

CSC 3210

# Computer Organization and Programming

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CHAPTER 5: PROCEDURES (STACK)

# Why Procedures? (Why stack?)

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- Procedures, Functions, Methods, or subroutine
  - ARE the most fundamental language feature for:
    - Code reuse.
    - Ease debugging
    - Abstraction
- They allow us to refer to some piece of code **by a name**
- All we need to know is
  - how many arguments are needed,
  - what type of arguments,
  - what the subroutine returns,
  - what the subroutine computes
    - it's not necessary to know how the subroutine does what it does

---

```
Main (){  
    pickMin( int x, int y, int z )  
}
```

```
int pickMin( int x, int y, int z )  
{  
    int min = x ;  
    if ( y < min ) min = y ;  
    if ( z < min ) min = z ;  
    return min ;  
}
```

# Why Procedures? (Why stack?)

---

- What happens in a subroutine call?

1. When a subroutine call is executed,
  - The arguments **need to be evaluated** to values
2. Then:
  - **control flow jumps** to the body of the subroutine (How?),
  - and code **begins executing**
3. Once a **return statement** has been encountered,
  - we're done with the subroutine,
  - and **return** back to the subroutine call (How?)

`pickMin( int x, int y, int z )`

```
int pickMin( int x, int y, int z )
{
    int min = x ;
    if ( y < min ) min = y ;
    if ( z < min ) min = z ;
    return min ;
}
```

- In order to understand subroutine calls, you need to understand **the stack**

# Outline

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- **Stack Operations**
- Defining and Using Procedures
- Linking to an External Library
- The Irvine32 Library

# Stack Operations

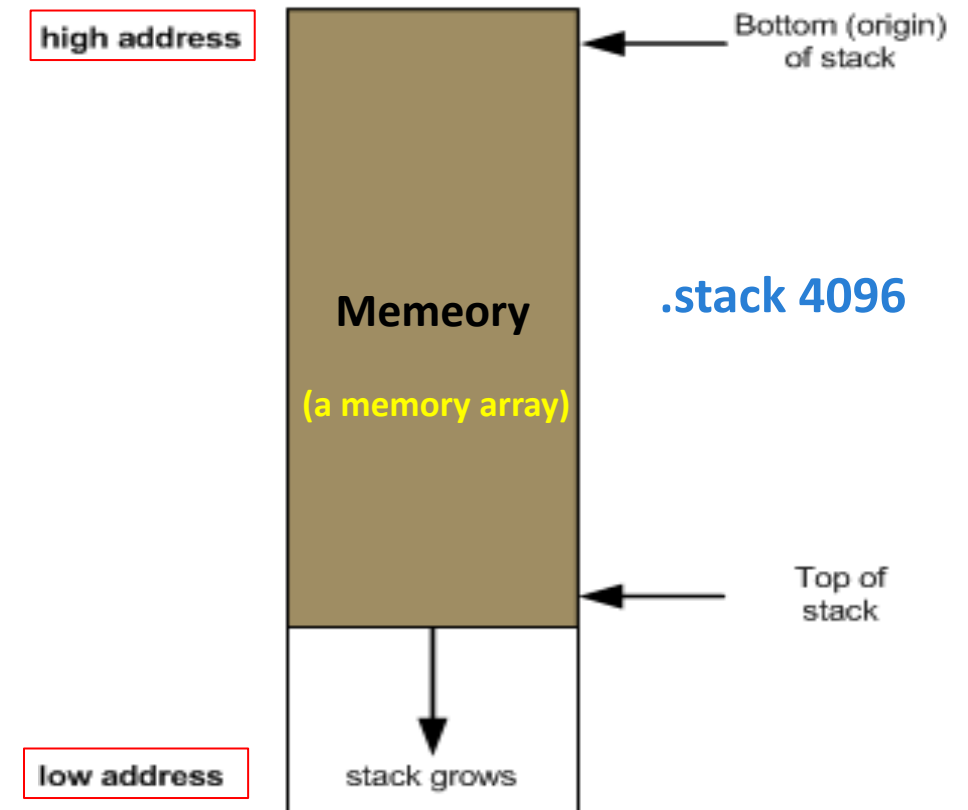
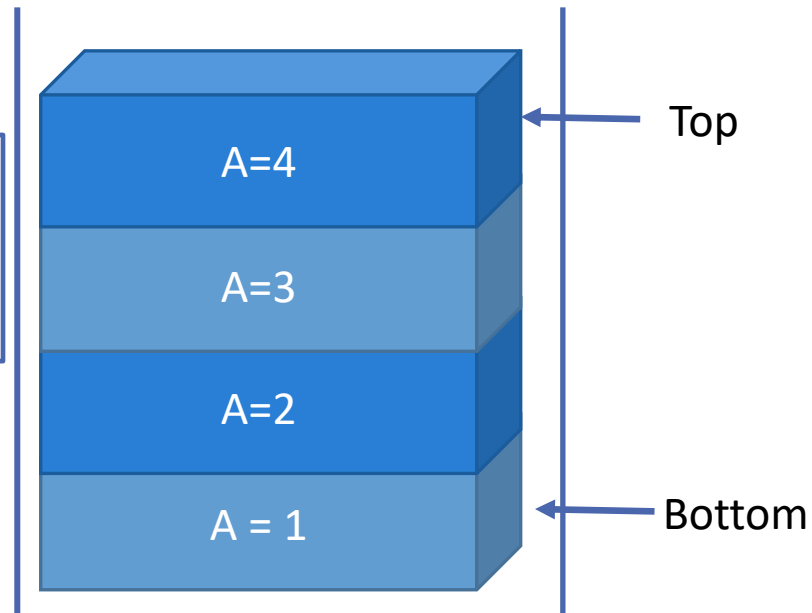
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- **Runtime Stack**
- PUSH Operation
- POP Operation
- PUSH and POP Instructions
- Using PUSH and POP
  - Example: Reversing a String
- Related Instructions

