Operating Systems 2021/2022

T Class 02 – Introduction to Operating Systems

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operating system
noun
the collection of software that directs a computer's operations,
controlling and scheduling the execution of other programs, and
managing storage, input/output, and communication resources.

Abbreviation:  OS

                                          Source: Dictionary.com
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Updated on
20 February 2022

Disclaimer

- This slides and notes are based on the companion material [Silberschatz13]. The original material can be found at:
  - http://codex.cs.yale.edu/avi/os-book/OS9/slide-dir/
- In some cases, material from [Stallings15] may also be used. The original material can be found at:
  - http://williamstallings.com/OS/OS5e.html
  - http://williamstallings.com/OperatingSystems/
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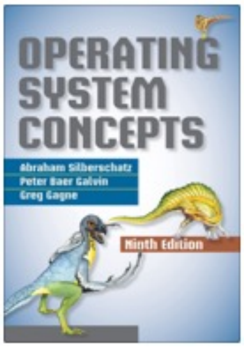
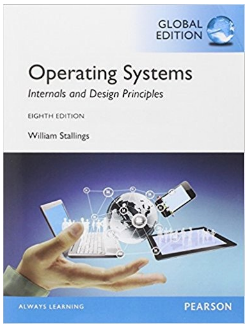


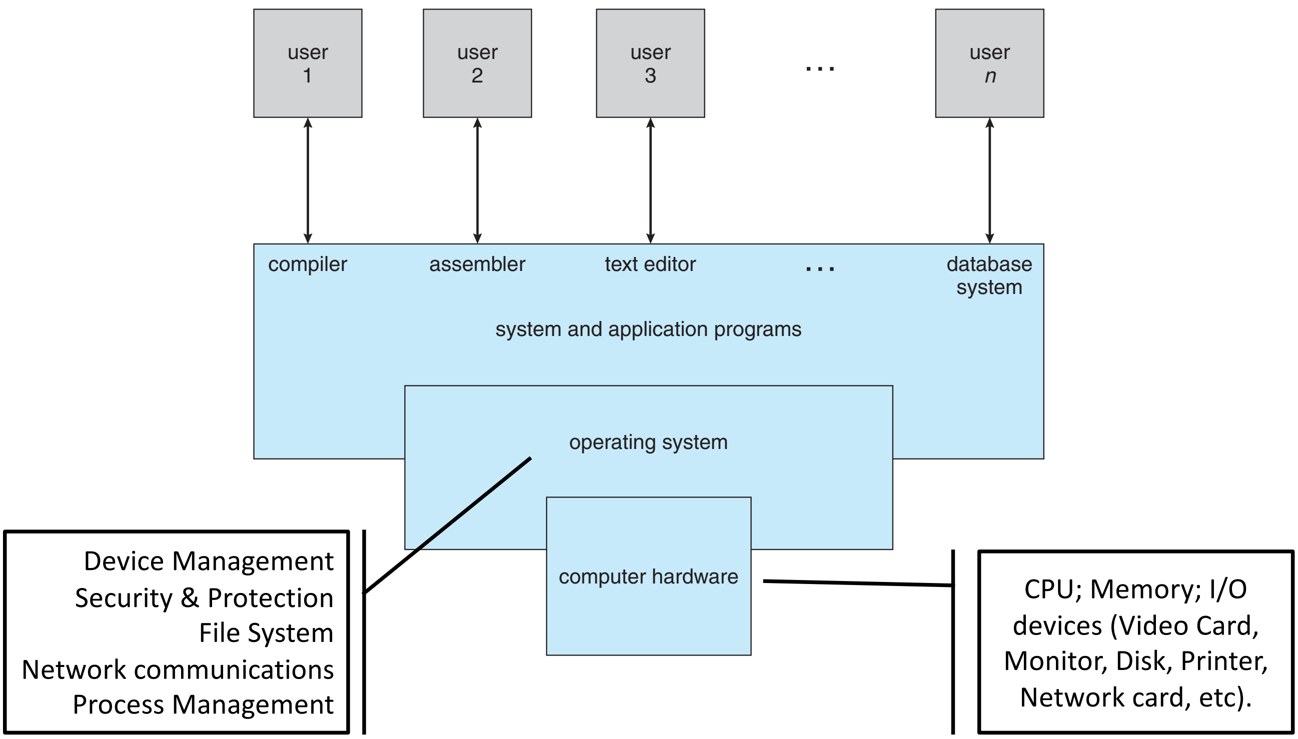


Note: Some slides are also based on previous versions from Bruno Cabral, Paulo Marques and Luis Silva (Operating Systems classes of DEI-FCTUC).

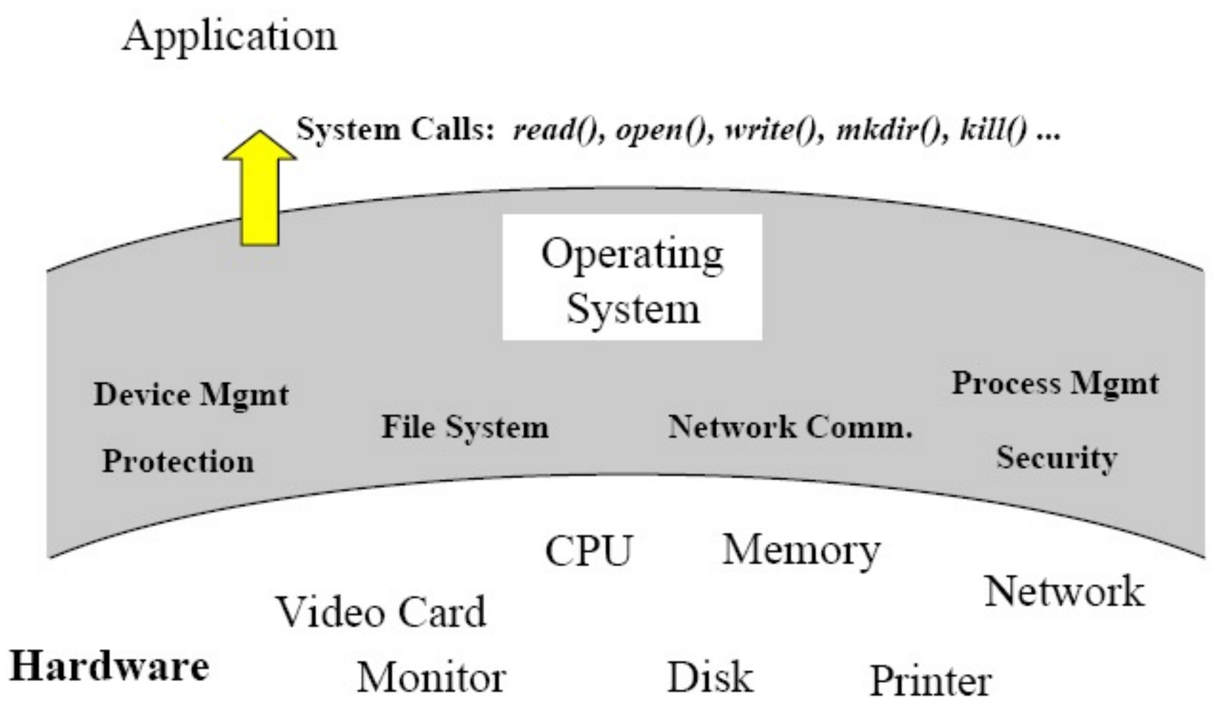
Introduction to Operating Systems
Outlines

