

Binary \rightarrow Unsigned :

$$\begin{array}{r} 0 \ 0 \ 0 \ 1 \\ \times \ 2^3 \ 2^2 \ 2^1 \ 2^0 \\ \hline 0 \ 0 \ 0 \ 1 \\ \hline \therefore 1 \end{array}$$

there's no -ve 0 (-0) in math so
binary 1000 \rightarrow unsigned is 0

* 2's complement is always 1 value smaller than 1's complement *

Represent the number -9210 in...

(i) 8-bit signed magnitude

signed = -9210

unsigned = 9210

$$\begin{array}{r} 2 \overline{) 92} \\ 2 \overline{) 46} \ 0 \\ 2 \overline{) 23} \ 0 \\ 2 \overline{) 11} \ 1 \\ 2 \overline{) 5} \ 1 \\ 2 \overline{) 2} \ 1 \\ 2 \overline{) 1} \ 0 \\ 2 \overline{) 0} \ 1 \end{array}$$

unsigned \Rightarrow 01011100

~~unsigned = 00100011~~ x

signed = 11011100 x ✓

(1 = negative)

(ii) 8-bit 2's complement

$$\begin{array}{r} 1's = 00100011 \\ 2's = 00100011 \\ + 1 \\ \hline 00100100 \end{array}$$

(10100011 + 1)
= 10100100

convert the following numbers to 6-bit

2's complement notation & add them :

(i) -16 + 11 (ii) 16 - 11

$$\begin{array}{l} 16 \rightarrow 010000 \\ 11 \rightarrow 001011 \\ -16 \rightarrow 101111 + 1 = 110000 \\ -11 \rightarrow 110100 + 1 = 110101 \end{array}$$

$$\begin{array}{r} (i) \ 110000 \\ + 001011 \\ \hline 111011 \end{array} \quad \begin{array}{r} (ii) \ 010000 \\ + 110101 \\ \hline 100101 \end{array}$$

* 2's complement can discard the extra bit

Because...

1's \rightarrow 2's need +1 already
save the process to create
memory to store the discarded 1

Binary arithmetic with 8-bit 2's complement :

(i) 33 + 92

$$\begin{array}{r} 2 \overline{) 33} \quad 2 \overline{) 92} \\ 2 \overline{) 16} \ 1 \quad 2 \overline{) 46} \ 0 \\ 2 \overline{) 8} \ 0 \quad 2 \overline{) 23} \ 0 \\ 2 \overline{) 4} \ 0 \quad 2 \overline{) 11} \ 1 \\ 2 \overline{) 2} \ 0 \quad 2 \overline{) 5} \ 1 \\ 2 \overline{) 1} \ 0 \quad 2 \overline{) 2} \ 1 \\ 2 \overline{) 0} \ 1 \quad 2 \overline{) 1} \ 0 \\ \quad \quad \quad 2 \overline{) 0} \ 1 \end{array}$$

100001

1011100

1's :

011110

0100011

2's :

011111

0100100

(ii) 33 - 92 \Rightarrow 33 + (-92)

33 2's = 1011111

92 2's = 0100100

-92 2's = 1100100

$$\begin{array}{r} 1011111 \\ + 1100100 \\ \hline 1100011 \end{array}$$

CORRECTION : 33 \Rightarrow 00100001

92 \Rightarrow 01011100

-92 \Rightarrow 10100011 + 1 \rightarrow 10100100

$$(i) \ 00100001 + 01011100 \\ = 01111101$$

$$(ii) \ 00100001 \\ + 10100100 \\ \hline 11000101$$

4-bit binary number	signed	1's	2's
0000	0	0	0
0001	1	1	1
0010	2	2	2
0011	3	3	3
0100	4	4	4
0101	5	5	5
0110	6	6	6
0111	7	7	7
1000	0	-7	-8
:	:	:	:
:	:	:	:
1111	-7	0	-1

	range	* n = number of bits
unsigned	0 → $2^n - 1$	
signed i's	$-(\frac{2^n}{2} - 1) < 0 < (\frac{2^n}{2} - 1)$	
2's	$-(\frac{2^n}{2}) < 0 < (\frac{2^n}{2} - 1)$	

* overflow = out-of-range

↳ How computer detect:

ANS: - + 2 positive number results in -ve number

- + 2 negative number results in +ve number

What are the ranges of number that can be represented in a computer if it's using...

(i) 4-bit 2's complement representation

ANS: $-8 < 0 < +7$

(ii) 6-bit 2's complement representation

ANS: $-32 < 0 < 31$

(iii) 8-bit 2's complement representation

ANS: $-128 < 0 < 127$

Use binary arithmetic with 8-bit 2's complement representation to calculate $92 + 92$. result of operation accepted?

ANS: $92 \rightarrow 01011100$

01011100
+ 01011100

10111000 (overflow)

↳ results of operation is NOT ACCEPTED.

$92 + 92 = 184$

184 is outside the range of numbers that can be represented in a computer using 8-bit 2's complement representation.

