

1. Using Bellman-Ford, find the shortest path tree from the node 3

- the shortest-path tree consists of edges 1-5, 2-4, 2-6, 3-1, 5-2

- the number of iterations of BF is 3+1 = 4

- renumber the nodes such that after renumbering BF needs only 2 iterations to find shortest path tree

3-1, 1-5, 5-2, 2-4, 2-6

old number	1	2	3	4	5	6
new number	<u>3</u>	<u>4</u>	<u>1</u>	<u>5</u>	<u>3</u>	<u>6</u>

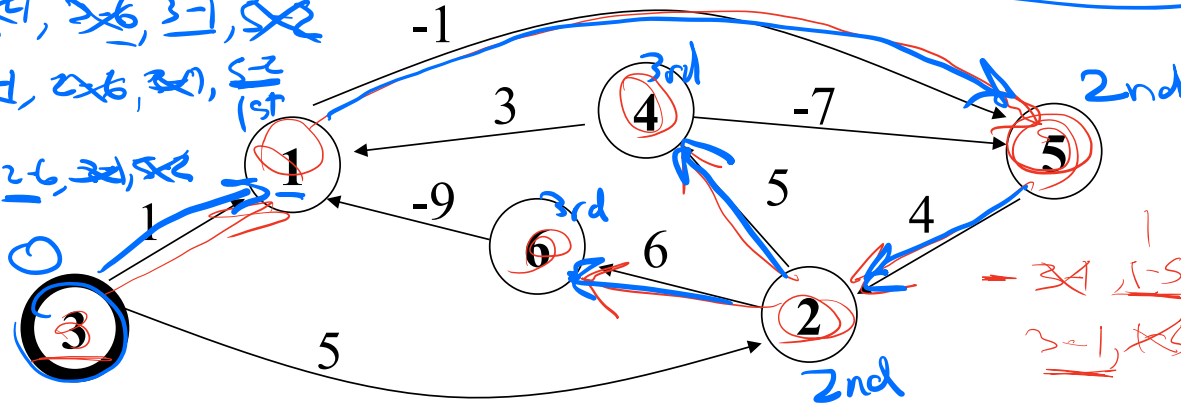
- the maximum number of iterations over all possible renumberings is 5?

#nodes - 1

~~- 1-5, 2-4, 2-6, 3-1, 5-2~~

~~- 1-5, 2-4, 2-6, 3-1, 5-2~~

~~- 1-5, 2-4, 2-6, 3-1, 5-2~~



~~- 3-1, 1-5, 5-2, 2-4, 2-6~~  
~~3-1, 1-5, 5-2, 2-4, 2-6~~  
 4 5

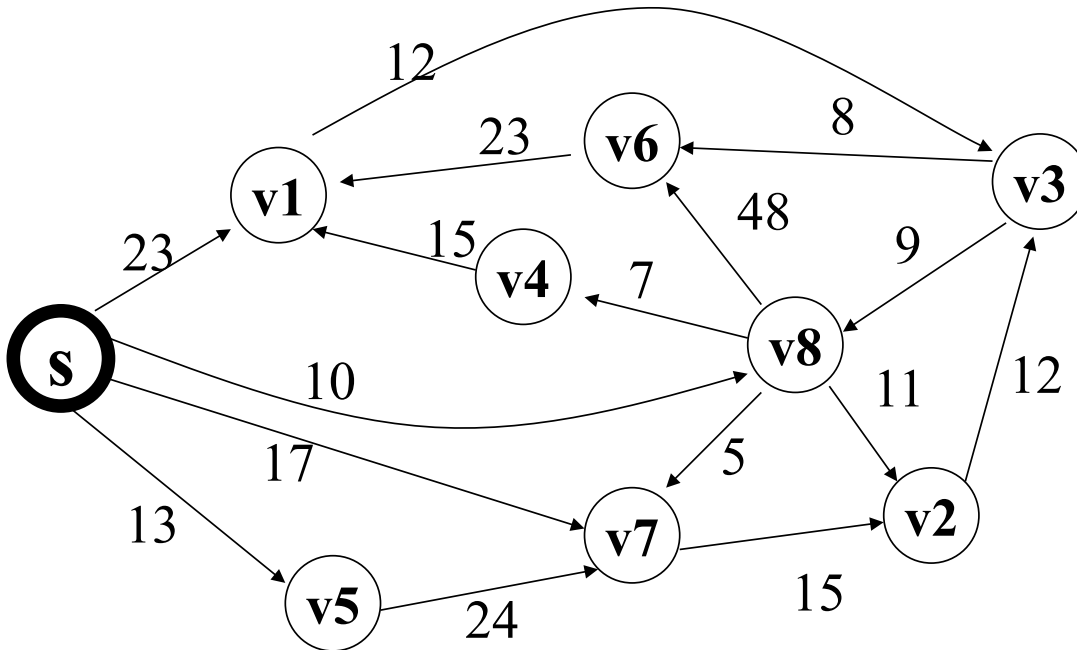
3 → 1 → 5 → 2 → 4 → 6  
 1 2 3 4 5 6

7. Write the content of the queue Q/ the set S/ keys  $d(v)$  after 5 iterations of the Dijkstra algorithm for the graph G below and source s (weights are on edges):