- Describe the basic organization of a computer system
- Overview of OS components
- OS boot
- Overview of different types of computing environments
- OS services to users, processes and other systems

Operating system in a computer system



What’s an Operating System?

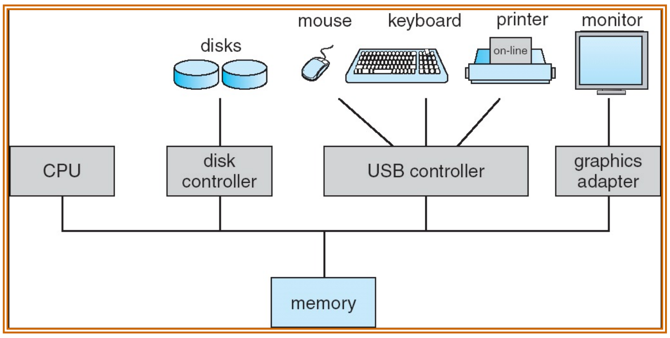


What’s an Operating System?

- OS is a control program
  - Controls execution of programs to prevent errors and improper use of the computer
- OS is a resource manager
  - Manages all resources
    - Process management
    - Memory management
    - File-system management
    - Mass-storage management
    - I/O management
  - Decides between conflicting requests for efficient and fair resource use

What Operating Systems Do

- Sharing resources among applications
  - scheduling
  - allocation
- Making efficient use of limited resources
  - improving utilization
  - minimizing overhead
  - improving throughput/goodput
- Protecting applications from each other
  - enforcement of boundaries



Classification of Operating Systems

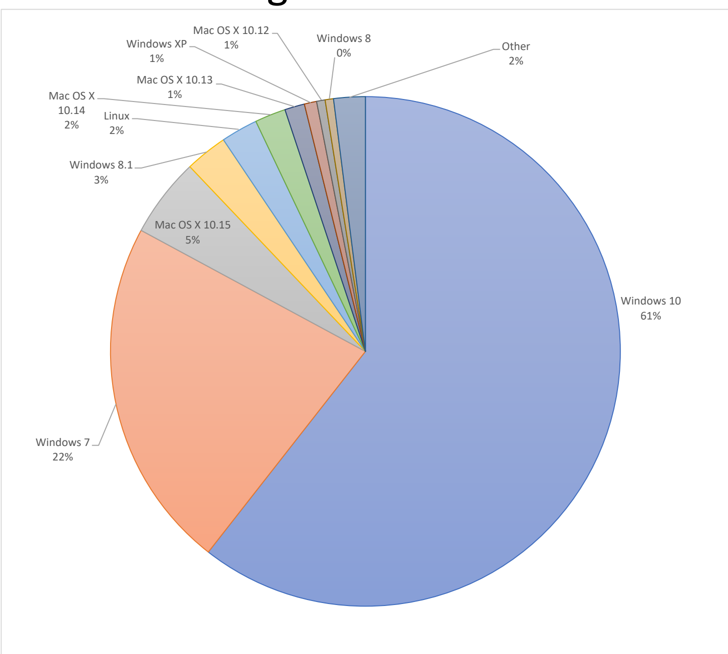
- Interface
  - CLI (Command-line interface)
  - Batch interface (commands are entered in files)
  - GUI (Graphical User Interface)
- Users
  - Multi-user
  - Mono-user
- Tasks
  - Single-task - can support only a single task at any time; tasks are handled sequentially (e.g. MS-DOS)
  - Multi-task - executes more than one task at a time, by sharing resources
- Real-time Operating Systems (RTOS)
  - Operating systems that have strict time bounds and have a predictable behavior (e.g. VxWorks, VRTX, pSOS, LynxOS, …)

Classification of Operating Systems (cont.)

- Open/Closed Operating Systems:
  - Proprietary
  - Open systems
    - Have standardized programming interfaces and peripheral interconnects
  - Open-source systems (GPL license)
    - Operating systems made available in source-code format rather than just binary closed-source
    - Started by Free Software Foundation (FSF), which has “copyleft” GNU Public License (GPL)
      - (GNU = GNU's Not Unix!)
    - Examples include GNU/Linux and BSD UNIX (included in the core of Mac OS X), and many more

