

Carnegie Mellon University

CMU 2

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$\underline{\text{Contest}}$ (1)

template.cpp

#include <bits/stdc++.h>

31 lines

```
using namespace std;
#define f first
#define s second
#define pb push_back
#define mp make pair
#define all(v) v.begin(), v.end()
#define sz(v) (int)v.size()
#define MOO(i, a, b) for(int i=a; i <b; i++)
#define M00(i, a) for(int i=0; i<a; i++)</pre>
#define MOOd(i,a,b) for(int i = (b)-1; i \ge a; i--)
#define M00d(i,a) for (int i = (a)-1; i >= 0; i--)
#define FAST ios::sync_with_stdio(0); cin.tie(0);
#define finish(x) return cout << x << '\n', 0;</pre>
#define dbg(x) cerr << ">>>> " << #x << " = " << x << "\n";</pre>
#define _ << " _ " <<
typedef long long 11;
typedef long double ld;
typedef vector<int> vi;
typedef pair<int,int> pi;
typedef pair<ld,ld> pd;
typedef complex<ld> cd;
int main() { FAST
```

set nocp bs=indent,eol,start nu ru si ts=4 sw=4 sts=0 sta \hookrightarrow et is hls sm mouse=a syntax on filetype plugin indent on

${\it cppreference.} txt$

```
atan(m) -> angle from -pi/2 to pi/2
atan2(y,x) -> angle from -pi to pi
acos(x) -> angle from 0 to pi
asin(y) -> angle from -pi/2 to pi/2
lower_bound -> first element >= val
upper_bound -> first element > val
```

<u>Data Structures</u> (2)

2.1 STL

MapComparator.h

CustomHash.h

 $\textbf{Description:} \ \ \text{faster than standard unordered map}$

```
struct chash {
  static uint64_t splitmix64(uint64_t x) {
    // http://xorshift.di.unimi.it/splitmix64.c
```

OrderStatisticTree.h

Description: A set (not multiset!) with support for finding the n'th element, and finding the index of an element.

Time: $\mathcal{O}(\log N)$

Rope.h

Description: insert element at n-th position, cut a substring and reinsert somewhere else

Time: $\mathcal{O}(\log N)$ per operation? not well tested

```
<ext/rope> 13 lines
using namespace __gnu_cxx;

void ropeExample() {
  rope<int> v(5, 0);
```

```
FOR(i, sz(v)) v.mutable reference at(i) = i+1; // or
rope<int> cur = v.substr(1,2); v.erase(1,2);
FOR(i,sz(v)) cout << v[i] << " "; // 1 4 5
cout << "\n";
v.insert(v.mutable_begin()+2,cur);
for (rope<int>::iterator it = v.mutable begin(); it != v
  cout << *it << " "; // 1 4 2 3 5
cout << "\n";
```

LineContainer.h

Description: Given set of lines, computes greatest y-coordinate for

```
Time: \mathcal{O}(\log N)
struct Line {
 mutable 11 k, m, p; // slope, y-intercept, last optimal
    \hookrightarrow x
 11 eval (11 x) { return k*x+m; }
 bool operator<(const Line& o) const { return k < o.k; }</pre>
 bool operator<(ll x) const { return p < x; }</pre>
struct LC : multiset<Line,less<>>> {
 // for doubles, use inf = 1/.0, div(a,b) = a/b
 const ll inf = LLONG_MAX;
 ll div(ll a, ll b) { return a/b-((a^b) < 0 \&\& a\%b); } //

→ floored division

 ll bet (const Line& x, const Line& y) { // last x such
    \hookrightarrowthat first line is better
    if (x.k == y.k) return x.m >= y.m? inf : -inf;
   return div(y.m-x.m,x.k-y.k);
 bool isect(iterator x, iterator y) { // updates x->p,
    \hookrightarrow determines if y is unneeded
   if (y == end()) { x->p = inf; return 0; }
   x->p = bet(*x,*y); return x->p >= y->p;
  void add(ll k, ll m) {
   auto z = insert(\{k, m, 0\}), y = z++, x = y;
   while (isect(y, z)) z = erase(z);
   if (x != begin() \&\& isect(--x, y)) isect(x, y = erase(
    while ((y = x) != begin() \&\& (--x)->p >= y->p) isect(x
       \hookrightarrow, erase(y));
  11 query(11 x) {
   assert(!empty());
   auto 1 = *lb(x);
   return l.k*x+l.m;
};
```

2.2 1D Range Queries

RMQ.h

Description: 1D range minimum query **Time:** $\mathcal{O}(N \log N)$ build, $\mathcal{O}(1)$ query

template<class T> struct RMQ { constexpr static int level(int x) { return 31-__builtin_clz(x); } // floor(log_2(x)) vector<vi> jmp; vector<T> v; int comb(int a, int b) { return v[a] == v[b] ? min(a,b) : (v[a] < v[b] ? a : b)} // index of minimum void init(const vector<T>& _v) { $v = _v; jmp = {vi(sz(v))}; iota(all(jmp[0]), 0);$ for (int j = 1; 1 << j <= sz(v); ++j) { jmp.pb(vi(sz(v)-(1<<j)+1));FOR(i,sz(jmp[j])) jmp[j][i] = comb(jmp[j-1][i],jmp[j-1][i+(1<<(j-1))]);int index(int 1, int r) { // get index of min element int d = level(r-l+1);return comb(jmp[d][1],jmp[d][r-(1<<d)+1]); T query(int 1, int r) { return v[index(1,r)]; }

};

Description: N-D range sum query with point update

Time: $\mathcal{O}\left((\log N)^D\right)$

```
19 lines
template <class T, int ...Ns> struct BIT {
 T val = 0:
 void upd(T v) { val += v; }
 T query() { return val; }
template <class T, int N, int... Ns> struct BIT<T, N, Ns
  BIT<T, Ns...> bit[N+1];
 template<typename... Args> void upd(int pos, Args...
   for (; pos <= N; pos += (pos&-pos)) bit[pos].upd(args</pre>
  template<typename... Args> T sum(int r, Args... args) {
   T res = 0; for (; r; r -= (r\&-r)) res += bit[r].query(
       \hookrightarrowargs...);
   return res;
 template<typename... Args> T query(int 1, int r, Args...
    return sum(r,args...)-sum(1-1,args...);
```

```
}; // BIT<int,10,10> gives a 2D BIT
```

BITrange.h

Description: 1D range increment and sum query Time: $\mathcal{O}(\log N)$

```
"BIT.h"
                                                         11 lines
template < class T, int SZ> struct BITrange {
 BIT<T,SZ> bit[2]; // piecewise linear functions
 // let cum[x] = sum_{i=1}^{x}a[i]
 void upd(int hi, T val) { // add val to a[1..hi]
   bit[1].upd(1,val), bit[1].upd(hi+1,-val); // if x \ll
       \hookrightarrowhi, cum[x] += val*x
    bit [0].upd (hi+1, hi*val); // if x > hi, cum[x] += val*
 void upd(int lo, int hi, T val) { upd(lo-1,-val), upd(hi
 T sum(int x) { return bit[1].sum(x) *x+bit[0].sum(x); }
    \hookrightarrow // get cum[x]
 T query(int x, int y) { return sum(y)-sum(x-1); }
```

SegTree.h

};

Description: 1D point update, range query

Time: $\mathcal{O}(\log N)$

```
33 lines
template<class T, int SZ> struct segtree {
   // modify these
   T identity = 0;
   T comb(T 1, T r) {
        return 1 + r;
   void updLeaf(T& 1, T val) {
       l = val;
   T tree[2*SZ+1];
   segtree() {
       M00(i, 2*SZ+1) tree[i] = identity;
   void upd(int pos, T val) {
       pos += SZ+1;
        updLeaf(tree[pos], val);
        for (pos >>= 1; pos >= 1; pos >>= 1) {
            tree[pos] = comb(tree[2*pos], tree[2*pos+1]);
   T query(int 1, int r) {
       1 += SZ+1;
       r += SZ+1;
       T res = identity;
        while (1 \le r)
           if(1&1) res = comb(res, tree[1++]);
           if(!(r&1)) res = comb(res, tree[r--]);
           1 >>= 1; r >>= 1;
        return res;
```

SegTreeBeats Lazy SegTree Sparse SegTree

SegTreeBeats.h

Description: supports modifications in the form ckmin(a_i,t) for all l < i < r, range max and sum queries

Time: $\mathcal{O}(\log N)$ 65 lines template<int SZ> struct SegTreeBeats { int N; 11 sum[2*SZ]; int mx[2*SZ][2], maxCnt[2*SZ]; void pull(int ind) { FOR(i,2) mx[ind][i] = max(mx[2*ind][i], mx[2*ind+1][i]) maxCnt[ind] = 0;FOR(i,2) { if (mx[2*ind+i][0] == mx[ind][0])maxCnt[ind] += maxCnt[2*ind+i]; else ckmax(mx[ind][1], mx[2*ind+i][0]); sum[ind] = sum[2*ind] + sum[2*ind+1];void build(vi& a, int ind = 1, int L = 0, int R = -1) { if (R == -1) { R = (N = sz(a))-1; } if (L == R) { mx[ind][0] = sum[ind] = a[L];maxCnt[ind] = 1; mx[ind][1] = -1;int M = (L+R)/2; build (a, 2*ind, L, M); build (a, 2*ind+1, M+1, R); pull (ind); void push (int ind, int L, int R) { if (L == R) return; FOR(i,2) if (mx[2*ind^i][0] > mx[ind][0]) $sum[2*ind^i] -= (11) maxCnt[2*ind^i]*$ (mx[2*ind^i][0]-mx[ind][0]); $mx[2*ind^i][0] = mx[ind][0];$ void upd(int x, int y, int t, int ind = 1, int L = 0, \hookrightarrow int R = -1) { if (R == -1) R += N;if (R < x || y < L || mx[ind][0] <= t) return;</pre> push (ind, L, R); if $(x \le L \&\& R \le y \&\& mx[ind][1] \le t)$ { sum[ind] -= (ll)maxCnt[ind]*(mx[ind][0]-t); mx[ind][0] = t;return; if (L == R) return; int M = (L+R)/2; upd(x,y,t,2*ind,L,M); upd(x,y,t,2*ind+1,M+1,R); pull(11 qsum(int x, int y, int ind = 1, int L = 0, int R =→-1) { if (R == -1) R += N;

if $(R < x \mid | y < L)$ return 0;

```
push (ind, L, R);
    if (x <= L && R <= y) return sum[ind];
    int M = (L+R)/2;
    return qsum(x,y,2*ind,L,M)+qsum(x,y,2*ind+1,M+1,R);
  int qmax(int x, int y, int ind = 1, int L = 0, int R =
     →-1) {
    if (R == -1) R += N;
    if (R < x \mid | y < L) return -1;
    push (ind, L, R);
    if (x <= L && R <= y) return mx[ind][0];</pre>
    int M = (L+R)/2;
    return max(qmax(x,y,2*ind,L,M), qmax(x,y,2*ind+1,M+1,R)
       \hookrightarrow));
};
```

Lazy SegTree.h

```
Description: 1D range update, range query
                                                        85 lines
template<class T> struct node {
   T val:
   T lazv;
    int 1, r;
    node* left;
    node* right;
    node(int 1, int r) {
        this -> 1 = 1:
        this \rightarrow r = r;
        this -> left = nullptr;
        this -> right = nullptr;
};
template<class T, int SZ> struct segtree {
    // modify these
    T combIdentity = 1e9;
    T comb(T 1, T r) {
        return min(l,r);
    T pushIdentity = 0;
    void push(node<T>* n) {
        n->val += n->lazy;
        if(n->1 != n->r)
            n->left->lazy += n->lazy;
            n->right->lazy += n->lazy;
        n->lazy = pushIdentity;
    node<T>* root;
    segtree() {
        int ub = 1;
        while (ub < SZ) ub \star= 2;
        root = new node < T > (0, ub-1);
        root->val = pushIdentity;
        root->lazy = pushIdentity;
    void propagate(node<T>* n) {
        if(n->1 != n->r) {
```

```
int mid = ((n->1) + (n->r))/2;
            if(n->left == nullptr) {
                 n\rightarrow left = new node < T > (n\rightarrow l, mid);
                 n->left->val = pushIdentity;
                 n->left->lazy = pushIdentity;
            if(n->right == nullptr) {
                 n->right = new node<T>(mid+1, n->r);
                 n->right->val = pushIdentity;
                 n->right->lazy = pushIdentity;
        push(n);
    void updN(node<T>* n, int i1, int i2, T val) {
        propagate(n);
        if(i2 < n->1 || i1 > n->r) return;
        if(i1 <= n->1 && i2 >= n->r) {
            n->lazy = val;
            push(n);
            return;
        updN(n->left, i1, i2, val);
        updN(n->right, i1, i2, val);
        n->val = comb(n->left->val, n->right->val);
    void upd(int i1, int i2, T val) {
        updN(root, i1, i2, val);
    T queryN(node<T>* n, int i1, int i2) {
        propagate(n);
        if(i2 < n->1 || i1 > n->r) return combIdentity;
        if(n->1 >= i1 \&\& n->r <= i2) return n->val;
        T a = combIdentity;
        if(n->left != nullptr) a = comb(a, queryN(n->left,
           \hookrightarrow i1, i2));
        if(n->right != nullptr) a = comb(a, queryN(n->
           \hookrightarrowright, i1, i2));
        return a;
    T query(int i1, int i2) {
        return queryN(root, i1, i2);
};
```

Sparse SegTree.h

Description: Does not allocate storage for nodes with no data 75 lines

```
template<class T> struct node {
   T val;
    int 1, r;
    node* left;
    node* right;
    node(int 1, int r) {
       this -> 1 = 1;
        this \rightarrow r = r;
        this -> left = nullptr;
```

PersSegTree Treap

```
this -> right = nullptr;
1:
template<class T, int SZ> struct segtree {
   // modify these
   T identity = 0;
   T comb(T l, T r) {
        return 1 + r;
   void updLeaf(T& 1, T val) {
        1 = val;
   node<T>* root:
   segtree() {
        int ub = 1;
        while (ub < SZ) ub \star= 2;
        root = new node < T > (0, ub-1);
        root->val = identity;
   void updN(node<T>* n, int pos, T val) {
        if (pos < n->1 \mid \mid pos > n->r) return;
        if(n->1 == n->r) {
            updLeaf(n->val, val);
            return;
        int mid = (n->1 + n->r)/2;
        if (pos > mid) {
            if(n->right == nullptr) {
                n->right = new node<T>(mid+1, n->r);
                n->right->val = identity;
            updN(n->right, pos, val);
        else {
            if(n->left == nullptr) {
                n->left = new node<T>(n->l, mid);
                n->left->val = identity;
            updN(n->left, pos, val);
        T lv = (n->left == nullptr) ? identity : n->left->
        T rv = (n->right == nullptr) ? identity : n->right
           \hookrightarrow->val;
        n->val = comb(lv, rv);
    void upd(int pos, T val) {
        updN(root, pos, val);
    T queryN(node<T>* n, int i1, int i2) {
        if (i2 < n->1 || i1 > n->r) return identity;
        if(n->1 == n->r) return n->val;
        if(n->1 >= i1 \&\& n->r <= i2) return n->val;
        T a = identity;
```

