# 33085625\_FIT3182\_Assignment\_1

April 9, 2024

## 1 FIT3182 - Assignment 1 - Questions Workbook

This assignment consists of **four questions** that are, taken together, worth 10% of the final marks. **Please treat this Assignment as a take-home test. This is an individual assignment and you should complete the questions on your own**. You should not attempt to post questions on EdForum seeking solutions to the answers. If you require clarification on the Assignment questions, you can post a private post on EdForum or seek consultation from the tutors.

- 1. The first question is related to Parallel Search Algorithms (2 Mark).
- 2. The second question is related to Parallel Join Algorithms (4 Marks).
- 3. The third question is related to Parallel Sort Algorithms (2 Marks).
- 4. The fourth question is related to Parallel GroupBy Algorithms (2 Marks).

Instructions: - You will be using Python 3. Answer all questions inside this Jupyter Notebook. Please use the provided Docker to load the Jupyter Notebook. - Read the instructions, skeleton code, and comments carefully. - There are code blocks that you need to complete yourself as a part of the assignment. - You are also required to answer the questions below. - Comment each line of code properly such that the tutor can easily understand what you are trying to do in the code. Marks may be deducted for insufficient or unclear comments. - Once completed, please rename the Jupyter Notebook to include your Student ID at the beginning of the filename (e.g., 12345678\_FIT3182\_Assignment\_1.ipynb). Submit this Jupyter Notebook into the Assignment 1 submission link in Moodle. Please refer to the Assignment 1 submission link (or page) in Moodle for information about the submission due date.

Please include your details as follows: - Student ID: 33085625 - Last Name: Foo - First (or Preferred) Name: Kai Yan - Monash Student Email: kfoo0012@student.monash.edu

```
[1]: from pprint import pprint
```

## 1.1 Dataset

For this test, we will use the following tables R and S to write solutions to three parallel algorithms.

## 1.2 1. Parallel Searching Algorithm (2 marks)

This section consist of 4 sub-questions.

In this task, you will build a **parallel search algorithm for range selection (continuous)** for a given query.

Implement a parallel search algorithm that uses local linear search (i.e., linear\_search()) and is able to work with the hash partitioning method (i.e., h\_partition()). Complete the code block between ### START CODE HERE ### and ### END CODE HERE ###.

#### 1.2.1 1.1 Local Search

```
[3]: # Linear search function
     def linear_search(data, key):
         Perform linear search on data for the given key.
         Arguments:
         data -- an input dataset (list OR numpy array)
         key -- an query attribute (number)
         Return:
         result -- list of matched records such as [('Adele', 8)]
                     or the position of searched record
                     whichever is more appropriate for your code
         11 11 11
         matches = []
         match record = None
         position = -1 # not found position
         ### START CODE HERE ###
         for record in data:
             # Because data is in tuple form so we extract data from tuple with this
             name, number = record
             if number == key: # If number is matched with key
                 match_record = (name, number)
                 position = data.index(record)
                 # Add the matched record as a tuple in a list
                 matches.append(match_record)
                 break
         ### END CODE HERE ###
```

```
return matches
```

```
[4]: results = linear_search(R, 20) print(results)
```

[('Lim', 20)]

#### 1.2.2 1.2 Hash Partition

```
[5]: # Define a simple hash function.
def s_hash(x, n):
    """
    Define a simple hash function for demonstration

Arguments:
    x -- an input record attribute (int)
    n -- the number of processors (int)

Return:
    result -- the hash value of x
    """
    return x % n
```

```
[6]: # Hash data partitioning function.
     # We will use the s_hash function defined above
     def h_partition(data, n):
         n n n
         Perform hash data partitioning on data
         Arguments:
         data -- an input dataset (list)
         n -- the number of processors (int)
         result -- the partitioned subsets of D
         partitions = {}
         ### START CODE HERE ###
         for record in data:
             # Because data is in tuple form so we extract data from tuple with this
             name, number = record
             h = s_hash(number, n) # Get the hash key of the input
             if h in partitions:
                 partitions[h].append((name, number))
             else:
```

```
partitions[h] = [(name, number)]
### END CODE HERE ###
return partitions
```

```
[7]: partitions = h_partition(R, 5)
pprint(partitions)
```

```
{0: [('Fung', 25), ('Lim', 20), ('Noor', 5)],
1: [('Clement', 16), ('Ed', 11), ('Kelly', 6), ('Meng', 1)],
2: [('Bob', 22), ('Harry', 17), ('Joanna', 2)],
3: [('Adele', 8), ('Dave', 23), ('Goel', 3)],
4: [('Irene', 14), ('Omar', 19)]}
```

#### 1.2.3 1.3 Parallel Search

```
[8]: from multiprocessing import Pool
     def collect_result(results, result):
         # Callback for when linear_search returns a result.
         # The results list is modified only by the main process, not the pool \Box
      →workers.
         results.append(result)
     # Parallel searching algorithm for range selection
     def parallel_search_range(data, query_range, n_processor):
         Perform parallel search for range selection on data for the given key.
         Arguments:
         data -- the input dataset (list)
         query_range -- a tuple with inclusive range end-points (e.g. [30, 50] means⊔
      \hookrightarrow the interval 30-50)
         n_processor -- the number of parallel processors (int)
         results -- the matched record information
         pool = Pool(processes=n_processor)
         ### START CODE HERE ###
         results = []
         # Perform data partitioning first
```

```
data_partitioning = h_partition(data, n_processor)

for query in range(query_range[0], query_range[1], 1):
    # Each element in data_partitioning has a pair (hash key: records)
    query_hash = s_hash(query, n_processor)
    partitioned_data = list(data_partitioning[query_hash])
    result = pool.apply(linear_search, [partitioned_data, query])
    collect_result(results, result)

### END CODE HERE ###

return results
```

```
[9]: n_processor = 3
results = parallel_search_range(R, [5, 15], n_processor)
pprint(results)
```

```
[[('Noor', 5)],
  [('Kelly', 6)],
  [],
  [('Adele', 8)],
  [],
  [('Ed', 11)],
  [],
  [],
  [('Irene', 14)]]
```