Q = \_\_\_\_\_ S = \_\_\_\_\_

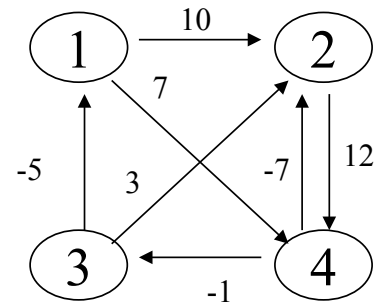
$d(s)$  = \_\_\_\_\_  $d(v1)$  = \_\_\_\_\_  $d(v2)$  = \_\_\_\_\_  $d(v3)$  = \_\_\_\_\_

$d(v4)$  = \_\_\_\_\_  $d(v5)$  = \_\_\_\_\_  $d(v6)$  = \_\_\_\_\_  $d(v7)$  = \_\_\_\_\_  $d(v8)$  = \_\_\_\_\_.



3. Find all shortest path weights with the matrix multiplication method for the graph on the right side.

- give all matrices that are obtained on the way,
- are there any negative cycles in the graph?



$M =$

$M^2 =$

$M^4 =$

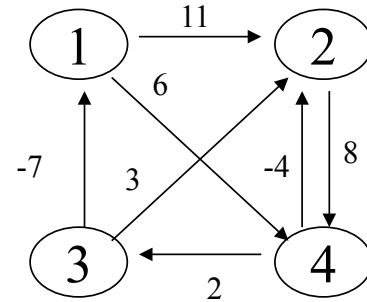
$M^8 =$

$M^{16} =$

$M^{32} =$

4. Find all shortest path weights with the Floyd-Warshall method for the graph on the right side.

- give all matrices that are obtained on the way



$D^0 =$


$D^1 =$


$D^2 =$

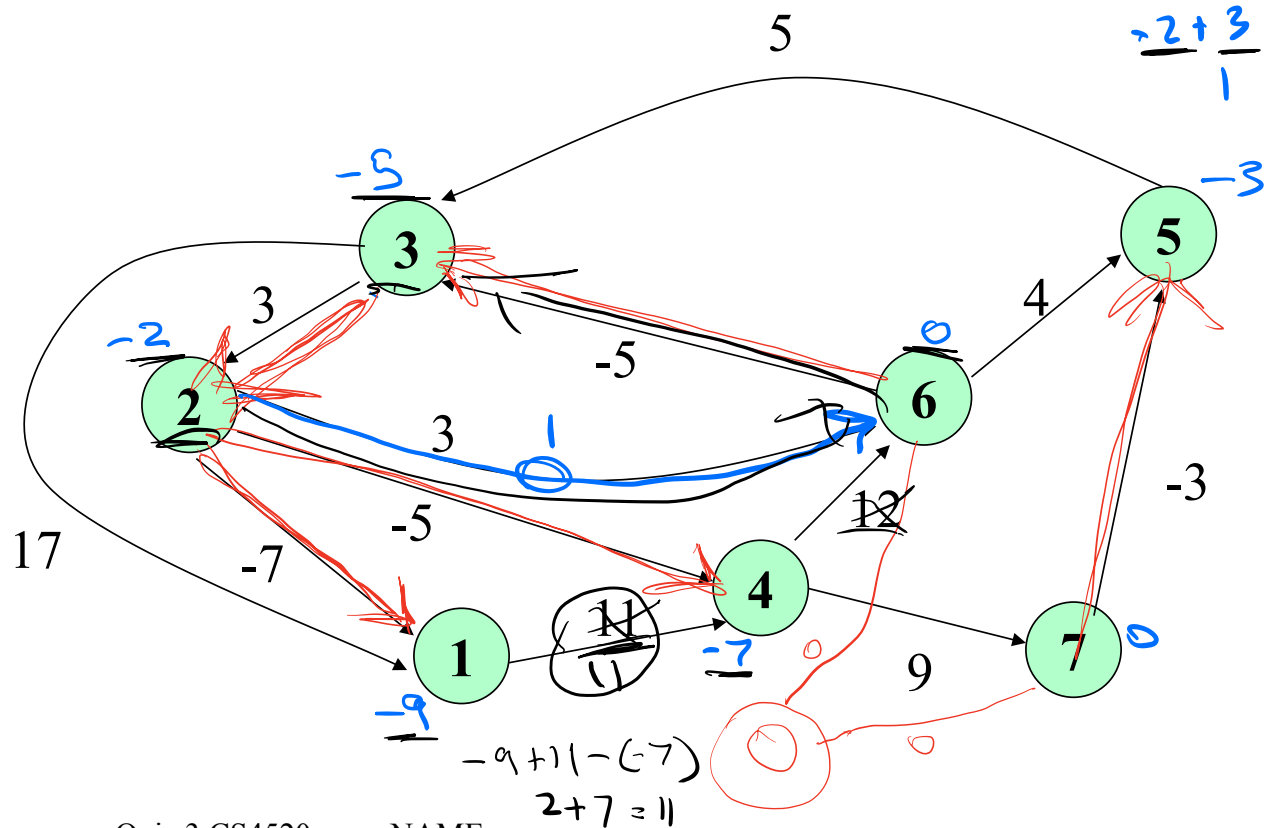

$D^3 =$


$D^4 =$

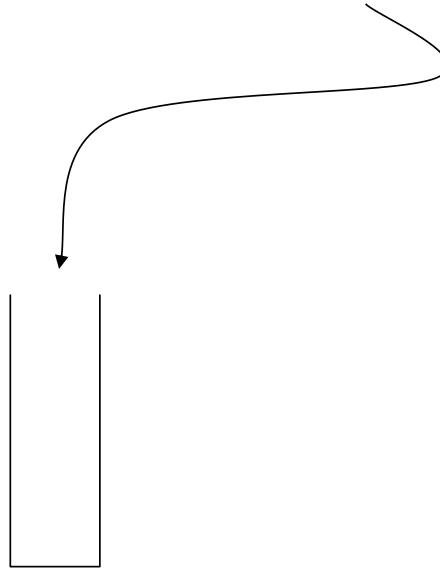

5. Johnson's algorithm is applied to the graph below. (20pts)

Give the **modified** weight of the edge (2,6) 1

The shortest path from 2 to 3 is 2-6, 6-3



