseudo inverse* of \mathbf{A} satisfies the following properties:

- 1) $\mathbf{A} \mathbf{A}^\dagger \mathbf{A} = \mathbf{A}$.
- 2) $\mathbf{A}^\dagger \mathbf{A} \mathbf{A}^\dagger = \mathbf{A}^\dagger$.
- 3) $(\mathbf{A} \mathbf{A}^\dagger)^T = \mathbf{A} \mathbf{A}^\dagger$.
- 4) $(\mathbf{A}^\dagger \mathbf{A})^T = \mathbf{A}^\dagger \mathbf{A}$.

Answer the following questions:

- 1) Prove that \mathbf{x}_{LS} is a solution to (2) and is of minimum 2-norm in X_{LS} , that is

$$\mathbf{x}_{\text{LS}} = \arg \min_{\mathbf{x} \in X_{\text{LS}}} \|\mathbf{x}\|_2 .$$

Hint. Notice that the orthogonal projection onto $\mathcal{N}(A)$ is given by

$$\mathbf{\Pi}_{\mathcal{N}(A)} = \mathbf{I} - \mathbf{A}^\dagger \mathbf{A}$$

- 2) Prove that $X_{\text{LS}} = \{\mathbf{x}_{\text{LS}}\}$ if and only if $\text{rank}(\mathbf{A}) = n$.

Solution. Please insert your solution here ...