## 1.2.4 1.4 Parallel Search Variants

#### Briefly answer the following.

1. How would the parallel search you implemented need to be changed if we switched to round-robin partition for the data? Provide a small snippet of code with the modified search loop.

**Your Answers**: First I would have specifically implement the round-robin partition in another function and perform the round-robin partition for the data then save it in data\_partitioning variable then only perform parallel search.

```
index_bin = (int) (name % n_processor)
    result[index_bin].append(number)

return result

def round_robin_partition_parallel_search(data, query_range, n_processor):
    pool = Pool(processes=n_processor)
    results = []

# Perform data partitioning first
    data_partitioning = round_robin_partition(data, n_processor)

for query in range(query_range[0], query_range[1], 1):
    # Each element in data_partitioning has a pair (hash key: records)
    query_hash = s_hash(query, n_processor)
    partitioned_data = list(data_partitioning[query_hash])
    result = pool.apply(linear_search, [partitioned_data, query])
    collect_result(results, result)

return result
```

```
[11]: n_processor = 3
results = round_robin_partition_parallel_search(R, [12, 15], n_processor)
pprint(results)
```

[('Irene', 14)]

## 1.3 2. Parallel Join Algorithm (4 marks)

This section consist of 5 sub-questions.

In this task, you will implement a **disjoint-partitioning based parallel join algorithm**. This algorithm consist of two stages. 1. Data disjoint-partitioning across processors. 1. Local join in each processor.

As a data partitioning method, use the range partitioning algorithm (i.e. range\_partition()). Assume that we have **3 parallel processors**, processor 1 will get records with join attribute value between 1 and 9, processor 2 between 10 and 19, and processor 3 between 20 and 29. Note that both tables R and S need to be partitioned using the same ranges described earlier.

As a joining technique, use the hash based join algorithm (i.e., HB\_join()). Complete the code block between ### START CODE HERE ### and ### END CODE HERE ###.

#### 1.3.1 2.1 Range Partition

```
[12]: # Range data partitionining function (Need to modify as instructed above)

def range_partition(data, range_indices):

"""

Perform range data partitioning on data based on the join attribute
```

```
Arguments:
          data -- an input dataset (list)
          range_indices -- the range end-points (e.g., [5, 10, 15] means
                               4 intervals: x < 5, 5 <= x < 10, 10 <= x < 15, 15 <= x)
          Return:
          result -- the partitioned subsets of D
          result = []
          ### START CODE HERE ###
          # First, we sort the dataset according their values
          sort_by_value = sorted(data, key=lambda x: x[1])
          # Calculate the number of bins
          bins_number = len(range_indices)
          # For each bin, perform the following
          for i in range(bins_number):
              # Find elements to be belonging to each range
              partitioned_list = [x for x in sort_by_value if x[1] < range_indices[i]]</pre>
              # Add the partitioned list to the result
              result.append(partitioned list)
              # Find the last element in the previous partition
              last_element = partitioned_list[len(partitioned_list)-1]
              # Find the index of of the last element
              last = sort_by_value.index(last_element)
              # Remove the partitioned list from the dataset
              sort_by_value = sort_by_value[int(last)+1:]
              # Append the last remaining data list
          result.append([x for x in sort_by_value if x[1] \geq=_\(\text{\text{\text{or}}}
       →range_indices[bins_number-1]])
          ### END CODE HERE ###
          return result
[13]: pprint(range_partition(R, [10, 20]))
     [[('Meng', 1),
       ('Joanna', 2),
       ('Goel', 3),
```

('Noor', 5), ('Kelly', 6), ('Adele', 8)],

```
[('Ed', 11), ('Irene', 14), ('Clement', 16), ('Harry', 17), ('Omar', 19)],
[('Lim', 20), ('Bob', 22), ('Dave', 23), ('Fung', 25)]]
```

```
1.3.2 2.2 Hash Join
[14]: def H(r):
          We define a hash function 'H' that is used in the hashing process works
          by summing the first and second digits of the hashed attribute, which
          in this case is the join attribute.
          Arguments:
          r -- a record where hashing will be applied on its join attribute
          Return:
          result -- the hash index of the record r
          # Convert the value of the join attribute into the digits
          digits = [int(d) for d in str(r[1])]
          # Calulate the sum of elemenets in the digits
          return sum(digits)
[15]: def HB_join(T1, T2):
          Perform the hash-based join algorithm.
          The join attribute is the numeric attribute (in second position) in the \Box
       →input tables T1 & T2.
          Arguments:
          T1 & T2 -- Tables to be joined
          Return:
          result -- the joined table
          11 11 11
          result = []
```

dic = {} # We will use a dictionary

# For each record in table T2

for s in T2:

*into hash table* 

else:

s\_key = H(s)
if s\_key in dic:

# Hash the record based on join attribute value using hash function  $H_{f L}$ 

dic[s\_key].add(s) # If there is an entry

```
dic[s_key] = {s}
print(dic)
# For each record in table T1 (probing)
for r in T1:
    # Hash the record based on join attribute value using H
    r_key = H(r)
    # If an index entry is found Then
    if r_key in dic:
        # Compare each record on this index entry with the record of table_u

>T1

for item in dic[r_key]:
    if item[1] == r[1]:
        # Put the rsult
        result.append((r[0], r[1], item[0]))

return result
```

```
[16]: # print the partitioned result
HB_join(R,S)
```

```
{8: {('Arts', 8)}, 6: {('Business', 15)}, 2: {('Health', 11), ('CompSc', 2)}, 3: {('Finance', 21), ('Dance', 12)}, 7: {('Engineering', 7)}, 1: {('Geology', 10)}, 9: {('IT', 18)}}

[16]: [('Adele', 8, 'Arts'), ('Ed', 11, 'Health'), ('Joanna', 2, 'CompSc')]
```

#### [10]. [( Addie , 0, Artis ), ( Ed , II, Reditin ), ( Johnson , 2, Johnson )

#### 1.3.3 2.3 Parallel Join

```
[17]: # Include this package for parallel processing import multiprocessing as mp

def consolidate_result(results, result):
    # This is called whenever HB_Join(T1, T2) returns a result.
    # The results list is modified only by the main process, not the pool
    →workers.
    results.append(result)

def DPBP_join(T1, T2, n_processor):
    """

Perform a disjoint partitioning-based parallel join algorithm.
    The join attribute is the numeric attribute in the input tables T1 & T2

Arguments:
    T1 & T2 -- Tables to be joined
    n_processor -- the number of parallel processors

Return:
```

```
### START CODE HERE ###

# Partition T1 & T2 into sub-tables using range_partition().
# The number of the sub-tables must be the equal to the n_processor
T1_subsets = range_partition(T1, [10, 20])
T2_subsets = range_partition(T2, [10, 20])

# Apply local join for each processor
pool = mp.Pool(processes = n_processor)
for i in range(len(T1_subsets)):
    result = pool.apply_async(HB_join, [T1_subsets[i], T2_subsets[i]])
    output = result.get()
    results.append(output)

### END CODE HERE ###
return results
```

## 1.3.4 2.4 Duplication Outer Join Algorithm (DOJA)

In this task, you will implement a **Duplication Outer Join Algorithm (DOJA)**. This algorithm consist of four main steps. 1. (Replication): Duplicate the small table 2. (Local Inner Join): Perform inner join in each processor 3. (Distribution): Reshuffle the inner join result based on R.x 4. (Local Outer Join): Perform outer join between the initial table R and the inner join result

Assume n\_processor=3 and the dataset is initially partition using the hash method. For joining technique, please use the provided outer join function for the inner join and outer join operation. Complete the code block between ### START CODE HERE ### and ### END CODE HERE ###.

```
[20]: import numpy as np
      def H(r):
          We define a hash function 'H' that is used in the hashing process works
          by summing the first and second digits of the hashed attribute, which
          in this case is the join attribute.
          Arguments:
          r -- a record where hashing will be applied on its join attribute
          result -- the hash index of the record r
          # Convert the value of the join attribute into the digits
          digits = [int(d) for d in str(r[1])]
          # Calulate the sum of elemenets in the digits
          return sum(digits)
      def outer_join(L, R, join="left"):
          """outer join using Hash-based join algorithm"""
          if join == "right":
              L, R = R, L
          # inner join
          if join == "inner":
              h dic = \{\}
              for r in R:
                  h r = H(r)
                  if h_r in h_dic.keys():
                      h_dic[h_r].add(r)
                  else:
                      h_dic[h_r] = \{r\}
              result = []
              for 1 in L:
                  h_1 = H(1)
                  if h_l in h_dic.keys():
                      for item in h_dic[h_l]:
```

```
if item[1] == 1[1]:
                    result.append([1[0], item[1], item[0]])
    return result
elif join in ["left", "right"]:
    h_dic = {}
    for r in R:
        print(r)
        h_r = H(r)
        if h_r in h_dic.keys():
            h_dic[h_r].add(r)
            h_dic[h_r] = \{r\}
    print(h_dic)
    result = []
    for l in L:
        isFound = False # to check whether there is a match found.
        h_1 = H(1)
        if h_l in h_dic.keys():
            for item in h_dic[h_l]:
                if item[1] == 1[1]:
                    result.append([1[0], item[1], item[2]])
                    isFound = True
                    break
        if not isFound:
            result.append([1[0], 1[1],str(np.nan)])
    return result
else:
    raise AttributeError('join should be in {left, right, inner}.')
```

```
[21]: def hash_distribution(T, n):
    """distribute data using hash partitioning"""

# Define a simple hash function for demonstration
    def s_hash(x, n):
        h = x%n
        return h

result = {}
    for t in T:
```

```
h_key = s_hash(t[1], n)
if h_key in result.keys():
    result[h_key].add(tuple(t))

else:
    result[h_key] = {tuple(t)}

return result
```

```
[22]: import multiprocessing as mp
      def DOJA(L, R, n):
          """left outer join using DOJA"""
          ### Start CODE
          output = []
          \# Compare to see which is the shorter and longer list between R and S
          # Temporary place inputted list then may change later
          longer_list = R # S list
          shorter_list = L # R list
          if len(L) < len(R):</pre>
              longer_list = R
              shorter_list = L
              print("List L is shorter than List R")
          elif len(L) > len(R):
              longer_list = L
              shorter_list = R
              print("List R is shorter than List L")
              print("Both lists are equal length") # Can randomly assign either of □
       ⇔list to be the 'shorter' side
          # Distribute longer list to n number of processors using hash distribution
          left_distribution = hash_distribution(longer_list, n)
          print("left distribution", left_distribution)
          \# Create a new distribution with the shorter list added to each processor \sqcup
       → like in a form of [duplicated short list][left distribution]
          duplicated_list = []
          for dup in range(n):
              duplicated_list.append(shorter_list)
          # INNER JOIN
          pool = mp.Pool(n)
```

```
results = []
  for i in left_distribution.keys():
      result = pool.apply_async(outer_join, [left_distribution[i],__

¬duplicated_list[i], "inner"])
      # use async cus you want them to perform outer join at the same time
      results.append(result)
  # REDISTRIBUTE
  inner_peace = []
  for x in results:
      inner_peace.extend(x.get())
  # Reshuffle the inner join result
  resuffled = hash_distribution(inner_peace, n)
  print("Resuffled: ", resuffled)
  # Get the key
  the key = 0
  for key in resuffled:
      the_key = key
  # local outer join and combine
  outputed = []
  for i in left_distribution.keys():
      outputs = pool.apply_async(outer_join, [left_distribution[i],_
→resuffled[the_key], "left"])
      print("left_distribution", left_distribution)
      print("resuffled", resuffled)
      # use async cus you want them to perform outer join at the same time in_
\hookrightarrow parallel
      outputed.append(outputs)
  for y in outputed:
      output.extend(y.get())
  ### End CODE
  return output
```