# Runtime Stack

- Imagine a stack of plates . . .
  - plates are only **added to the top**
  - plates are only **removed from the top**
  - **LIFO** structure

PUSH operation  
Stores item in stack  
POP operation  
Retrieve item from stack

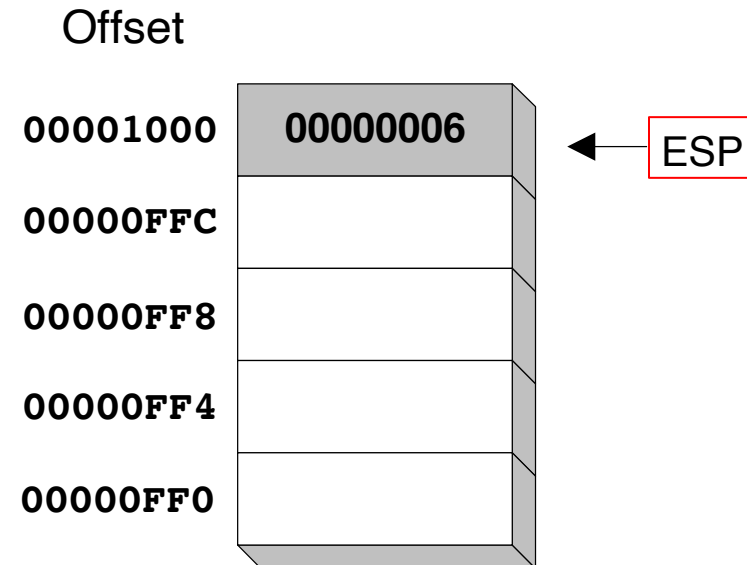


# Runtime Stack

- The runtime stack is **a memory array**
- Managed by the CPU, using two registers
  - **SS** (stack **segment**)
  - **ESP** (stack pointer) \*

- Stack **starts at some address**
- Then **grows down to a lower address.**

**.stack 4096**



\* SP in Real-address mode



# Stack Operations

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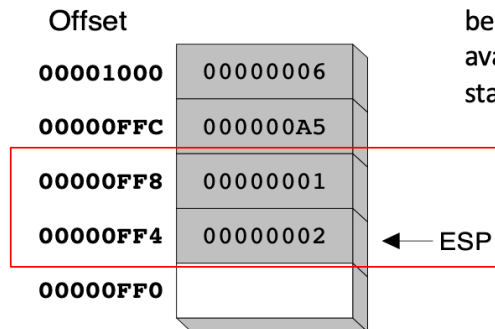
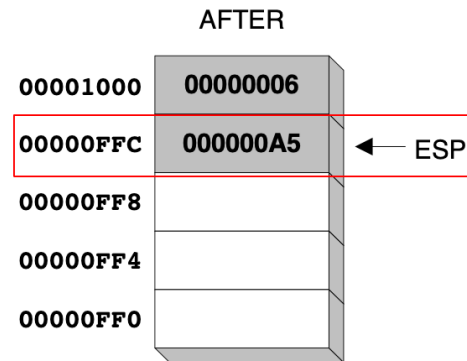
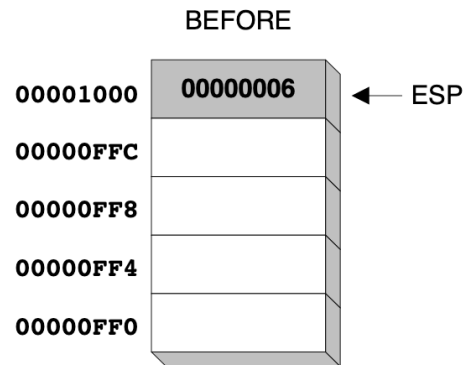
- Runtime Stack
- **PUSH Operation**
- POP Operation
- Using PUSH and POP
  - Example: Reversing a String
- Related Instructions

# PUSH Operation

- A 32-bit push operation

- **decrements** the **stack pointer** by 4
- and **copies a value** into the location **pointed to by the stack pointer**.

```
PUSH reg/mem16  
PUSH reg/mem32  
PUSH imm32
```



The stack **grows downward**. The area below ESP is always available (unless the stack has overflowed).

ESP = stack pointer (Top)  
EBP = Base Pointer (Bottom)

# Stack Operations

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- Runtime Stack
- PUSH Operation
- **POP Operation**
- Using PUSH and POP
  - Example: Reversing a String
- Related Instructions

# POP Operation

- Copies value at stack[ESP] into a register or variable.
- Adds ***n*** to ESP, where ***n*** is either **2 bytes** or **4 bytes**
  - value of ***n*** depends on the attribute of the operand receiving the data