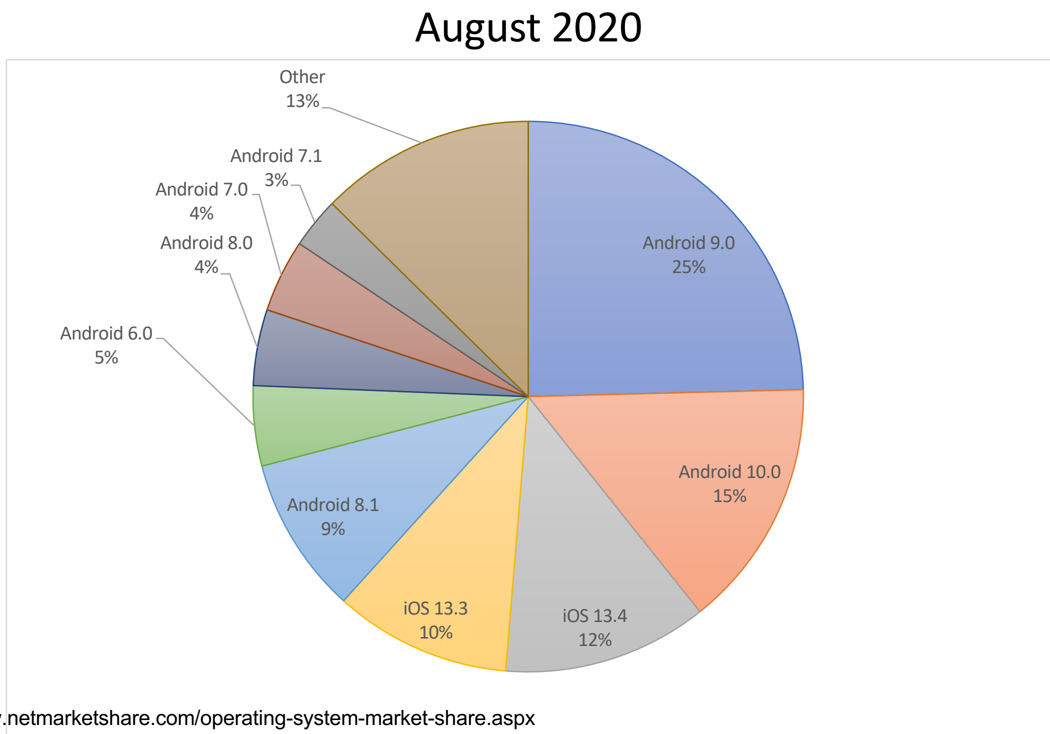
Desktop Operating System Market Share

August 2020



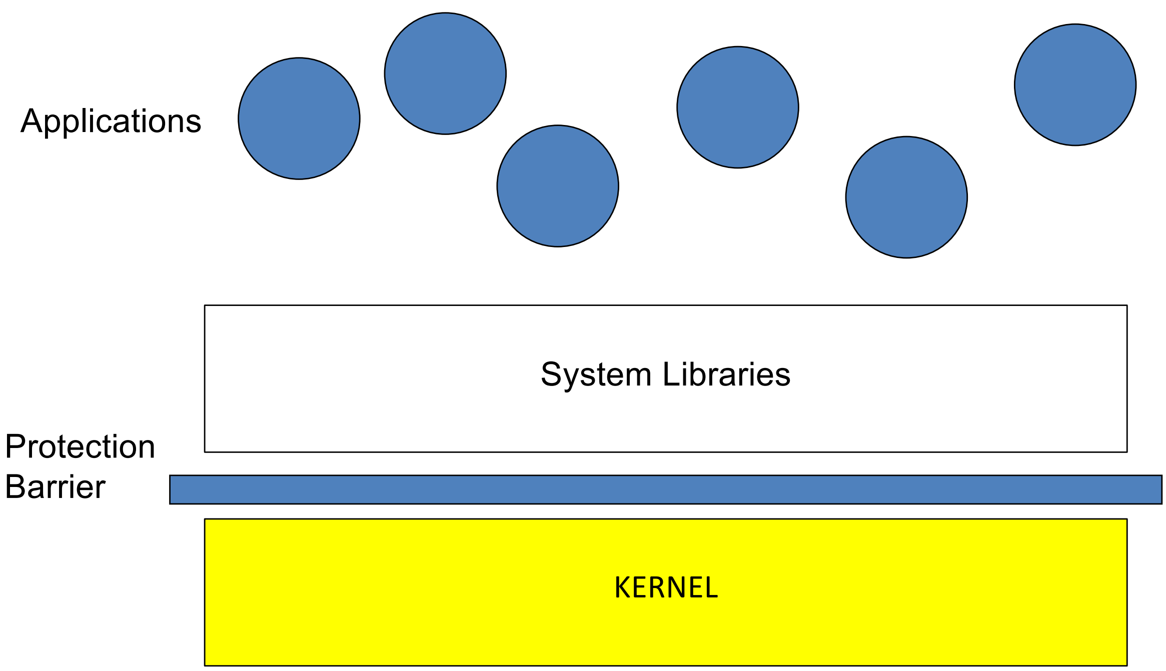
Source:
http://www.netmarketshare.com/operating-system-market-share.aspx

Mobile/tablet Operating System Market Share



Source:
http://www.netmarketshare.com/operating-system-market-share.aspx

Simple Organization of an Operating System



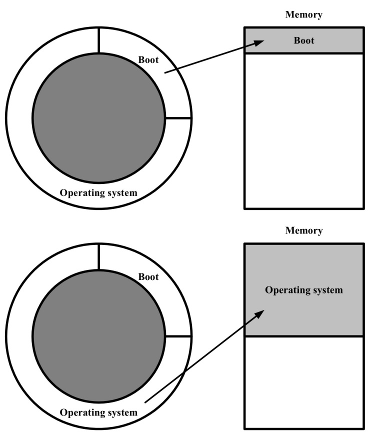
OS: Startup

- How does the hw know where the kernel is and how to load it?
- Starting a computer and loading the kernel is called booting.
- Most computers have a multistage boot

- Power up the computer (IBM - compatible PC)
  - Test operation (power-up diagnostics) :
    - In the power-up (and at reset) the CPU executes a jump instruction to a pre-defined address where an initial boot loader located in nonvolatile firmware (BIOS or UEFI) is run;
      - BIOS = Basic Input Output System
      - UEFI = Unified Extensible Firmware Interface
    - The initial boot loader program executes tests for verifying the correct operation of the CPU (subset of instructions) and RAM (basic tests)
      - Power-on self-test, also known as POST
    - If OK , a second boot loader is started ...

OS: Startup (2)

- Sequence startup (boot) OS :
  - A second boot loader is started, normally from the first sector of the boot device. This second boot loader may be able to start the OS, but normally it only knows the location of the remainder of the **bootstrap loader** program, which in turn locates the kernel, loads it into main memory, and starts its execution. .
    - Common bootstrap loader for Linux, GRUB, allows selection of kernel from multiple disks, versions, kernel options



BIOS vs UEFI

- UEFI (Unified Extensible Firmware Interface)
  - Appeared in 2007
  - Supports all the functionalities of the BIOS
  - Can be saved in any support (disk, flash memory, …)
  - Faster than BIOS
  - Allows for bigger disks
  - Better graphic interface
  - Allows Secure Boot

