

CONVERSIONS ($B_2 \leftrightarrow B_{10} \leftrightarrow B_{16}$)

65₁₀ to binary...

$$\begin{array}{r} 2 \overline{) 65} \\ 2 \overline{) 32} \ 1 \\ 2 \overline{) 16} \ 0 \\ 2 \overline{) 8} \ 0 \\ 2 \overline{) 4} \ 0 \\ 2 \overline{) 2} \ 0 \\ 2 \overline{) 1} \ 0 \\ 0 \ 1 \end{array}$$
 ANS: 1000001 ✓

binary \rightarrow hexadecimal
 (2) \rightarrow (16)
 4 bits \rightarrow 1 bit

group binary to 4 x 4 bits

11 0010 1110₂

11	0010	1110
3	2	14
3	2	E

$\Rightarrow 32E_{16}$

hexadecimal \rightarrow binary

C4D₁₆

C	4	D
12	4	13
1100	0100	1101

hexadecimal \rightarrow decimal
 16 \rightarrow 10

C4D₁₆

op		
x16	C	12 = 192
+4	4	4 = 196
x16		= 3136
+13	D	13 = 3149 #

ANS: 3149₁₀

1019₁₀

$$\begin{array}{r} 16 \overline{) 1019} \\ 16 \overline{) 63} \ 11 \\ 16 \overline{) 3} \ 15 \\ 0 \ 3 \end{array}$$
 ANS: 3FB₁₆ #

SIGN - MAGNITUDE / 1's & 2's complement

X calculation

Waste single bit for sign

\rightarrow waste space

Drawbacks
 (sign-magnitude)

decimal number = 5

unsigned = 0101

signed = 0101

1's = 0101

2's = 0101

decimal number = -5

unsigned =

signed = 1101

1's = 0101 \rightarrow 1010

2's = 0101 \rightarrow +1 \rightarrow 1011

2's complement (4-bit)

Q1) 3+4

$$3_{10} = 0011_2$$

$$4_{10} = 0100_2$$

$$\begin{array}{r} 0011 \\ + 0100 \\ \hline 0111 \end{array} \rightarrow 2^0 + 2^1 + 2^2 = 1 + 2 + 4 = 7 \neq$$

Q2) 5-2 = 5+(-2)

$$5_{10} = 0101$$

$$-2_{10} = 0010 \rightarrow 1101 + 1 \rightarrow 1110$$

$$\begin{array}{r} 0101 \\ + 1110 \\ \hline 10011 \end{array} \rightarrow 0011 \#$$

\hookrightarrow it's 4-bit so...
the 10011
 \hookrightarrow this can be removed

HOW / WHEN will an overflow occur

- (+ve) + (+ve) = (-ve)
- (-ve) + (-ve) = (+ve)

WHY?

- no enough to represent real answer in 4 bits

SINGLE PRECISION (IEEE 754 FLOATING-POINT STANDARD)



1) convert to binary exponent form:

$$1593_{10} = 110\ 0011\ 1001_2$$

$$= 1.1000111001 \times 2^{10}$$

2) compute sign-bit positive = 0; negative = 1:

0

3) mantissa is whatever after the decimal point

4) Compute exponent using excess-k method:

$$k = 2^4 - 1 = 127 \quad e = 10$$

8-bit exponent

$$= k + e$$

$$= 127 + 10$$

$$= 137_{10}$$

Error detection

↳ send signal to one computer to another

3 ways:

- Parity bits

↳ add a single bit

↳ drawbacks:

cannot detect multiple bits error

even parity $\Rightarrow 0$

odd parity $\Rightarrow 1$

- Checksum

- Cyclic Redundancy Checks (CRC)