

# Numerical Analysis homework # 1

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## I. Analysis of the Bisection Method on the Interval $[1.5, 3.5]$

### Width of the interval

Given the initial interval  $[a_0, b_0] = [1.5, 3.5]$ , the initial width is  $W_0 = b_0 - a_0 = 2$ . The interval width is halved at each step, so the width at step  $n$ ,  $W_n$ , is:

$$W_n = \frac{W_0}{2^n} = \frac{2}{2^n} = \frac{1}{2^{n-1}}$$

### Supremum of the error

At step  $n$ , let the interval be  $[a_n, b_n]$  and the midpoint  $c_n = (a_n + b_n)/2$ . The root  $r$  satisfies  $r \in [a_n, b_n]$ . The distance  $|r - c_n|$  is maximized when  $r$  is located at an endpoint of the interval. This gives the supremum of the error as  $1/2$  the interval width.

$$\sup |r - c_n| = c_n - a_n = b_n - c_n = \frac{b_n - a_n}{2} = \frac{W_n}{2^1} = \frac{W_{n-1}}{2^2} = \frac{1}{2^n}$$

## II. Proof for the Number of Steps for a Given Relative Error

Denote the root by  $r$ , we need to find the number of steps  $n$  such that  $|r - c_n|/|r|$  is no greater than  $\epsilon$ .

$$\frac{|r - c_n|}{|r|} \leq \epsilon$$

First, the absolute error is bounded by half the interval width at step  $n$ .

$$|r - c_n| \leq \frac{b_n - a_n}{2} = \frac{b_0 - a_0}{2^{n+1}}$$

Given that  $r \in [a_0, b_0]$  and  $a_0 > 0$ , the magnitude of the root is bounded below by  $a_0$ .

$$|r| \geq a_0$$

Therefore, we can conclude that:

$$\frac{|r - c_n|}{|r|} \leq \frac{(b_0 - a_0)/2^{n+1}}{a_0} = \frac{b_0 - a_0}{a_0 \cdot 2^{n+1}}$$

To guarantee the desired accuracy, we enforce this upper bound to be less than or equal to  $\epsilon$ .

$$\frac{b_0 - a_0}{a_0 \cdot 2^{n+1}} \leq \epsilon$$

$$\frac{b_0 - a_0}{\epsilon \cdot a_0} \leq 2^{n+1}$$

Taking the logarithm of both sides yields:

$$\log \left( \frac{b_0 - a_0}{\epsilon \cdot a_0} \right) \leq \log(2^{n+1})$$

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$$\log(b_0 - a_0) - \log \epsilon - \log a_0 \leq (n + 1) \log 2$$

Rearranging terms :

$$n \geq \frac{\log(b_0 - a_0) - \log \epsilon - \log a_0}{\log 2} - 1$$

This completes the proof.

### III. Application of Newton's Method

The problem is to find a root for the polynomial equation  $p(x) = 4x^3 - 2x^2 + 3 = 0$  with the starting point  $x_0 = -1$ . We have the derivative of  $p(x)$ :

$$p'(x) = 12x^2 - 4x$$

The iteration formula for this specific problem is:

$$x_{k+1} = x_k - \frac{4x_k^3 - 2x_k^2 + 3}{12x_k^2 - 4x_k}$$

Starting with  $x_0 = -1$ , we perform 4 iterations. The results are organized in the following table.