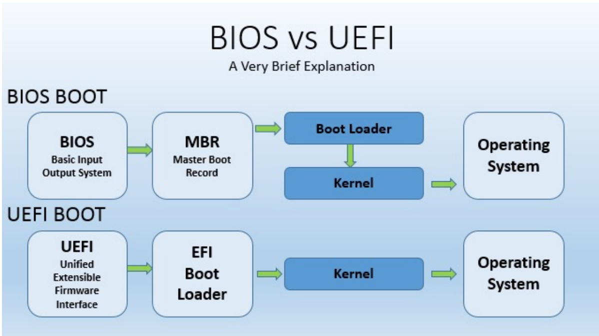


Image from:
https://passmoz.com/

OS: Initialization

- OS performs its boot “routines”:
  - hardware - Initialization of registers of the controllers , the system interrupts, the interrupt vector ;
  - software - Initialization of data structures that represent the various system resources and algorithms that support the management of these resources .
- Creation of the first cases :
  - The system auxiliary processes
    - login , file system , network ...
  - the interpreter commands
    - command line (sh , bash, csh, command.com) or windows
- "Wait for Interrupt "

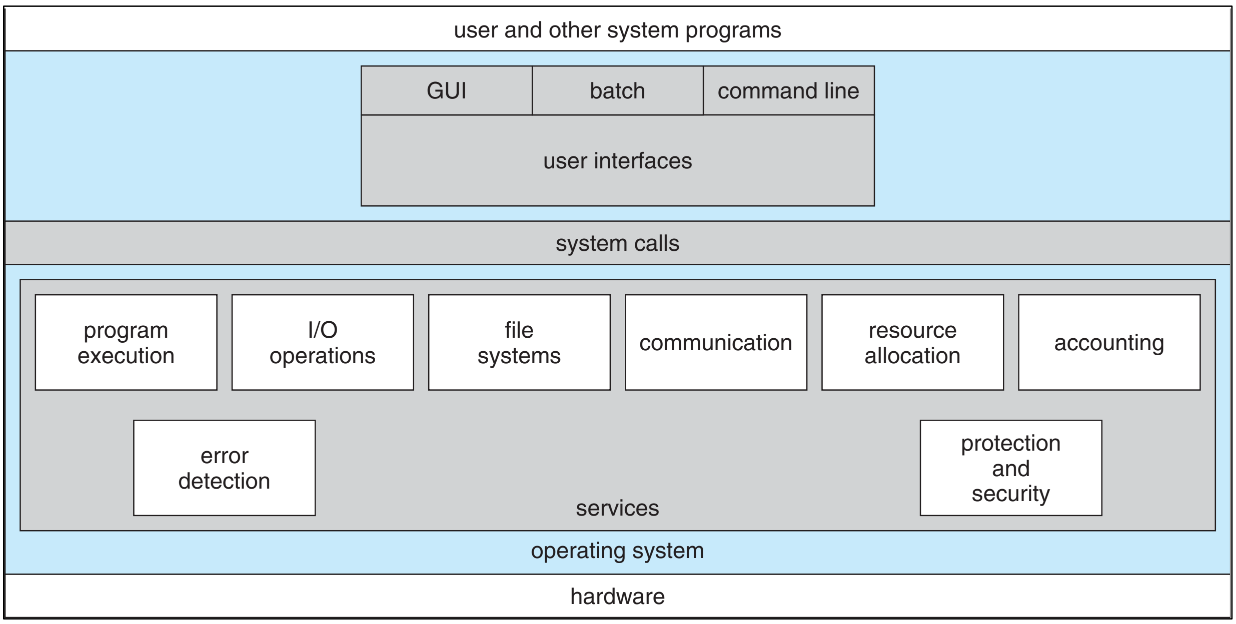
Wait for interrupt

- After booting, the OS is “quiet" waiting for work and ... there is a hardware interrupt
  - OS performs a task that handles the interruption .
    - Ex : CPU clock interrupts , OS updates the time
- The OS is " quiet" and ... a program requests a service - through a software interrupt or trap:
  - OS triggers a task that runs the service (this interruption)
    - Ex : Print a string on the screen (in MS -DOS)

```
mov dx , <address of string>
mov ah, 09H
int 21H
```

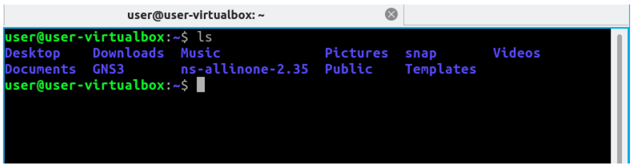

Operating System Provided Services

- The specific services provided to depend on the specific OS...



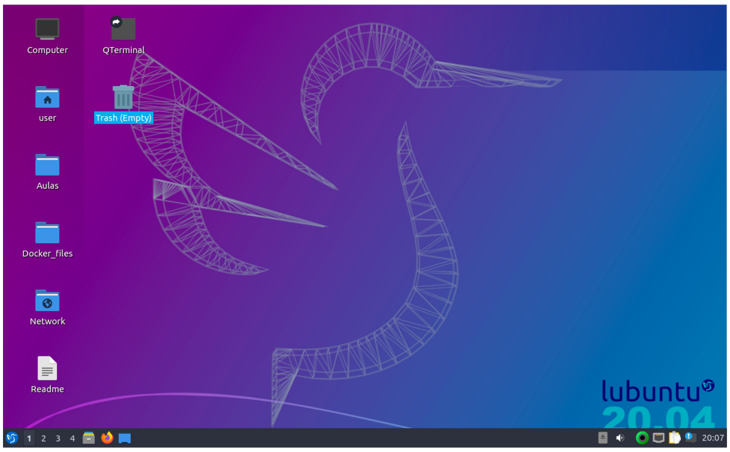
User and OS interface

- Command interpreters
  - Shells



Qterminal,
Lubuntu 20.04

- Graphical User Interface (GUI)