1. Give the final content of the stack in Graham's algorithm for convex hull for these 3 points A, B and C (check the order!):  
A=(190,291), B=(300,400) and C=(200,300)

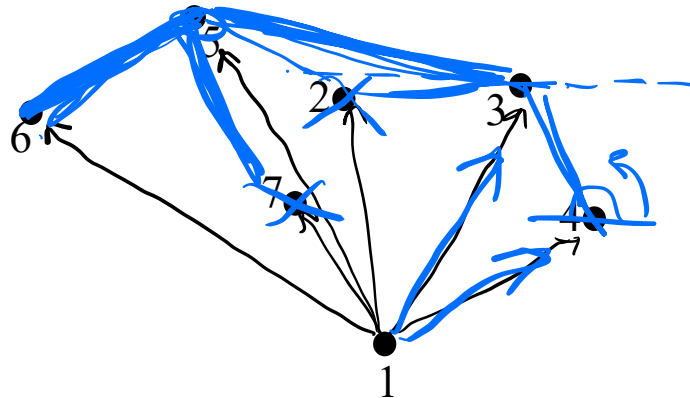


2. For Graham's scan finding convex hull of the point set given below:

- Give the sorted sequence of points for Graham scan 1, 4, 3, 2, 5, 7, 6
- Show the content of the stack after each change

empty	<del>1</del>	4 1	3 4 1	2 3 4 1	3 4 1	5 3 4 1	7 5 3 4 1	5 3 4 1	6 5 3 4 1	
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- Give the convex hull of this point set 1, 4, 3, 5, 6



3. Below given a point set in the **rectilinear** metric (the height/width of any cell=1) where the closest pair of points should be found using divide and conquer. Show

1) the first partition of the point set (draw a line)

2) the closest pair in the left part (connect solid),  $\delta_{\text{left}} = \underline{1}$ , 2-3

and the right part (connect solid),  $\delta_{\text{right}} = \underline{2}$  8-9