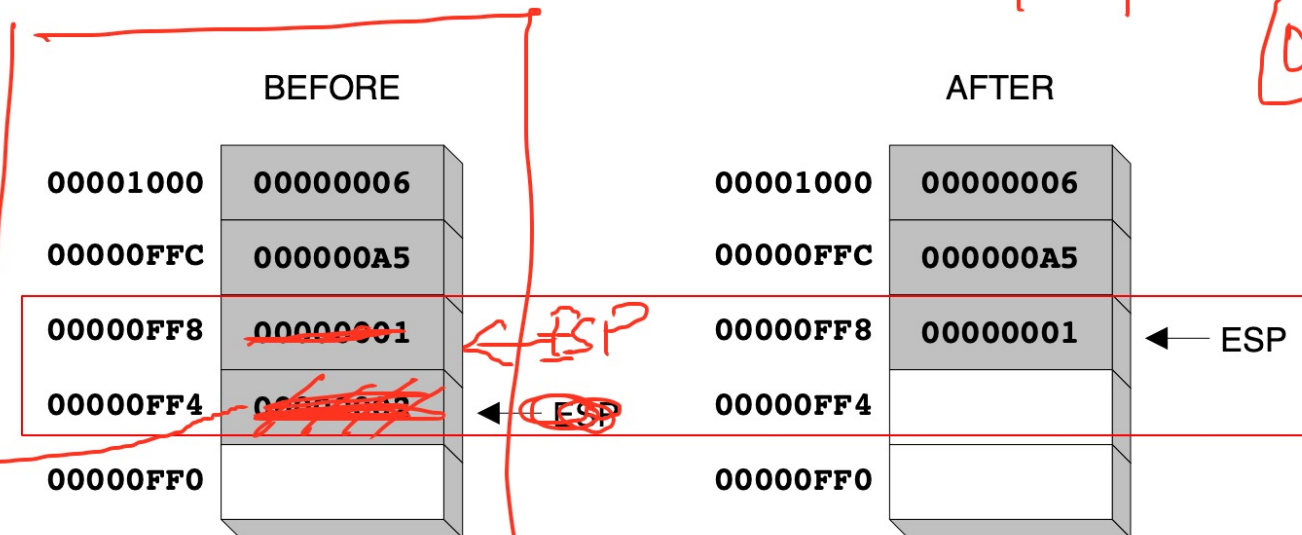
POP reg/mem16  
POP reg/mem32

pop EAX

EAX [00000002]

pop EBX

[00000001]



Stack

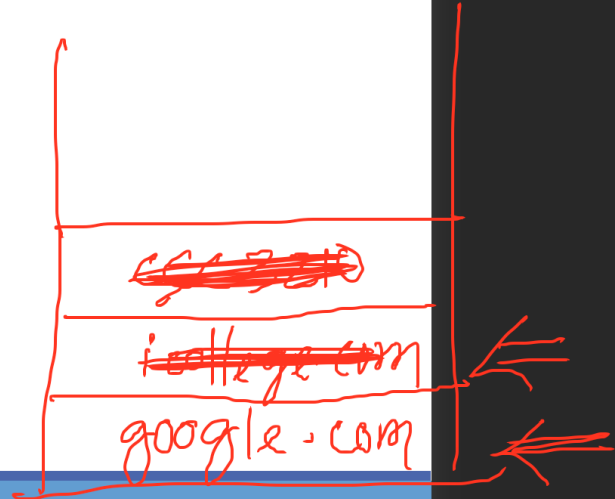
## Stack Applications

- A **stack** makes
  - a convenient temporary save area for registers when they (ex. ECX) are used for more than one purpose
  - After they are modified, they can be **restored** to their original values
- When the **CALL** instruction executes,
  - the **CPU** saves the current subroutine's return address on the stack
- When calling a subroutine,
  - you **pass input values** called arguments by **pushing** them on the stack
- The stack provides temporary storage for local variables inside subroutines

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# Stack Operations

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- Runtime Stack
- PUSH Operation
- POP Operation
- **Using PUSH and POP**
  - Example: Reversing a String
- Related Instructions

# Using PUSH and POP

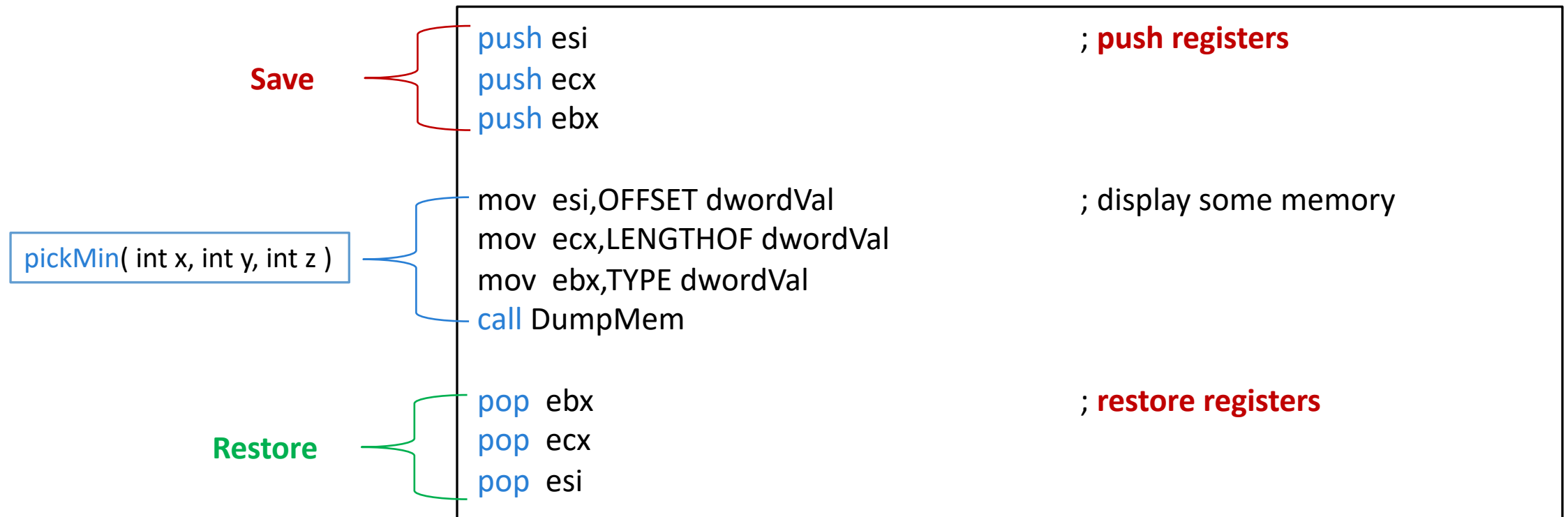
- **Save** and **restore** registers **when** they contain important values.
- PUSH and POP instructions occur in the opposite order

```
.data
array DWORD 1,2,3,4,5,6,7,8,9,0Ah,0Bh
.code
main PROC
    mov     esi,OFFSET array           ; starting OFFSET
    mov     ecx,LENGTHOF array        ; number of units
    mov     ebx,TYPE array             ; doubleword format
    call    DumpMem
```

