

# Wavelets Homework 1

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- Mathematical foundation: Haar transform theory
- Implementation: Fast Haar Transform algorithms
- Analysis: Understanding coefficient patterns
- Applications: Compression and approximation

- Haar basis functions and their properties
- Matrix construction for discrete Haar transform
- Complexity analysis:  $\mathcal{O}(n \log n)$  vs  $\mathcal{O}(n)$
- Theoretical bounds and properties

## Exercise 1: Take $T_2$ as a start

We compute  $T_2$ :

$$\begin{bmatrix} s & s & 0 & 0 \\ s & -s & 0 & 0 \\ 0 & 0 & s & s \\ 0 & 0 & s & -s \end{bmatrix} \begin{pmatrix} c_{20} \\ c_{21} \\ c_{22} \\ c_{23} \end{pmatrix} = \begin{pmatrix} c_{10} \\ h_{10} \\ c_{11} \\ h_{11} \end{pmatrix}$$

$$\begin{bmatrix} s & 0 & s & 0 \\ s & 0 & -s & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} \begin{pmatrix} c_{10} \\ h_{10} \\ c_{11} \\ h_{11} \end{pmatrix} = \begin{pmatrix} c_{00} \\ h_{00} \\ h_{10} \\ h_{11} \end{pmatrix}$$

$$\begin{bmatrix} s & 0 & s & 0 \\ s & 0 & -s & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} s & s & 0 & 0 \\ s & -s & 0 & 0 \\ 0 & 0 & s & s \\ 0 & 0 & s & -s \end{bmatrix} = \begin{bmatrix} s^2 & s^2 & s^2 & s^2 \\ s^2 & s^2 & -s^2 & -s^2 \\ s & -s & 0 & 0 \\ 0 & 0 & s & -s \end{bmatrix}$$

# Exercise 1: Construction of $T_j$

When computing  $T_3$ , we get the following matrix construction:

$$\begin{bmatrix} * & * & * & * & * & * & * & * \\ * & * & * & * & * & * & * & * \\ * & * & * & * & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & * & * & * & * \\ * & * & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & * & * & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & * & * & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & * & * \end{bmatrix}$$

With this idea in mind, we can compute the amount of nonzero entries for all  $T_j$ :

- For every box in the matrix  $T_j$ , there are  $2^j$  nonzero entries;
- There are  $j + 1$  boxes;
- The total amount of nonzero entries is  $(j + 1) \cdot 2^j$ .

## Exercise 2: Complexity of $T_j$

- Take the amount of nonzeros of  $T_j := (j + 1) \cdot 2^j$ ;
- Assume  $2^j = n$ ;
- Note that we apply also  $(j + 1) \cdot 2^j$  multiplications;
- Then the complexity of  $T_j$  is  $\mathcal{O}(n \log n)$ .

## Exercise 3: How about $T_j^{-1}$

As we've seen in the lecture,  $T_j^{-1} = T_j^T, \forall j \in \mathbb{N}$ , thus:

$$T_2^{-1} = \begin{bmatrix} s^2 & s^2 & s & 0 \\ s^2 & s^2 & -s & 0 \\ s^2 & -s^2 & 0 & s \\ s^2 & -s^2 & 0 & -s \end{bmatrix} \quad T_3^{-1}: \left[ \begin{array}{c|c|c|c|c|c|c|c} * & * & * & 0 & * & 0 & 0 & 0 \\ * & * & * & 0 & * & 0 & 0 & 0 \\ * & * & * & 0 & 0 & * & 0 & 0 \\ * & * & * & 0 & 0 & * & 0 & 0 \\ * & * & 0 & * & 0 & 0 & * & 0 \\ * & * & 0 & * & 0 & 0 & * & 0 \\ * & * & 0 & * & 0 & 0 & 0 & * \\ * & * & 0 & * & 0 & 0 & 0 & * \end{array} \right]$$

Furthermore, The amount of nonzeros entries hasn't changed, thus this is still equal to  $(j+1) \cdot 2^j$ .