#### [23]: DOJA(R,S,3)

```
List R is shorter than List L
left distribution {2: {('Dave', 23), ('Adele', 8), ('Ed', 11), ('Noor', 5),
    ('Harry', 17), ('Joanna', 2), ('Irene', 14), ('Lim', 20)}, 1: {('Omar', 19),
    ('Fung', 25), ('Meng', 1), ('Clement', 16), ('Bob', 22)}, 0: {('Kelly', 6),
    ('Goel', 3)}}
```

```
('Joanna', 2, 'CompSc')('Joanna', 2, 'CompSc')
     ('Ed', 11, 'Health')('Ed', 11, 'Health')
     {8: {('Adele', 8, 'Arts')}, 2: {('Joanna', 2, 'CompSc'), ('Ed', 11,
     'Health')}}{8: {('Adele', 8, 'Arts')}, 2: {('Joanna', 2, 'CompSc'), ('Ed', 11,
     'Health')}}
     ('Adele', 8, 'Arts')
     ('Joanna', 2, 'CompSc')
     ('Ed', 11, 'Health')
     {8: {('Adele', 8, 'Arts')}, 2: {('Joanna', 2, 'CompSc'), ('Ed', 11, 'Health')}}
     Resuffled: {2: {('Adele', 8, 'Arts'), ('Joanna', 2, 'CompSc'), ('Ed', 11,
     'Health')}}
     left_distribution {2: {('Dave', 23), ('Adele', 8), ('Ed', 11), ('Noor', 5),
     ('Harry', 17), ('Joanna', 2), ('Irene', 14), ('Lim', 20)}, 1: {('Omar', 19),
     ('Fung', 25), ('Meng', 1), ('Clement', 16), ('Bob', 22)}, 0: {('Kelly', 6),
     ('Goel', 3)}}
     resuffled {2: {('Adele', 8, 'Arts'), ('Joanna', 2, 'CompSc'), ('Ed', 11,
     'Health')}}
     left_distribution {2: {('Dave', 23), ('Adele', 8), ('Ed', 11), ('Noor', 5),
     ('Harry', 17), ('Joanna', 2), ('Irene', 14), ('Lim', 20)}, 1: {('Omar', 19),
     ('Fung', 25), ('Meng', 1), ('Clement', 16), ('Bob', 22)}, 0: {('Kelly', 6),
     ('Goel', 3)}}
     resuffled {2: {('Adele', 8, 'Arts'), ('Joanna', 2, 'CompSc'), ('Ed', 11,
     'Health')}}
     left_distribution {2: {('Dave', 23), ('Adele', 8), ('Ed', 11), ('Noor', 5),
     ('Harry', 17), ('Joanna', 2), ('Irene', 14), ('Lim', 20)}, 1: {('Omar', 19),
     ('Fung', 25), ('Meng', 1), ('Clement', 16), ('Bob', 22)}, 0: {('Kelly', 6),
     ('Goel', 3)}}
     resuffled {2: {('Adele', 8, 'Arts'), ('Joanna', 2, 'CompSc'), ('Ed', 11,
     'Health')}}
[23]: [['Dave', 23, 'nan'],
       ['Adele', 8, 'Arts'],
       ['Ed', 11, 'Health'],
       ['Noor', 5, 'nan'],
       ['Harry', 17, 'nan'],
       ['Joanna', 2, 'CompSc'],
       ['Irene', 14, 'nan'],
       ['Lim', 20, 'nan'],
       ['Omar', 19, 'nan'],
       ['Fung', 25, 'nan'],
       ['Meng', 1, 'nan'],
       ['Clement', 16, 'nan'],
```

('Adele', 8, 'Arts')('Adele', 8, 'Arts')

```
['Bob', 22, 'nan'],
['Kelly', 6, 'nan'],
['Goel', 3, 'nan']]
```

#### 1.3.5 2.5 Parallel Join Variants

Briefly answer the following question.