pickMin( int x, int y, int z )

**DumpMem** Writes a block of memory to the console window in hexadecimal

# Using PUSH and POP



- **DumpMem** Writes a block of memory to the console window in hexadecimal



"Good morning"



← Top

GNINROM - DOOG



# Example1: Nested Loop

---

- When creating a nested loop, push the outer loop counter before entering the inner loop:

	mov ecx,100	; set outer loop count
L1:		; begin the outer loop
	push ecx	; save outer loop count
	mov ecx,20	; set inner loop count
L2:		; begin the inner loop
	;	
	;	
	loop L2	; repeat the inner loop
	pop ecx	; restore outer loop count
	loop L1	; repeat the outer loop



## Example2: Reversing a String

Why **must** each character be put in **EAX** before it is pushed?

Because only **doubleword** (32-bit) or **word** (16-bit) values can be pushed on the stack.

```
; Reversing a String                                (RevStr.asm)

.386
.model flat,stdcall
.stack 4096
ExitProcess PROTO,dwExitCode:DWORD

.data
aName BYTE "Abraham Lincoln",0
nameSize = ($ - aName) - 1

.code
main PROC
; Push the name on the stack.
    mov     ecx,nameSize
    mov     esi,0
L1:  movzx   eax,aName[esi]           ; get character
    push    eax                     ; push on stack
    inc     esi
    loop    L1

; Pop the name from the stack, in reverse,
; and store in the aName array.
    mov     ecx,nameSize
    mov     esi,0
L2:  pop     eax                     ; get character
    mov     aName[esi],al           ; store in string
    inc     esi
    loop    L2
```

# Attendance!

---

# Stack Operations

---

- Runtime Stack
- PUSH Operation
- POP Operation
- Using PUSH and POP
  - Example: Reversing a String
- **Related Instructions**

# Related Instructions

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- **PUSHFD** and **POPFD**
  - push and pop the **EFLAGS register**
- **PUSHAD** pushes the **32-bit general-purpose registers on the stack**
  - order: **EAX, ECX, EDX, EBX, ESP, EBP, ESI, EDI**
- **POPAD** pops **the same registers off the stack in reverse order**
  - **PUSHA** and **POPA** do the same for **16-bit** registers

# Outline

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- Stack Operations
- **Defining and Using Procedures**
- Linking to an External Library
- The Irvine32 Library



# Defining and Using Procedures

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- **Creating Procedures**
- **Documenting Procedures**
- **Example: SumOf Procedure**
- CALL and RET Instructions
- Nested Procedure Calls
- Local and Global Labels
- Procedure Parameters

# Creating Procedures

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- A **procedure** is the ASM equivalent of a Java or C++ **function/method**
- A procedure is **declared** using the **PROC** and **ENDP** directives
- It must be **assigned** a **name** (a valid identifier)
- **When** you **create** a **procedure** other than your **program's startup procedure (main)**:
  - End it with a **RET** instruction
  - **RET** forces the CPU to return to the location from where the procedure was called (address of call next instruction)

sample **PROC**

•

•

**ret**

sample **ENDP**

main **PROC**

•

•

main **ENDP**

# Documenting Procedures

---

- Suggested **documentation** for each procedure:
  - A description of all tasks accomplished by the procedure
  - Receives: A list of input parameters; state their usage and requirements
  - Returns: A description of values returned by the procedure
  - Requires: Optional list of requirements called **preconditions** that must be satisfied before the procedure is called

If a procedure is called **without its preconditions satisfied**, it will probably **not produce the expected output**.

# Example: SumOf Procedure

---

```
;-----  
; Calculates and returns the sum of three 32-bit integers.  
; Receives: EAX, EBX, ECX, the three integers. May be  
; signed or unsigned.  
; Returns: EAX = sum, and the status flags (Carry,  
; Overflow, etc.) are changed.  
; Requires: nothing  
;-----
```

```
SumOf PROC  
    add eax,ebx  
    add eax,ecx  
    ret  
SumOf ENDP
```

```
sample PROC  
    .  
    .  
    ret  
sample ENDP
```

# Defining and Using Procedures

---

- Creating Procedures
- Documenting Procedures
- Example: SumOf Procedure
- **CALL and RET Instructions**
- Nested Procedure Calls
- Local and Global Labels
- Procedure Parameters

# CALL and RET Instructions

- The **CALL** instruction calls a procedure:
  - 1) pushes offset of next instruction on the **stack**
  - 2) copies the address of the called procedure into **EIP**
- The **RET** instruction **returns** from a procedure:
  - **pops** top of **stack** into **EIP** 😞

```
main PROC
    call MySub
    mov eax, ebx
    .
    .
main ENDP

MySub PROC
    mov eax, edx
    .
    .
    ret
MySub ENDP
```