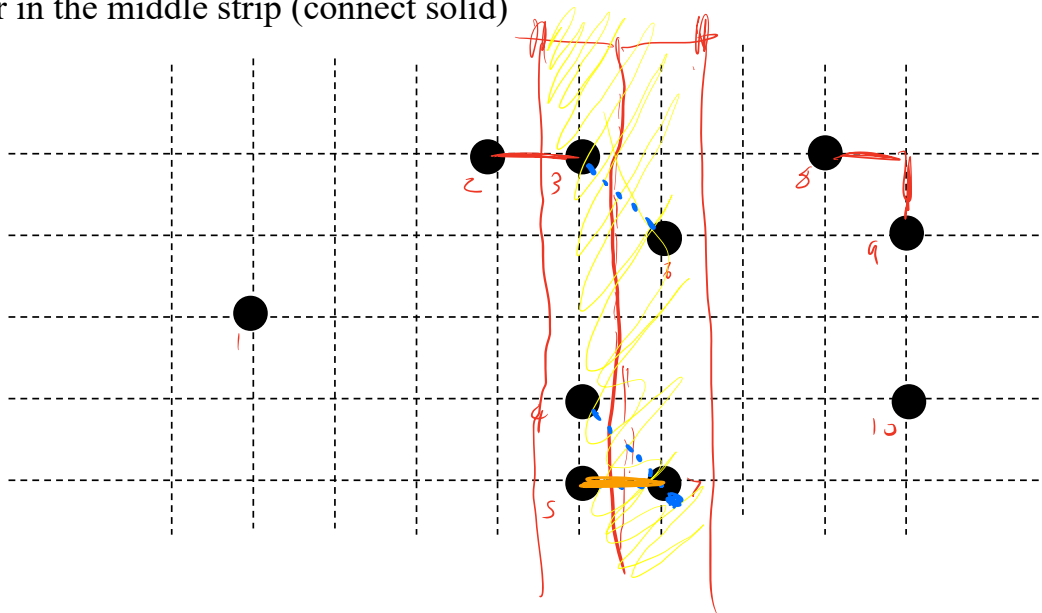
*Smaller of two is Delta*

3) the middle strip (shade) *- go out delta from first partition*

- pairs in the middle strip for which distances should be computed (connect dashed)

- closest pair in the middle strip (connect solid)