1. For each partitioning algorithm besides range-partitioning, state and justify whether we can use it with the code for disjoint-partitioning based parallel join above.

Your answer: Yes. Other partitioning algorithm like hash-partitioning and round-robin partitioning can be used with disjoint-partitioning based parallel joins. There is another partitioning algorithm which is the random-unequal data partitionining method but this method should not be used with disjoint-partitioning based parallel joins.

For hash partitioning method, it distributes the tuples by implementing a hash function on the join key which is appropriate for parallel joins as it evenly spreads data among partitions, reducing the possibility of skewing. Disjoint-partitioning method can be used as long as the hash functions used in both hash partitioning and joining phases are compatible.

For round-robin partitioning method, it evenly distributes tuples among partitions by distributing them in a sequential manner whereby there is an equal spread of tuples among partitions and which is good for balancing workloads in parallel joins. It can work with disjoint-partitioning method because it is not dependent on the data value, which reduced the possibility of skewing.

For random-unequal data partitionining method, it does not evenly distributes tuples among partitions into disjoint steps which is the main requirement for disjoint-partitioning based parallel joins. In addition to that, if random-unequal data partitionining method is used, there would be a high possibility of skewing as the data are randomly partitioned and the partitions are most likely unequal partitions.

#### 1.4 3. Parallel Sorting Algorithm (2 marks)

In this task, you will implement **parallel binary-merge sort** method. It has two phases same as the parallel merge-all sort that you learnt in the labs. 1. Local sort 2. Final merge.

The first phase is similar to the parallel merge-all sort. The second phase, the merging phase, is pipelined instead of concentrating on one processor. In this phase, we take the results from two processors and then merge the two in one processor, called binary merging. The result of the merge between two processors is passed on to the next level until one processor (the host) is left.

Complete the code block between ### START CODE HERE ### and ### END CODE HERE ###. Assume that we use the round robin partitioning method (i.e. rr\_partition()).

#### 1.4.1 3.1 Round-robin Partition

```
[24]: # Round-robin data partitionining function

def rr_partition(data, n):

"""

Perform data partitioning on data
```

```
Arguments:
  data -- an input dataset which is a list
  n -- the number of processors
  Return:
  result -- the paritioned subsets of D
  result = []
  for i in range(n):
       result.append([])
  ### START CODE HERE ###
  # For each bin,
  for name, number in enumerate(data):
       # Calculate the index of the bin that the current data point will be ...
\hookrightarrow assigned
       index_bin = (int) (name % n)
       result[index_bin].append(number)
  ### END CODE HERE ###
  return result
```

```
[25]: # Test the round-robin partitioning function
result = rr_partition(R, 3)
pprint(result)
```

```
[[('Adele', 8), ('Dave', 23), ('Goel', 3), ('Joanna', 2), ('Meng', 1)],
[('Bob', 22), ('Ed', 11), ('Harry', 17), ('Kelly', 6), ('Noor', 5)],
[('Clement', 16), ('Fung', 25), ('Irene', 14), ('Lim', 20), ('Omar', 19)]]
```