# CALL-RET Example

0000025 is the **offset** of the instruction immediately following the **CALL** instruction

00000040 is the offset of the first instruction inside **MySub**

```
main PROC
    00000020 call MySub
    00000025 mov  eax,ebx
    .
    .
main ENDP

MySub PROC
    00000040 mov  eax,edx
    .
    .
    ret
MySub ENDP
```

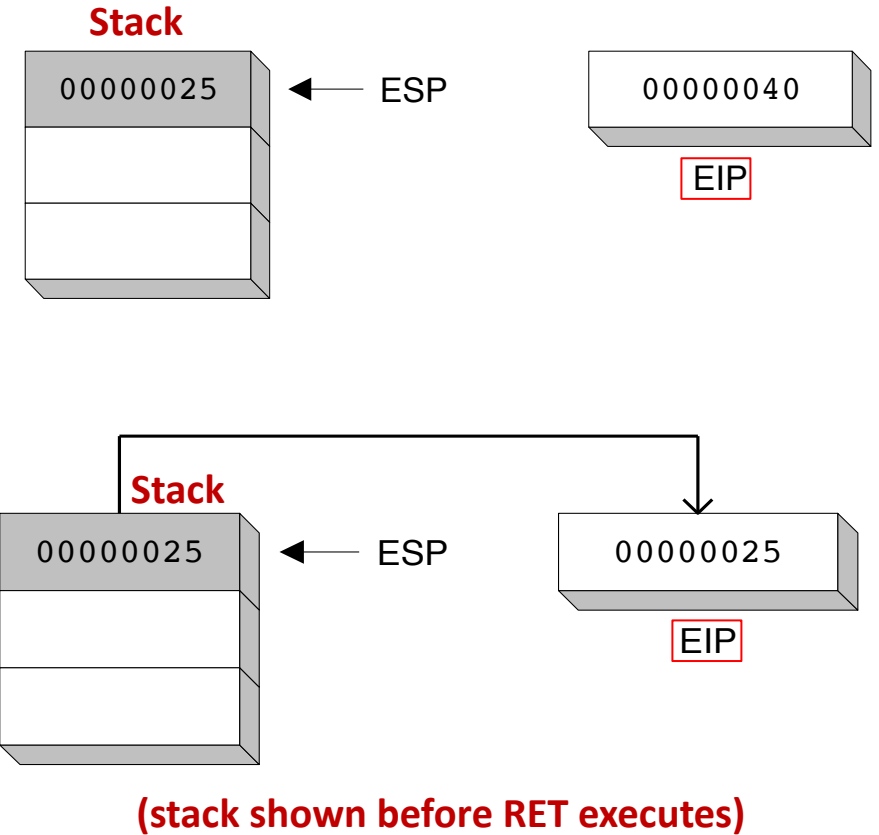
# CALL-RET Example

The **CALL** instruction pushes **00000025** onto the **stack**, and loads **00000040** into **EIP**

```
main PROC
    00000020 call MySub
    00000025 mov eax,ebx
    .
    .
main ENDP

MySub PROC
    00000040 mov eax,edx
    .
    .
    ret
MySub ENDP
```

The **RET** instruction pops **00000025** from the **stack** into **EIP**





# Defining and Using Procedures

---

- Creating Procedures
- Documenting Procedures
- Example: SumOf Procedure
- CALL and RET Instructions
- **Nested Procedure Calls**
- **Local and Global Labels**
- Procedure Parameters

# Nested Procedure Calls

1. By the time **Sub3** is called, **the stack** contains all **three** return addresses (**main**, **sub1**, **sub2**):
2. After the **return**,
  - **ESP** points to the next-highest stack entry.
3. Finally, when **Sub1** returns, **stack[ESP]** is **popped** into the instruction pointer, and execution resumes in **main**:

**Stack**

```
main PROC  
.  
.  
    call Sub1  
    exit  
main ENDP
```

```
Sub1 PROC  
.  
.  
    call Sub2  
    ret  
Sub1 ENDP
```

```
Sub2 PROC  
.  
.  
    call Sub3  
    ret  
Sub2 ENDP
```

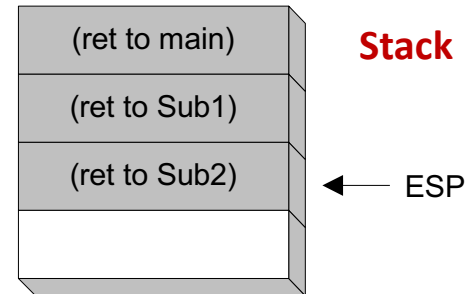
```
Sub3 PROC  
.  
.  
    ret  
Sub3 ENDP
```



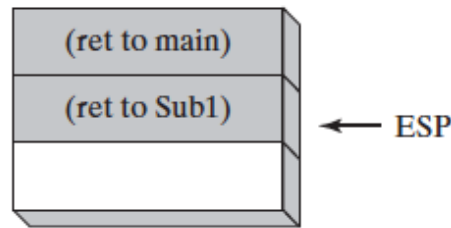


# Nested Procedure Calls

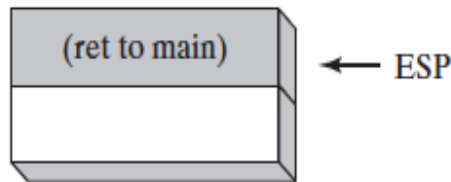
1. By the time **Sub3** is called, **the stack** contains all **three** return addresses (**main**, **sub1**, **sub2**):