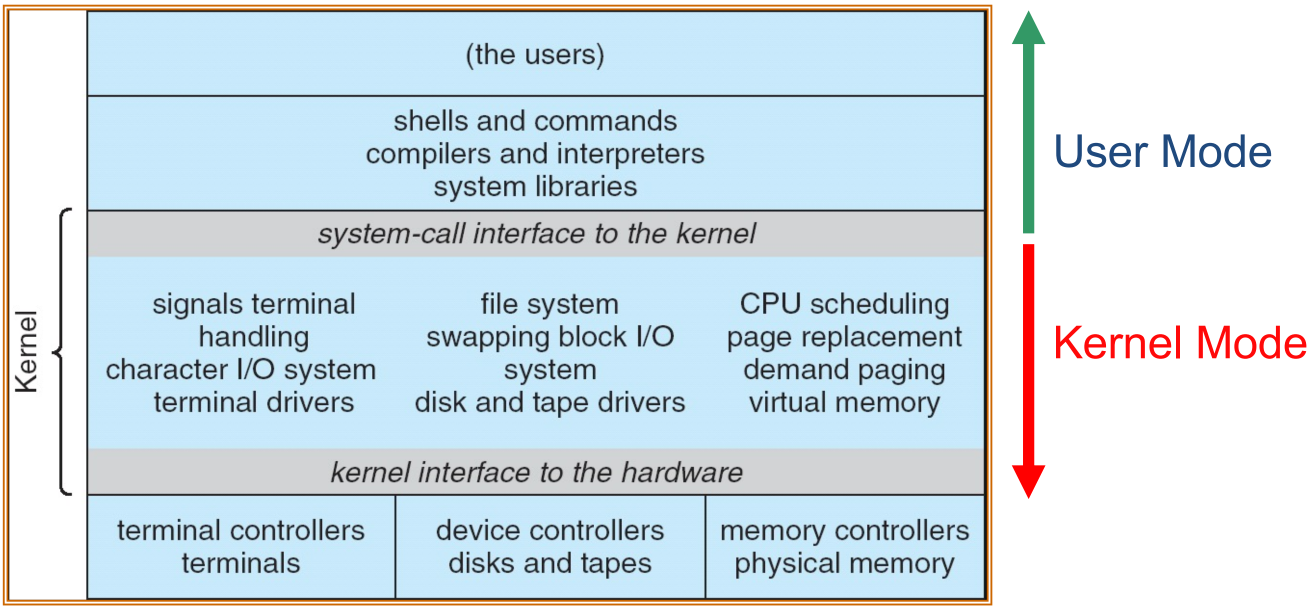


LUbuntu 20.04

Protection provided by the OS

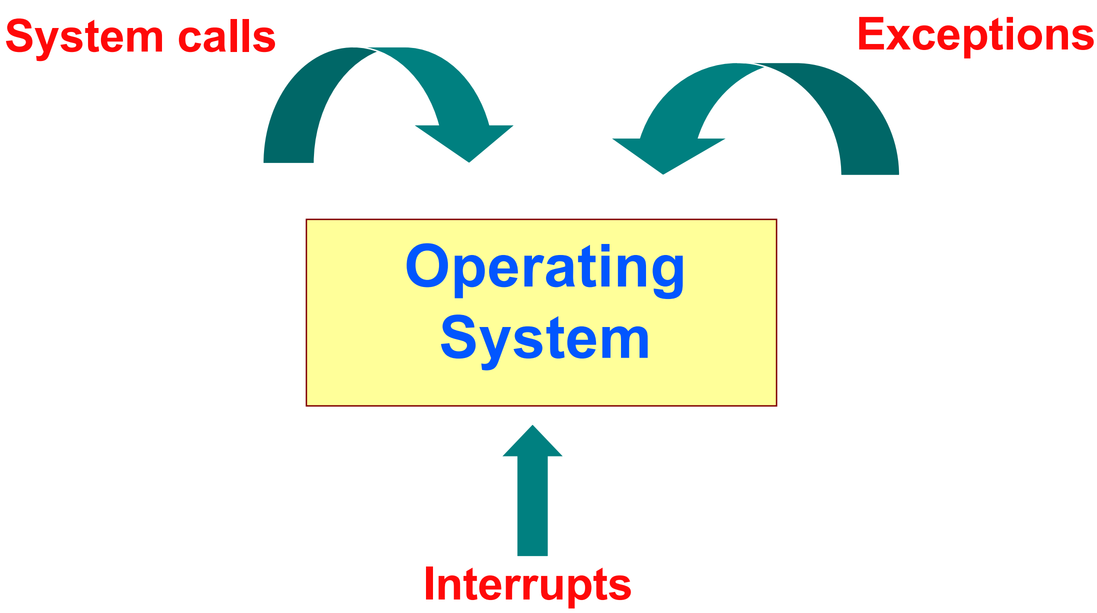
- Dual Mode of Operation
  - “User Mode” and “Kernel Mode”
  - In “User Mode” only common instructions can be performed (e.g. arithmetic and logic). In “Kernel Mode” anything can be done.
  - A mode bit contains the current mode of operation
  - RTI: Return from interruption
  - Kernel mode assumed on traps and on interrupts
- I/O Protection
  - All I/O Instructions are privileged.
    For doing I/O a trap into the operating system must be generated.
- Memory Protection
  - Each program can only access its memory
- CPU Protection
  - The CPU must remain in control of the Operating System, even if the applications does not want to let go.