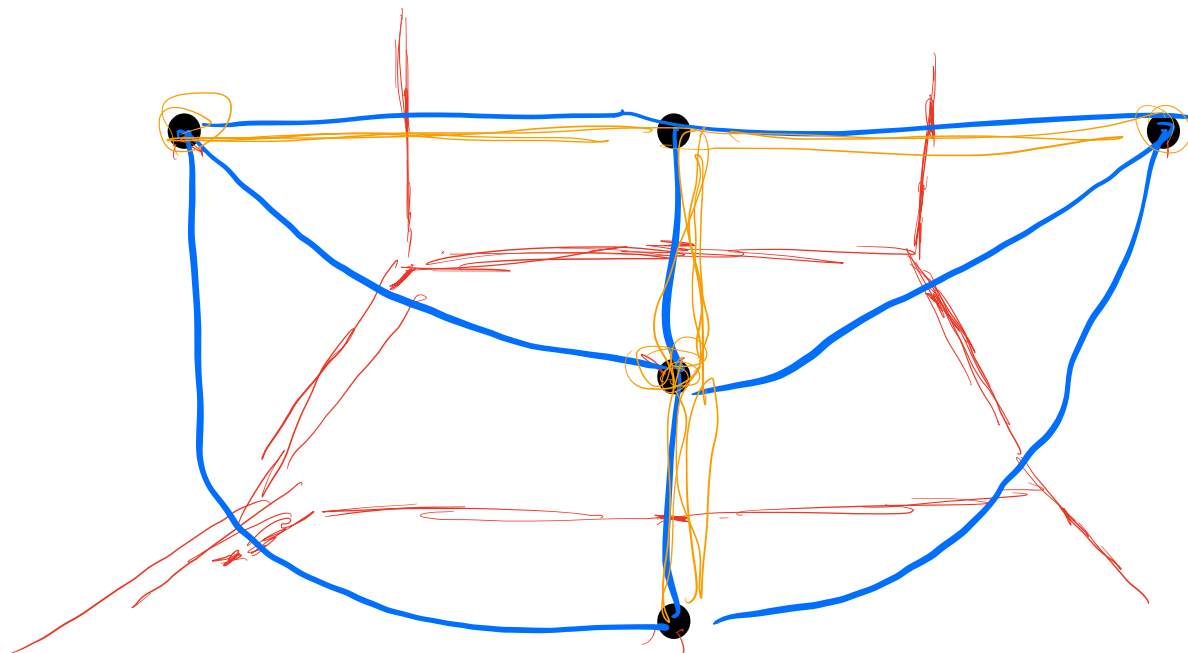
*least  
vert*





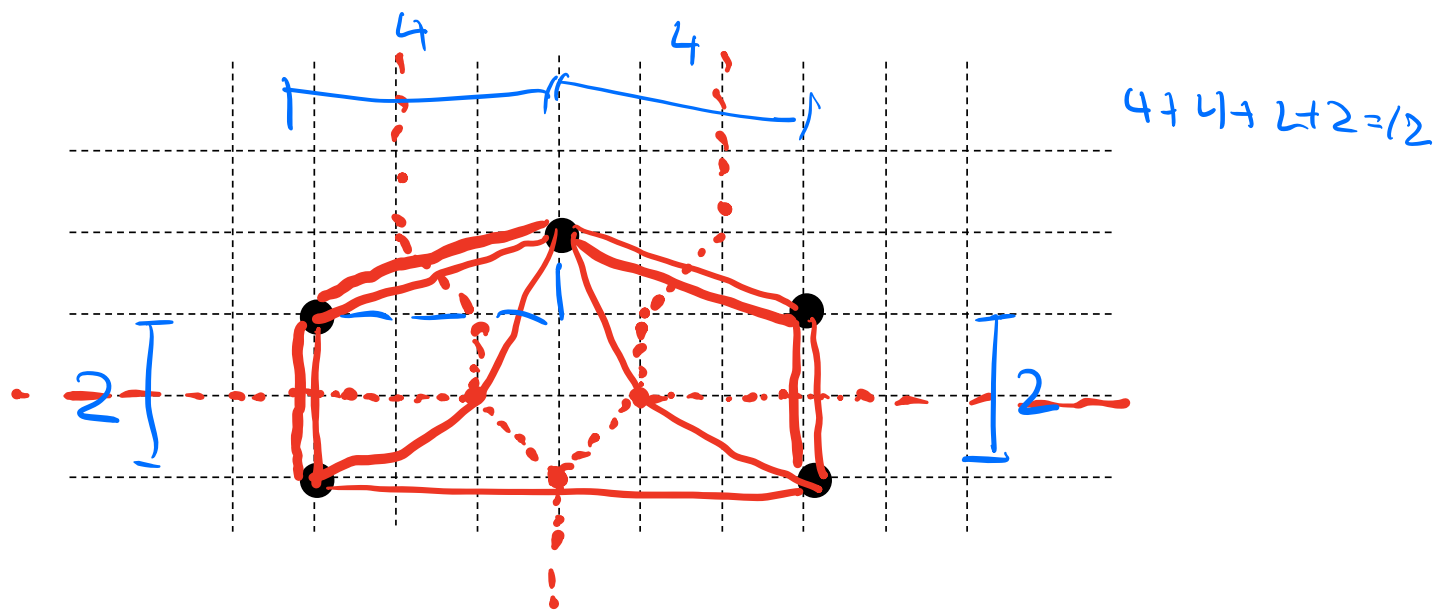
4. Below given a point set in the **Euclidean** metric. Draw

- Voronoi regions (dashed edges) — barriers between points
- Voronoi graph / Delanau triangulation (solid edges) — connections with one barrier
- minimum spanning tree (double edges) — closest path



5. Below given a point set in the **rectilinear** metric (the height/width of any cell=1) . Draw

- Voronoi regions(dashed lines) - barriers
- Voronoi graph / Delanau triangulation (solid edges) - on barrier connection
- minimum spanning tree (double edges) - shortest path
- the length of the minimum spanning tree is 12



6. Prove, that the k-tree problem (finding minimum weight tree subgraph with k vertices) is in class NP

- a) Optimization formulation
- b) Decision formulation
- c) Polynomial-size certificate
- d) Polynomial time verification algorithm

## 7. For the 3-CNF

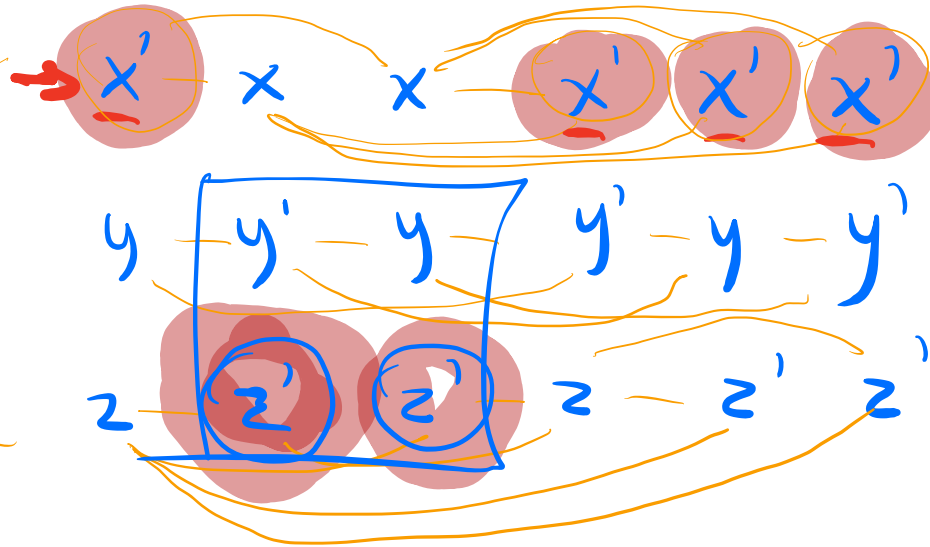
$$f = (x' + y + z) \& (x + y' + z') \& \underline{(x + y + z')} \& (x' + y' + z) \& (x' + y + z') \& (x' + y' + z')$$

