



书面作业 第5次 参考解答或提示

第1部分 基础

无

第2部分 理论

无

第3部分 综合应用 (T2 可选做)

T1. 形式化并证明下列自然语言推理是有效的. 个体域均为全总域.

- (1) 所有的玫瑰和蔷薇都是芳香带刺的, 因此, 所有的玫瑰都是带刺的.
- (2) 三角函数都是周期函数; 一些三角函数是连续函数; 所以, 一些周期函数是连续函数.
- (3) 每个科学工作者都是勤奋的, 每个既勤奋又聪明的人在他的事业中都将获得成功, 王大志是科学工作者并且是聪明的, 所以, 王大志在他的事业中将获得成功.
- (4) 如果一个人长期吸烟或酗酒, 那么他身体绝不会健康; 如果一个人身体不健康, 那么他就不能参加体育比赛; 有人参加了体育比赛. 所以, 有人不长期酗酒.
- (5) 小王是一年级生理科生; 每个非文科的一年级生都有辅导员; 凡小王的辅导员都是理科生; 所有的理科生都不是文科生. 所以, 至少有一个不是文科生的辅导员.

(1) 设 $S(x)$: x 是三角函数; $T(x)$: x 是周期函数; $P(x)$: x 是连续函数, 则推理可以形式化为:
 $\forall x(S(x) \rightarrow T(x)), \exists x(S(x) \wedge P(x)) \Rightarrow \exists x(T(x) \wedge P(x)).$

证明:

- | | |
|---|------------|
| (1) $\exists x(S(x) \wedge P(x))$ | P |
| (2) $S(c) \wedge P(c)$ | ES(1) |
| (3) $\forall x (S(x) \rightarrow T(x))$ | P |
| (4) $S(c) \rightarrow T(c)$ | US(3) |
| (5) $S(c)$ | T, I(2) |
| (6) $T(c)$ | T, I(4)(5) |
| (7) $P(c)$ | T, I(2) |
| (8) $T(c) \wedge P(c)$ | T, I(6)(7) |
| (9) $\exists x(T(x) \wedge P(x))$ | EG(8) |

(2) 设 $P(x)$: x 是玫瑰; $Q(x)$: x 是蔷薇; $R(x)$: x 是芳香的; $S(x)$: x 是带刺的, 则推理可以形式化为:
 $\forall x((P(x) \vee Q(x)) \rightarrow (R(x) \wedge S(x))) \Rightarrow \forall x(P(x) \rightarrow S(x))$

证明: 可以直接证明, 也可以用CP规则, 以及反证法.



直接法

- (1) $\forall x((P(x) \vee Q(x)) \rightarrow (R(x) \wedge S(x)))$ P
- (2) $P(x) \vee Q(x) \rightarrow R(x) \wedge S(x)$ US(1)
- (3) $(\neg P(x) \wedge \neg Q(x)) \vee (R(x) \wedge S(x))$ T,I(2)
- (4) $(\neg P(x) \vee S(x)) \wedge (\neg P(x) \vee R(x)) \wedge (\neg Q(x) \vee S(x)) \wedge (\neg Q(x) \vee S(x))$ T,I(3)
- (5) $\neg P(x) \vee S(x)$ T,I(4)
- (6) $P(x) \rightarrow S(x)$ T,I(5)
- $\forall x(P(x) \rightarrow S(x))$ UG,5

CP规则

- (1) $\forall x((P(x) \vee Q(x)) \rightarrow (R(x) \wedge S(x)))$ P
- (2) $P(x) \vee Q(x) \rightarrow R(x) \wedge S(x)$ US(1)
- (3) $P(x)$ P(附加, for: $P(x) \rightarrow S(x)$)
- (4) $P(x) \vee Q(x)$ T,I(3)
- (5) $R(x) \wedge S(x)$ T,I(2)(4)
- (6) $S(x)$ T,I(5)
- (7) $P(x) \rightarrow S(x)$ CP(3)(6)
- (8) $\forall x(P(x) \rightarrow S(x))$ UG(7)

反证法(略).

(3) 令 $M(x)$: x 是人; $K(x)$: x 是科学工作者; $Q(x)$: x 勤奋; $T(x)$: x 聪明; $S(x)$: x 将获得成功; a : 王壮志, 则推理可以形式化为:

$$\forall x((M(x) \wedge K(x)) \rightarrow Q(x)), \forall x((M(x) \wedge Q(x) \wedge T(x)) \rightarrow S(x)), M(a) \wedge K(a) \wedge T(a) \Rightarrow S(a)$$

证明:

- 1) $M(a) \wedge K(a) \wedge T(a)$ P
- 2) $\forall x((M(x) \wedge K(x)) \rightarrow Q(x))$ P
- 3) $(M(a) \wedge K(a)) \rightarrow Q(a)$ US,(2)
- 4) $M(a) \wedge K(a)$ T,I,(1)
- 5) $Q(a)$ T,I,(2),(4)
- 6) $M(a) \wedge T(a)$ T,I,(1)
- 7) $M(a) \wedge Q(a) \wedge T(a)$ T,I,(5),(6)
- 8) $\forall x((M(x) \wedge Q(x) \wedge T(x)) \rightarrow S(x))$ P
- 9) $(M(a) \wedge Q(a) \wedge T(a)) \rightarrow S(a)$ US,(8)
- 10) $S(a)$ T,I,(7),(9)

(4) 令 $M(x)$: x 是人; $C(x)$: x 长期吸烟; $K(x)$: x 长期酗酒; $J(x)$: x 身体健康; $P(x)$: x 能参加体育比赛, 则推理可以形式化为:



$\forall x((M(x) \wedge (C(x) \vee K(x))) \rightarrow \neg J(x)), \forall x((M(x) \wedge \neg J(x)) \rightarrow \neg P(x)), \exists x(M(x) \wedge P(x)) \Rightarrow \exists x(M(x) \wedge \neg K(x))$

证明:

- | | |
|---|--------------|
| 1) $\exists x(M(x) \wedge P(x))$ | P |
| 2) $M(c) \wedge P(c)$ | ES,(1) |
| 3) $\forall x((M(x) \wedge \neg J(x)) \rightarrow \neg P(x))$ | P |
| 4) $(M(c) \wedge \neg J(c)) \rightarrow \neg P(c)$ | US,(3) |
| 5) $P(c)$ | T,I,(2) |
| 6) $\neg(M(c) \wedge \neg J(c))$ | T,I,(4),(5) |
| 7) $\neg M(c) \vee J(c)$ | R,E,(6) |
| 8) $M(c)$ | T,I,(2) |
| 9) $J(c)$ | T,I,(7),(8) |
| 10) $\forall x((M(x) \wedge (C(x) \vee K(x))) \rightarrow \neg J(x))$ | P |
| 11) $(M(c) \wedge (C(c) \vee K(c))) \rightarrow \neg J(c)$ | US,(10) |
| 12) $\neg(M(c) \wedge (C(c) \vee K(c)))$ | T,I,(9),(11) |
| 13) $\neg M(c) \vee (\neg C(c) \wedge \neg K(c))$ | R,E,(12) |
| 14) $\neg C(c) \wedge \neg K(c)$ | T,I,(8),(13) |
| 15) $\neg K(c)$ | T,I,(14) |
| 16) $M(c) \wedge \neg K(c)$ | T,I,(8),(15) |
| 17) $\exists x(M(x) \wedge \neg K(x))$ | EG,(16) |

(5) 令a: 小王; $S(x)$: x是一年级生; $L(x)$: x是理科生; $W(x)$: x是文科生; $F(x,y)$: x是y的辅导员, 则推理可以形式化为:

$\forall x(S(x) \wedge \neg W(x) \rightarrow \exists y F(y,x)), S(a), L(a), \forall x(F(x,a) \rightarrow L(x)), \forall x(L(x) \rightarrow \neg W(x)) \Rightarrow \exists x \exists y (\neg W(x) \wedge F(x,y)).$

证明:

- | | |
|---|-------------|
| (1) $\forall x(L(x) \rightarrow \neg W(x))$ | P |
| (2) $L(a) \rightarrow \neg W(a)$ | US,(1) |
| (3) $L(a)$ | P |
| (4) $\neg W(a)$ | T,I,(2),(3) |
| (5) $S(a)$ | P |
| (6) $S(a) \wedge \neg W(a)$ | T,I,(4),(5) |
| (7) $\forall x(S(x) \wedge \neg W(x) \rightarrow \exists y F(y,x))$ | P |
| (8) $S(a) \wedge \neg W(a) \rightarrow \exists y F(y,a)$ | US,(7) |
| (9) $\exists y F(y,a)$ | T,I,(6),(8) |
| (10) $F(c,a)$ | ES,(9) |
| (11) $\forall x(F(x,a) \rightarrow L(x))$ | P |



(12) $F(c, a) \rightarrow L(c)$	US, (11)
(13) $L(c)$	T, I, (10), (12)
(14) $L(c) \rightarrow \neg W(c)$	US, (1)
(15) $\neg W(c)$	T, I, (13), (14)
(16) $\neg W(c) \wedge F(c, a)$	T, I, (10), (15)
(17) $\exists y (\neg W(c) \wedge F(c, y))$	EG, (16)
(18) $\exists x \exists y (\neg W(x) \wedge F(x, y))$	EG, (17)

因此，该推理是有效的。

T2. Consider the following problem. We know that horses are faster than dogs and that there is a greyhound that is faster than every rabbit. We know that Harry is a horse and that Ralph is a rabbit. Our job is to derive the fact that Harry is faster than Ralph (可以应用直接证明方法与归结法. 注意要增加两个公认的事实: Greyhound是Dog, 速度可以传递比较的).