PersSegTree.h

Description: persistent segtree with lazy updates, assumes that lazy[cur] is included in val[cur] before propagating cur **Time:** $\mathcal{O}(\log N)$

```
60 lines
template<class T, int SZ> struct pseg {
 static const int LIMIT = 10000000; // adjust
  int l[LIMIT], r[LIMIT], nex = 0;
 T val[LIMIT], lazv[LIMIT];
 int copy(int cur) {
   int x = nex++;
   val[x] = val[cur], l[x] = l[cur], r[x] = r[cur], lazy[
       \hookrightarrow x] = lazy[cur];
   return x;
 T comb(T a, T b) { return min(a,b); }
  void pull(int x) { val[x] = comb(val[l[x]],val[r[x]]); }
  void push(int cur, int L, int R) {
   if (!lazv[cur]) return;
   if (L != R) {
      1[cur] = copy(1[cur]);
      val[l[cur]] += lazy[cur];
      lazy[l[cur]] += lazy[cur];
      r[cur] = copy(r[cur]);
      val[r[cur]] += lazy[cur];
      lazy[r[cur]] += lazy[cur];
   lazy[cur] = 0;
 T query(int cur, int lo, int hi, int L, int R) {
   if (lo <= L && R <= hi) return val[cur];</pre>
   if (R < lo || hi < L) return INF;
   int M = (L+R)/2;
    return lazy[cur]+comb(query(l[cur],lo,hi,L,M), query(r
       \hookrightarrow [cur], lo, hi, M+1, R));
  int upd(int cur, int lo, int hi, T v, int L, int R) {
   if (R < lo || hi < L) return cur;
   int x = copy(cur);
    if (lo <= L && R <= hi) { val[x] += v, lazy[x] += v;
       →return x; }
   push(x, L, R);
    int M = (L+R)/2;
```

```
l[x] = upd(l[x], lo, hi, v, L, M), r[x] = upd(r[x], lo, hi, v, L, M)
        \rightarrowM+1,R);
    pull(x); return x;
  int build(vector<T>& arr, int L, int R) {
    int cur = nex++;
    if (L == R) {
      if (L < sz(arr)) val[cur] = arr[L];</pre>
      return cur;
    int M = (L+R)/2;
    l[cur] = build(arr,L,M), r[cur] = build(arr,M+1,R);
    pull(cur); return cur;
  vi loc;
  void upd(int lo, int hi, T v) { loc.pb(upd(loc.back(), lo
     \hookrightarrow, hi, v, 0, SZ-1)); }
  T query (int ti, int lo, int hi) { return query (loc[ti],
     \hookrightarrowlo, hi, 0, SZ-1); }
  void build(vector<T>& arr) { loc.pb(build(arr, 0, SZ-1));
};
```

Treap.h

Description: easy BBST, use split and merge to implement insert and delete

Time: $\mathcal{O}(\log N)$

```
77 lines
typedef struct tnode* pt;
struct thode {
 int pri, val; pt c[2]; // essential
 int sz; 11 sum; // for range queries
 bool flip; // lazy update
 tnode (int _val) {
   pri = rand()+(rand()<<15); val = _val; c[0] = c[1] =</pre>
    sz = 1; sum = val;
    flip = 0;
};
int getsz(pt x) { return x?x->sz:0; }
11 getsum(pt x) { return x?x->sum:0; }
pt prop(pt x) {
 if (!x || !x->flip) return x;
 swap(x->c[0], x->c[1]);
  x \rightarrow flip = 0;
 FOR(i,2) if (x->c[i]) x->c[i]->flip ^= 1;
 return x;
pt calc(pt x) {
 assert(!x->flip);
 prop(x->c[0]), prop(x->c[1]);
  x->sz = 1+qetsz(x->c[0])+qetsz(x->c[1]);
```

SqrtDecomp Mo MaxQueue 2D Sumtree

```
x->sum = x->val+getsum(x->c[0])+getsum(x->c[1]);
 return x;
void tour(pt x, vi& v) {
 if (!x) return;
 prop(x);
 tour (x->c[0],v); v.pb(x->val); tour(x->c[1],v);
pair<pt,pt> split(pt t, int v) { // >= v goes to the right
 if (!t) return {t,t};
 prop(t);
 if (t->val >= v) {
   auto p = split(t->c[0], v); t->c[0] = p.s;
   return {p.f, calc(t)};
   auto p = split(t->c[1], v); t->c[1] = p.f;
   return {calc(t), p.s};
pair<pt, pt> splitsz(pt t, int sz) { // leftmost sz nodes
  \hookrightarrowgo to left
 if (!t) return {t,t};
 prop(t);
 if (getsz(t->c[0]) >= sz) {
   auto p = splitsz(t->c[0], sz); t->c[0] = p.s;
   return {p.f, calc(t)};
 } else {
   auto p = splitsz(t->c[1], sz-getsz(t->c[0])-1); t->c
      \hookrightarrow [1] = p.f;
    return {calc(t), p.s};
pt merge(pt 1, pt r) {
 if (!1 || !r) return 1 ? 1 : r;
 prop(l), prop(r);
 pt t;
 if (1->pri > r->pri) 1->c[1] = merge(1->c[1],r), t = 1;
 else r - c[0] = merge(1, r - c[0]), t = r;
  return calc(t);
pt ins(pt x, int v) { // insert v
 auto a = split(x,v), b = split(a.s,v+1);
  return merge(a.f, merge(new tnode(v), b.s));
pt del(pt x, int v) { // delete v
 auto a = split(x, v), b = split(a.s, v+1);
  return merge(a.f,b.s);
SqrtDecomp.h
```

Description: 1D point update, range query

Time: $\mathcal{O}\left(\sqrt{N}\right)$

struct sqrtDecomp { const static int blockSZ = 10; //change this int val[blockSZ*blockSZ]; int lazy[blockSZ];

44 lines

int 1 = 0, r = -1, cans = 0;

void modify (int x, int y = 1) {

// if condition: cans --;

```
sqrtDecomp() {
        M00(i, blockSZ*blockSZ) val[i] = 0;
        M00(i, blockSZ) lazy[i] = 0;
    void upd(int 1, int r, int v) {
        int ind = 1;
        while(ind%blockSZ && ind <= r) {
            val[ind] += v;
            lazy[ind/blockSZ] += v;
            ind++;
        while(ind + blockSZ <= r) {
            lazy[ind/blockSZ] += v*blockSZ;
            ind += blockSZ;
        while(ind <= r) {
            val[ind] += v;
            lazy[ind/blockSZ] += v;
            ind++;
    int query(int 1, int r) {
        int res = 0;
        int ind = 1;
        while(ind%blockSZ && ind <= r) {</pre>
            res += val[ind];
            ind++;
        while (ind + blockSZ <= r) {
            res += lazy[ind/blockSZ];
            ind += blockSZ;
        while(ind <= r) {
            res += val[ind];
            ind++;
        return res;
};
Mo.h
Description: Answers queries offline in (N+Q)sqrt(N) Also see Mo's
int N, A[MX];
int ans[MX], oc[MX], BLOCK;
vector<array<int,3>> todo; // store left, right, index of
  \hookrightarrowans
bool cmp(array<int,3> a, array<int,3> b) { // sort queries
  if (a[0]/BLOCK != b[0]/BLOCK) return a[0] < b[0];</pre>
  return a[1] < b[1];
```

```
oc[x] += y;
 // if condition: cans ++;
int answer(int L, int R) { // modifyjust interval
 while (1 > L) modify (--1);
 while (r < R) modify(++r);
 while (1 < L) \mod (1++,-1);
 while (r > R) modify (r--,-1);
 return cans;
void solve() {
 BLOCK = sqrt(N); sort(all(todo),cmp);
 trav(x,todo) {
    answer(x[0],x[1]);
    ans[x[2]] = cans;
MaxQueue.h
Description: queue, but get() returns max element
Time: \mathcal{O}(1)
                                                        16 lines
struct maxQueue {
    queue<int> q;
    deque<int> dq;
    void push(int v) {
        q.push(v);
        if(q.empty()) {dq.push_back(v); return;}
        while(!dq.empty() && dq.back() < v) dq.pop_back();</pre>
        dq.push_back(v);
    void pop() {
        if(q.front() == dq.front()) dq.pop_front();
    int get() {return dg.front();}
    int size() {return (int)q.size();}
```

2.3 2D Range Queries

2D Sumtree.h

Description: Lawrence's 2d sum segment tree

104 lines

```
struct sumtreenode{
    node* root;
    sumtreenode* left;
    sumtreenode* right;
    int 1, r;
    sumtreenode(int 1, int r, int SZ) {
        int ub = 1;
         while (ub < SZ) ub \star= 2;
         root = new node(0, ub-1);
        this \rightarrow 1 = 1;
         this \rightarrow r = r;
         this->left = nullptr;
         this->right = nullptr;
```

```
void updN(node* n, int pos, int val) {
        if(pos < n->1 || pos > n->r) return;
        if(n->1 == n->r) {
            n->val = val;
            return;
        int mid = (n->1 + n->r)/2;
        if (pos > mid) {
            if (n->right == nullptr) n->right = new node(
               \hookrightarrowmid+1, n->r);
            updN(n->right, pos, val);
        else (
            if (n->left == nullptr) n->left = new node(n->l
               \hookrightarrow, mid);
            updN(n->left, pos, val);
        int s = 0;
        if(n->right != nullptr) s += n->right->val;
        if(n->left != nullptr) s += n->left->val;
        n->val = s;
    void upd(int pos, int val) {
        updN(root, pos, val);
    int queryN(node* n, int i1, int i2) {
        if(i2 < n->1 || i1 > n->r) return 0;
        if (n->1 == n->r) return n->val;
        if(n->1 >= i1 \&\& n->r <= i2) return n->val;
        int s = 0:
        if (n->left != nullptr) s += queryN(n->left, i1, i2
        if (n->right != nullptr) s += queryN(n->right, i1,
           \hookrightarrowi2);
        return s;
    int query(int i1, int i2) {
        return queryN(root, i1, i2);
};
template<int w, int h> struct sumtree2d{
    sumtreenode* root;
    sumtree2d() {
        int ub = 1:
        while (ub < w) ub \star= 2;
        this->root = new sumtreenode(0, ub-1, h);
        root->left = nullptr;
        root->right = nullptr;
    void updN(sumtreenode* n, int x, int y, int val) {
        if (x < n->1 \mid | x > n->r) return;
        if(n->1 == n->r) {
            n->upd(y, val);
            return;
```

```
int mid = (n->1 + n->r)/2;
    if(x > mid)  {
        if(n->right == nullptr) n->right = new
           \hookrightarrow sumtreenode (mid+1, n->r, h);
        updN(n->right, x, y, val);
        if(n->left == nullptr) n->left = new
           ⇒sumtreenode(n->1, mid, h);
        updN(n->left, x, y, val);
    int s = 0:
    if(n->left != nullptr) s += n->left->query(y, y);
    if(n->right != nullptr) s += n->right->query(y, y)
       \hookrightarrow ;
    n->upd(y, s);
void upd(int x, int y, int val) {
    updN(root, x, y, val);
int queryN(sumtreenode* n, int x1, int y1, int x2, int
   if (x2 < n->1 | | x1 > n->r) return 0;
    if (n->1 == n->r) return n->query(y1, y2);
    if(n->1) = x1 \&\& n->r <= x2) return n->query(y1,
       \hookrightarrow v2);
    int s = 0;
    if(n->left != nullptr) s += queryN(n->left, x1, y1
       \hookrightarrow, x2, y2);
    if (n->right != nullptr) s += queryN(n->right, x1,
       \hookrightarrowy1, x2, y2);
    return s;
int query(int x1, int y1, int x2, int y2) {
    return queryN(root, x1, y1, x2, y2);
```

Number Theory (3)

if (val < 0) val += MOD;</pre>

3.1 Modular Arithmetic

Modular.h

};

Description: modular arithmetic operations

```
template<class T> struct modular {
  T val;
  explicit operator T() const { return val; }
  modular() { val = 0; }
  modular(const 11& v) {
```

val = (-MOD <= v && v <= MOD) ? v : v % MOD;</pre>

```
friend bool operator!=(const modular& a, const modular&
    \hookrightarrowb) { return ! (a == b); }
 friend bool operator < (const modular& a, const modular& b
    modular operator-() const { return modular(-val); }
 modular& operator+=(const modular& m) { if ((val += m.
    →val) >= MOD) val -= MOD; return *this; }
 modular& operator = (const modular& m) { if ((val -= m.
    ⇔val) < 0) val += MOD; return *this; }</pre>
 modular& operator *= (const modular& m) { val = (11) val *m.
    →val%MOD; return *this; }
 friend modular pow(modular a, 11 p) {
   modular ans = 1; for (; p; p /= 2, a \star= a) if (p&1)
       \hookrightarrowans *= a;
   return ans:
 friend modular inv(const modular& a) {
   assert(a != 0); return exp(a, MOD-2);
 modular& operator/=(const modular& m) { return (*this)
    \hookrightarrow \star = inv(m); }
 friend modular operator+(modular a, const modular& b) {
    →return a += b; }
 friend modular operator-(modular a, const modular& b) {
     friend modular operator*(modular a, const modular& b) {
    \hookrightarrowreturn a *= b; }
 friend modular operator/(modular a, const modular& b) {
     →return a /= b; }
typedef modular<int> mi;
typedef pair<mi, mi> pmi;
typedef vector<mi> vmi;
typedef vector<pmi> vpmi;
```

// friend ostream& operator << (ostream& os, const modular

friend void re(modular& a) { ll x; re(x); a = modular(x)

friend bool operator == (const modular& a, const modular&

friend void pr(const modular& a) { pr(a.val); }

 $\hookrightarrow \&$ a) { return os << a.val; }

⇒b) { return a.val == b.val; }

ModFact.h

Description: pre-compute factorial mod inverses for MOD, assumes MOD is prime and SZ < MOD

Time: $\mathcal{O}\left(SZ\right)$

10 lines

```
ifac[i] = ifac[i-1]*inv[i]%MOD;
}
```

ModMulLL.h

Description: multiply two 64-bit integers mod another if 128-bit is not available works for $0 \le a, b < mod < 2^{63}$

typedef unsigned long long ul;

// equivalent to (ul) (__int128(a) *b\$mod)
ul modMul(ul a, ul b, const ul mod) {
 l1 ret = a*b-mod*(ul)((ld)a*b/mod);
 return ret+((ret<0)-(ret>=(l1)mod)) *mod;
}
ul modPow(ul a, ul b, const ul mod) {
 if (b == 0) return 1;
 ul res = modPow(a,b/2,mod);
 res = modMul(res,res,mod);
 if (b&l) return modMul(res,a,mod);
 return res;
}

ModSart.h

Description: find sqrt of integer mod a prime

Time: ?

```
template<class T> T sgrt(modular<T> a) {
 auto p = pow(a, (MOD-1)/2); if (p != 1) return p == 0 ? 0
    \hookrightarrow : -1; // check if zero or does not have sqrt
 T s = MOD-1, e = 0; while (s % 2 == 0) s /= 2, e ++;
 modular < T > n = 1; while (pow(n, (MOD-1)/2) == 1) n = (T) (
    \hookrightarrown)+1; // find non-square residue
 auto x = pow(a, (s+1)/2), b = pow(a, s), g = pow(n, s);
  int r = e:
  while (1) {
   auto B = b; int m = 0; while (B != 1) B *= B, m ++;
   if (m == 0) return min((T)x, MOD-(T)x);
   FOR(i,r-m-1) g *= g;
   x *= g; g *= g; b *= g; r = m;
/* Explanation:
* Initially, x^2=ab, ord(b) = 2^m, ord(g) = 2^r where m < r
 * q = q^{2^{r-m-1}} -> ord(q) = 2^{m+1}
 * if x'=x*g, then b'=b*g^2
    (b')^{2^{m-1}} = (b*g^2)^{2^{m-1}}
             = b^{2^{m-1}} *q^{2^m}
             = -1 * -1
             = 1
 -> ord(b') | ord(b) /2
 * m decreases by at least one each iteration
```

ModSum.h

3.2 Primality

PrimeSieve.h

Description: tests primality up to SZ

Time: $\mathcal{O}\left(SZ\log\log SZ\right)$

11 lines

```
template<int SZ> struct Sieve {
  bitset<SZ> isprime;
  vi pr;
  Sieve() {
   isprime.set(); isprime[0] = isprime[1] = 0;
   for (int i = 4; i < SZ; i += 2) isprime[i] = 0;
   for (int i = 3; i*i < SZ; i += 2) if (isprime[i])
      for (int j = i*i; j < SZ; j += i*2) isprime[j] = 0;
   FOR(i,2,SZ) if (isprime[i]) pr.pb(i);
  }
};</pre>
```