Dual Mode of Operation (at least) - to assure security

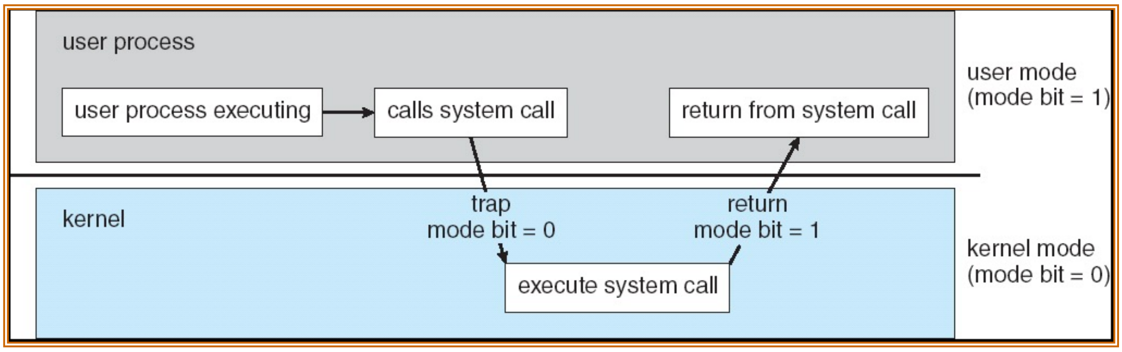


UNIX Operating System Structure

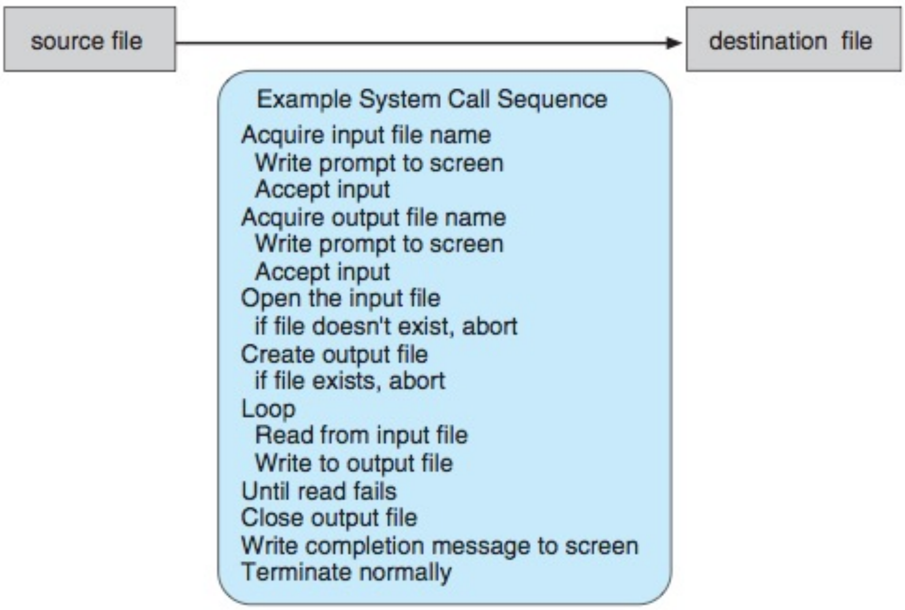
System Calls, Exceptions and Interrupts



System Calls
Dual Mode of Operation

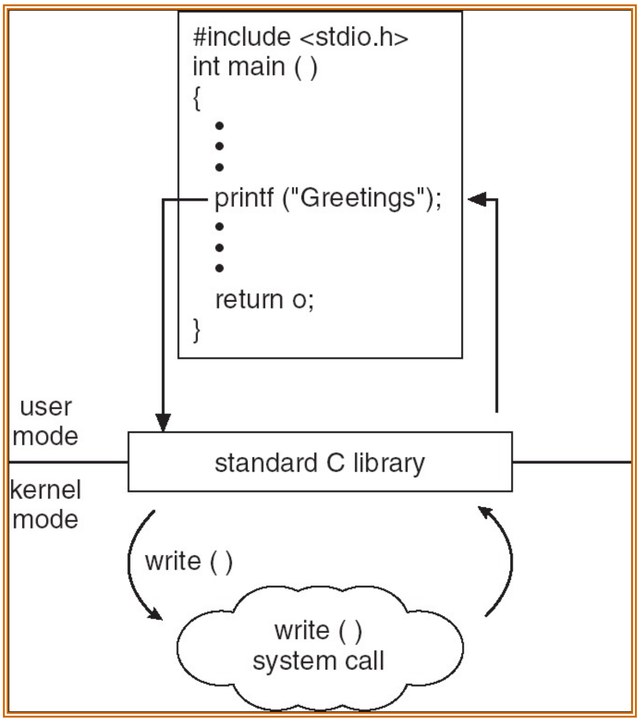


System Calls
A file copy operation

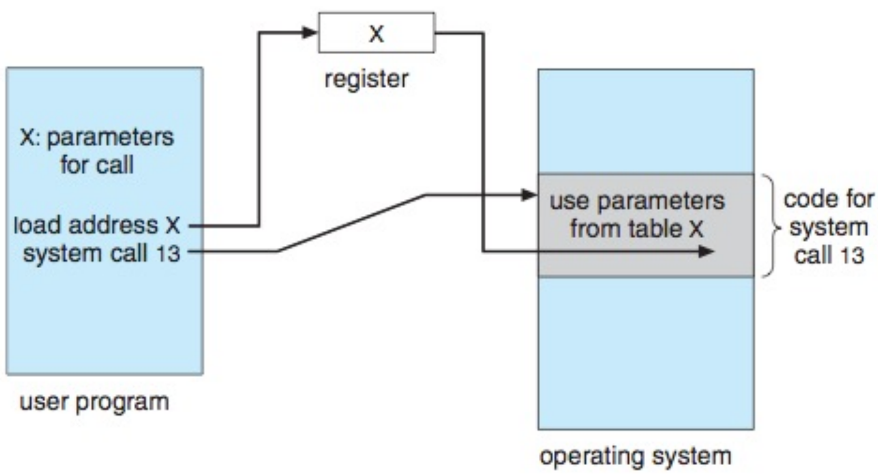


Even a simple program creates multiple system calls!

System Calls
Structure



System Calls
Passing of parameters as a table



System Calls

Example with a "Hello World" Application!

```
#include <stdio.h>

int main() {
  printf("Hello World\n");
  return 0;
}
```

- The "strace" program is very useful:
  - It allows you to see what system calls are made and with what parameters!

System Calls

"strace" command

```
$strace ./hello
```

```
[vasco@student2]$ ./hello
Hello world
[vasco@student2]$ strace ./hello
execve("./hello", ["./hello"], [/* 28 vars */]) = 0
brk(0)                                  = 0x1bf9000
mmap(NULL, 4096, PROT_READ|PROT_WRITE, MAP_PRIVATE|MAP_ANONYMOUS, -1, 0) = 0x7fe594bb5000
access("/etc/ld.so.preload", R_OK)      = -1 ENOENT (No such file or directory)
open("/etc/ld.so.cache", O_RDONLY)      = 3
fstat(3, {st_mode=S_IFREG|0644, st_size=53205, ...}) = 0
mmap(NULL, 53205, PROT_READ, MAP_PRIVATE, 3, 0) = 0x7fe594ba8000
close(3)                                = 0
open("/lib64/libc.so.6", O_RDONLY)      = 3
read(3, "\177ELF\2\1\1\3\0\0\0\0\0\0\0\0\3\0>\0\1\0\0\0000\356\201f=\0\0\0"..., 832) = 832
fstat(3, {st_mode=S_IFREG|0755, st_size=1928936, ...}) = 0
mmap(0x3d66800000, 3750184, PROT_READ|PROT_EXEC, MAP_PRIVATE|MAP_DENYWRITE, 3, 0) = 0x3d66800000
mprotect(0x3d6698a000, 2097152, PROT_NONE) = 0

(…)

write(1, "Hello world\n", 12Hello world
)           = 12
exit_group(0)                           = ?
+++ exited with 0 +++
[vasco@student2]$
```