# Code Inspection: DHT Implementation

- DHT algorithm: sparse matrix multiplication
- Inverse transform:  $\text{iDHT}$



# Code Inspection: DHT Implementation

- DHT algorithm: sparse matrix multiplication
- Inverse transform: iDHT

$$\begin{bmatrix}
 s^3 & s^3 & s^3 & s^3 & s^3 & s^3 & s^3 & s^3 \\
 \hline
 s^3 & s^3 & s^3 & s^3 & -s^3 & -s^3 & -s^3 & -s^3 \\
 \hline
 s^2 & s^2 & -s^2 & -s^2 & 0 & 0 & 0 & 0 \\
 0 & 0 & 0 & 0 & s^2 & s^2 & -s^2 & -s^2 \\
 \hline
 s & -s & 0 & 0 & 0 & 0 & 0 & 0 \\
 0 & 0 & s & -s & 0 & 0 & 0 & 0 \\
 0 & 0 & 0 & 0 & s & -s & 0 & 0 \\
 0 & 0 & 0 & 0 & 0 & 0 & s & -s
 \end{bmatrix}
 \begin{matrix}
 (0) \\
 (0) \\
 (0)
 \end{matrix}$$

We want a loop that accesses exactly these entries from a virtual matrix.

# Code Inspection: DHT Implementation

Within and outside of the intervals [**start**, **mid**) and [**mid**, **end**), the values of the matrix are constant.

- DHT algorithm: sparse matrix multiplication
- Inverse transform: iDHT

$$\begin{bmatrix}
 s^3 & s^3 & s^3 & s^3 & s^3 & s^3 & s^3 & s^3 \\
 \hline
 \textcolor{blue}{s}^3 & s^3 & s^3 & s^3 & \textcolor{orange}{-s}^3 & -s^3 & -s^3 & -s^3 \\
 \hline
 \textcolor{blue}{s}^2 & s^2 & \textcolor{orange}{-s}^2 & -s^2 & 0 & 0 & 0 & 0 \\
 0 & 0 & 0 & 0 & \textcolor{blue}{s}^2 & s^2 & \textcolor{orange}{-s}^2 & -s^2 \\
 \hline
 \textcolor{blue}{s} & \textcolor{orange}{-s} & 0 & 0 & 0 & 0 & 0 & 0 \\
 0 & 0 & \textcolor{blue}{s} & \textcolor{orange}{-s} & 0 & 0 & 0 & 0 \\
 0 & 0 & 0 & 0 & \textcolor{blue}{s} & \textcolor{orange}{-s} & 0 & 0 \\
 0 & 0 & 0 & 0 & 0 & 0 & \textcolor{blue}{s} & \textcolor{orange}{-s}
 \end{bmatrix}
 \begin{matrix}
 (0) \\
 (0) \\
 (0)
 \end{matrix}$$

We want a loop that accesses exactly these entries from a virtual matrix.

## Exercise 4: Multiplications of the FHT

- Take  $j \in \mathbb{N}$ . The first matrix is of size  $2^j \times 2^j$ ;
- The amount of multiplications in the first step of the FHT is  $2 \cdot 2^j$  (the amount of nonzero elements);
- Going one layer lower leaves a matrix of size  $2^{j-1} \times 2^{j-1}$ ;
- This gives us  $2 \cdot 2^{j-1}$  multiplications;
- the scalar difference of multiplications of going from  $j$  to  $j - 1$  is a factor of 2.

As for  $T_3$ :

$$\begin{bmatrix} S & 0 & 0 & 0 \\ 0 & S & 0 & 0 \\ 0 & 0 & S & 0 \\ 0 & 0 & 0 & S \end{bmatrix}$$

## Exercise 4: Multiplications and complexity

What is the total amount of multiplications needed?

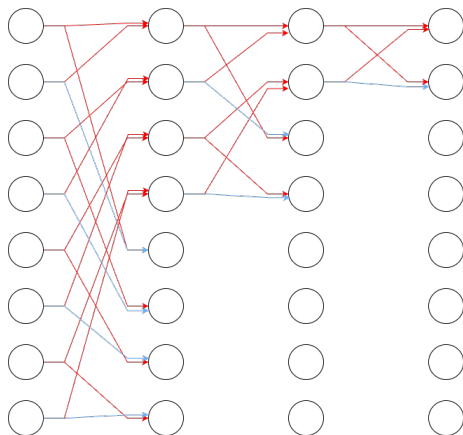
- We want to add up all the multiplications  $2 \cdot 2^k$ , where  $1 \leq k \leq j$ ;
- This is equal to  $\sum_{k=1}^j 2^{k+1} = 2^{j+2} - 2$ .