#### 1.4.2 3.2 Serial Sort

```
[26]: def qsort(arr):
    """
    Quicksort a list

Arguments:
    arr -- the input list to be sorted

Return:
    result -- the sorted arr
    """
    if len(arr) <= 1:
        return arr</pre>
```

```
else:
              # take the first element as the pivot
              pivot = arr[0]
              left_arr = [x for x in arr[1:] if x[1] < pivot[1]]</pre>
              right_arr = [x for x in arr[1:] if x[1] >= pivot[1]]
              # uncomment this to see what to print
              # print("Left:" + str(left_arr)+" Pivot : "+ str(pivot)+" Right: " +
       ⇔str(right_arr))
              sorted_arr = qsort(left_arr) + [pivot] + qsort(right_arr)
              return sorted_arr
[27]: qsort(R)
[27]: [('Meng', 1),
       ('Joanna', 2),
       ('Goel', 3),
       ('Noor', 5),
       ('Kelly', 6),
       ('Adele', 8),
       ('Ed', 11),
       ('Irene', 14),
       ('Clement', 16),
       ('Harry', 17),
       ('Omar', 19),
       ('Lim', 20),
       ('Bob', 22),
       ('Dave', 23),
       ('Fung', 25)]
[28]: | # Let's first look at 'k-way merging algorithm' that will be used
      # to merge sub-record sets in our external sorting algorithm.
      import sys
      # Find the smallest record
      def find_min(records):
          Find the smallest record
          Arguments:
          records -- the input record set
          Return:
          result -- the smallest record's index
          m = records[0]
          index = 0
```

```
for i in range(len(records)):
        if(records[i][1] < m[1]):</pre>
            index = i
            m = records[i]
    return index
def k_way_merge(record_sets):
    11 11 11
    K-way merging algorithm
    Arguments:
    record_sets -- the set of multiple sorted sub-record sets
    Return:
    result -- the sorted and merged record set
    # indexes will keep the indexes of sorted records in the given buffers
    indexes = []
    for _ in record_sets:
        indexes.append(0) # initialisation with 0
    # final result will be stored in this variable
    result = []
    while(True):
        merged_result = [] # the merging unit (i.e. # of the given buffers)
        # This loop gets the current position of every buffer
        for i in range(len(record_sets)):
            if(indexes[i] >= len(record_sets[i])):
                merged_result.append((None, sys.maxsize))
            else:
                merged_result.append(record_sets[i][indexes[i]])
        # find the smallest record
        smallest = find_min(merged_result)
        # if we only have sys.maxsize on the tuple, we reached the end of everyu
 ⇔record set
        if(merged_result[smallest][1] == sys.maxsize):
            break
        # This record is the next on the merged list
        result.append(record_sets[smallest][indexes[smallest]])
        indexes[smallest] +=1
    return result
```

```
[29]: # Test k-way merging method
      buffers = rr_partition(R, 3)
      print(buffers)
      result = k_way_merge([qsort(b) for b in buffers])
      pprint(result)
     [[('Adele', 8), ('Dave', 23), ('Goel', 3), ('Joanna', 2), ('Meng', 1)], [('Bob',
     22), ('Ed', 11), ('Harry', 17), ('Kelly', 6), ('Noor', 5)], [('Clement', 16),
     ('Fung', 25), ('Irene', 14), ('Lim', 20), ('Omar', 19)]]
     [('Meng', 1),
      ('Joanna', 2),
      ('Goel', 3),
      ('Noor', 5),
      ('Kelly', 6),
      ('Adele', 8),
      ('Ed', 11),
      ('Irene', 14),
      ('Clement', 16),
      ('Harry', 17),
      ('Omar', 19),
      ('Lim', 20),
      ('Bob', 22),
      ('Dave', 23),
      ('Fung', 25)]
[30]: def serial_sorting(dataset, buffer_size):
          Perform a serial external sorting method based on sort-merge
          The buffer size determines the size of each sub-record set
          Arguments:
          dataset -- the entire record set to be sorted
          buffer_size -- the buffer size determining the size of each sub-record set
          Return:
          result -- the sorted record set
          11 11 11
          if (buffer_size <= 2):</pre>
              print("Error: buffer size should be greater than 2")
              return
          result = []
          ### START CODE HERE ###
          # Length of dataset
```

```
n = len(dataset)
  # --- Sort Phase ---
  # Sort and divide the dataset into subsets
  start_position = 0
  sorted_set = []
  while True:
      if (n - start_position) > buffer_size:
           subset = dataset[start_position: start_position + buffer_size]
           sorted_subset = qsort(subset)
           sorted set.append(sorted subset)
          start_position += buffer_size
      else:
          subset = dataset[start_position:]
          sorted_subset = qsort(subset)
           sorted_set.append(sorted_subset)
          break
  # --- Merge Phase ---
  merge_buffer_size = buffer_size - 1
  while True:
      merged_set = []
      n = len(sorted_set)
      start position = 0
      while True:
          if (n - start_position) > merge_buffer_size:
               subset = sorted_set[start_position: start_position +__
→merge_buffer_size]
               merged_set.append(k_way_merge(subset))
               start_position += merge_buffer_size
          else:
               subset = sorted_set[start_position:]
              merged_set.append(k_way_merge(subset))
              break
      sorted_set = merged_set
      if len(sorted_set) <= 1:</pre>
          result = merged_set
          break
  ### END CODE HERE ###
  return result
```

```
[31]: print(R)
```

```
[('Adele', 8), ('Bob', 22), ('Clement', 16), ('Dave', 23), ('Ed', 11), ('Fung',
```