2. After the **return**,
  - **ESP** points to the next-highest stack entry.



3. Finally, when **Sub1** returns, **stack[ESP]** is **popped** into the instruction pointer, and execution resumes in **main**:



```
main PROC
.  
.  
    call Sub1  
    exit  
main ENDP
```

```
Sub1 PROC
.  
.  
    call Sub2  
    ret  
Sub1 ENDP
```

```
Sub2 PROC
.  
.  
    call Sub3  
    ret  
Sub2 ENDP
```

```
Sub3 PROC
.  
.  
    ret  
Sub3 ENDP
```

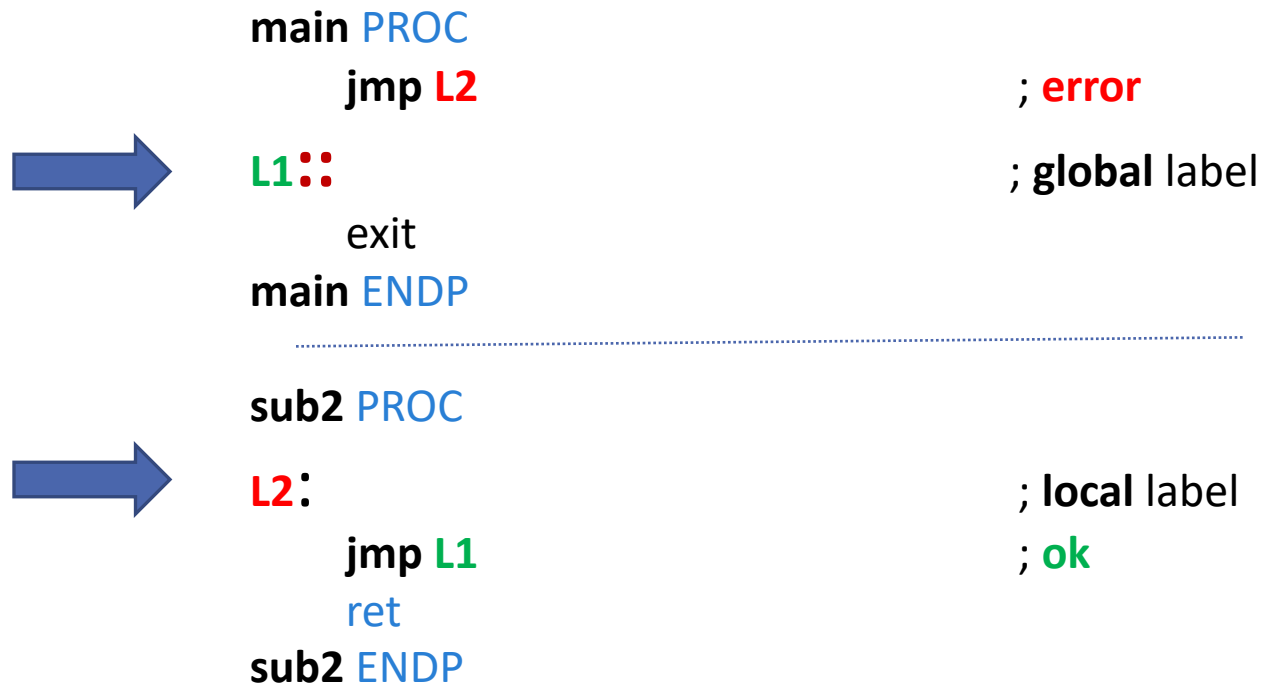
# Nested Procedure Calls

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- **The stack** proves itself a useful device for
  - **remembering information**, including nested procedure calls
- **Stack structures** are used in situations where programs **must retrace their steps in a specific order**

# Local and Global Labels

- A **local label** is visible only to statements inside the same procedure
- A **global label** is **visible everywhere**



```
main PROC
    jmp L2                                ; error
    L1::                                  ; global label
    exit
main ENDP

.....

sub2 PROC
    L2:                                  ; local label
    jmp L1                                ; ok
    ret
sub2 ENDP
```

# Defining and Using Procedures

---

- Creating Procedures
- Documenting Procedures
- Example: SumOf Procedure
- CALL and RET Instructions
- Nested Procedure Calls
- Local and Global Labels
- **Procedure Parameters**



# Procedure Parameters

- Why parameters?

- A good procedure might be usable in many different programs

- but not if it refers to specific variable names

Example: calculating the sum of an integer array:

- it's **not a good idea to include references to specific variable names** inside the procedure
- If you did, the procedure could only be used with one array

- **Parameters** help to make procedures flexible because **parameter values can change at runtime**

ArraySum PROC

```
myarray DWORD 10000h,20000h,30000h,40000h,50000h
mov esi,0                ; array index
mov eax,0                ; set the sum to zero
mov ecx,LENGTHOF myarray ; set number of elements
```

```
L1:  add eax,myArray[esi] ; add each integer to sum
     add esi,4           ; point to next integer
     loop L1             ; repeat for array size
```

```
mov theSum,eax           ; store the sum
ret
```

ArraySum ENDP

What if you wanted to calculate the sum of two or three arrays within the same program?