FactorFast.h

Description: Factors integers up to 2⁶⁰

Time: ?

```
"PrimeSieve.h"
                                                       46 lines
Sieve<1<<20> S = Sieve<1<<20>(); // should take care of
  \hookrightarrowall primes up to n^(1/3)
bool millerRabin(ll p) { // test primality
 if (p == 2) return true;
 if (p == 1 || p % 2 == 0) return false;
  11 s = p - 1; while (s % 2 == 0) s /= 2;
  FOR(i,30) { // strong liar with probability <= 1/4
   11 a = rand() % (p - 1) + 1, tmp = s;
    11 mod = mod_pow(a, tmp, p);
    while (tmp != p - 1 \&\& mod != 1 \&\& mod != p - 1) {
      mod = mod_mul(mod, mod, p);
      tmp *= 2;
    if (mod != p - 1 && tmp % 2 == 0) return false;
  return true;
```

```
11 f(11 a, 11 n, 11 &has) { return (mod_mul(a, a, n) + has
vpl pollardsRho(ll d) {
 vpl res:
 auto& pr = S.pr;
  for (int i = 0; i < sz(pr) && pr[i]*pr[i] <= d; i++) if
    \hookrightarrow (d % pr[i] == 0) {
    int co = 0; while (d % pr[i] == 0) d /= pr[i], co ++;
    res.pb({pr[i],co});
 if (d > 1) { // d is now a product of at most 2 primes.
    if (millerRabin(d)) res.pb({d,1});
   else while (1) {
     11 \text{ has} = \text{rand}() % 2321 + 47;
      11 x = 2, y = 2, c = 1;
      for (; c == 1; c = \_gcd(abs(x-y), d)) {
       x = f(x, d, has);
        y = f(f(y, d, has), d, has);
      } // should cycle in ~sqrt(smallest nontrivial
         ⇒divisor) turns
      if (c != d) {
        d \neq c; if (d > c) swap(d,c);
        if (c == d) res.pb({c,2});
        else res.pb({c,1}), res.pb({d,1});
        break:
 return res;
```

3.3 Divisibility

Euclid.h

Description: Euclidean Algorithm

0.1:-

CRT.h

$\textbf{Description:} \ \ \textbf{Chinese} \ \ \textbf{Remainder} \ \ \textbf{Theorem}$

IntPerm MatroidIntersect PermGroup

Combinatorial (4)

IntPerm.h

Description: convert permutation of $\{0, 1, ..., N-1\}$ to integer in

Usage: assert (encode (decode (5, 37)) == 37);

Time: $\mathcal{O}(N)$

20 lines

```
vi decode(int n, int a) {
 vi el(n), b; iota(all(el),0);
 FOR(i,n) {
   int z = a sz(e1):
   b.pb(el[z]); a \neq sz(el);
   swap(el[z],el.back()); el.pop_back();
 return b;
int encode(vi b) {
 int n = sz(b), a = 0, mul = 1;
 vi pos(n); iota(all(pos),0); vi el = pos;
 FOR(i,n) {
   int z = pos[b[i]]; a += mul*z; mul *= sz(el);
   swap(pos[el[z]],pos[el.back()]);
   swap(el[z],el.back()); el.pop_back();
 return a;
```

MatroidIntersect.h

Description: computes a set of maximum size which is independent in both graphic and colorful matroids, aka a spanning forest where no two edges are of the same color

Time: $\mathcal{O}\left(GI^{1.5}\right)$ calls to oracles, where G is the size of the ground set and I is the size of the independent set

```
"DSU.h"
                                                      108 lines
int R:
map<int, int> m;
struct Element {
 pi ed;
 int col;
 bool in_independent_set = 0;
  int independent_set_position;
 Element (int u, int v, int c) { ed = \{u,v\}; col = c; }
vi independent_set;
vector<Element> ground_set;
bool col used[300];
struct GBasis {
  void reset() { D.init(sz(m)); }
  void add(pi v) { assert(D.unite(v.f,v.s)); }
  bool independent_with(pi v) { return !D.sameSet(v.f,v.s)
};
```

```
GBasis basis, basis wo[300];
bool graph oracle(int inserted) {
  return basis.independent_with(ground_set[inserted].ed);
bool graph_oracle(int inserted, int removed)
  int wi = ground set[removed].independent set position;
  return basis wo[wi].independent_with(ground_set[inserted
void prepare_graph_oracle() {
  basis.reset();
  FOR(i,sz(independent set)) basis wo[i].reset();
  FOR(i,sz(independent_set)) {
    pi v = ground_set[independent_set[i]].ed; basis.add(v)
    FOR(j,sz(independent_set)) if (i != j) basis_wo[j].add
      \hookrightarrow (v);
bool colorful_oracle(int ins) {
  ins = ground set[ins].col;
  return !col_used[ins];
bool colorful_oracle(int ins, int rem) {
  ins = ground_set[ins].col;
  rem = ground_set[rem].col;
  return !col_used[ins] || ins == rem;
void prepare_colorful_oracle() {
 FOR(i,R) col\_used[i] = 0;
  trav(t,independent_set) col_used[ground_set[t].col] = 1;
bool augment() {
  prepare_graph_oracle();
  prepare_colorful_oracle();
  vi par(sz(ground_set),MOD);
  queue<int> q;
  FOR(i,sz(ground_set)) if (colorful_oracle(i)) {
    assert(!ground_set[i].in_independent_set);
    par[i] = -1; q.push(i);
  int lst = -1;
  while (sz(q)) {
    int cur = q.front(); q.pop();
    if (ground_set[cur].in_independent_set) {
      FOR(to,sz(ground_set)) if (par[to] == MOD) {
        if (!colorful_oracle(to,cur)) continue;
        par[to] = cur; q.push(to);
    } else {
      if (graph_oracle(cur)) { lst = cur; break; }
      trav(to,independent set) if (par[to] == MOD) {
        if (!graph_oracle(cur,to)) continue;
        par[to] = cur; q.push(to);
```

```
if (lst == -1) return 0;
   ground_set[lst].in_independent_set ^= 1;
   lst = par[lst];
 } while (lst != -1);
 independent set.clear();
 FOR(i,sz(ground_set)) if (ground_set[i].
    →in independent set) {
   ground_set[i].independent_set_position = sz(
       →independent_set);
   independent set.pb(i);
 return 1;
void solve() {
 re(R); if (R == 0) exit(0);
 m.clear(); ground_set.clear(); independent_set.clear();
 FOR (i, R) {
   int a,b,c,d; re(a,b,c,d);
   ground_set.pb(Element(a,b,i));
   ground_set.pb(Element(c,d,i));
   m[a] = m[b] = m[c] = m[d] = 0;
 int co = 0;
 trav(t,m) t.s = co++;
 trav(t,ground_set) t.ed.f = m[t.ed.f], t.ed.s = m[t.ed.s
 while (augment());
 ps(2*sz(independent_set));
```

PermGroup.h

Description: Schreier-Sims, count number of permutations in group and test whether permutation is a member of group Time: ?

51 lines

```
const int N = 15:
int n;
vi inv(vi v) { vi V(sz(v)); FOR(i,sz(v)) V[v[i]] = i;}
  →return V; }
vi id() { vi v(n); iota(all(v),0); return v; }
vi operator*(const vi& a, const vi& b) {
 vi c(sz(a)); FOR(i,sz(a)) c[i] = a[b[i]];
  return c;
struct Group {
 bool flag[N];
  vi sigma[N]; // sigma[t][k] = t, sigma[t][x] = x if x >
    \hookrightarrow k
  vector<vi> gen;
  void clear(int p) {
    memset(flag,0, sizeof flag);
    flag[p] = 1; sigma[p] = id();
    gen.clear();
} q[N];
```

Matrix MatrixInv MatrixTree VecOp

```
bool check(const vi& cur, int k) {
 if (!k) return 1;
 int t = cur[k];
 return g[k].flag[t] ? check(inv(g[k].sigma[t])*cur,k-1)
void updateX(const vi& cur, int k);
void ins(const vi& cur, int k) {
 if (check(cur,k)) return;
 g[k].gen.pb(cur);
 FOR(i,n) if (g[k].flag[i]) updateX(cur*g[k].sigma[i],k);
void updateX(const vi& cur, int k) {
 int t = cur[k];
 if (g[k].flag[t]) ins(inv(g[k].sigma[t])*cur,k-1); //
    \hookrightarrow fixes k \rightarrow k
 else {
    g[k].flag[t] = 1, g[k].sigma[t] = cur;
    trav(x,g[k].gen) updateX(x*cur,k);
ll order (vector<vi> gen) {
 assert(sz(gen)); n = sz(gen[0]); FOR(i,n) g[i].clear(i);
 trav(a,gen) ins(a,n-1); // insert perms into group one
     \hookrightarrowby one
 11 \text{ tot} = 1;
 FOR(i,n) {
   int cnt = 0; FOR(j, i+1) cnt += g[i].flag[j];
    tot *= cnt;
 return tot;
```

Numerical (5)

5.1 Matrix

Matrix.h

Description: 2D matrix operations

36 lines

```
FOR(i,r) FOR(i,c) d[i][i] -= m.d[i][i];
    return *this;
  Mat operator* (const Mat& m) {
    assert(c == m.r); Mat x(r,m.c);
    FOR(i,r) FOR(j,c) FOR(k,m.c) x.d[i][k] += d[i][j]*m.d[
    return x;
  Mat operator+(const Mat& m) { return Mat(*this)+=m; }
  Mat operator-(const Mat& m) { return Mat(*this)-=m; }
  Mat& operator *= (const Mat& m) { return *this = (*this) *m
    \hookrightarrow; }
  friend Mat pow(Mat m, ll p) {
    assert(m.r == m.c);
    Mat r(m.r,m.c);
    FOR(i, m.r) r.d[i][i] = 1;
    for (; p; p /= 2, m \star= m) if (p&1) r \star= m;
    return r;
};
```

MatrixInv.h

Description: calculates determinant via gaussian elimination **Time:** $\mathcal{O}(N^3)$

```
template < class T > T gauss (Mat < T > & m) { // determinant of
  \hookrightarrow1000x1000 Matrix in \sim1s
  int n = m.r;
 T prod = 1; int nex = 0;
 FOR(i,n) {
    int row = -1; // for 1d use EPS rather than 0
    FOR(j,nex,n) if (m.d[j][i] != 0) { row = j; break; }
    if (row == -1) { prod = 0; continue; }
    if (row != nex) prod *= -1, swap(m.d[row], m.d[nex]);
    prod *= m.d[nex][i];
    auto x = 1/m.d[nex][i]; FOR(k,i,m.c) m.d[nex][k] *= x;
    FOR(j,n) if (j != nex) {
      auto v = m.d[j][i];
      if (v != 0) FOR(k,i,m.c) m.d[j][k] -= v*m.d[nex][k];
    nex ++;
  return prod;
template<class T> Mat<T> inv(Mat<T> m) {
 int n = m.r;
 Mat < T > x(n, 2*n);
 FOR(i,n) {
   x.d[i][i+n] = 1;
    FOR(j,n) \times d[i][j] = m.d[i][j];
  if (gauss(x) == 0) return Mat < T > (0,0);
 Mat < T > r(n,n);
 FOR(i,n) FOR(j,n) r.d[i][j] = x.d[i][j+n];
```

MatrixTree.h

Description: Kirchhoff's Matrix Tree Theorem: given adjacency matrix, calculates # of spanning trees

5.2 Polynomials

VecOp.h

Description: arithmetic + misc polynomial operations with vectors

```
namespace VecOp {
 template<class T> vector<T> rev(vector<T> v) { reverse(
     ⇒all(v)); return v; }
 template<class T> vector<T> shift(vector<T> v, int x) {
    →v.insert(v.begin(),x,0); return v; }
 template<class T> vector<T> integ(const vector<T>& v) {
   vector<T> res(sz(v)+1);
   FOR(i, sz(v)) res[i+1] = v[i]/(i+1);
   return res;
 template<class T> vector<T> dif(const vector<T>& v) {
   if (!sz(v)) return v;
   vector<T> res(sz(v)-1); FOR(i,1,sz(v)) res[i-1] = i*v[

→i];

   return res;
 template<class T> vector<T>& remLead(vector<T>& v) {
   while (sz(v) && v.back() == 0) v.pop_back();
   return v;
 template < class T > T eval(const vector < T > & v, const T & x)
   T res = 0; ROF(i,sz(v)) res = x*res+v[i];
   return res;
 template<class T> vector<T>& operator+=(vector<T>& 1,
    ⇔const vector<T>& r) {
   1.rsz(max(sz(1),sz(r))); FOR(i,sz(r)) l[i] += r[i];
      →return 1;
 template<class T> vector<T>& operator == (vector<T>& 1,
```

PolyRoots Karatsuba FFT

```
1.rsz(max(sz(1),sz(r))); FOR(i,sz(r)) l[i] -= r[i];
     \hookrightarrowreturn 1:
template<class T> vector<T>& operator *= (vector<T>& 1,
   \rightarrowconst T& r) { trav(t,1) t *= r; return 1; }
template<class T> vector<T>& operator/=(vector<T>& 1,
   \hookrightarrowconst T& r) { trav(t,1) t /= r; return 1; }
template<class T> vector<T> operator+(vector<T> 1, const
   \hookrightarrow vector<T>& r) { return 1 += r; }
template<class T> vector<T> operator-(vector<T> 1, const

    vector<T>& r) { return 1 -= r; }

template<class T> vector<T> operator* (vector<T> 1, const
   → T& r) { return 1 *= r; }
template<class T> vector<T> operator*(const T& r, const
   →vector<T>& 1) { return 1*r; }
template<class T> vector<T> operator/(vector<T> 1, const
  template<class T> vector<T> operator*(const vector<T>& 1
  \hookrightarrow, const vector<T>& r) {
 if (min(sz(l),sz(r)) == 0) return {};
 vector < T > x(sz(1) + sz(r) - 1); FOR(i, sz(1)) FOR(j, sz(r))
     \hookrightarrow x[i+j] += l[i] *r[j];
 return x;
template<class T> vector<T>& operator *= (vector<T>& 1,
   \hookrightarrowconst vector<T>& r) { return 1 = 1*r; }
template < class T > pair < vector < T > , vector < T > > qr (vector < T >
   \hookrightarrow a, vector<T> b) { // quotient and remainder
 assert(sz(b)); auto B = b.back(); assert(B != 0);
 B = 1/B; trav(t,b) t *= B;
  remLead(a); vector<T> q(max(sz(a)-sz(b)+1,0));
 while (sz(a) >= sz(b)) {
    q[sz(a)-sz(b)] = a.back();
    a = a.back()*shift(b,sz(a)-sz(b));
    remLead(a);
 trav(t,q) t *= B;
  return {q,a};
template<class T> vector<T> quo(const vector<T>& a,

→const vector<T>& b) { return gr(a,b).f; }

template<class T> vector<T> rem(const vector<T>& a,
   \rightarrowconst vector<T>& b) { return gr(a,b).s; }
template<class T> vector<T> interpolate(vector<pair<T,T</pre>
  >>> v) {
  vector<T> ret, prod = {1};
  FOR(i, sz(v)) prod *= vector<T>({-v[i].f,1});
 FOR(i,sz(v)) {
   T todiv = 1; FOR(j,sz(v)) if (i != j) todiv *= v[i].
    ret += qr(prod, \{-v[i].f, 1\}).f*(v[i].s/todiv);
  return ret;
```

```
using namespace VecOp;
PolyRoots.h
Description: Finds the real roots of a polynomial.
Usage: poly_roots(\{\{2,-3,1\}\},-1e9,1e9\}) // solve x^2-3x+2=
Time: \mathcal{O}\left(N^2\log(1/\epsilon)\right)
"VecOp.h"
                                                          19 lines
vd polyRoots(vd p, ld xmin, ld xmax)
  if (sz(p) == 2) \{ return \{-p[0]/p[1]\}; \}
  auto dr = polyRoots(dif(p),xmin,xmax);
  dr.pb(xmin-1); dr.pb(xmax+1); sort(all(dr));
  vd ret;
  FOR(i,sz(dr)-1) {
    auto l = dr[i], h = dr[i+1];
    bool sign = eval(p, 1) > 0;
    if (sign ^ (eval(p,h) > 0)) {
      FOR(it, 60) { // while (h - 1 > 1e-8)
        auto m = (1+h)/2, f = eval(p,m);
        if ((f \le 0) \hat{sign}) l = m;
        else h = m;
      ret.pb((1+h)/2);
  return ret;
Karatsuba.h
Description: multiply two polynomials
Time: \mathcal{O}\left(N^{\log_2 3}\right)
int size(int s) { return s > 1 ? 32-_builtin_clz(s-1) :
  →0; }
void karatsuba(ll *a, ll *b, ll *c, ll *t, int n) {
  int ca = 0, cb = 0; FOR(i,n) ca += !!a[i], cb += !!b[i];
  if (min(ca, cb) <= 1500/n) { // few numbers to multiply</pre>
    if (ca > cb) swap(a, b);
```

```
karatsuba(&a[0], &b[0], &c[0], &t[0], n);
c.rsz(sa+sb-1); return c;
}
```

