



CSC 355 Database Systems

Lecture 9

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Today:

- ◆ Transactions

Problem: Interruptions

- ◆ Two SQL statements are written to transfer money from one bank account to another...
- ◆ ...one executes, then the server crashes
 - What happens to the money?

Problem: Concurrency

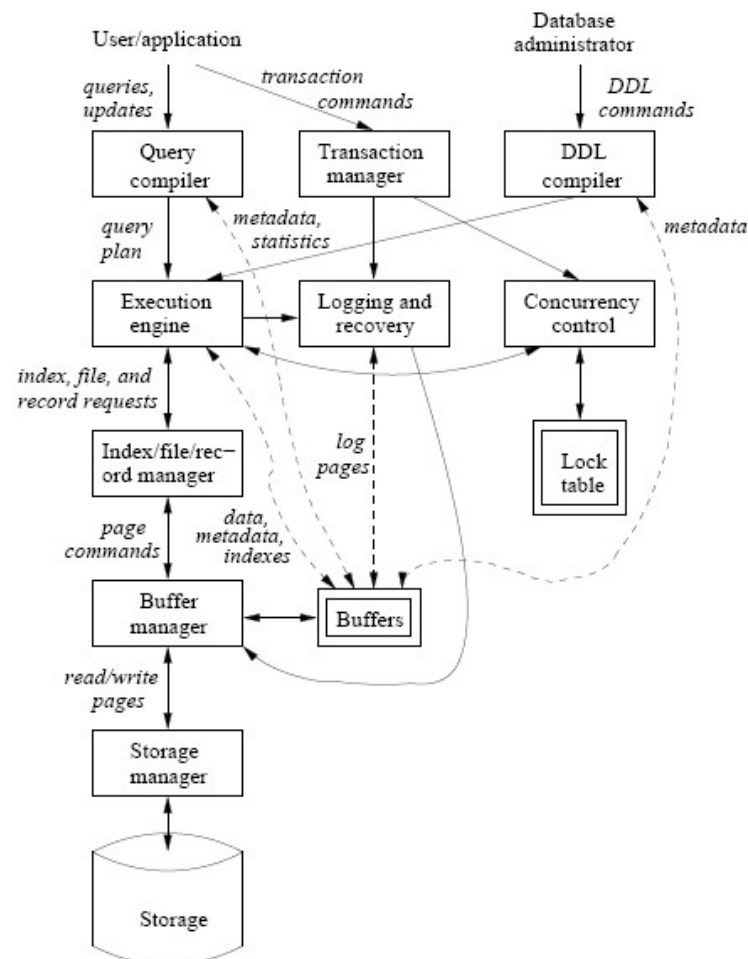
- ◆ Two users check the balance of the same bank account...
- ◆ ...then both try to transfer money out of it
 - Who gets it?

Solution: Transactions

- ◆ A *transaction* is a collection of SQL statements that must be executed as a unit
- ◆ The Transaction manager in the DBMS must handle
 - Interruptions: Logging and recovery, Buffers
 - Concurrency: Concurrency control, Lock tables

Components of a DBMS

(From Ullman/Widom)



Transactions in Oracle

- ◆ Any operation that changes the database state starts an implicit transaction in Oracle
- ◆ Can also start a transaction explicitly with a `SET TRANSACTION` statement
- ◆ End a transaction with a `COMMIT` statement
 - Transaction can also be ended with `ROLLBACK`, usually done by system rather than user...

ACID Properties

- ◆ A transaction should satisfy the following properties:
 - Atomicity: Executes completely or not at all
 - Consistency: Satisfies all database constraints
 - Isolation: Executes “separately” from others
 - Durability: Once executed, results are permanent

Atomicity

- ◆ Transaction operations are kept in a local store, not applied to the database immediately
 - Transaction can see its own changes, others can't
- ◆ When a transaction is completed, COMMIT applies changes to the shared database in their entirety (“executes completely...”)
- ◆ ROLLBACK during a transaction undoes any partial results (“...or not at all”)



Durability

- ◆ After COMMIT, the changes applied to the shared database are permanent, and cannot be rolled back later

Consistency

- ◆ Constraints can be “deferred”, so that they are only checked when a transaction commits, not for each individual statement
 - Add DEFERRABLE INITIALLY DEFERRED to the constraint definition
 - If a constraint is violated when you COMMIT, a ROLLBACK of the entire transaction is done!

Isolation

- ◆ DBMS maintains “separation” among transactions that access data concurrently
 - Various different levels of isolation
 - Transactions might modify, or just read, data
 - Tradeoff between performance and data integrity

Serializable Isolation

- ◆ Transactions must behave as though they were run serially (first one, then the other)
- ◆ Usually implemented by “locking” the tables (or parts of tables) used by a transaction
 - Other transactions using the same tables will have to wait until they are released
 - Other transactions using other tables could run at the same time

Read Committed Isolation

- ◆ Transaction operations can be interleaved
- ◆ If one transaction tries to read data that were written by another, it can only see the changes that have been committed
 - Can't be rolled back, but could be changed later
 - Multiple queries of the same table might not yield the same results... “non-repeatable reads”, including “phantoms”...

Read Uncommitted Isolation

- ◆ Transaction operations can be interleaved
- ◆ If one transaction tries to read data that were written by another, it can see all changes, even if they have not been committed
 - Could be rolled back by the other transaction!
 - Transaction might make decisions based on values that are later rolled back and so were never really there ... “dirty reads”...

Isolation Levels

- ◆ **SERIALIZABLE:** Transactions must appear to run serially – cannot read any changes from others
- ◆ **[REPEATABLE READ:** Can read committed changes that only add data (allows only phantoms)]
- ◆ **READ COMMITTED:** Can read all committed changes (allows all non-repeatable reads)
- ◆ **READ UNCOMMITTED:** Can read all changes (allows all non-repeatable reads and dirty reads)

Transactions in Oracle

- ◆ SET TRANSACTION statement can specify ISOLATION LEVEL
 - READ COMMITTED (default)
 - SERIALIZABLE
 - Oracle does not support REPEATABLE READ and READ UNCOMMITTED
- ◆ COMMIT or ROLLBACK ends transaction



Next:

- ◆ Finish Transactions
- ◆ Review for Midterm Exam
- ◆ Introduction to Relational Database Design