# Procedure Parameters

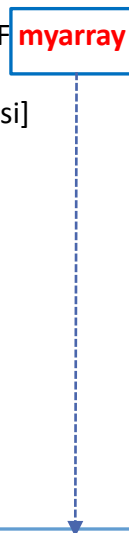
## Does not use parameters

ArraySum PROC

```
myarray DWORD 10000h,20000h,30000h,40000h,50000h
mov esi,0 ; array index
mov eax,0 ; set the sum to zero
mov ecx,LENGTHOF myarray ; set number of elements

L1: add eax,myArray[esi] ; add each integer to sum
    add esi,4 ; point to next integer
    loop L1 ; repeat for array size

mov theSum,eax ; store the sum
ret
ArraySum ENDP
```



Not a good idea to include references to specific variable names inside the procedure

## Uses parameters

ArraySum PROC

; Receives: **ESI** points to an array of doublewords,  
; **ECX** = number of array elements.  
; Returns: **EAX** = sum  
;-----

```
mov eax,0 ; set the sum to zero

L1: add eax,[esi] ; add each integer to sum
    add esi,4 ; point to next integer
    loop L1 ; repeat for array size

ret
ArraySum ENDP
```

- This version of ArraySum returns the sum of any doubleword array whose address is in ESI.
- The sum is returned in EAX:

# Procedure Parameters: Passing Register Arguments to Procedures

## Main

```
; Testing the ArraySum procedure (TestArraySum.asm)

.386
.model flat, stdcall
.stack 4096
ExitProcess PROTO, dwExitCode:DWORD

.data
array DWORD 10000h,20000h,30000h,40000h,50000h
theSum DWORD ?

.code
main PROC
    mov     esi,OFFSET array    ; ESI points to array
    mov     ecx,LENGTHOF array ; ECX = array count
    call    ArraySum           ; calculate the sum
    mov     theSum,eax          ; returned in EAX

    INVOKE ExitProcess,0
main ENDP
```

Subroutines **must preserve all registers**, except for **eax**, **ecx**, and **edx**, which **can be changed** across Subroutines **call**, and **esp**, which must be updated according to the calling convention.

## Subroutine

```
-----
; ArraySum
;
; Calculates the sum of an array of 32-bit integers.
; Receives: ESI = the array offset
;           ECX = number of elements in the array
; Returns:  EAX = sum of the array elements
-----
ArraySum PROC
    push    esi                ; save ESI, ECX
    push    ecx
    mov     eax,0              ; set the sum to zero

L1: add     eax,[esi]           ; add each integer to sum
    add     esi,TYPE DWORD     ; point to next integer
    loop    L1                 ; repeat for array size

    pop     ecx                ; restore ECX, ESI
    pop     esi
    ret                     ; sum is in EAX
ArraySum ENDP
```



## Procedure Parameters: Passing Register Arguments to Procedures

After the **CALL** statement, we have the option of **copying the sum in EAX** to a variable.

```
.data
theSum  DWORD  ?
.code
main PROC
    mov     eax,10000h    ; argument
    mov     ebx,20000h    ; argument
    mov     ecx,30000h    ; argument
    call    Sumof         ; EAX = (EAX + EBX + ECX)
    mov     theSum,eax    ; save the sum

; -----
; sumof
;
; Calculates and returns the sum of three 32-bit integers.
; Receives: EAX, EBX, ECX, the three integers. May be
;           signed or unsigned.
; Returns:  EAX = sum
; -----

SumOf PROC
    add     eax,ebx
    add     eax,ecx
    ret
SumOf ENDP
```

# When not to push a register

- The **sum** of the three registers is stored in **EAX** on line (3),
- **but** the **POP** instruction replaces it with the starting value of **EAX** on line (4):

```
SumOf PROC                                ; sum of three integers
    push eax                               ; 1
    add eax,ebx                             ; 2
    add eax,ecx                             ; 3
    pop eax                                ; 4
    ret
SumOf ENDP
```

causing the procedure's return value to be lost

# Outline

---

- Stack Operations
- Defining and Using Procedures
- **Linking to an External Library**
- The Irvine32 Library

# Linking to an External Library

---

- What is a Link Library?
- How the Linker Works



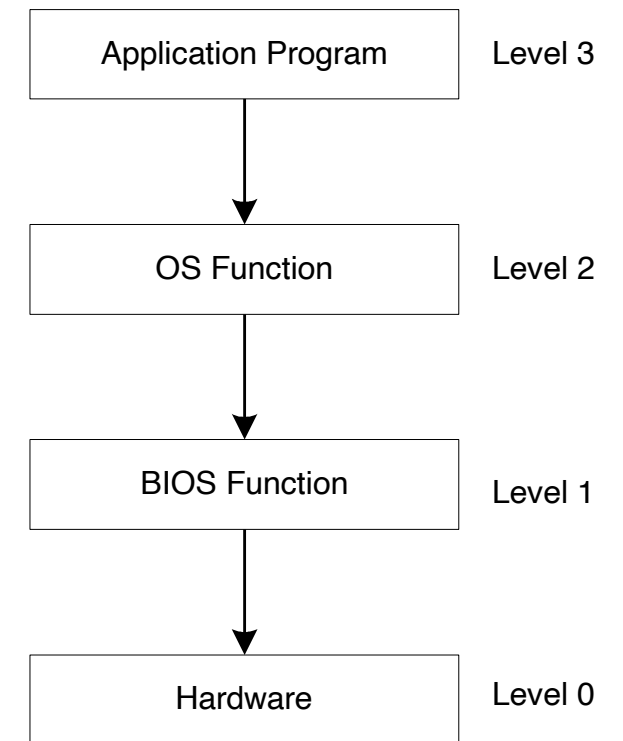
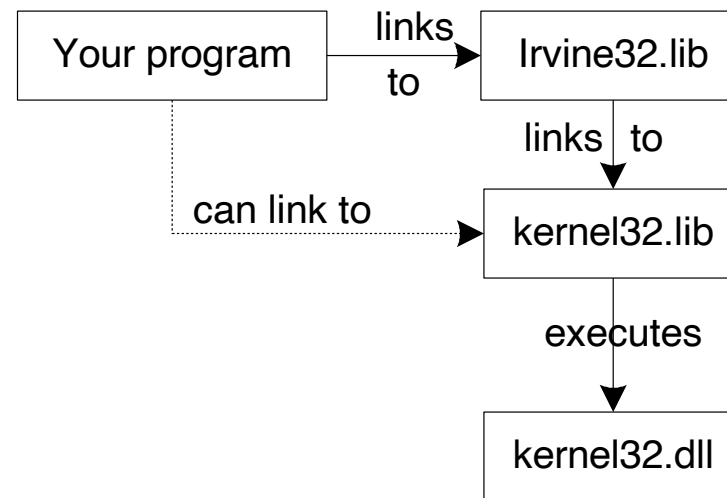
# What is a Link Library?