- give 0-1 assignment to variables such that  $f=1$   $x'=1, z'=1$

1. - give 0-1 assignment to variables such that  $f=0$   $x=y=z'=0$

- Draw the corresponding graph and mark the maximum independent set

1 2 3 4 5 6



Prime  $\rightarrow$  Not prime

8. In the following graph find

- Maximum Independent Set 2, 4, 6, 9, b

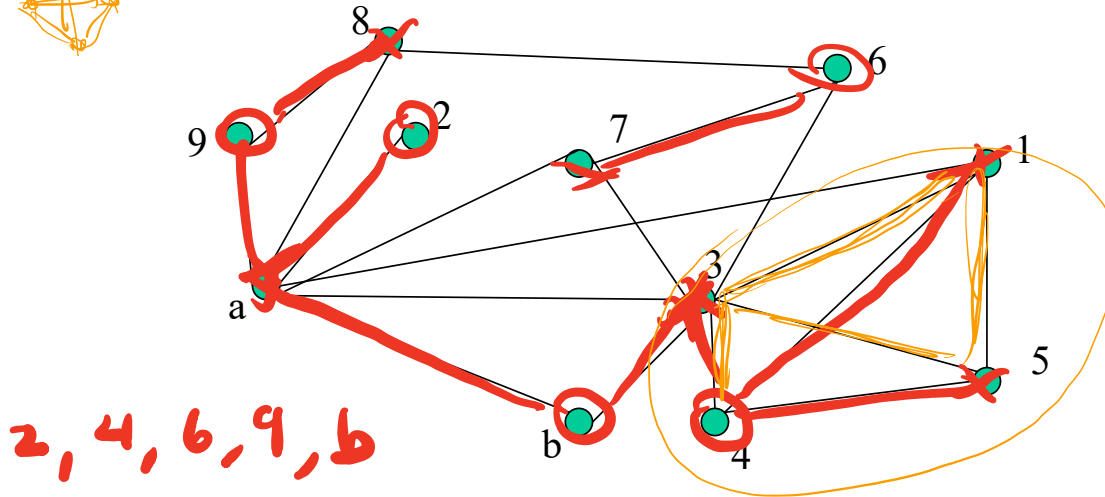
- Minimum Vertex Cover 1, 3, 5, 7, 8, a

- Maximum Clique 1, 3, 4, 5

5 coms



1, 3, 4, 5

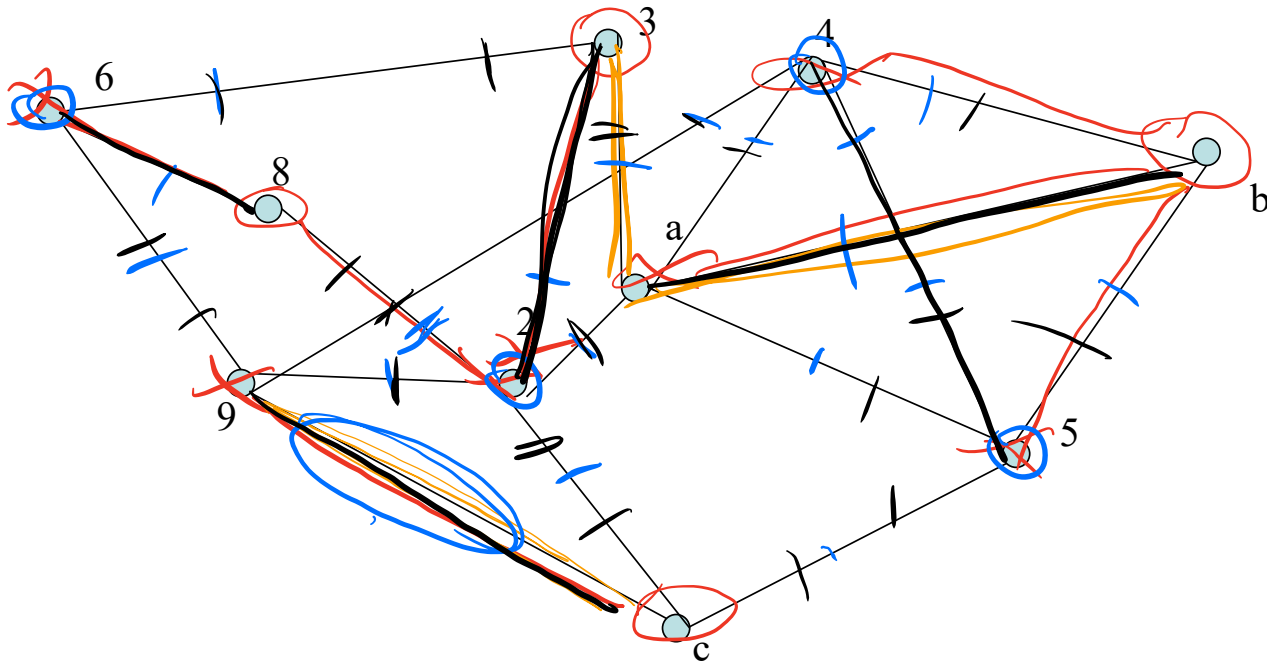


1. For minimum vertex cover problem in the following graph give

2. greedy solution = nodes 2, 4, 5, 6, a, 9

3. 2-VC solution = nodes 2-3, 4-5, 6-8, 9-c, a-b