What is the corresponding complexity?

- Assume again  $2^j = n$ ;
- This gives  $4n - 2$ , which has complexity  $\mathcal{O}(n)$ .

# Code Inspection: FHT Implementation

- FHT function implementation
- Inverse transform:  $iFHT$



# Live Demo: Function Reconstruction

- Sample a test function
- Apply FHT / DHT to get Haar coefficients
- Reconstruct using inverse FHT / DHT

## Exercise 7: Hölder continuity

We eventually want to show that  $|\langle f, H_{jk} \rangle| \leq C 2^{-j(\alpha+1/2)}$ ,  $C \in \mathbb{R}$ .

- We are given that  $f$  is Hölder continuous with coefficient  $\alpha \in [0, 1]$ ;
- $\exists c \in \mathbb{R}$  such that  $|f(x) - f(y)| \leq c|x - y|^\alpha$ ,  $\forall x, y \in \mathbb{R}$ ;
- We also know that  $H_{jk}$  is non-zero on the interval  $[2^{-j}k, 2^{-j}(k+1)]$  and  $|h_{jk}| = |\langle f, H_{jk} \rangle|$

Throughout the computation, we apply the substitution  $u = t - 2^{-j-1}$ , which we denote as 'sub.'

## Exercise 7: Proof of the inequality

$$\begin{aligned} |\langle f, H_{jk} \rangle| &= \left| \int_{2^{-j}k}^{2^{-j}(k+1)} f(t) H_{jk}(t) dt \right| \\ &= 2^{j/2} \left| \int_{2^{-j}k}^{2^{-j}(k+1/2)} f(t) dt - \int_{2^{-j}(k+1/2)}^{2^{-j}(k+1)} f(t) dt \right| \\ &\stackrel{\text{sub.}}{=} 2^{j/2} \left| \int_{2^{-j}(k+1/2)}^{2^{-j}(k+1)} f(u + 2^{-j-1}) - f(u) du \right| \\ &\leq 2^{j/2} \int_{2^{-j}(k+1/2)}^{2^{-j}(k+1)} |f(u + 2^{-j-1}) - f(u)| du \\ &\stackrel{\text{H\"ol.}}{\leq} C \cdot 2^{j/2} 2^{(-j-1)\alpha} \\ &= C \cdot 2^{-j(\alpha+1/2)}. \end{aligned}$$



# Analysis: Understanding Coefficient Patterns

- Multi-resolution structure of Haar coefficients
- Coefficient decay across scales
- Theoretical bounds for Hölder continuous functions
- Empirical verification and sampling methods

$$|\langle f, H_{jk} \rangle| \leq C 2^{-j(\alpha+1/2)}$$

- Visualize coefficients at different layers
- Compute ratios of consecutive layers
- Compare theoretical bounds with empirical results
- Effect of sampling methods on bounds

## Exercise 10: Invariance under translation and dialation

Define  $I$  such that  $\forall 0 \leq q \leq I$ ,  $\int_{\mathbb{R}} x^q W(x) dx = 0$ . Does a translation or dialation on  $W(x)$  gives back zero?

- We start by writing  $2^{j/2} W(2^{-j}x - k)$  for  $W(x)$  as the translation and dialation;
- This results in  $2^{j/2} \int_{\mathbb{R}} x^q W(2^{-j}x - k) dx$ , which we want to solve.