---

- A **file containing procedures** that have been compiled into machine code
  - constructed from one or more **OBJ files**
- To **build** a **library**, . . .
  - start with **one or more ASM source files**
  - **assemble** each into an **OBJ** file
  - **create** an empty library file (extension .LIB)
  - add the **OBJ file(s)** to the library file, using the Microsoft LIB utility

# How The Linker Works

- Your programs link to Irvine32.lib using the linker command inside a batch file named [make32.bat](#).
- Notice the two **LIB files**: [Irvine32.lib](#), and [kernel32.lib](#)
  - the latter is part of the Microsoft *Win32 Software Development Kit (SDK)*



**Chapter2**

# Outline

---

- Stack Operations
- Defining and Using Procedures
- Linking to an External Library
- **The Irvine32 Library**

# Calling Irvine32 Library Procedures

---

- Call each procedure using the **CALL instruction**
- Some procedures **require input arguments**
- The **INCLUDE directive** copies in the procedure prototypes (**declarations**)
- The following example **displays "1234" on the console**:

```
INCLUDE Irvine32.inc
```

```
.code
```

```
    mov    eax, 1234h        ; input argument  
    call   WriteHex          ; show hex number  
    call   Crlf              ; end of line
```

**WriteHex** **PROTO**

Next, a **Call** instruction executes **WriteHex**:

Call **WriteHex**

# Some Library Procedures

---

- Clrscr - **Clears console**, locates cursor at upper left corner
- Crlf - **Writes** end of line sequence to standard output
- DumpMem - **Writes** **block of memory** to standard output in hex
- DumpRegs – **Displays** **general-purpose registers** and **flags** (hex)
- ReadChar - **Reads** a **single character** from standard input
- ReadHex - **Reads** 32-bit **hexadecimal integer** from keyboard
- ReadDec - **Reads** 32-bit **unsigned decimal integer** from keyboard
- ReadInt - **Reads** 32-bit **signed decimal integer** from keyboard
- ReadString - **Reads** **string** from ***stdin***, terminated by [Enter]

# Some Library Procedures

---

- **WriteChar** - **Writes** a single character to standard output
- **WriteDec** - **Writes** unsigned 32-bit integer in decimal format
- **WriteHex** - **Writes** an unsigned 32-bit integer in hexadecimal format
- **WriteHexB** – **Writes** byte, word, or doubleword in hexadecimal format
- **WriteInt** - **Writes** signed 32-bit integer in decimal format
- **WriteString** - **Writes** null-terminated string to console window

# Some Library Procedures: Example 1

---

- **Clear** the screen, **delay** the program for 500 milliseconds, and **dump** the registers and flags

```
.code
    call Clrscr
    mov  eax, 500
    call Delay
    call DumpRegs
```

**DumpRegs** – Displays  
general-purpose registers  
and flags (hex)

Sample output:

```
EAX=00000613 EBX=00000000 ECX=000000FF EDX=00000000
ESI=00000000 EDI=00000100 EBP=0000091E ESP=000000F6
EIP=00401026 EFL=00000286 CF=0 SF=1 ZF=0 OF=0
```

# Some Library Procedures: Example 2

---

- **Display** a null-terminated string and **move** the cursor to the **beginning** of the next screen line

```
.data
str1 BYTE "Assembly language is easy!",0

.code
    mov     edx,OFFSET str1
    call WriteString
    call Crlf
```

**WriteString** - Writes  
null-terminated  
string to console  
window



# Some Library Procedures: Example 3

---

- **Display** an unsigned integer in **binary**, **decimal**, and **hexadecimal**, each on a separate line.

`IntVal = 35`

`.code`

`mov eax, IntVal`

`call WriteBin` ; display binary

`call Crlf`

`call WriteDec` ; display decimal

`call Crlf`

`call WriteHex` ; display hexadecimal

`call Crlf`

Sample output:

```
0000 0000 0000 0000 0000 0000 0010 0011
35
23
```

# Some Library Procedures: Example 4

---

- **Input a string** from the user
- **EDX** points to the string and **ECX** specifies the maximum number of characters the user is permitted to enter.

```
.data
fileName BYTE 80 DUP(0)

.code
    mov     edx, OFFSET fileName
    mov     ecx, SIZEOF fileName - 1
    call ReadString
```

**ReadString** - Reads  
string from stdin,  
terminated by [Enter]

A null byte is automatically appended to the string.

## Example 5

; Adding Three Numbers

Include Irvine32.inc

;WriteInt PROTO ; Irvine32 library

; Crlf PROTO

.386

.model flat,stdcall

.stack 4096

ExitProcess proto,dwExitCode:dword

.code

**main PROC**

mov ebx,10 ; pass three parameters, in order

mov ecx,20

mov edx,30

call AddThree ; look for return value in EAX

call WriteInt ; display the number

call Crlf ; output a CR/LF

invoke ExitProcess,0

**main ENDP**

**AddThree PROC**

Mov eax,0

add eax,ebx

add eax,ecx

add eax,edx

ret

**AddThree ENDP**

**end main**