1. Optimal solution = nodes 2, 3, 4, 5, 6, 9, a



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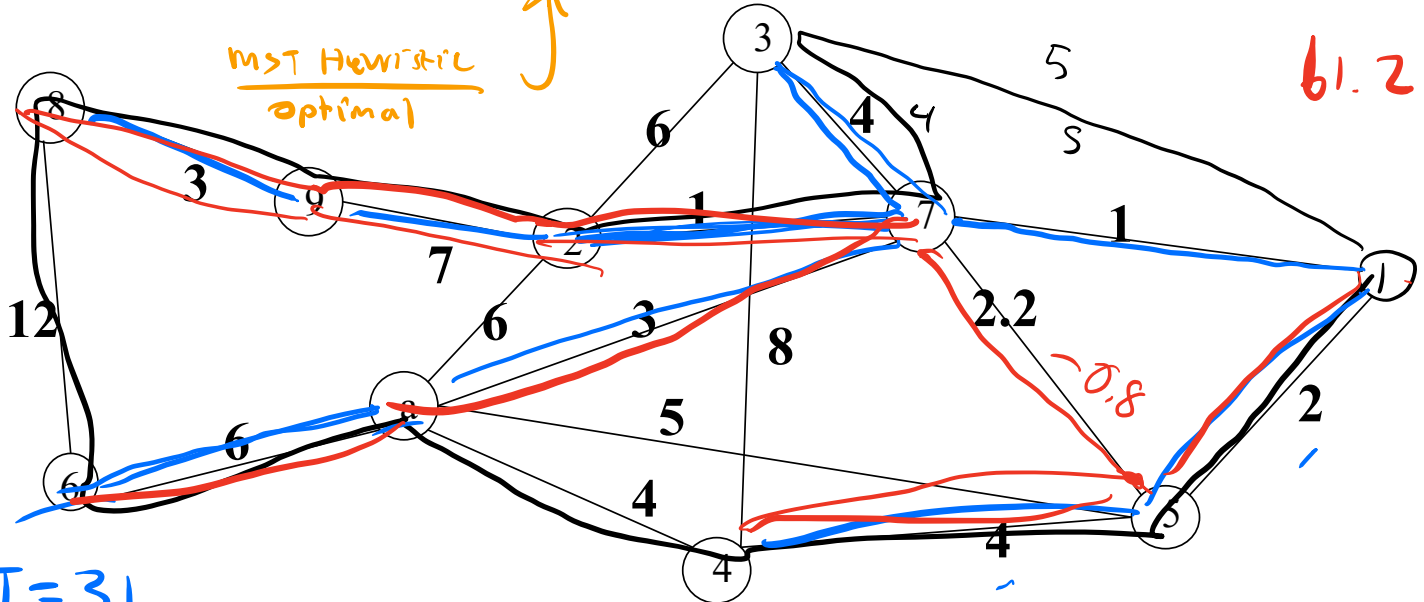
2. For the following graph, find

1. Optimal TSP tour 1, 3, 7, 2, 9, 8, 6, a, 4, 5, 1 length = 48

2. Double MST tour 1, 5, 4, 5, 1, 7, 2, 9, 8, 9, 2, 7, 3, 7, a, 6, a, 7, 1 length = 62

3. MST-heuristic tour (with shortcuts) 1, 5, 4, 7, 2, 9, 8, 3, a, 6, 1 length = 61.2

4. The error of the MST-heuristic is 27.5%  $(\frac{61.2}{48} - 1) \times 100$  62 - 0.8  
61.2



mst=31

3x2=62

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3. For the following graph, find