We start by applying the substitution:  $x \mapsto x + \frac{k}{2^{-j}}$ . This gives us:

$$2^{j/2} \int_{\mathbb{R}} \left( x + \frac{k}{2^{-j}} \right)^q W(x) dx.$$

## Exercise 10: Invariance under translation and dialation

To go further, we see:

$$\left(x + \frac{k}{2^{-j}}\right)^q = x^q + \frac{k}{2^{-j}}x^{q-1} + \cdots + \frac{k^q}{2^{-qj}}$$

As integrals are linear, we can rewrite the previous integral as:

$$2^{j/2} \int_{\mathbb{R}} x^q W(x) dx + \frac{2^{j/2}k}{2^{-j}} \int_{\mathbb{R}} x^{q-1} W(x) dx + \cdots + \frac{2^{j/2}k^q}{2^{-qj}} \int_{\mathbb{R}} W(x) dx.$$

We already knew that  $\int_{\mathbb{R}} x^q W(x) dx = 0, \forall 0 \leq q \leq l$ . Which implies:

$$2^{j/2} \int_{\mathbb{R}} \left(x + \frac{k}{2^{-j}}\right)^q W(x) dx = 0.$$

# Applications: Compression and Approximation

- N-term approximation in Haar basis
- Coefficient sorting and thresholding
- Error analysis and compression efficiency
- Practical applications to discontinuous functions

# Live Demo: Compression

- Apply FHT to discontinuous function
- Sort coefficients by magnitude
- Compute best N-term approximation error
- Visualize compressed reconstruction

# Thank You



# Presentation Script: Part 1 - Foundation

- **Mathematical Foundation:** Introduce Haar basis functions, matrix construction, and complexity analysis as the theoretical foundation for our implementation work.
- **Exercise 1: Matrix construction  $T_j$ :** Show how  $T_2$  matrix construction reveals the recursive pattern that leads to the general  $T_j$  structure with  $(j+1) \cdot 2^j$  nonzero entries.
- **Exercise 2: Complexity analysis  $O(n \log n)$ :** Derive  $O(n \log n)$  complexity for the direct matrix approach by counting multiplications in the  $T_j$  matrix operations.
- **Exercise 3: Inverse transform properties:** Demonstrate that  $T_j^{-1} = T_j^T$  due to orthogonality, maintaining the same complexity but enabling efficient inverse transforms.
- **Code Inspection: DHT Implementation:** Show the DHT algorithm implementation with visual matrix pattern, demonstrating how the  $T_j$  structure maps to sparse matrix multiplication with colored intervals.



- **Exercise 4: Multiplications and complexity:** Complete the FHT complexity analysis showing total multiplications and  $O(n)$  complexity, bridging from DHT theory to FHT implementation.
- **Code Inspection: FHT Implementation:** Show the FHT algorithm implementation with visual diagram, demonstrating the iterative approach and comparing with DHT matrix-based method.
- **Live Demo: Function Reconstruction:** Sample a test function and apply FHT/DHT to get Haar coefficients, then reconstruct using inverse transforms. Display visual comparison showing how the mathematical theory translates into accurate function approximation.

# Presentation Script: Part 3 - Analysis

- **Exercise 7: Hölder continuity:** Establish theoretical bounds  $|\langle f, H_{jk} \rangle| \leq C 2^{-j(\alpha+1/2)}$  for Hölder continuous functions, providing the foundation for understanding coefficient decay across scales.
- **Exercise 7: Proof of the inequality:** Complete the mathematical proof of the Hölder continuity bounds, showing the step-by-step derivation and substitution techniques.
- **Analysis: Understanding Coefficient Patterns:** Bridge from theoretical bounds to empirical analysis by introducing multi-resolution coefficient patterns and explaining how the Hölder continuity bounds provide the theoretical foundation.
- **Live Demo: Coefficient Analysis:** Visualize coefficients at different layers using `draw_coefficients_at_layer()` and compute ratios using `compute_ratios()`. Compare theoretical bounds with empirical results, showing how sampling methods affect the bounds.
- **Exercise 10: Invariance properties:** Prove invariance properties of  

- **Applications: Compression and Approximation:** Introduce N-term approximation and compression as practical applications of the theoretical foundation, explaining how the coefficient decay properties enable effective compression strategies.
- **Live Demo: Compression:** Apply FHT to a discontinuous function, sort coefficients by magnitude, and compute best N-term approximation error. Visualize compressed reconstruction showing how the mathematical theory enables practical compression with quantifiable error bounds.
- **Thank You:** Conclude the presentation and open for questions.