FFT.h Description: multiply two polynomials

```
Time: \mathcal{O}(N \log N)
"Modular.h"
                                                          40 lines
typedef complex<db> cd;
const int MOD = (119 << 23) + 1, root = 3; // = 998244353
// NTT: For p < 2^30 there is also e.g. (5 << 25, 3), (7
  \hookrightarrow << 26, 3),
// (479 << 21, 3) and (483 << 21, 5). The last two are >
  \hookrightarrow 10^9.
constexpr int size(int s) { return s > 1 ? 32-
    \rightarrow __builtin_clz(s-1) : 0; }
void genRoots(vmi& roots) { // primitive n-th roots of
   \hookrightarrowunity
 int n = sz(roots); mi r = pow(mi(root), (MOD-1)/n);
 roots[0] = 1; FOR(i,1,n) roots[i] = roots[i-1]*r;
void genRoots(vcd& roots) { // change cd to complex<double</pre>
  \hookrightarrow> instead?
  int n = sz(roots); double ang = 2*PI/n;
 FOR(i,n) roots[i] = cd(cos(ang*i), sin(ang*i)); // is
     ⇒there a way to do this more quickly?
template<class T> void fft(vector<T>& a, const vector<T>&
   \hookrightarrowroots, bool inv = 0) {
  int n = sz(a);
  for (int i = 1, j = 0; i < n; i++) { // sort by reverse
     ⇒bit representation
    int bit = n >> 1;
    for (; j&bit; bit >>= 1) j ^= bit;
    j ^= bit; if (i < j) swap(a[i], a[j]);</pre>
  for (int len = 2; len <= n; len <<= 1)
    for (int i = 0; i < n; i += len)
      FOR(j,len/2) {
        int ind = n/len*j; if (inv && ind) ind = n-ind;
        auto u = a[i+j], v = a[i+j+len/2]*roots[ind];
        a[i+j] = u+v, a[i+j+len/2] = u-v;
 if (inv) { T i = T(1)/T(n); trav(x,a) x *= i; }
template<class T> vector<T> mult(vector<T> a, vector<T> b)
 int s = sz(a) + sz(b) - 1, n = 1 << size(s);
 vector<T> roots(n); genRoots(roots);
 a.rsz(n), fft(a,roots);
 b.rsz(n), fft(b,roots);
 FOR(i,n) a[i] *= b[i];
 fft(a,roots,1); return a;
```

FFTmod.h

Description: multiply two polynomials with arbitrary MOD ensures precision by splitting in half

```
"FFT.h"
                                                           27 lines
vl multMod(const vl& a, const vl& b) {
 if (!min(sz(a),sz(b))) return {};
  int s = sz(a) + sz(b) - 1, n = 1 < size(s), cut = sqrt(MOD);
 vcd roots(n); genRoots(roots);
  vcd ax(n), bx(n);
  FOR(i,sz(a)) ax[i] = cd((int)a[i]/cut, (int)a[i]%cut);
     \hookrightarrow // ax(x) = a1(x) + i * a0(x)
  FOR(i,sz(b)) bx[i] = cd((int)b[i]/cut, (int)b[i]%cut);
     \hookrightarrow // bx (x) =b1 (x) +i *b0 (x)
  fft(ax, roots), fft(bx, roots);
  vcd v1(n), v0(n);
  FOR(i,n) {
   int j = (i ? (n-i) : i);
    v1[i] = (ax[i]+conj(ax[j]))*cd(0.5,0)*bx[i]; // v1 =
       \hookrightarrow a1*(b1+b0*cd(0,1));
    v0[i] = (ax[i]-conj(ax[j]))*cd(0,-0.5)*bx[i]; // v0 =
       \hookrightarrow a0*(b1+b0*cd(0,1));
  fft(v1, roots, 1), fft(v0, roots, 1);
  vl ret(n);
  FOR(i,n) {
   11 V2 = (11) round(v1[i].real()); // a1*b1
    11 V1 = (11) round(v1[i].imag()) + (11) round(v0[i].real()
       \hookrightarrow); // a0*b1+a1*b0
    11 V0 = (11) round(v0[i].imag()); // a0*b0
    ret[i] = ((V2%MOD*cut+V1)%MOD*cut+V0)%MOD;
  ret.rsz(s); return ret;
// \sim 0.8s when sz(a) = sz(b) = 1 << 19
```

PolyInv.h Description: ?

Time: ?

PolyDiv.h

Description: divide two polynomials

PolySqrt.h

Description: find sqrt of polynomial

Time: $\mathcal{O}(N \log N)$?

5.3 Misc

LinRec.h

Description: Berlekamp-Massey: computes linear recurrence of order n for sequence of 2n terms

```
Time: ?
                                                         35 lines
using namespace vecOp;
struct LinRec {
 vmi x; // original sequence
  vmi C, rC;
  void init(const vmi& _x) {
   x = x; int n = sz(x), m = 0;
    vmi B; B = C = \{1\}; // B is fail vector
    mi b = 1; // B gives 0,0,0,...,b
    FOR(i,n) {
      mi d = x[i]; FOR(\dot{j}, 1, sz(C)) d += C[\dot{j}] *x[i-\dot{j}];
      if (d == 0) continue; // recurrence still works
      auto B = C; C.rsz(max(sz(C), m+sz(B)));
      mi coef = d/b; FOR(j,m,m+sz(B)) C[j] -= coef*B[j-m];
         \hookrightarrow // recurrence that gives 0,0,0,...,d
      if (sz(B) < m+sz(B)) \{ B = B; b = d; m = 0; \}
    rC = C; reverse(all(rC)); // polynomial for getPo
    C.erase(begin(C)); trav(t,C) t *=-1; // x[i]=sum_{i}
       \hookrightarrow =0} \{sz(C)-1\}C[j]*x[i-j-1]
  vmi getPo(int n) {
   if (n == 0) return {1};
```

```
vmi x = getPo(n/2); x = rem(x*x,rC);
if (n&l) { vmi v = {0,1}; x = rem(x*v,rC); }
  return x;
}
mi eval(int n) {
  vmi t = getPo(n);
  mi ans = 0; FOR(i,sz(t)) ans += t[i]*x[i];
  return ans;
};
```

Integrate.h Description: ?

```
// db f(db x) { return x*x+3*x+1; }

db quad(db (*f)(db), db a, db b) {
  const int n = 1000;
  db dif = (b-a)/2/n, tot = f(a)+f(b);
  FOR(i,1,2*n) tot += f(a+i*dif)*(i&1?4:2);
  return tot*dif/3;
}
```

IntegrateAdaptive.h

Description: ?

19 lines

8 lines

Simplex.h

Time: ?

Description: Simplex algorithm for linear programming, maximize $c^T x$ subject to $Ax \leq b, x \geq 0$

```
struct LPSolver {
 int m, n;
 vi N, B;
 vvd D;
 LPSolver(const vvd& A, const vd& b, const vd& c) :
   m(sz(b)), n(sz(c)), N(n+1), B(m), D(m+2), vd(n+2)) {
      FOR(i,m) FOR(j,n) D[i][j] = A[i][j];
      FOR(i,m) \{ B[i] = n+i; D[i][n] = -1; D[i][n+1] = b[i] \}
         \hookrightarrow]; } // B[i] -> basic variables, col n+1 is for
         \hookrightarrow constants, why D[i][n]=-1?
      FOR(j,n) \{ N[j] = j; D[m][j] = -c[j]; \} // N[j] ->

→non-basic variables, all zero

      N[n] = -1; D[m+1][n] = 1;
  void print() {
   ps("D");
   trav(t,D) ps(t);
   ps();
   ps("B",B);
   ps("N",N);
   ps();
  void pivot(int r, int s) { // row, column
   T *a = D[r].data(), inv = 1/a[s]; // eliminate col s
       \hookrightarrowfrom consideration
    FOR(i,m+2) if (i != r \&\& abs(D[i][s]) > eps) {
     T *b = D[i].data(), inv2 = b[s]*inv;
      FOR(j,n+2) b[j] -= a[j]*inv2;
     b[s] = a[s] * inv2;
    FOR(j,n+2) if (j != s) D[r][j] *= inv;
    FOR(i,m+2) if (i != r) D[i][s] *= -inv;
   D[r][s] = inv; swap(B[r], N[s]); // swap a basic and
       →non-basic variable
  bool simplex(int phase) {
   int x = m + phase - 1;
    for (;;) {
      int s = -1; FOR(j, n+1) if (N[j] != -phase) ltj(D[x])
         \hookrightarrow; // find most negative col
      if (D[x][s] >= -eps) return true; // have best
         \hookrightarrowsolution
      int r = -1;
      FOR(i,m) {
        if (D[i][s] <= eps) continue;</pre>
        if (r == -1 \mid | mp(D[i][n+1] / D[i][s], B[i])
               < mp(D[r][n+1] / D[r][s], B[r])) r = i; //

→ find smallest positive ratio

      if (r == -1) return false; // unbounded
     pivot(r, s);
 T solve(vd &x) {
```

```
int r = 0; FOR(i,1,m) if (D[i][n+1] < D[r][n+1]) r = i
    if (D[r][n+1] < -eps)  { // x=0 is not a solution
      pivot(r, n); // -1 is artificial variable, initially
          \hookrightarrow set to smth large but want to get to 0
      if (!simplex(2) || D[m+1][n+1] < -eps) return -inf;</pre>
         \hookrightarrow // no solution
      // D[m+1][n+1] is max possible value of the negation
         \hookrightarrow of artificial variable, starts negative but
         \hookrightarrowshould get to zero
      FOR(i,m) if (B[i] == -1) {
        int s = 0; FOR(j, 1, n+1) ltj(D[i]);
        pivot(i,s);
    bool ok = simplex(1); x = vd(n);
    FOR(i,m) if (B[i] < n) x[B[i]] = D[i][n+1];
    return ok ? D[m][n+1] : inf;
};
```

Graphs (6)

6.1 Fundamentals

DSU.h

};

Description: ? Time: $\mathcal{O}(N\alpha(N))$

```
29 lines
template<int SZ> struct DSU {
    int par[SZ];
    int size[SZ]:
       M00(i, SZ) par[i] = i, size[i] = 1;
    int get(int node) {
        if (par[node] != node) par[node] = get (par[node]);
        return par[node];
    bool connected(int n1, int n2) {
        return (get(n1) == get(n2));
    int sz(int node) {
        return size[get(node)];
    void unite(int n1, int n2) {
       n1 = get(n1);
        n2 = get(n2);
        if(n1 == n2) return;
        if(rand()%2) {
            par[n1] = n2;
            size[n2] += size[n1];
        } else {
            par[n2] = n1;
            size[n1] += size[n2];
```

ManhattanMST.h

 $\bf Description:$ Compute minimum spanning tree of points where edges are manhattan distances

Time: $\mathcal{O}(N \log N)$

```
"MST.h"
                                                          60 lines
int N;
vector<arrav<int,3>> cur;
vector<pair<ll,pi>> ed;
vi ind;
struct {
 map<int,pi> m;
 void upd(int a, pi b) {
    auto it = m.lb(a);
    if (it != m.end() && it->s <= b) return;
    m[a] = b; it = m.find(a);
    while (it != m.begin() && prev(it)->s >= b) m.erase(
       \hookrightarrowprev(it));
 pi query(int y) { // for all a > y find min possible
     \hookrightarrowvalue of b
    auto it = m.ub(y);
    if (it == m.end()) return {2*MOD,2*MOD};
    return it->s;
} S;
void solve() {
 sort(all(ind),[](int a, int b) { return cur[a][0] > cur[
     \hookrightarrowbl[0]; });
  S.m.clear();
  int nex = 0:
  trav(x, ind) { // cur[x][0] <= ?, cur[x][1] < ?}
    while (nex < N \&\& cur[ind[nex]][0] >= cur[x][0])  {
      int b = ind[nex++];
      S.upd(cur[b][1], {cur[b][2],b});
    pi t = S.query(cur[x][1]);
    if (t.s != 2*MOD) ed.pb({(11)t.f-cur[x][2],{x,t.s}});
ll mst(vpi v) {
 N = sz(v); cur.resz(N); ed.clear();
  ind.clear(); FOR(i,N) ind.pb(i);
  sort(all(ind),[&v](int a, int b) { return v[a] < v[b];</pre>
  FOR(i, N-1) if (v[ind[i]] == v[ind[i+1]]) ed.pb(\{0, \{ind[i]\}\})
     \hookrightarrow],ind[i+1]}});
  FOR(i,2) { // it's probably ok to consider just two
     \hookrightarrow quadrants?
    FOR(i,N) {
      auto a = v[i];
      cur[i][2] = a.f+a.s;
    FOR(i,N) { // first octant
      auto a = v[i];
      cur[i][0] = a.f-a.s;
      cur[i][1] = a.s;
```

Dijkstra FloydWarshall LCAjumps LCArmq

```
solve();
 FOR(i,N) { // second octant
   auto a = v[i];
   cur[i][0] = a.f;
   cur[i][1] = a.s-a.f;
 trav(a,v) a = {a.s,-a.f}; // rotate 90 degrees, repeat
return kruskal (ed):
```

Diikstra.h

Description: Dijkstra's algorithm for shortest path

Time: $\mathcal{O}\left(E\log V\right)$

```
template<int SZ> struct dijkstra {
   vector<pair<int, ll>> adj[SZ];
   bool vis[SZ];
   11 d[SZ];
    void addEdge(int u, int v, ll l) {
        adj[u].PB(MP(v, 1));
   11 dist(int v) {
        return d[v];
   void build(int u) {
       M00(i, SZ) vis[i] = 0;
        priority_queue<pair<ll, int>, vector<pair<ll, int
          ⇔>>, greater<pair<11, int>>> pg;
        M00(i, SZ) d[i] = 1e17;
        d[u] = 0;
        pq.push(MP(0, u));
        while(!pq.empty()) {
           pair<11, int> t = pq.top(); pq.pop();
            while(!pq.empty() && vis[t.S]) t = pq.top(),
               \hookrightarrowpq.pop();
            vis[t.S] = 1;
            for(auto& v: adj[t.S]) if(!vis[v.F]) {
                if(d[v.F] > d[t.S] + v.S) {
                    d[v.F] = d[t.S] + v.S;
                    pq.push(MP(d[v.F], v.F));
};
```

FlovdWarshall.h

Description: Floyd Warshall's algorithm for all pairs shortest path Time: $\mathcal{O}(V^3)$

let dist be a |V| * |V| array of minimum distances \hookrightarrow initialized to inf for each edge (u, v) do

```
dist[u][v] \leftarrow w(u, v) // The weight of the edge (u, v
for each vertex v do
   dist[v][v] <- 0
for k from 1 to |V|
   for i from 1 to IVI
       for j from 1 to |V|
           if dist[i][j] > dist[i][k] + dist[k][j]
                dist[i][j] \leftarrow dist[i][k] + dist[k][j]
            end if
```

6.2 Trees

LCAjumps.h

Description: calculates least common ancestor in tree with binary jumping

```
Time: \mathcal{O}(N \log N)
                                                        44 lines
template<int SZ> struct tree {
    vector<pair<int, ll>> adj[SZ];
    const static int LGSZ = 32-__builtin_clz(SZ-1);
    pair<int, 11> ppar[SZ][LGSZ];
    int depth[SZ];
    11 distfromroot[SZ];
    void addEdge(int u, int v, int d) {
        adj[u].PB(MP(v, d));
        adj[v].PB(MP(u, d));
    void dfs(int u, int dep, ll dis) {
        depth[u] = dep;
        distfromroot[u] = dis;
        for(auto& v: adj[u]) if(ppar[u][0].F != v.F) {
            ppar[v.F][0] = MP(u, v.S);
            dfs(v.F, dep + 1, dis + v.S);
    void build() {
        ppar[0][0] = MP(0, 0);
        M00(i, SZ) depth[i] = 0;
        dfs(0, 0, 0);
        MOO(i, 1, LGSZ) MOO(j, SZ) {
            ppar[j][i].F = ppar[ppar[j][i-1].F][i-1].F;
            ppar[j][i].S = ppar[j][i-1].S + ppar[ppar[j][i]
               \hookrightarrow-1].F][i-1].S;
    int lca(int u, int v) {
        if(depth[u] < depth[v]) swap(u, v);</pre>
        M00d(i, LGSZ) if(depth[ppar[u][i].F] >= depth[v])
           \hookrightarrowu = ppar[u][i].F;
        if(u == v) return u;
        M00d(i, LGSZ) {
            if(ppar[u][i].F != ppar[v][i].F) {
                u = ppar[u][i].F;
                v = ppar[v][i].F;
        return ppar[u][0].F;
```

```
11 dist(int u, int v) {
      return distfromroot[u] + distfromroot[v] - 2*
        };
```