1. Christofides heuristic matching (3,4) (6,8)

length = 20

MST + matching tour 1, 5, 4, 3, 7, 2, 9, 8, 6, a, 7, 1

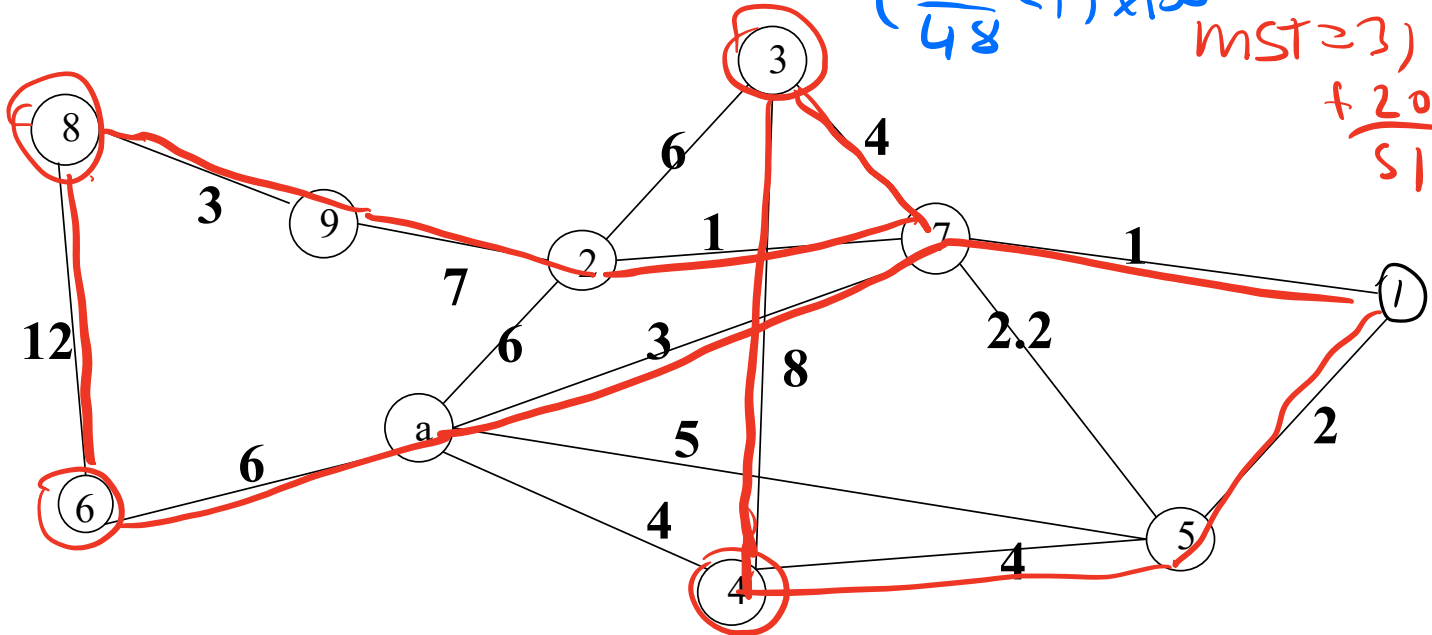
length = 51

Christofides heuristic tour (w/shorts) 1, 5, 4, 3, 7, 2, 9, 8, 6, a, 1 length = 51

The error of the Christofides heuristic is 6.25 %

$$\left( \frac{51}{48} - 1 \right) \times 100$$

MST = 31  
 $+ \frac{20}{51}$

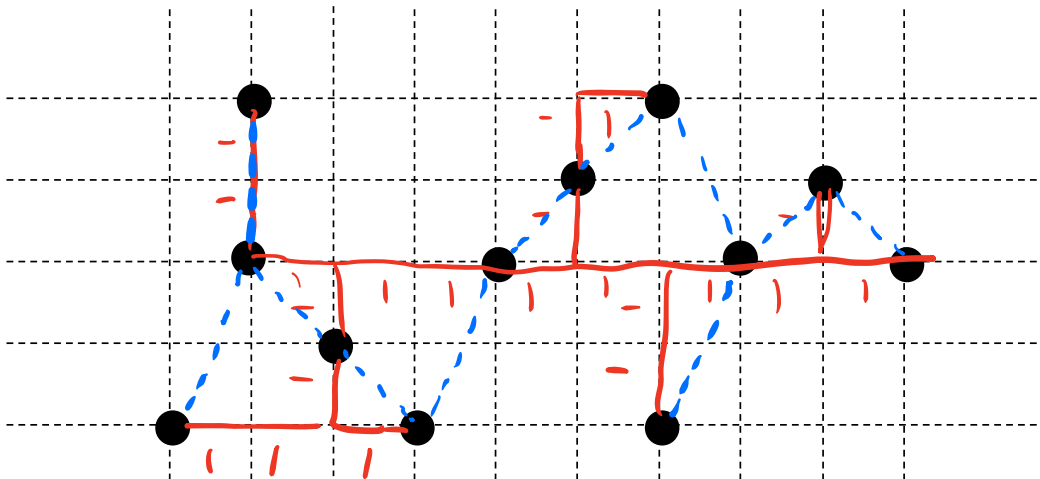




4 . In the rectilinear metric for points given below find

- The minimum Steiner tree (bold), its length is  $9+12 = 21$
- The minimum spanning tree (dashed), its length is  $14+12 = 26$
- the approximation ratio of the MST heuristic in this case is  $\frac{26}{21}$
- Approximation error in % 24%

←  
span  
Stein



8. Given the following stable marriage instance:

B1's preferences: 1st choice=G1, 2nd=G3, next – G4, last – G2

B2: G2,G3,G4,G1

B3: G4,G2,G1,G3

B4: G4,G1,G3,G2

G1: B2,B1,B3,B4

G2: B4,B3,B1,B2

G3: B3,B2,B1,B4

G4: B1,B2,B3,B4

Boys' best stable marriage is \_\_\_\_\_

Girls' best stable marriage is \_\_\_\_\_