## 1.4.3 3.3 Parallel Binary-merge Sort

```
[33]: # Include this package for parallel processing
      import multiprocessing as mp
      def parallel_binary_merge_sorting(dataset, n_processor, buffer_size):
          Perform a parallel binary-merge sorting method
          Arguments:
          dataset -- entire record set to be sorted
          n_processor -- number of parallel processors
          buffer_size -- buffer size determining the size of each sub-record set
          Return:
          result -- the merged record set
          11 11 11
          if (buffer_size <= 2):</pre>
              print("Error: buffer size should be greater than 2")
              return
          result = []
          ### START CODE HERE ###
          # Call rr partition to do round-robin partitioning
          round_robin = rr_partition(dataset, n_processor)
          # pool for parallel processing
          pool = mp.Pool(processes = n_processor)
          # Sort phase
          sorted_dataset = []
          for elem in round_robin:
              # call serial sorting method
              print("elem: ", elem)
```

```
sorted_dataset.append(*pool.apply_async(serial_sorting, [elem,_
⇒buffer_size]).get())
       print("sorted_dataset: ", sorted_dataset)
  pool.close()
  # Merge phase
  # For loop for binary merge
  while len(sorted_dataset) > 1:
       updated = []
       for i in range(0, len(sorted_dataset), 2):
           if i + 1 < len(sorted_dataset):</pre>
               merged = k_way_merge([sorted_dataset[i], sorted_dataset[i + 1]])
               updated.append(merged)
           else:
               # If there's an odd number of sorted datasets, just append the
→last one
               updated.append(sorted_dataset[i])
       sorted_dataset = updated
  # The final merged result
  result = sorted_dataset[0]
  ### END CODE HERE ###
  return result
```

```
[34]: result = parallel_binary_merge_sorting(R, 10, 20)
print("final result:" + str(result))
```

```
elem: [('Adele', 8), ('Kelly', 6)]
sorted_dataset: [[('Kelly', 6), ('Adele', 8)]]
elem: [('Bob', 22), ('Lim', 20)]
sorted_dataset: [[('Kelly', 6), ('Adele', 8)], [('Lim', 20), ('Bob', 22)]]
      [('Clement', 16), ('Meng', 1)]
sorted_dataset: [[('Kelly', 6), ('Adele', 8)], [('Lim', 20), ('Bob', 22)],
[('Meng', 1), ('Clement', 16)]]
elem: [('Dave', 23), ('Noor', 5)]
sorted_dataset: [[('Kelly', 6), ('Adele', 8)], [('Lim', 20), ('Bob', 22)],
[('Meng', 1), ('Clement', 16)], [('Noor', 5), ('Dave', 23)]]
elem: [('Ed', 11), ('Omar', 19)]
sorted_dataset: [[('Kelly', 6), ('Adele', 8)], [('Lim', 20), ('Bob', 22)],
[('Meng', 1), ('Clement', 16)], [('Noor', 5), ('Dave', 23)], [('Ed', 11),
('Omar', 19)]]
elem: [('Fung', 25)]
sorted_dataset: [[('Kelly', 6), ('Adele', 8)], [('Lim', 20), ('Bob', 22)],
[('Meng', 1), ('Clement', 16)], [('Noor', 5), ('Dave', 23)], [('Ed', 11),
('Omar', 19)], [('Fung', 25)]]
```

```
[('Goel', 3)]
elem:
                 [[('Kelly', 6), ('Adele', 8)], [('Lim', 20), ('Bob', 22)],
sorted_dataset:
[('Meng', 1), ('Clement', 16)], [('Noor', 5), ('Dave', 23)], [('Ed', 11),
('Omar', 19)], [('Fung', 25)], [('Goel', 3)]]
elem: [('Harry', 17)]
sorted_dataset: [[('Kelly', 6), ('Adele', 8)], [('Lim', 20), ('Bob', 22)],
[('Meng', 1), ('Clement', 16)], [('Noor', 5), ('Dave', 23)], [('Ed', 11),
('Omar', 19)], [('Fung', 25)], [('Goel', 3)], [('Harry', 17)]]
elem: [('Irene', 14)]
sorted_dataset: [[('Kelly', 6), ('Adele', 8)], [('Lim', 20), ('Bob', 22)],
[('Meng', 1), ('Clement', 16)], [('Noor', 5), ('Dave', 23)], [('Ed', 11),
('Omar', 19)], [('Fung', 25)], [('Goel', 3)], [('Harry', 17)], [('Irene', 14)]]
elem: [('Joanna', 2)]
sorted_dataset: [[('Kelly', 6), ('Adele', 8)], [('Lim', 20), ('Bob', 22)],
[('Meng', 1), ('Clement', 16)], [('Noor', 5), ('Dave', 23)], [('Ed', 11),
('Omar', 19)], [('Fung', 25)], [('Goel', 3)], [('Harry', 17)], [('Irene', 14)],
[('Joanna', 2)]]
final result: [('Meng', 1), ('Joanna', 2), ('Goel', 3), ('Noor', 5), ('Kelly',
6), ('Adele', 8), ('Ed', 11), ('Irene', 14), ('Clement', 16), ('Harry', 17),
('Omar', 19), ('Lim', 20), ('Bob', 22), ('Dave', 23), ('Fung', 25)]
```

#### 1.4.4 3.4 Parallel Binary-merge Behavior

Briefly answer the following.

1. Does parallel binary-merge utilize all available processors? State and justify with respect to each phase: sort, merge.

Your answer: Yes, parallel binary-merge do utilize all available processors whereby all the available processors can be used to sort different data subsets in parallel during the sorting phase and during the merge phase, the binary-merge algorithm merges two sorted lists at a time but when merging, the number of lists to merge reduces so less and less processors will be used so in conclusion, although all processors can be utilized at the sorting and merging phase, the merging phase may not fully use all the processors as it goes on as the workload has decreased as it goes on.

## 1.5 4 Parallel GroupBy (2 marks)

#### 1.5.1 4.1 Parallel GroupBy with Merge-All

In this task, you will implement **Parallel GroupBy with Merge-All** method. It involves the following two steps: 1. Local aggregate in each processor 2. Global aggregation