LCArmq.h

```
Description: Euler Tour LCA w/ O(1) query
                                                        58 lines
template<int SZ> struct tree +
    vector<pair<int, ll>> adi[SZ];
    pair<int, ll> par[SZ];
    const static int LGSZ = 33-__builtin_clz(SZ-1);
    11 distfromroot[SZ];
    int depth[SZ], t, tin[SZ], RMQ[2*SZ-1][LGSZ], oldToNew
       \hookrightarrow [SZ], newToOld[SZ], numNodes;
    void addEdge(int u, int v, int d) {
        adj[u].PB(MP(v, d));
        adj[v].PB(MP(u, d));
    void dfs(int u, int dep, ll dis) {
        depth[u] = dep;
        distfromroot[u] = dis;
        for(auto& v: adj[u]) if(par[u].F != v.F) {
            par[v.F] = MP(u, v.S);
            dfs(v.F, dep + 1, dis + v.S);
    void buildtarr(int u) {
        RMQ[t][0] = oldToNew[u], tin[oldToNew[u]] = t++;
        for(auto& v: adj[u]) if(par[u].F != v.F) {
            buildtarr(v.F);
            RMQ[t++][0] = oldToNew[u];
    void build(int n) {
        this->numNodes = n;
        par[0] = MP(0, 0);
        M00(i, numNodes) depth[i] = 0;
        dfs(0, 0, 0);
        t = 0;
        queue<int> q;
        q.push(0);
        while(!q.empty()) {
            int u = q.front(); q.pop();
            oldToNew[u] = t++;
            for(auto& v: adj[u]) if(par[u].F != v.F) q.
               \hookrightarrow push (v.F);
        M00(i, numNodes) newToOld[oldToNew[i]] = i;
        t = 0;
        buildtarr(0);
        MOO(j, 1, LGSZ) MOO(i, 2*numNodes-1) if(i+(1<<(j
           \hookrightarrow-1)) < 2*numNodes-1)
            RMQ[i][j] = min(RMQ[i][j-1], RMQ[i+(1<<(j-1))
               →][j-1]);
    int lca(int u, int v) {
```

CentroidDecomp HLD SCC TopoSort

```
u = oldToNew[u], v = oldToNew[v];
        if(tin[u] > tin[v]) swap(u, v);
        int 1 = tin[u];
        int r = tin[v];
        int len = r-1+1;
        int hl = 31-__builtin_clz(len-1);
        return newToOld[min(RMO[1][h1], RMO[r-(1<<h1)+1][
    11 dist(int u, int v) {
        return distfromroot[u]+distfromroot[v]-2*

→distfromroot[lca(u, v)];
};
```

CentroidDecomp.h

Description: can support tree path queries and updates

Time: $\mathcal{O}(N \log N)$

```
47 lines
template<int SZ> struct centroidDecomp {
    vi neighbor[SZ];
    int subsize[SZ];
    bool vis[SZ];
    int p[SZ];
    int par[SZ];
    vi child[SZ];
    int numNodes;
    centroidDecomp(int num) {
        this->numNodes = num;
    void addEdge(int u, int v) {
        neighbor[u].PB(v);
        neighbor[v].PB(u);
    void build() {
        M00(i, numNodes) vis[i] = 0, par[i] = -1;
        solve(0);
        M00(i, numNodes) if (par[i] != -1) child[par[i]].PB
           \hookrightarrow (i);
    void getSizes(int node) {
        subsize[node] = 1;
        for(int ch: neighbor[node]) if(!vis[ch] && ch != p
           \hookrightarrow [node]) {
            p[ch] = node;
            getSizes(ch);
            subsize[node] += subsize[ch];
    int getCentroid(int root) {
        p[root] = -1;
        getSizes(root);
        int cur = root;
        while(1) {
            pi hi = MP(subsize[root]-subsize[cur], cur);
            for(int v: neighbor[cur]) if(!vis[v] && v != p
                \hookrightarrow [cur]) hi = max(hi, MP(subsize[v], v));
            if(hi.F <= subsize[root]/2) return cur;</pre>
            cur = hi.S:
```

```
int solve(int node) {
        node = getCentroid(node);
        vis[node] = 1;
         for(int ch: neighbor[node]) if(!vis[ch]) par[solve
            \hookrightarrow (ch) | = node;
         return node:
};
```

HLD.h

Description: Heavy Light Decomposition

```
Time: \mathcal{O}(\log^2 N) per path operations
                                                       50 lines
template<int SZ, bool VALUES_IN_EDGES> struct HLD {
  int N; vi adi[SZ];
  int par[SZ], sz[SZ], depth[SZ];
  int root[SZ], pos[SZ];
  LazySegTree<11,SZ> tree;
  void addEdge(int a, int b) { adj[a].pb(b), adj[b].pb(a);
  void dfs_sz(int v = 1) {
   if (par[v]) adj[v].erase(find(all(adj[v]),par[v]));
    sz[v] = 1;
    trav(u,adj[v]) {
      par[u] = v; depth[u] = depth[v]+1;
      dfs_sz(u); sz[v] += sz[u];
      if (sz[u] > sz[adj[v][0]]) swap(u, adj[v][0]);
  void dfs_hld(int v = 1) {
    static int t = 0;
    pos[v] = t++;
    trav(u,adj[v]) {
      root[u] = (u == adj[v][0] ? root[v] : u);
      dfs_hld(u);
  void init(int _N) {
   N = N; par[1] = depth[1] = 0; root[1] = 1;
   dfs sz(); dfs hld();
  template <class BinaryOperation>
  void processPath(int u, int v, BinaryOperation op) {
   for (; root[u] != root[v]; v = par[root[v]]) {
      if (depth[root[u]] > depth[root[v]]) swap(u, v);
      op(pos[root[v]], pos[v]);
    if (depth[u] > depth[v]) swap(u, v);
    op(pos[u]+VALUES_IN_EDGES, pos[v]);
  void modifyPath(int u, int v, int val) { // add val to
     →vertices/edges along path
    processPath(u, v, [this, &val](int 1, int r) { tree.
       \hookrightarrowupd(1, r, val); });
```

```
void modifySubtree(int v, int val) { // add val to

→vertices/edges in subtree

    tree.upd(pos[v]+VALUES IN EDGES,pos[v]+sz[v]-1,val);
  11 queryPath(int u, int v) { // query sum of path
    11 res = 0; processPath(u, v, [this, &res](int 1, int
       \hookrightarrowr) { res += tree.gsum(1, r); });
    return res;
};
```

DFS Algorithms

SCC.h

Description: Kosaraju's Algorithm: DFS two times to generate SCCs in topological order

Time: $\mathcal{O}(N+M)$

24 lines

```
template<int SZ> struct SCC {
 int N. comp[SZ];
  vi adj[SZ], radj[SZ], todo, allComp;
  bitset<SZ> visit;
  void addEdge(int a, int b) { adj[a].pb(b), radj[b].pb(a)
    \hookrightarrow; }
  void dfs(int v) {
    visit[v] = 1:
    trav(w,adj[v]) if (!visit[w]) dfs(w);
    todo.pb(v);
  void dfs2(int v, int val) {
    comp[v] = val;
    trav(w, radj[v]) if (comp[w] == -1) dfs2(w, val);
 void init(int N) { // fills allComp
    N = N;
    FOR(i,N) comp[i] = -1, visit[i] = 0;
    FOR(i,N) if (!visit[i]) dfs(i);
    reverse(all(todo)); // now todo stores vertices in
       →order of topological sort
    trav(i,todo) if (comp[i] == -1) dfs2(i,i), allComp.pb(
       \hookrightarrowi);
};
```

TopoSort.h

Description: sorts vertices such that if there exists an edge x->v, then x goes before v

```
template<int SZ> struct TopoSort {
   int N, in[SZ];
   vi res, adj[SZ];
   void ae(int x, int y) { adj[x].pb(y), in[y] ++; }
   bool sort(int _N) {
       N = _N; queue<int> todo;
       FOR(i,1,N+1) if (!in[i]) todo.push(i);
        while (sz(todo)) {
            int x = todo.front(); todo.pop(); res.pb(x);
            trav(i,adj[x]) if (!(--in[i])) todo.push(i);
```

2SAT EulerPath BCC Dinic

```
return sz(res) == N;
};
```

2SAT.h Description: ?

38 lines template<int SZ> struct TwoSat { SCC<2*SZ> S: bitset<SZ> ans; int N = 0; int addVar() { return N++; } void either(int x, int y) { $x = \max(2*x, -1-2*x), y = \max(2*y, -1-2*y);$ S.addEdge(x^1, y); S.addEdge(y^1, x); void implies (int x, int y) { either $(\sim x, y)$; } void setVal(int x) { either(x,x); } void atMostOne(const vi& li) { if (sz(li) <= 1) return; int cur = \sim li[0]; FOR(i,2,sz(li)) { int next = addVar(); either(cur,~li[i]); either(cur,next); either(~li[i],next); cur = ~next; either(cur,~li[1]); bool solve(int N) { if (N != -1) N = N;S.init(2*N); for (int i = 0; i < 2*N; i += 2) if (S.comp[i] == S.comp[i^1]) return 0; reverse(all(S.allComp)); vi tmp(2*N); trav(i, S.allComp) if (tmp[i] == 0) $tmp[i] = 1, tmp[S.comp[i^1]] = -1;$ FOR(i,N) if (tmp[S.comp[2*i]] == 1) ans[i] = 1; return 1; };

EulerPath.h

Description: Eulerian Path for both directed and undirected graphs Time: $\mathcal{O}(N+M)$

```
template<int SZ, bool directed> struct Euler {
 int N, M = 0;
 vpi adj[SZ];
 vpi::iterator its[SZ];
 vector<bool> used;
 void addEdge(int a, int b) {
   if (directed) adj[a].pb({b,M});
```

```
else adj[a].pb({b,M}), adj[b].pb({a,M});
    used.pb(0); M ++;
  vpi solve(int _N, int src = 1) {
    N = N;
    FOR(i,1,N+1) its[i] = begin(adj[i]);
    vector<pair<pi, int>> ret, s = \{\{\{src, -1\}, -1\}\};
    while (sz(s)) {
      int x = s.back().f.f;
      auto& it = its[x], end = adj[x].end();
      while (it != end && used[it->s]) it ++;
      if (it == end) {
        if (sz(ret) && ret.back().f.s != s.back().f.f)
           →return {}; // path isn't valid
        ret.pb(s.back()), s.pop_back();
      } else { s.pb({{it->f,x},it->s}); used[it->s] = 1; }
    if (sz(ret) != M+1) return {};
    vpi ans; trav(t,ret) ans.pb({t.f.f,t.s});
    reverse(all(ans)); return ans;
};
```

BCC.h

return low;

Description: computes biconnected components Time: $\mathcal{O}(N+M)$

```
37 lines
template<int SZ> struct BCC {
  int N:
  vpi adj[SZ], ed;
  void addEdge(int u, int v) {
    adj[u].pb({v,sz(ed)}), adj[v].pb({u,sz(ed)});
    ed.pb({u,v});
  int disc[SZ];
  vi st; vector<vi> fin;
  int bcc(int u, int p = -1) { // return lowest disc
    static int ti = 0;
    disc[u] = ++ti; int low = disc[u];
    int child = 0;
    trav(i,adj[u]) if (i.s != p)
      if (!disc[i.f]) {
        child ++; st.pb(i.s);
        int LOW = bcc(i.f,i.s); ckmin(low,LOW);
        // disc[u] < LOW -> bridge
        if (disc[u] <= LOW) {</pre>
          // if (p != -1 || child > 1) -> u is
             \hookrightarrowarticulation point
          vi tmp; while (st.back() != i.s) tmp.pb(st.back
             \hookrightarrow ()), st.pop_back();
          tmp.pb(st.back()), st.pop_back();
          fin.pb(tmp);
      } else if (disc[i.f] < disc[u]) {</pre>
        ckmin(low,disc[i.f]);
        st.pb(i.s);
```

```
void init(int N) {
   N = N; FOR(i,N) disc[i] = 0;
   FOR(i,N) if (!disc[i]) bcc(i); // st should be empty

→after each iteration

};
```

6.4 Flows

Dinic.h

Description: faster flow

Time: $\mathcal{O}(N^2M)$ flow, $\mathcal{O}(M\sqrt{N})$ bipartite matching

```
45 lines
template<int SZ> struct Dinic {
 typedef 11 F; // flow type
 struct Edge { int to, rev; F flow, cap; };
 int N.s.t:
 vector<Edge> adi[SZ];
 typename vector<Edge>::iterator cur[SZ];
 void addEdge(int u, int v, F cap) {
   assert(cap >= 0); // don't try smth dumb
   Edge a\{v, sz(adj[v]), 0, cap\}, b\{u, sz(adj[u]), 0, 0\};
   adj[u].pb(a), adj[v].pb(b);
 int level[SZ];
 bool bfs() { // level = shortest distance from source
   // after computing flow, edges {u,v} such that level[u
      \hookrightarrow \neq -1, level[v] = -1 are part of min cut
   M00(i,N) level[i] = -1, cur[i] = begin(adj[i]);
   queue<int> q({s}); level[s] = 0;
   while (sz(q)) {
     int u = q.front(); q.pop();
            for(Edge e: adj[u]) if (level[e.to] < 0 && e.</pre>
               \hookrightarrowflow < e.cap)
        q.push(e.to), level[e.to] = level[u]+1;
   return level[t] >= 0;
 F sendFlow(int v, F flow) {
   if (v == t) return flow;
   for (; cur[v] != end(adj[v]); cur[v]++) {
     Edge& e = *cur[v];
     if (level[e.to] != level[v]+1 || e.flow == e.cap)
         →continue;
      auto df = sendFlow(e.to,min(flow,e.cap-e.flow));
     if (df) { // saturated at least one edge
        e.flow += df; adj[e.to][e.rev].flow -= df;
        return df;
   return 0;
 F maxFlow(int _N, int _s, int _t) {
   N = N, s = s, t = t; if (s == t) return -1;
   F tot = 0:
```

MCMF GomoryHu DFSmatch Hungarian

```
while (bfs()) while (auto df = sendFlow(s,
   return tot:
};
```

MCMF.h

Description: Min-Cost Max Flow, no negative cycles allowed Time: $\mathcal{O}(NM^2 \log M)$