Problem translated in FOPL:

$\forall x \forall y ((\text{Horse}(x) \wedge \text{Dog}(y)) \rightarrow \text{Faster}(x, y))$

$\exists y (\text{Greyhound}(y) \wedge (\forall z \text{Rabbit}(z) \rightarrow \text{Faster}(y, z)))$

$\text{Horse}(\text{Harry})$

$\text{Rabbit}(\text{Ralph})$

Derive the following fact:

$\text{Faster}(\text{Harry}, \text{Ralph})$

Added axioms to represent commonsense knowledge:

$\forall y (\text{Greyhound}(y) \rightarrow \text{Dog}(y))$

$\forall x \forall y \forall z ((\text{Faster}(x, y) \wedge \text{Faster}(y, z)) \rightarrow \text{Faster}(x, z))$

Proving using Proof Theory and a set of inference rules

1. $\forall x \forall y \text{Horse}(x) \wedge \text{Dog}(y) \rightarrow \text{Faster}(x, y)$	Premise
2. $\exists y \text{Greyhound}(y) \wedge (\forall z \text{Rabbit}(z) \rightarrow \text{Faster}(y, z))$	Premise
3. $\forall y \text{Greyhound}(y) \rightarrow \text{Dog}(y)$	Premise
4. $\forall x \forall y \forall z \text{Faster}(x, y) \wedge \text{Faster}(y, z) \rightarrow \text{Faster}(x, z)$	Premise
5. $\text{Horse}(\text{Harry})$	Premise
6. $\text{Rabbit}(\text{Ralph})$	Premise
7. $\text{Greyhound}(\text{Greg}) \wedge (\forall z \text{Rabbit}(z) \rightarrow \text{Faster}(\text{Greg}, z))$	ES (2)
8. $\text{Greyhound}(\text{Greg})$	T, I (7)
9. $\forall z \text{Rabbit}(z) \rightarrow \text{Faster}(\text{Greg}, z)$	T, I (7)
10. $\text{Rabbit}(\text{Ralph}) \rightarrow \text{Faster}(\text{Greg}, \text{Ralph})$	US (9)
11. $\text{Faster}(\text{Greg}, \text{Ralph})$	T, I (6) (10)
12. $\text{Greyhound}(\text{Greg}) \rightarrow \text{Dog}(\text{Greg})$	US (3)
13. $\text{Dog}(\text{Greg})$	T, I (12) (8)
14. $\text{Horse}(\text{Harry}) \wedge \text{Dog}(\text{Greg}) \rightarrow \text{Faster}(\text{Harry}, \text{Greg})$	US (1)
15. $\text{Horse}(\text{Harry}) \wedge \text{Dog}(\text{Greg})$	T, I (5) (13)
16. $\text{Faster}(\text{Harry}, \text{Greg})$	T, I (14) (15)
17. $\text{Faster}(\text{Harry}, \text{Greg}) \wedge \text{Faster}(\text{Greg}, \text{Ralph}) \rightarrow \text{Faster}(\text{Harry}, \text{Ralph})$	US (4)



18. $\text{Faster}(\text{Harry}, \text{Greg}) \wedge \text{Faster}(\text{Greg}, \text{Ralph})$ T, I (11) (16)
19. $\text{Faster}(\text{Harry}, \text{Ralph})$ T, I (17) (19)
QED.

Using Resolution(归结法) to determine logical entailment

1. $\{\neg \text{Horse}(x), \neg \text{Dog}(y), \text{Faster}(x, y)\}$ Premise
 2. $\{\text{Greyhound}(\text{Greg})\}$ Premise
 3. $\{\neg \text{Rabbit}(z), \text{Faster}(\text{Greg}, z)\}$ Premise
 4. $\{\neg \text{Greyhound}(y), \text{Dog}(y)\}$ Premise
 5. $\{\neg \text{Faster}(x, y), \neg \text{Faster}(y, z), \text{Faster}(x, z)\}$ Premise
 6. $\{\text{Horse}(\text{Harry})\}$ Premise
 7. $\{\text{Rabbit}(\text{Ralph})\}$ Premise
 8. $\{\neg \text{Faster}(\text{Harry}, \text{Ralph})\}$ Negated Goal
 9. $\{\text{Dog}(\text{Greg})\}$ (2) (4) 归结
 10. $\{\neg \text{Dog}(y), \text{Faster}(\text{Harry}, y)\}$ (6) (1) 归结
 11. $\{\text{Faster}(\text{Harry}, \text{Greg})\}$ (9) (10) 归结
 12. $\{\text{Faster}(\text{Greg}, \text{Ralph})\}$ (7) (3) 归结
 13. $\{\neg \text{Faster}(\text{Greg}, z), \text{Faster}(\text{Harry}, z)\}$ (11) (5) 归结
 14. $\{\text{Faster}(\text{Harry}, \text{Ralph})\}$ (12) (13) 归结
 15. $\{\}$ (14) (8) 归结
- QED.