Iteration, $k$	$x_k$	$p(x_k)$	$p'(x_k)$
0	-1.000000	-3.000000	16.000000
1	-0.812500	-0.465820	11.171875
2	-0.770804	-0.020138	10.212886
3	-0.768832	-0.000044	10.168568
4	-0.768828	-0.000000	10.168472

After four iterations, the approximation of the root is  $x_4$ .

$$x_4 \approx -0.768828$$

### IV. Convergence Analysis of a Modified Newton's Method

Let  $r$  be the root, with  $f(r) = 0$ , and let the error be  $e_n = x_n - r$ . The error recurrence relation derived from the iteration formula is:

$$e_{n+1} = x_{n+1} - r = e_n - \frac{f(x_n)}{f'(x_0)}$$

Expanding  $f(x_n) = f(r + e_n)$  as a Taylor series around  $r$  gives:

$$\begin{aligned} f(x_n) &= f'(r)e_n + \frac{f''(r)}{2}e_n^2 + O(e_n^3) \\ e_{n+1} &= e_n - \frac{f'(r)e_n + \frac{f''(r)}{2}e_n^2 + O(e_n^3)}{f'(x_0)} \\ e_{n+1} &= \left(1 - \frac{f'(r)}{f'(x_0)}\right)e_n - \left(\frac{f''(r)}{2f'(x_0)}\right)e_n^2 + O(e_n^3) \end{aligned}$$

By comparing with the form  $e_{n+1} = Ce_n^s$ , we identify the exponent of the dominant term as  $s$  and its coefficient as  $C$ . Therefore, this iteration formula is **linearly convergent**.

$$s = 1$$

$$C = 1 - \frac{f'(r)}{f'(x_0)}$$

## V. Convergence of the Iteration $x_{n+1} = \tan^{-1} x_n$

The iteration  $x_{n+1} = \tan^{-1} x_n$  has a unique fixed point  $\alpha$  that solves  $\tan^{-1} \alpha = \alpha$ , which is  $\alpha = 0$ . If  $x_0 = 0$ , the sequence converges trivially. Thus, we only consider the case where  $x_0 \neq 0$ . To analyze convergence, we consider the absolute value of the terms. For any  $x \neq 0$ , it is a known property that:

$$|\tan^{-1} x| < |x|$$

Applying this property to the iteration for any  $x_n \neq 0$  gives a strictly **contractive** relationship.

$$|x_{n+1}| = |\tan^{-1} x_n| < |x_n|$$

This shows that the sequence of absolute values,  $\{|x_n|\}$ , is **strictly decreasing** and **bounded below by 0**. Therefore, the sequence of absolute values must converge to 0.

$$\lim_{n \rightarrow \infty} |x_n| = 0$$

This implies that the sequence  $\{x_n\}$  itself converges to 0 for all initial values  $x_0 \in (-\pi/2, \pi/2)$ .

## VI. Analysis of a Continued Fraction

Assuming the sequence converges to a value  $x$ , this value must be a fixed point satisfying the relation:

$$x_n = \frac{1}{p + \frac{1}{p + \dots}} \implies x = \frac{1}{p + x}$$

The solutions are trivial.

$$x = \frac{-p \pm \sqrt{p^2 + 4}}{2}$$

Since the sequence of convergents  $x_n$  is always positive for  $p > 1$ , its limit  $x$  must be positive. We thus select the positive root.

$$x = \frac{-p + \sqrt{p^2 + 4}}{2}$$

To prove convergence, we analyze the iteration  $x_{n+1} = g(x_n)$  with the function  $g(x) = 1/(p + x)$ . We will show that  $g(x)$  is a contraction mapping for  $x \geq 0$ . First, we find the derivative.

$$g'(x) = -\frac{1}{(p + x)^2}$$

The magnitude of the derivative is **bounded**. For  $p > 1$  and any  $x \geq 0$ , the denominator  $(p + x)^2 > p^2$ . This provides a uniform bound on the magnitude of the derivative.

$$|g'(x)| = \frac{1}{(p + x)^2} < \frac{1}{p^2}$$

Since  $p > 1$ , we have  $p^2 > 1$ , which implies that  $1/p^2 < 1$  and  $k = \frac{1}{p^2} < 1$ . Because  $|g'(x)| \leq k < 1$  for all  $x \geq 0$ , the function  $g(x)$  is a contraction mapping on  $[0, \infty)$ . Therefore, by the **Contraction Mapping Theorem**, the sequence converges for any initial  $x_0 \geq 0$ .