53 lines

```
template < class T > using pgg = priority_queue < T, vector < T >,
   template<class T> T poll(pqg<T>& x) {
 T y = x.top(); x.pop();
 return y;
template<int SZ> struct mcmf {
 typedef 11 F; typedef 11 C;
 struct Edge { int to, rev; F flow, cap; C cost; int id;
    \hookrightarrow};
  vector<Edge> adi[SZ];
  void addEdge(int u, int v, F cap, C cost) {
    assert(cap >= 0);
    Edge a\{v, sz(adj[v]), 0, cap, cost\}, b\{u, sz(adj[u]),
       \hookrightarrow0, 0, -cost};
    adj[u].pb(a), adj[v].pb(b);
 int N, s, t;
 pi pre[SZ]; // previous vertex, edge label on path
 pair<C,F> cost[SZ]; // tot cost of path, amount of flow
  C totCost, curCost; F totFlow;
  void reweight() { // makes all edge costs non-negative
    // all edges on shortest path become 0
   FOR(i,N) trav(p,adj[i]) p.cost += cost[i].f-cost[p.to
       \hookrightarrow1.f;
  bool spfa() { // reweight ensures that there will be
    \hookrightarrownegative weights
    // only during the first time you run this
    FOR(i,N) cost[i] = {INF,0}; cost[s] = {0,INF};
    pgg<pair<C, int>> todo; todo.push({0,s});
    while (sz(todo)) {
      auto x = poll(todo); if (x.f > cost[x.s].f) continue
         \hookrightarrow ;
      trav(a,adj[x.s]) if (x.f+a.cost < cost[a.to].f && a.
         \hookrightarrowflow < a.cap) {
        // if costs are doubles, add some EPS to ensure
           \hookrightarrow that
        // you do not traverse some 0-weight cycle
           \hookrightarrowrepeatedly
        pre[a.to] = {x.s,a.rev};
        cost[a.to] = \{x.f+a.cost, min(a.cap-a.flow, cost[x.s])\}
        todo.push({cost[a.to].f,a.to});
    curCost += cost[t].f; return cost[t].s;
```

```
void backtrack() {
   F df = cost[t].s; totFlow += df, totCost += curCost*df
    for (int x = t; x != s; x = pre[x].f) {
      adj[x][pre[x].s].flow -= df;
     adj[pre[x].f][adj[x][pre[x].s].rev].flow += df;
  pair<F,C> calc(int _N, int _s, int _t) {
   N = N; s = s, t = t; totFlow = totCost = curCost = s
    while (spfa()) reweight(), backtrack();
   return {totFlow, totCost};
};
```

GomorvHu.h

Description: Compute max flow between every pair of vertices of undirected graph

```
"Dinic.h"
                                                        56 lines
template<int SZ> struct GomorvHu {
 int N;
  vector<pair<pi,int>> ed;
  void addEdge(int a, int b, int c) { ed.pb({{a,b},c}); }
  vector<vi> cor = {{}}; // groups of vertices
  map<int,int> adj[2*SZ]; // current edges of tree
  int side[SZ];
  int gen(vector<vi> cc) {
   Dinic<SZ> D = Dinic<SZ>();
    vi comp(N+1); FOR(i,sz(cc)) trav(t,cc[i]) comp[t] = i;
    trav(t,ed) if (comp[t.f.f] != comp[t.f.s]) {
      D.addEdge(comp[t.f.f],comp[t.f.s],t.s);
      D.addEdge(comp[t.f.s],comp[t.f.f],t.s);
    int f = D.maxFlow(0,1);
    FOR(i,sz(cc)) trav(j,cc[i]) side[j] = D.level[i] >= 0;
       \hookrightarrow // min cut
    return f;
  void fill(vi& v, int a, int b) {
   trav(t,cor[a]) v.pb(t);
    trav(t,adj[a]) if (t.f != b) fill (v,t.f,a);
  void addTree(int a, int b, int c) { adj[a][b] = c, adj[b
  void delTree(int a, int b) { adj[a].erase(b), adj[b].
     \hookrightarrowerase(a); }
  vector<pair<pi,int>> init(int _N) { // returns edges of
     \hookrightarrow Gomory-Hu Tree
    N = N;
    FOR(i,1,N+1) cor[0].pb(i);
    queue<int> todo; todo.push(0);
    while (sz(todo)) {
      int x = todo.front(); todo.pop();
```

```
vector<vi> cc; trav(t,cor[x]) cc.pb({t});
      trav(t,adj[x]) {
        : ({})dq.pp
        fill(cc.back(),t.f,x);
      int f = gen(cc); // run max flow
      cor.pb({}), cor.pb({});
      trav(t, cor[x]) cor[sz(cor)-2+side[t]].pb(t);
      FOR(i,2) if (sz(cor[sz(cor)-2+i]) > 1) todo.push(sz(
         \hookrightarrowcor)-2+i);
      FOR(i,sz(cor)-2) if (i != x \&\& adj[i].count(x)) {
        addTree(i,sz(cor)-2+side[cor[i][0]],adj[i][x]);
        delTree(i,x);
      } // modify tree edges
      addTree(sz(cor)-2,sz(cor)-1,f);
    vector<pair<pi, int>> ans;
    FOR(i,sz(cor)) trav(j,adj[i]) if (i < j.f)
      ans.pb({{cor[i][0],cor[j.f][0]},j.s});
    return ans;
};
```

6.5 Matching

DFSmatch.h

Description: naive bipartite matching

Time: $\mathcal{O}(NM)$

26 lines

```
template<int SZ> struct MaxMatch {
 int N, flow = 0, match[SZ], rmatch[SZ];
  bitset<SZ> vis;
  vi adi[SZ];
 MaxMatch() {
    memset (match, 0, sizeof match);
    memset (rmatch, 0, sizeof rmatch);
  void connect(int a, int b, bool c = 1) {
    if (c) match[a] = b, rmatch[b] = a;
    else match[a] = rmatch[b] = 0;
 bool dfs(int x) {
    if (!x) return 1;
    if (vis[x]) return 0:
    vis[x] = 1;
    trav(t,adj[x]) if (t != match[x] && dfs(rmatch[t]))
     return connect(x,t),1;
    return 0:
  void tri(int x) { vis.reset(); flow += dfs(x); }
  void init(int N) {
    N = N; FOR(i, 1, N+1) if (!match[i]) tri(i);
};
```

Hungarian.h

Description: finds min cost to complete n jobs w/ m workers each worker is assigned to at most one job $(n \le m)$

21 lines

UnweightedMatch MaximalCliques LCT

```
\mathbf{Time:}\ ?
int HungarianMatch(const vector<vi>& a) { // cost array,

→negative values are ok

  int n = sz(a)-1, m = sz(a[0])-1; // jobs 1..., workers
  vi u(n+1), v(m+1), p(m+1); // p[j] -> job picked by
    →worker j
  FOR(i,1,n+1) { // find alternating path with job i
   p[0] = i; int j0 = 0;
   vi dist(m+1, MOD), pre(m+1,-1); // dist, previous
      vector<bool> done(m+1, false);
     done[j0] = true;
     int i0 = p[j0], j1; int delta = MOD;
     FOR(j,1,m+1) if (!done[j]) {
        auto cur = a[i0][j]-u[i0]-v[j];
       if (cur < dist[j]) dist[j] = cur, pre[j] = j0;</pre>
       if (dist[j] < delta) delta = dist[j], j1 = j;</pre>
      FOR(j,m+1) // just dijkstra with potentials
        if (done[i]) u[p[i]] += delta, v[i] -= delta;
        else dist[j] -= delta;
      j0 = j1;
    } while (p[i0]);
   do { // update values on alternating path
     int j1 = pre[j0];
     p[j0] = p[j1];
     j0 = j1;
   } while (j0);
 return -v[0]; // min cost
```

UnweightedMatch.h

Description: general unweighted matching

pv = par[v]; nv = match[pv];

```
Time: ?
                                                       79 lines
template<int SZ> struct UnweightedMatch {
 int vis[SZ], par[SZ], orig[SZ], match[SZ], aux[SZ], t, N
    \hookrightarrow; // 1-based index
 vi adj[SZ];
 queue<int> Q;
 void addEdge(int u, int v) {
   adj[u].pb(v); adj[v].pb(u);
  void init(int n) {
   N = n; t = 0;
   FOR(i,N+1) {
      adj[i].clear();
      match[i] = aux[i] = par[i] = 0;
  void augment(int u, int v) {
   int pv = v, nv;
```

```
match[v] = pv; match[pv] = v;
   v = nv;
  } while(u != pv);
int lca(int v, int w) {
 ++t;
 while (1) {
   if (v) {
     if (aux[v] == t) return v; aux[v] = t;
     v = orig[par[match[v]]];
    swap(v, w);
void blossom(int v, int w, int a) {
 while (orig[v] != a) {
   par[v] = w; w = match[v];
   if (vis[w] == 1) Q.push(w), vis[w] = 0;
   orig[v] = orig[w] = a;
   v = par[w];
bool bfs(int u) {
 fill(vis+1, vis+1+N, -1); iota(orig + 1, orig + N + 1,
 Q = queue < int > (); Q.push(u); vis[u] = 0;
 while (sz(Q)) {
   int v = Q.front(); Q.pop();
   trav(x,adj[v]) {
     if (vis[x] == -1) {
        par[x] = v; vis[x] = 1;
        if (!match[x]) return augment(u, x), true;
        Q.push(match[x]); vis[match[x]] = 0;
      } else if (vis[x] == 0 && orig[v] != orig[x]) {
        int a = lca(orig[v], orig[x]);
        blossom(x, v, a); blossom(v, x, a);
  return false;
int match() {
 int ans = 0;
  // find random matching (not necessary, constant
     →improvement)
 vi V(N-1); iota(all(V), 1);
  shuffle(all(V), mt19937(0x94949));
 trav(x, V) if(!match[x])
   trav(v,adj[x]) if (!match[v]) {
     match[x] = y, match[y] = x;
      ++ans; break;
 FOR(i,1,N+1) if (!match[i] && bfs(i)) ++ans;
  return ans;
```

6.6 Misc

MaximalCliques.h

Description: Ûsed only once. Finds all maximal cliques. **Time:** $\mathcal{O}\left(3^{N/3}\right)$

```
typedef bitset<128> B;
int N;
B adj[128];

// possibly in clique, not in clique, in clique
void cliques(B P = ~B(), B X={}, B R={}) {
   if (!P.any()) {
      if (!X.any()) {
        // do smth with R
      }
      return;
   }
   int q = (P|X)._Find_first();
   // clique must contain q or non-neighbor of q
   B cands = P&~adj[q];
   FOR(i,N) if (cands[i]) {
      R[i] = 1;
      cliques(P&adj[i],X&adj[i],R);
```

LCT.h

Description: Link-Cut Tree, use vir for subtree size queries **Time:** $\mathcal{O}(\log N)$

R[i] = P[i] = 0; X[i] = 1;

96 lines

```
typedef struct snode* sn;
struct snode {
 sn p, c[2]; // parent, children
  int val; // value in node
  int sum, mn, mx; // sum of values in subtree, min and
    \hookrightarrow max prefix sum
 bool flip = 0;
  // int vir = 0; stores sum of virtual children
  snode(int v) {
   p = c[0] = c[1] = NULL;
    val = v; calc();
  friend int getSum(sn x) { return x?x->sum:0; }
  friend int getMn(sn x) { return x?x->mn:0;
  friend int getMx(sn x) { return x?x->mx:0; }
  void prop() {
    if (!flip) return;
    swap(c[0],c[1]); tie(mn,mx) = mp(sum-mx,sum-mn);
    FOR(i,2) if (c[i]) c[i]->flip ^= 1;
    flip = 0;
  void calc() {
```

DirectedMST DominatorTree

```
FOR(i,2) if (c[i]) c[i]->prop();
  int s0 = getSum(c[0]), s1 = getSum(c[1]); sum = s0+val
     →+s1: // +vir
 mn = min(getMn(c[0]), s0+val+getMn(c[1]));
  mx = max(qetMx(c[0]), s0+val+qetMx(c[1]));
int dir() {
 if (!p) return -2;
  FOR(i,2) if (p->c[i] == this) return i;
  return -1; // p is path-parent pointer, not in current
     \hookrightarrow splay tree
bool isRoot() { return dir() < 0; }</pre>
friend void setLink(sn x, sn y, int d) {
 if (y) y->p = x;
 if (d >= 0) x -> c[d] = y;
void rot() { // assume p and p->p propagated
 assert(!isRoot()); int x = dir(); sn pa = p;
  setLink(pa->p, this, pa->dir());
  setLink(pa, c[x^1], x);
  setLink(this, pa, x^1);
  pa->calc(); calc();
void splay() {
  while (!isRoot() && !p->isRoot()) {
    p->p->prop(), p->prop(), prop();
    dir() == p->dir() ? p->rot() : rot();
    rot();
  if (!isRoot()) p->prop(), prop(), rot();
  prop();
void access() { // bring this to top of tree
  for (sn v = this, pre = NULL; v; v = v->p) {
    v->splay();
    // if (pre) v->vir -= pre->sz;
    // if (v->c[1]) v->vir += v->c[1]->sz;
    v->c[1] = pre; v->calc();
    pre = v;
    // v->sz should remain the same if using vir
  splay(); assert(!c[1]); // left subtree of this is now
     \hookrightarrow path to root, right subtree is empty
void makeRoot() { access(); flip ^= 1; }
void set(int v) { splay(); val = v; calc(); } // change
   ⇒value in node, splay suffices instead of access
   ⇒because it doesn't affect values in nodes above it
friend sn lca(sn x, sn y) {
 if (x == y) return x;
  x->access(), y->access(); if (!x->p) return NULL; //
     \hookrightarrowaccess at y did not affect x, so they must not be
     \hookrightarrow connected
  x->splay(); return x->p ? x->p : x;
```

```
friend bool connected(sn x, sn y) { return lca(x,y); }
  friend int balanced(sn x, sn y) {
    x->makeRoot(); v->access();
    return v->sum-2*v->mn;
  friend bool link(sn x, sn y) { // make x parent of y
    if (connected(x,y)) return 0; // don't induce cycle
    y->makeRoot(); y->p = x;
    // x->access(); x->sz += y->sz; x->vir += y->sz;
    return 1; // success!
  friend bool cut(sn x, sn y) { // x is originally parent
     \hookrightarrow of V
    x->makeRoot(); y->access();
    if (y->c[0] != x || x->c[0] || x->c[1]) return 0; //
       \hookrightarrow splay tree with y should not contain anything
       \hookrightarrowelse besides x
    x->p = y->c[0] = NULL; y->calc(); return 1; // calc is
       \hookrightarrow redundant as it will be called elsewhere anyways
       \hookrightarrow ?
};
```

DirectedMST.h

Description: computes minimum weight directed spanning tree, edge from $inv[i] \to i$ for all $i \neq r$

Time: $\mathcal{O}(M \log M)$

```
"DSUrb.h"
                                                   64 lines
struct Edge { int a, b; ll w; };
struct Node {
 Edge kev:
 Node *1, *r;
 ll delta;
 void prop() {
   kev.w += delta;
   if (1) 1->delta += delta;
   if (r) r->delta += delta;
   delta = 0;
 Edge top() { prop(); return key; }
Node *merge(Node *a, Node *b) {
 if (!a || !b) return a ?: b;
 a->prop(), b->prop();
 if (a->key.w > b->key.w) swap(a, b);
 swap(a->1, (a->r = merge(b, a->r)));
  return a;
void pop(Node*\& a) { a->prop(); a = merge(a->1, a->r); }
pair<ll, vi> dmst(int n, int r, const vector<Edge>& g) {
 DSUrb dsu; dsu.init(n); // DSU with rollback if need to

→ return edges

  vector<Node*> heap(n); // store edges entering each
    trav(e,g) heap[e.b] = merge(heap[e.b], new Node{e});
 ll res = 0; vi seen(n,-1); seen[r] = r;
  vpi in(n, \{-1, -1\});
```

```
vector<pair<int, vector<Edge>>> cycs;
FOR(s,n) {
  int u = s, w;
  vector<pair<int, Edge>> path;
  while (seen[u] < 0) {</pre>
    if (!heap[u]) return {-1,{}};
    seen[u] = s;
    Edge e = heap[u] \rightarrow top(); path.pb(\{u,e\});
    heap[u]->delta -= e.w, pop(heap[u]);
    res += e.w, u = dsu.get(e.a);
    if (seen[u] == s) { // compress verts in cycle
      Node * cyc = 0; cycs.pb(\{u, \{\}\}\);
        cyc = merge(cyc, heap[w = path.back().f]);
        cycs.back().s.pb(path.back().s);
        path.pop_back();
      } while (dsu.unite(u, w));
      u = dsu.get(u); heap[u] = cyc, seen[u] = -1;
  trav(t,path) in[dsu.get(t.s.b)] = {t.s.a,t.s.b}; //
     \hookrightarrow found path from root
while (sz(cycs)) { // expand cycs to restore sol
  auto c = cycs.back(); cycs.pop_back();
  pi inEdge = in[c.f];
  trav(t,c.s) dsu.rollback();
  trav(t,c.s) in[dsu.get(t.b)] = \{t.a,t.b\};
  in[dsu.get(inEdge.s)] = inEdge;
vi inv;
FOR(i,n)
  assert(i == r ? in[i].s == -1 : in[i].s == i);
  inv.pb(in[i].f);
return {res,inv};
```