Assume n\_processor=4 and the dataset has already been pre-partitioned. For local aggregation, please use the provided local\_groupby function. Complete the code block between ### START CODE HERE ### and ### END CODE HERE ###.

```
[35]: D1 = [('A', 1), ('B', 2), ('C', 3), ('A', 10), ('B', 20), ('C', 30)]
D2 = [('A', 2), ('B', 3), ('C', 4), ('A', 20), ('B', 30), ('C', 40)]
D3 = [('A', 3), ('B', 4), ('C', 5), ('A', 30), ('B', 40), ('C', 50)]
D4 = [('A', 4), ('B', 5), ('C', 6), ('A', 40), ('B', 50), ('C', 60)]
```

```
[36]: # The first step in the merge-all groupby method
      def local_groupby(dataset):
          Perform a local groupby method
          Arguments:
          dataset -- entire record set to be merged
          result -- the aggregated record set according to the group_by attribute\sqcup
       \hookrightarrow index
          11 11 11
          dict = \{\}
          for index, record in enumerate(dataset):
              key = record[0]
              val = record[1]
              if key not in dict:
                   dict[key] = 0
              dict[key] += val
          return dict
[37]: result = local_groupby (D1)
      print(result)
     {'A': 11, 'B': 22, 'C': 33}
[38]: import multiprocessing as mp
      def parallel_merge_all_groupby(dataset):
          Perform a parallel merge_all groupby method
          Arguments:
          dataset -- entire record set to be merged
          result -- the aggregated record dictionary according to the group by_{\sqcup}
       \ominusattribute index
          11 11 11
          result = {}
          ### START CODE HERE ###
          # Define the number of parallel processors: the number of sub-datasets.
          n_processor = len(dataset)
```

```
pool = mp.Pool(processes=n_processor)
# ---- Local aggregation step ----
local_result = []
for s in dataset:
    # call the local aggregation method
    local_result.append(pool.apply(local_groupby, [s]))
pool.close()
# ---- Global aggregation step ----
# Assume that the global operator is sum
for r in local result:
    for key, val in r.items():
        if key not in result:
            result[key] = 0
        result[key] += val
### END CODE HERE ###
return result
```

```
[39]: E = [D1, D2, D3, D4]
result = parallel_merge_all_groupby (E)
print(result)
```

{'A': 110, 'B': 154, 'C': 198}

#### 1.5.2 4.2 Redistribution Method

In this task, you will implement **Parallel GroupBy with Redistribution Method** method. It invloves the following two steps: 1. (Partitioning phase): Redistribute raw records to all processor 2. (Aggregation phase): Each processor performs a local aggregation

Assume n\_processor=4 and the dataset should be partitioned into 4 groups using range partition Where first group with A & B, second with C & D, third with E & F and fourth with G & H. For local aggregation, please use the provided local\_groupby function. Complete the code block between ### START CODE HERE ### and ### END CODE HERE ###.

```
[40]: D = [('A', 1), ('B', 2), ('C', 3), ('D', 10), ('E', 20), ('F', 30), ('G', 20), \( \text{'H'}, 30), \( \text{'H'}, 30), \( \text{'A'}, 2), \text{('B'}, 3), \( \text{'C'}, 4), \text{('D'}, 20), \( \text{'E'}, 30), \text{('F'}, 40), \( \text{'G'}, 30), \( \text{'H'}, 40), \( \text{'A'}, 3), \( \text{'B'}, 4), \( \text{'C'}, 5), \( \text{'D'}, 30), \( \text{'E'}, 40), \( \text{'F'}, 50), \( \text{'G'}, 40), \) \( \text{'H'}, 50), \( \text{'A'}, 4), \( \text{'B'}, 5), \( \text{'C'}, 6), \( \text{'D'}, 40), \( \text{'E'}, 50), \( \text{'F'}, 60), \( \text{'G'}, 50), \) \( \text{'H'}, 60)]
```

## Create range partition function based on GroupBy attribute

```
[41]: # Range data partitionining function (Need to modify as instructed above)
      def range_partition(data, range_indices):
          Perform range data partitioning on data based on the join attribute
          Arguments:
           data -- an input dataset (list of tuple)
          range_indices -- the range end-points (e.q., ['C', 'E', 'G'] means
                                4 intervals: x < 'C', 'C' \le x < 'E', 'E' \le x < 'G',
       \hookrightarrow 'G' <= x)
          Return:
          result -- the partitioned subsets of D
          result = []
          ### START CODE HERE ###
          # First, we sort the dataset according to their values
          new data = sorted(data, key=lambda x: x[1])
          # Calculate the number of bins
          n_bin = len(range_indices) + 1
          # For each bin, perform the following
          for i in range(n_bin):
               # Find elements belonging to each range
               if i == 0:
                   s = [x \text{ for } x \text{ in new_data if } x[0] \text{ in } ('A', 'B')]
               elif i == n_bin - 1:
                   s = [x \text{ for } x \text{ in new_data if } x[0] \text{ in } ('G', 'H')]
               else:
                   s = [x for x in new_data if range_indices[i-1] <= x[0] <__
        →range_indices[i]]
               # Add the partitioned list to the result
               result.append(s)
          ### END CODE HERE ###
          return result
```

```
[42]: import multiprocessing as mp

def parallel_redistributed_groupby(dataset, range_indices):
```

```
Perform a parallel redistributed groupby method
          Arguments:
          dataset -- entire record set to be merged
          Return:
          result -- the aggregated record dictionary according to the group_by_{\sqcup}
       \ominusattribute index
          11 11 11
          result = []
          ### START CODE HERE ###
          # call range partition
          ranged = range_partition(dataset, range_indices)
          # Define the number of parallel processors: the number of sub-datasets.
          n_processor = len(dataset)
          pool = mp.Pool(processes=n_processor)
          # ---- Local aggregation step ----
          result = []
          for each in ranged:
              # call the local aggregation method
              result.append(pool.apply(local_groupby, [each]))
          pool.close()
          ### END CODE HERE ###
          return result
[43]: result = parallel_redistributed_groupby (D, ['C', 'E', 'G'])
      print(result)
     [{'A': 10, 'B': 14}, {'C': 18, 'D': 100}, {'E': 140, 'F': 180}, {'G': 140, 'H':
     180}]
 []:
```