## VII. Bisection Method Analysis when the Interval Contains Zero

### Derivation of the Inequality for the Number of Steps

The upper bound for the relative error is:

$$\frac{|r - c_n|}{|r|} \leq \frac{b_0 - a_0}{|r| \cdot 2^{n+1}}$$

Under the condition  $a_0 < 0 < b_0$ , the root  $r$  can be arbitrarily close to 0. The supremum of the error bound is therefore not finite.

$$\sup_{r \in [a_0, b_0]} \left( \frac{b_0 - a_0}{|r| \cdot 2^{n+1}} \right) \rightarrow \infty \quad \text{as } r \rightarrow 0$$

Thus, no general inequality for  $n$  exists that can guarantee the relative error is bounded by a given  $\epsilon$ .

## Appropriateness of Relative Error

Relative error is an inappropriate measure when the root  $r$  is near zero. For a small but fixed absolute error  $|r - c_n| = \delta > 0$ , the relative error diverges as  $r \rightarrow 0$ .

$$\lim_{r \rightarrow 0} \frac{|r - c_n|}{|r|} = \lim_{r \rightarrow 0} \frac{\delta}{|r|} = \infty$$

Absolute error,  $|r - c_n|$ , is the proper measure in this case.

## VIII. Newton's Method for Roots of Multiplicity $k$

### Detection of a Multiple Zero

A multiple zero can be detected by two primary observations during the iteration process. First, the convergence of the sequence  $\{x_n\}$  to the root  $r$  is slow. It degrades from quadratic to linear. Second, as the iterates  $x_n$  approach the root  $r$ , both the function values  $f(x_n)$  and the derivative values  $f'(x_n)$  converge to 0. For a simple root,  $f'(x_n)$  would converge to a non-zero constant,  $f'(r)$ .

### Proof of Restored Quadratic Convergence

We are given the modified Newton's iteration for a root of multiplicity  $k$ .

$$x_{n+1} = g(x_n) \quad \text{where} \quad g(x) = x - k \frac{f(x)}{f'(x)}$$

To prove quadratic convergence, we must show that  $g'(r) = 0$ . The derivative of  $g(x)$  is:

$$\begin{aligned} g'(x) &= 1 - k \left[ \frac{f'(x)f'(x) - f(x)f''(x)}{[f'(x)]^2} \right] = 1 - k \left[ 1 - \frac{f(x)f''(x)}{[f'(x)]^2} \right] \\ g'(x) &= 1 - k + k \frac{f(x)f''(x)}{[f'(x)]^2} \end{aligned}$$

Since  $r$  is a root of multiplicity  $k$ , we can write  $f(x) = (x - r)^k h(x)$  for some function  $h(x)$  where  $h(r) \neq 0$ . The leading terms of the derivatives near  $x = r$  are:

$$\begin{aligned} f(x) &\approx (x - r)^k h(r) \\ f'(x) &\approx k(x - r)^{k-1} h(r) \\ f''(x) &\approx k(k-1)(x - r)^{k-2} h(r) \\ \lim_{x \rightarrow r} \frac{f(x)f''(x)}{[f'(x)]^2} &= \frac{(x - r)^k h(r) \cdot k(k-1)(x - r)^{k-2} h(r)}{[k(x - r)^{k-1} h(r)]^2} \\ &= \frac{k(k-1)(x - r)^{2k-2} h(r)^2}{k^2(x - r)^{2k-2} h(r)^2} = \frac{k-1}{k} \end{aligned}$$

Substituting this back for  $g'(r)$ :

$$g'(r) = 1 - k + k \left( \frac{k-1}{k} \right) = 1 - k + (k-1) = 0$$

Since  $g'(r) = 0$  (and assuming  $g''(r) \neq 0$ ), the fixed-point iteration for  $g(x)$  converges quadratically.

### Acknowledgement

Honestly, the **translation** was done using LLM tools, but the answers and solutions were written by myself.