DominatorTree.h

Description: a dominates b iff every path from 1 to b passes through

```
Time: \mathcal{O}\left(M\log N\right)
```

46 line

5 lines

EdgeColor Point AngleCmp LineDist

```
void dfs(int x) { // create DFS tree
   label[x] = ++co; rlabel[co] = x;
   sdom[co] = par[co] = bes[co] = co;
   trav(y,adj[x]) {
     if (!label[y]) {
       dfs(y);
       child[label[x]].pb(label[y]);
     radj[label[y]].pb(label[x]);
 void init() {
   dfs(root);
   ROF(i,1,co+1)
     trav(j,radj[i]) ckmin(sdom[i],sdom[get(j)]);
     if (i > 1) sdomChild[sdom[i]].pb(i);
     trav(j,sdomChild[i]) {
       int k = qet(j);
       if (sdom[j] == sdom[k]) dom[j] = sdom[j];
       else dom[j] = k;
     trav(j,child[i]) par[j] = i;
   FOR(i,2,co+1) {
     if (dom[i] != sdom[i]) dom[i] = dom[dom[i]];
     ans[rlabel[dom[i]]].pb(rlabel[i]);
};
```

EdgeColor.h

Description: naive implementation of Misra & Gries edge coloring, by Vizing's Theorem a simple graph with max degree d can be edge colored with at most d+1 colors

```
Time: \mathcal{O}\left(MN^2\right)
template<int SZ> struct EdgeColor {
 int N = 0, maxDeg = 0, adj[SZ][SZ], deg[SZ];
 EdgeColor() {
   memset(adj,0,sizeof adj);
   memset(deg,0,sizeof deg);
  void addEdge(int a, int b, int c) {
   adj[a][b] = adj[b][a] = c;
  int delEdge(int a, int b) {
   int c = adj[a][b];
   adj[a][b] = adj[b][a] = 0;
   return c;
  vector<bool> genCol(int x) {
   vector<bool> col(N+1); FOR(i,N) col[adj[x][i]] = 1;
    return col;
  int freeCol(int u) {
   auto col = genCol(u);
   int x = 1; while (col[x]) x ++; return x;
  void invert(int x, int d, int c) {
```

```
FOR(i,N) if (adi[x][i] == d)
      delEdge(x,i), invert(i,c,d), addEdge(x,i,c);
  void addEdge(int u, int v) { // follows wikipedia steps
    // check if you can add edge w/o doing any work
    assert(N); ckmax(maxDeg,max(++deg[u],++deg[v]));
    auto a = genCol(u), b = genCol(v);
    FOR(i,1,maxDeg+2) if (!a[i] && !b[i]) return addEdge(u
       \hookrightarrow, v, i);
    // 2. find maximal fan of u starting at v
    vector<bool> use(N); vi fan = {v}; use[v] = 1;
      auto col = genCol(fan.back());
      if (sz(fan) > 1) col[adj[fan.back()][u]] = 0;
      int i = 0; while (i < N && (use[i] || col[adj[u][i</pre>
         →]])) i ++;
      if (i < N) fan.pb(i), use[i] = 1;
      else break:
    // 3/4. choose free cols for endpoints of fan, invert
       \hookrightarrowcd u path
    int c = freeCol(u), d = freeCol(fan.back()); invert(u,
       \hookrightarrowd,c);
    // 5. find i such that d is free on fan[i]
    int i = 0; while (i < sz(fan) && genCol(fan[i])[d]</pre>
      && adj[u][fan[i]] != d) i ++;
    assert (i != sz(fan));
    // 6. rotate fan from 0 to i
    FOR(j,i) addEdge(u,fan[j],delEdge(u,fan[j+1]));
    // 7. add new edge
    addEdge(u,fan[i],d);
1:
```

Geometry (7)

7.1 Primitives

Point.h

```
Description: Easy Geo
```

```
P conj(P x) \{ return P(x.f,-x.s); \}
  P operator+(const P& 1, const P& r) { return P(1.f+r.f.1
     \hookrightarrow.s+r.s); }
  P operator-(const P& 1, const P& r) { return P(1.f-r.f,1
     \hookrightarrow.s-r.s); }
  P operator*(const P& 1, const T& r) { return P(1.f*r,1.s
     \hookrightarrow \star r): }
  P operator*(const T& 1, const P& r) { return r*1; }
  P operator/(const P& 1, const T& r) { return P(1.f/r,1.s
  P operator*(const P& 1, const P& r) { return P(1.f*r.f-1
     \hookrightarrow.s*r.s,l.s*r.f+l.f*r.s); }
  P operator/(const P& 1, const P& r) { return 1*conj(r)/
     \hookrightarrownorm(r); }
  P\& operator += (P\& l, const P\& r) { return l = l+r; }
  P\& operator = (P\& 1, const P\& r) { return 1 = 1-r; }
 P& operator*=(P& 1, const T& r) { return 1 = 1*r; }
 P& operator/=(P& 1, const T& r) { return 1 = 1/r; }
  P\& operator *= (P\& 1, const P\& r) { return 1 = 1*r; }
  P\& operator/=(P\& l, const P\& r) { return l = l/r; }
  P unit(P x) { return x/abs(x); }
 T dot(P a, P b) { return (conj(a)*b).f; }
 T cross(P a, P b) { return (conj(a)*b).s; }
 T cross(P p, P a, P b) { return cross(a-p,b-p); }
 P rotate(P a, T b) { return a*P(cos(b), sin(b)); }
  P reflect (P p, P a, P b) { return a+conj((p-a)/(b-a))*(b
 P foot(P p, P a, P b) { return (p+reflect(p,a,b))/(T)2;
 bool onSeg(P p, P a, P b) { return cross(a,b,p) == 0 &&
     \hookrightarrowdot (p-a, p-b) \langle = 0; \rangle
using namespace Point;
```

AngleCmp.h

Description: sorts points according to atan2

LineDist.h

Description: computes distance between P and line AB

```
T lineDist(P p, P a, P b) { return abs(cross(p,a,b))/abs(a ---b); }
```

SegDist.h

Description: computes distance between P and line segment AB

LineIntersect.h

Description: computes the intersection point(s) of lines AB, CD; returns -1,0,0 if infinitely many, 0,0,0 if none, 1,x if x is the unique point "Point.h" 8 lines

SegIntersect.h

Description: computes the intersection point(s) of line segments AB, CD

7.2 Polygons

Area.h

Description: computes area + the center of mass of a polygon with constant mass per unit area

Time: $\mathcal{O}\left(N\right)$

```
int j = (i+1)%sz(v); T a = cross(v[i],v[j]);
  cen += a*(v[i]+v[j]); area += a;
}
return cen/area/(T)3;
}
```

InPolv.

Description: tests whether a point is inside, on, or outside the perimeter of any polygon

Time: $\mathcal{O}\left(N\right)$

ConvexHull.h

Description: Top-bottom convex hull **Time:** $\mathcal{O}(N \log N)$

} else {

```
48 lines
struct convexHull {
    set<pair<ld,ld>> dupChecker;
    vector<pair<ld,ld>> points;
    vector<pair<ld,ld>> dn, up, hull;
    convexHull() {}
    bool cw(pd o, pd a, pd b) {
        return ((a.f-o.f) * (b.s-o.s) - (a.s-o.s) * (b.f-o.f) <=
           \hookrightarrow 0);
    void addPoint(pair<ld,ld> p) {
        if (dupChecker.count(p)) return;
        points.pb(p);
        dupChecker.insert(p);
    void addPoint(ld x, ld y) {
        addPoint(mp(x,y));
    void build() {
        sort(points.begin(), points.end());
        if(sz(points) < 3) {
            for (pair<ld, ld> p: points) {
                dn.pb(p);
                hull.pb(p);
            M00d(i, sz(points)) {
                up.pb(points[i]);
```

for(int i = 0; i < (int)points.size(); i++) {</pre>

while(dn.size() >= 2 && cw(dn[dn.size()

 \hookrightarrow -2], dn[dn.size()-1], points[i])) {

PolyDiameter.h

10 lines

Description: computes longest distance between two points in P **Time:** O(N) given convex hull

7.3 Circles

Circles.h

Description: misc operations with two circles

Circumcenter MinEnclosingCircle ClosestPair DelaunayFast

```
T intersectArea(circ x, circ y) { // not thoroughly tested
 T d = abs(x.f-v.f), a = x.s, b = v.s; if (a < b) swap(a, b)
 if (d >= a+b) return 0;
 if (d <= a-b) return PI*b*b;
  auto ca = (a*a+d*d-b*b)/(2*a*d), cb = (b*b+d*d-a*a)/(2*b)
 auto s = (a+b+d)/2, h = 2*sqrt(s*(s-a)*(s-b)*(s-d))/d;
 return a*a*acos(ca)+b*b*acos(cb)-d*h;
P tangent (P x, circ y, int t = 0) {
 y.s = abs(y.s); // abs needed because internal calls y.s
 if (y.s == 0) return y.f;
 T d = abs(x-v.f);
 P = pow(y.s/d, 2) * (x-y.f) + y.f;
 P b = sqrt(d*d-y.s*y.s)/d*y.s*unit(x-y.f)*dir(PI/2);
 return t == 0 ? a+b : a-b;
vector<pair<P,P>> external(circ x, circ y) { // external
  \hookrightarrowtangents
 vector<pair<P,P>> v;
 if (x.s == y.s) {
   P \text{ tmp} = \text{unit}(x.f-y.f)*x.s*dir(PI/2);
   v.pb(mp(x.f+tmp,y.f+tmp));
   v.pb(mp(x.f-tmp,y.f-tmp));
   P p = (y.s*x.f-x.s*y.f)/(y.s-x.s);
   FOR(i,2) v.pb({tangent(p,x,i),tangent(p,y,i)});
vector<pair<P,P>> internal(circ x, circ y) { // internal
  \hookrightarrowtangents
 x.s *= -1; return external(x,y);
```

Circumcenter.h

Description: returns {circumcenter,circumradius}

MinEnclosingCircle.h

Description: computes minimum enclosing circle **Time:** expected $\mathcal{O}(N)$

```
o = (ps[i]+ps[j])/2, r = abs(o-ps[i]);
FOR(k,j) if (abs(o-ps[k]) > r*EPS)
          tie(o,r) = ccCenter(ps[i],ps[j],ps[k]);
}
return {o,r};
}
```

7.4 Misc

ClosestPair.h

Description: line sweep to find two closest points **Time:** $\mathcal{O}(N \log N)$

```
21 lines
using namespace Point;
pair<P,P> solve(vP v) {
  pair<ld,pair<P,P>> bes; bes.f = INF;
  set < P > S; int ind = 0;
  sort(all(v));
  FOR(i,sz(v)) {
   if (i && v[i] == v[i-1]) return {v[i],v[i]};
    for (; v[i].f-v[ind].f >= bes.f; ++ind)
      S.erase({v[ind].s,v[ind].f});
    for (auto it = S.ub({v[i].s-bes.f,INF});
      it != end(S) && it->f < v[i].s+bes.f; ++it) {
      P t = \{it->s, it->f\};
      ckmin(bes, {abs(t-v[i]), {t,v[i]}});
    S.insert({v[i].s,v[i].f});
  return bes.s;
```

DelaunavFast.h

Description: Delaunay Triangulation, concyclic points are OK (but not all collinear) **Time:** $\mathcal{O}\left(N\log N\right)$

```
"Point.h"
                                                       94 lines
typedef 11 T;
typedef struct Quad* Q;
typedef __int128_t 111; // (can be 11 if coords are < 2e4)
P arb(LLONG_MAX, LLONG_MAX); // not equal to any other
  \hookrightarrowpoint
struct Quad {
 bool mark; Q o, rot; P p;
 P F() { return r()->p; }
  Q r() { return rot->rot; }
  Q prev() { return rot->o->rot;
  Q next() { return r()->prev(); }
// test if p is in the circumcircle
bool circ(P p, P a, P b, P c) {
 ll ar = cross(a,b,c); assert(ar); if (ar < 0) swap(a,b);
```

```
111 p2 = norm(p), A = norm(a) - p2,
    B = norm(b) - p2, C = norm(c) - p2;
 return cross(p,a,b)*C+cross(p,b,c)*A+cross(p,c,a)*B > 0;
Q makeEdge(P orig, P dest) {
 0 \text{ g}[] = \{\text{new Ouad}\{0, 0, 0, \text{orig}\}, \text{ new Ouad}\{0, 0, 0, \text{arb}\},
       new Quad{0,0,0,dest}, new Quad{0,0,0,arb}};
 FOR(i,4) q[i] \rightarrow o = q[-i \& 3], q[i] \rightarrow rot = q[(i+1) \& 3];
 return *q;
void splice(O a, O b) {
 swap(a->o->rot->o, b->o->rot->o); swap(a->o, b->o);
0 connect(0 a, 0 b) {
 Q = makeEdge(a->F(), b->p);
 splice(q, a->next());
 splice(q->r(), b);
 return q;
pair<Q,Q> rec(const vector<P>& s) {
 if (sz(s) \le 3) {
    Q = makeEdge(s[0], s[1]), b = makeEdge(s[1], s.back
    if (sz(s) == 2) return { a, a->r() };
    splice(a->r(), b);
    auto side = cross(s[0], s[1], s[2]);
    Q c = side ? connect(b, a) : 0;
    return {side < 0 ? c->r() : a, side < 0 ? c : b->r()
       \hookrightarrow };
\#define H(e) e \rightarrow F(), e \rightarrow p
#define valid(e) (cross(e->F(),H(base)) > 0)
 Q A, B, ra, rb;
 int half = sz(s) / 2;
 tie(ra, A) = rec({all(s) - half});
 tie(B, rb) = rec({sz(s) - half + all(s)});
  while ((cross(B->p,H(A)) < 0 \&& (A = A->next())) | |
       (cross(A->p,H(B)) > 0 && (B = B->r()->o)));
  O base = connect(B->r(), A);
  if (A->p == ra->p) ra = base->r();
 if (B->p == rb->p) rb = base;
#define DEL(e, init, dir) Q e = init->dir; if (valid(e)) \
    while (circ(e->dir->F(), H(base), e->F())) {
      0 t = e -> dir; \setminus
      splice(e, e->prev()); \
      splice(e->r(), e->r()->prev()); \
      e = t; \
  for (;;) {
    DEL(LC, base->r(), o); DEL(RC, base, prev());
    if (!valid(LC) && !valid(RC)) break;
    if (!valid(LC) || (valid(RC) && circ(H(RC), H(LC))))
      base = connect(RC, base->r());
      base = connect(base->r(), LC->r());
```

19 lines

Point3D Hull3D KMP Z

$7.5 \quad 3D$

Point3D.h

Description: Basic 3D Geometry

45 lines

```
typedef ld T;
namespace Point3D {
 typedef array<T,3> P3;
 typedef vector<P3> vP3;
 T norm(const P3& x) {
   T sum = 0; FOR(i,sz(x)) sum += x[i]*x[i];
   return sum;
 T abs(const P3& x) { return sqrt(norm(x)); }
 P3& operator+=(P3& 1, const P3& r) { F0R(i,3) 1[i] += r[
    \hookrightarrowi]; return 1; }
 P3& operator-=(P3& 1, const P3& r) { F0R(i,3) 1[i] -= r[
    \hookrightarrowil; return l; }
 P3& operator *= (P3& 1, const T& r) { F0R(i,3) 1[i] *= r;
     \hookrightarrowreturn 1; }
 P3& operator/=(P3& 1, const T& r) { F0R(i,3) 1[i] /= r;
    →return 1; }
 P3 operator+(P3 1, const P3& r) { return 1 += r; }
  P3 operator-(P3 1, const P3& r) { return 1 -= r; }
  P3 operator*(P3 1, const T& r) { return 1 *= r; }
  P3 operator*(const T& r, const P3& 1) { return 1*r; }
 P3 operator/(P3 1, const T& r) { return 1 /= r; }
 T dot(const P3& a, const P3& b) {
   T sum = 0; FOR(i,3) sum += a[i]*b[i];
    return sum;
 P3 cross(const P3& a, const P3& b) {
   return {a[1]*b[2]-a[2]*b[1],
```