System Calls

Examples of WINDOWS and UNIX

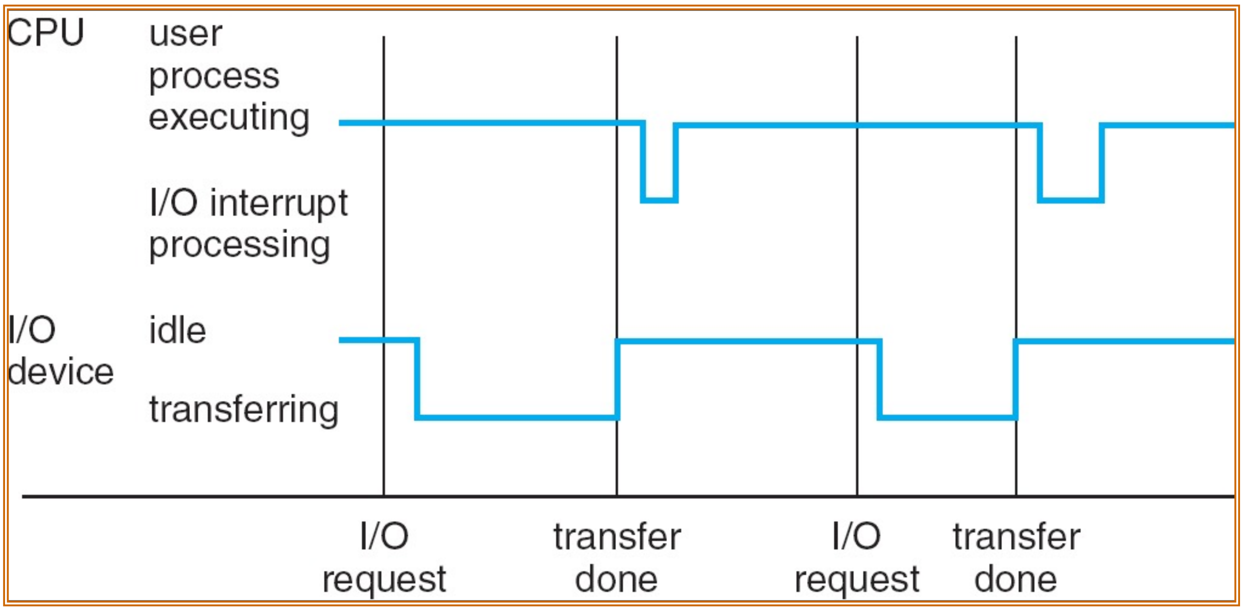
6 categories of system calls

	Windows	Unix
Process Control	CreateProcess() ExitProcess() WaitForSingleObject()	fork() exit() wait()
File Manipulation	CreateFile() ReadFile() WriteFile() CloseHandle()	open() read() write() close()
Device Manipulation	SetConsoleMode() ReadConsole() WriteConsole()	ioctl() read() write()
Information Maintenance	GetCurrentProcessID() SetTimer() Sleep()	getpid() alarm() sleep()
Communication	CreatePipe() CreateFileMapping() MapViewOfFile()	pipe() shm_open() mmap()
Protection	SetFileSecurity() InitlializeSecurityDescriptor() SetSecurityDescriptorGroup()	chmod() umask() chown()

Input/Output versus CPU Usage

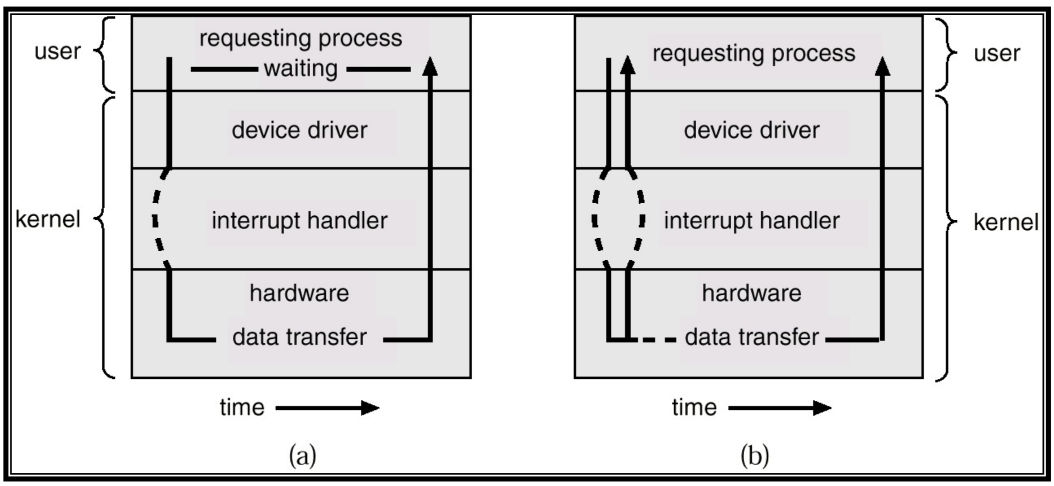
- I/O devices and the CPU can execute concurrently
- Each device controller is in charge of a particular device
- Each device controller has a local buffer
- CPU moves data from/to main memory to/from local buffers
- I/O is from the device to local buffer of controller
- Device controller informs CPU that it has finished its operation by causing an interrupt

Interrupt Timeline



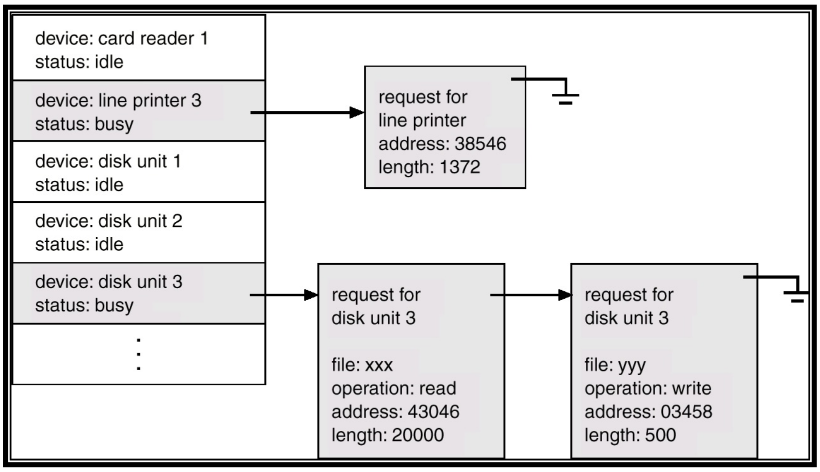
Different Types of I/O

- Synchronous (a) versus asynchronous I/O (b)
  - In (a) the application blocks on the system call
  - In (b) the application resumes execution after issuing the system call and later receives the answer