Hull3D.h

Description: 3D Convex Hull + Polyedron Volume **Time:** $\mathcal{O}\left(N^2\right)$

```
"Point3D.h"
                                                        48 \underline{\text{lines}}
struct ED {
 void ins(int x) { (a == -1 ? a : b) = x; }
  void rem(int x) { (a == x ? a : b) = -1; }
  int cnt() { return (a !=-1) + (b !=-1); }
 int a, b;
};
struct F { P3 q; int a, b, c; };
vector<F> hull3d(const vP3& A) {
  assert(sz(A) >= 4);
  vector < vector < ED >> E(sz(A), vector < ED > (sz(A), {-1, -1}))
  #define E(x,y) E[f.x][f.y]
  vector<F> FS; // faces
  auto mf = [\&] (int i, int j, int k, int l) { // make face}
   P3 q = cross(A[j]-A[i],A[k]-A[i]);
    if (dot(q,A[1]) > dot(q,A[i])) q *= -1; // make sure q

→ points outward

    F f{q, i, j, k};
    E(a,b).ins(k); E(a,c).ins(j); E(b,c).ins(i);
   FS.pb(f);
  FOR(i,4) FOR(j,i+1,4) FOR(k,j+1,4) mf(i, j, k, 6-i-j-k);
  FOR(i, 4, sz(A)) {
   FOR(j,sz(FS)) {
      F f = FS[j];
      if (dot(f.q,A[i]) > dot(f.q,A[f.a]))  { // face is
         \hookrightarrow visible, remove edges
        E(a,b).rem(f.c), E(a,c).rem(f.b), E(b,c).rem(f.a);
        swap(FS[j--], FS.back());
        FS.pop_back();
    FOR(j,sz(FS)) { // add faces with new point
```

Strings (8)

8.1 Lightweight

KMP.h

Time: $\mathcal{O}(N)$

Description: f[i] equals the length of the longest proper suffix of the *i*-th prefix of s that is a prefix of s

Z.h

Description: for each index i, computes the the maximum len such that s.substr(0,len) == s.substr(i,len) **Time:** $\mathcal{O}(N)$

```
vi z(string s) {
  int N = sz(s); s += '#';
  vi ans(N); ans[0] = N;
  int L = 1, R = 0;
  FOR(i,1,N) {
   if (i <= R) ans[i] = min(R-i+1,ans[i-L]);
}</pre>
```

Manacher.h

 $\bf Description:$ Calculates length of largest palindrome centered at each character of string

Time: $\mathcal{O}\left(N\right)$

MinRotation.h

Description: minimum rotation of string **Time:** $\mathcal{O}(N)$

. . . (11)

LyndonFactorization.h

Description: A string is "simple" if it is strictly smaller than any of its own nontrivial suffixes. The Lyndon factorization of the string s is a factorization $s=w_1w_2\dots w_k$ where all strings w_i are simple and $w_1\geq w_2\geq \dots \geq w_k$

 $\begin{array}{ll} w_1 \geq w_2 \geq \cdots \geq w_k \\ \hline \textbf{Time: } \mathcal{O}(N) & 20 \text{ lines} \\ \hline \text{vector} < \text{string} > \text{duval} (\text{const string} \& s) & \{\\ & \text{int n = sz(s); vector} < \text{string} > \text{factors;} \\ & \text{for (int i = 0; i < n;) } & \{\\ \end{array}$

RabinKarp.h

18 lines

8 lines

Description: generates hash values of any substring in O(1), equal strings have same hash value

Time: $\mathcal{O}(N)$ build, $\mathcal{O}(1)$ get hash value of a substring

25 lines

```
template<int SZ> struct rabinKarp {
    const 11 mods[3] = {1000000007, 999119999,
       \hookrightarrow1000992299};
    11 p[3][SZ];
    11 h[3][SZ];
    const 11 base = 1000696969;
    rabinKarp() {}
    void build(string a) {
        M00(i, 3) {
             p[i][0] = 1;
             h[i][0] = (int)a[0];
             MOO(j, 1, (int)a.length()) {
                 p[i][j] = (p[i][j-1] * mods[i]) % base;
                 h[i][j] = (h[i][j-1] * mods[i] + (int)a[j]
                     \hookrightarrow]) % base;
    tuple<11, 11, 11> hsh(int a, int b) {
         if(a == 0) return make_tuple(h[0][b], h[1][b], h
            \hookrightarrow [2] [b]);
        tuple<11, 11, 11> ans;
         get<0>(ans) = (((h[0][b] - h[0][a-1]*p[0][b-a+1])
            \hookrightarrow% base) + base) % base;
         get<1>(ans) = (((h[1][b] - h[1][a-1]*p[1][b-a+1])
            \hookrightarrow% base) + base) % base;
```

Trie.h

```
Description: trie
```

25 lines

```
struct tnode
    char c:
    bool used;
    tnode* next[26];
    tnode() {
        c = ' ';
        used = 0;
        M00(i, 26) next[i] = nullptr;
};
tnode* root;
void addToTrie(string s) {
    tnode* cur = root;
    for(char ch: s) {
        int idx = ch - 'a';
        if(cur->next[idx] == nullptr) {
            cur->next[idx] = new tnode();
        cur = cur->next[idx];
        cur->c = ch;
    cur->used = 1;
```

8.2 Suffix Structures

ACfixed.h

Description: for each prefix, stores link to max length suffix which is also a prefix

Time: $\mathcal{O}(N \sum)$

36 lines

```
struct ACfixed { // fixed alphabet
    struct node {
        array<int,26> to;
        int link;
    };
    vector<node> d;
    ACfixed() { d.eb(); }

    int add(string s) { // add word
        int v = 0;
        trav(C,s) {
            int c = C-'a';
            if (!d[v].to[c]) {
                  d[v].to[c] = sz(d);
                  d.eb();
            }
            v = d[v].to[c];
}
```

PalTree SuffixArray ReverseBW SuffixAutomaton

PalTree.h

 $\bf Description:$ palindromic tree, computes number of occurrences of each palindrome within string

Time: $\mathcal{O}(N \sum)$

25 lines

```
template<int SZ> struct PalTree {
  static const int sigma = 26;
  int s[SZ], len[SZ], link[SZ], to[SZ][sigma], oc[SZ];
  int n, last, sz;
 PalTree() { s[n++] = -1; link[0] = 1; len[1] = -1; sz = -1
    \hookrightarrow2; }
  int getLink(int v) {
   while (s[n-len[v]-2] != s[n-1]) v = link[v];
   return v;
  void addChar(int c) {
   s[n++] = c;
   last = getLink(last);
   if (!to[last][c]) {
     len[sz] = len[last]+2;
     link[sz] = to[getLink(link[last])][c];
     to[last][c] = sz++;
   last = to[last][c]; oc[last] ++;
 void numOc() {
   vpi v; FOR(i,2,sz) v.pb({len[i],i});
   sort(rall(v)); trav(a,v) oc[link[a.s]] += oc[a.s];
};
```

SuffixArray.h Description: ?

Time: $\mathcal{O}(N \log N)$

```
template<int SZ> struct suffixArray {
   const static int LGSZ = 33-__builtin_clz(SZ-1);
   pair<pi, int> tup[SZ];
```

```
int sortIndex[LGSZ][SZ];
    int res[SZ];
    int len:
    suffixArray(string s) {
        this->len = (int)s.length();
        M00(i, len) tup[i] = MP(MP((int)s[i], -1), i);
        sort(tup, tup+len);
        int temp = 0:
        tup[0].F.F = 0;
        MOO(i, 1, len) {
            if(s[tup[i].S] != s[tup[i-1].S]) temp++;
            tup[i].F.F = temp;
        M00(i, len) sortIndex[0][tup[i].S] = tup[i].F.F;
        MOO(i, 1, LGSZ) {
            M00(j, len) tup[j] = MP(MP(sortIndex[i-1][j],
               \hookrightarrow (j+(1<<(i-1))<len)?sortIndex[i-1][j+(1<<(
               \hookrightarrowi-1))]:-1), j);
            sort(tup, tup+len);
            int temp2 = 0;
            sortIndex[i][tup[0].S] = 0;
            MOO(j, 1, len) {
                if(tup[j-1].F != tup[j].F) temp2++;
                sortIndex[i][tup[j].S] = temp2;
        M00(i, len) res[sortIndex[LGSZ-1][i]] = i;
    int LCP(int x, int y) {
        if(x == y) return len - x;
        int ans = 0;
        M00d(i, LGSZ) {
            if (x \ge len | | y \ge len) break;
            if(sortIndex[i][x] == sortIndex[i][y]) {
                x += (1 << i);
                y += (1 << i);
                ans += (1 << i);
        return ans;
};
```

ReverseBW.h

Description: The Burrows-Wheeler Transform appends # to a string, sorts the rotations of the string in increasing order, and constructs a new string that contains the last character of each rotation. This function reverses the transform.

Time: $\mathcal{O}(N \log N)$

```
string reverseBW(string s) {
  vi nex(sz(s));
  vector<pair<char,int>> v; FOR(i,sz(s)) v.pb({s[i],i});
  sort(all(v)); FOR(i,sz(v)) nex[i] = v[i].s;
  int cur = nex[0]; string ret;
  for (; cur; cur = nex[cur]) ret += v[cur].f;
  return ret;
}
```

SuffixAutomaton.h

 $\bf Description:$ constructs minimal DFA that recognizes all suffixes of a string

```
Time: \mathcal{O}\left(N\log\sum\right)
```

73 lines

```
struct SuffixAutomaton {
 struct state {
   int len = 0, firstPos = -1, link = -1;
   bool isClone = 0;
   map<char, int> next;
   vi invLink;
 vector<state> st;
 int last = 0;
 void extend(char c) {
   int cur = sz(st); st.eb();
   st[cur].len = st[last].len+1, st[cur].firstPos = st[
       \hookrightarrowcurl.len-1;
   int p = last;
   while (p != -1 \&\& !st[p].next.count(c)) {
      st[p].next[c] = cur;
      p = st[p].link;
   if (p == -1) {
      st[cur].link = 0;
    } else {
      int q = st[p].next[c];
      if (st[p].len+1 == st[q].len) {
        st[cur].link = q;
      } else {
        int clone = sz(st); st.pb(st[q]);
        st[clone].len = st[p].len+1, st[clone].isClone =
           \hookrightarrow1:
        while (p != -1 \&\& st[p].next[c] == q) {
          st[p].next[c] = clone;
          p = st[p].link;
        st[q].link = st[cur].link = clone;
   last = cur;
 void init(string s) {
   st.eb(); trav(x,s) extend(x);
   FOR(v,1,sz(st)) st[st[v].link].invLink.pb(v);
 // APPLICATIONS
 void getAllOccur(vi& oc, int v) {
   if (!st[v].isClone) oc.pb(st[v].firstPos);
   trav(u, st[v].invLink) getAllOccur(oc, u);
 vi allOccur(string s) {
   int cur = 0;
   trav(x,s) {
      if (!st[cur].next.count(x)) return {};
      cur = st[cur].next[x];
```

vi oc; qetAllOccur(oc, cur); trav(t, oc) t += 1-sz(s);

SuffixTree TandemRepeats

```
sort(all(oc)); return oc;
 vl distinct;
  11 getDistinct(int x) {
    if (distinct[x]) return distinct[x];
    distinct[x] = 1;
    trav(y,st[x].next) distinct[x] += getDistinct(y.s);
    return distinct[x];
  11 numDistinct() { // # of distinct substrings,
    \hookrightarrowincluding empty
    distinct.rsz(sz(st));
    return getDistinct(0);
  ll numDistinct2() { // another way to get # of distinct
    \hookrightarrow substrings
    11 \text{ ans} = 1;
    FOR(i, 1, sz(st)) ans += st[i].len-st[st[i].link].len;
    return ans;
};
```

SuffixTree.h

Description: Ukkonen's algorithm for suffix tree

Time: $\mathcal{O}(N \log \Sigma)$

```
61 lines
struct SuffixTree
 string s; int node, pos;
  struct state {
   int fpos, len, link = -1;
   map<char,int> to;
   state(int fpos, int len) : fpos(fpos), len(len) {}
 };
  vector<state> st;
  int makeNode(int pos, int len) {
    st.pb(state(pos,len)); return sz(st)-1;
 void goEdge() {
   while (pos > 1 && pos > st[st[node].to[s[sz(s)-pos]]].
      \hookrightarrowlen) {
     node = st[node].to[s[sz(s)-pos]];
     pos -= st[node].len;
  void extend(char c) {
   s += c; pos ++; int last = 0;
   while (pos) {
     goEdge();
     char edge = s[sz(s)-pos];
     int& v = st[node].to[edge];
      char t = s[st[v].fpos+pos-1];
     if (v == 0) {
        v = makeNode(sz(s)-pos,MOD);
        st[last].link = node; last = 0;
      } else if (t == c) {
        st[last].link = node;
```

int u = makeNode(st[v].fpos,pos-1);

```
st[u].to[c] = makeNode(sz(s)-1,MOD); st[u].to[t] =
        st[v].fpos += pos-1; st[v].len -= pos-1;
        v = u; st[last].link = u; last = u;
      if (node == 0) pos --;
      else node = st[node].link;
  void init(string s) {
   makeNode(0,MOD); node = pos = 0;
    trav(c,_s) extend(c);
  bool isSubstr(string _x) {
    string x; int node = 0, pos = 0;
    trav(c,_x) {
      x += c; pos ++;
      while (pos > 1 && pos > st[st[node].to[x[sz(x)-pos]]
         \hookrightarrow111.len) {
        node = st[node].to[x[sz(x)-pos]];
        pos -= st[node].len;
      char edge = x[sz(x)-pos];
      if (pos == 1 && !st[node].to.count(edge)) return 0;
      int& v = st[node].to[edge];
      char t = s[st[v].fpos+pos-1];
      if (c != t) return 0;
    return 1;
};
```

Misc 8.3

TandemRepeats.h

Description: Main-Lorentz algorithm, finds all (x, y) such that s.substr(x,y-1) == s.substr(x+y,y-1)Time: $\mathcal{O}(N \log N)$

```
"Z.h"
                                                          54 lines
struct StringRepeat {
  string S;
  vector<array<int,3>> al;
  // (t[0],t[1],t[2]) -> there is a repeating substring
     \hookrightarrowstarting at x
  // with length t[0]/2 for all t[1] \le x \le t[2]
  vector<array<int,3>> solveLeft(string s, int m) {
   vector<array<int,3>> v;
    vi v2 = getPrefix(string(s.begin()+m+1, s.end()), string
        \hookrightarrow (s.begin(),s.begin()+m+1));
    string V = string(s.begin(),s.begin()+m+2); reverse(
       \hookrightarrowall(V)); vi v1 = z(V); reverse(all(v1));
    FOR(i, m+1) if (v1[i]+v2[i] >= m+2-i) {
      int lo = \max(1, m+2-i-v2[i]), hi = \min(v1[i], m+1-i);
      lo = i-lo+1, hi = i-hi+1; swap(lo,hi);
      v.pb({2*(m+1-i),lo,hi});
```

```
return v;
  void divi(int 1, int r) {
    if (1 == r) return;
    int m = (1+r)/2; divi(1, m); divi(m+1, r);
    string t = string(S.begin()+1,S.begin()+r+1);
    m = (sz(t)-1)/2;
    auto a = solveLeft(t,m);
    reverse(all(t));
    auto b = solveLeft(t, sz(t) - 2 - m);
    trav(x,a) al.pb(\{x[0],x[1]+1,x[2]+1\});
    trav(x,b) {
      int ad = r-x[0]+1;
      al.pb(\{x[0],ad-x[2],ad-x[1]\});
  void init(string _S) {
    S = _S; divi(0, sz(S)-1);
  vi genLen() { // min length of repeating substring
     \hookrightarrowstarting at each index
    priority_queue<pi, vpi, greater<pi>>> m; m.push({MOD, MOD
    vpi ins[sz(S)]; trav(a,al) ins[a[1]].pb({a[0],a[2]});
    vi len(sz(S));
    FOR(i,sz(S)) {
      trav(j,ins[i]) m.push(j);
      while (m.top().s < i) m.pop();</pre>
      len[i] = m.top().f;
    return len;
};
```