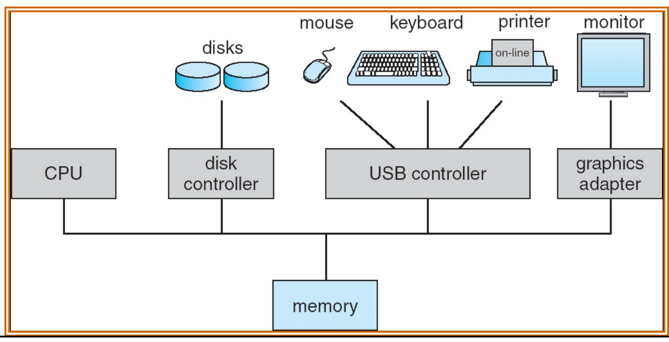
Device Status Table

- Each table entry indicates the device's type, address, and state (not functioning, idle, or busy). If the device is busy with a request, the type of request and other parameters will be stored in the table entry for that device.
- If the kernel supports asynchronous I/O, it also keeps track of many I/O requests at the same time.

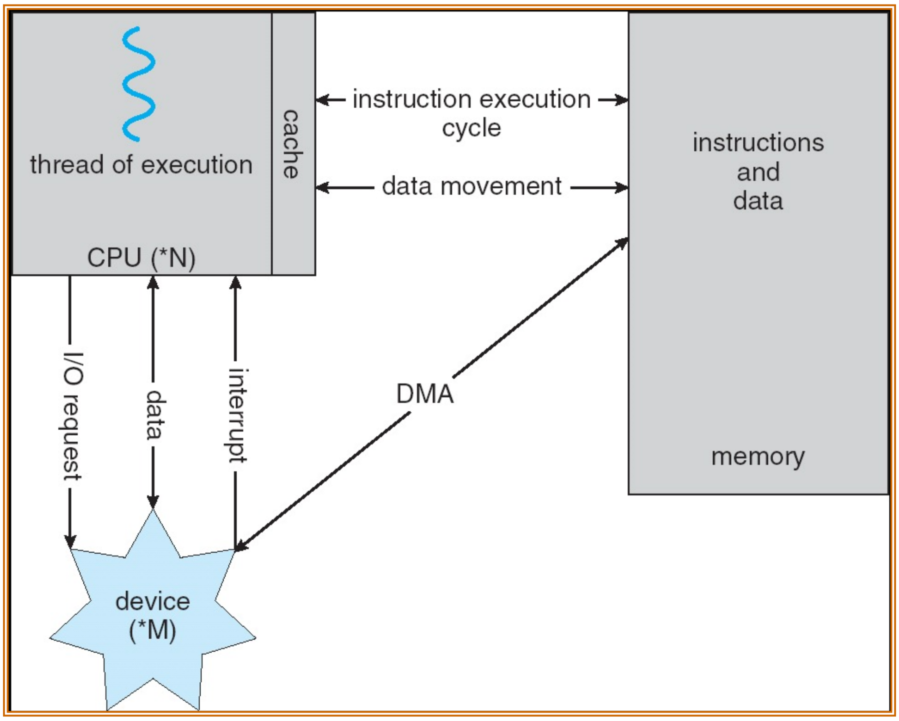


Direct Memory Access (DMA)

- Used for high-speed I/O devices able to transmit information at close to memory speeds
- Device controller transfers blocks of data from buffer storage directly to main memory without CPU intervention
- Implies the capability of arbitrating the system bus
- Only one interrupt is generated per block, rather than one interrupt per byte

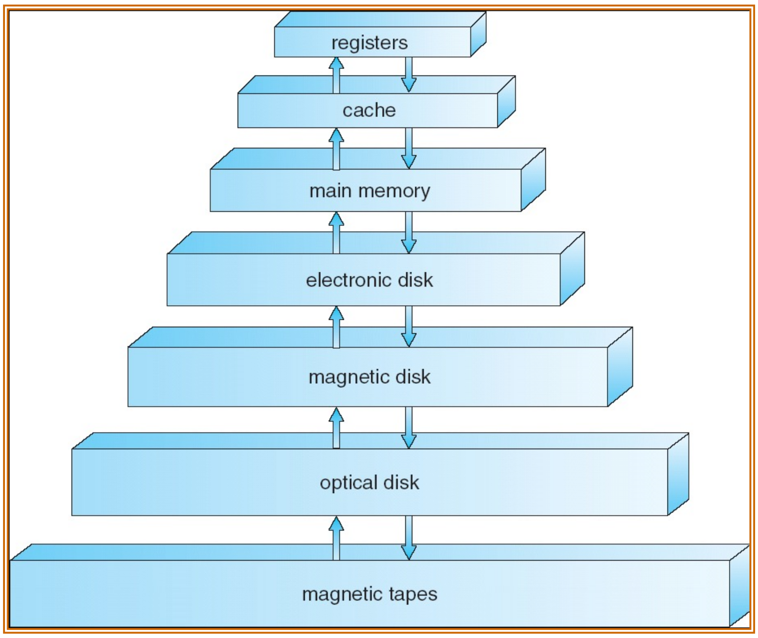


DMA (2)



Storage Hierarchy

- Caches, Buffers and a Storage Hierarchy are Essential



Storage Hierarchy

- A great part of this hierarchy is controlled by the Operating System
  - Cache Management Policies
  - Buffer Management Policies
  - Memory Management Policies

Principles of Locality:
- Temporal Locality
- Spatial Locality

Level	1	2	3	4
Name	registers	cache	main memory	disk storage
Typical size	< 1 KB	> 16 MB	> 16 GB	> 100 GB
Implementation technology	custom memory with multiple ports, CMOS	on-chip or off-chip CMOS SRAM	CMOS DRAM	magnetic disk
Access time (ns)	0.25 – 0.5	0.5 – 25	80 – 250	5,000.000
Bandwidth (MB/sec)	20,000 – 100,000	5000 – 10,000	1000 – 5000	20 – 150
Managed by	compiler	hardware	operating system	operating system
Backed by	cache	main memory	disk	CD or tape

Storage Hierarchy – a human perspective

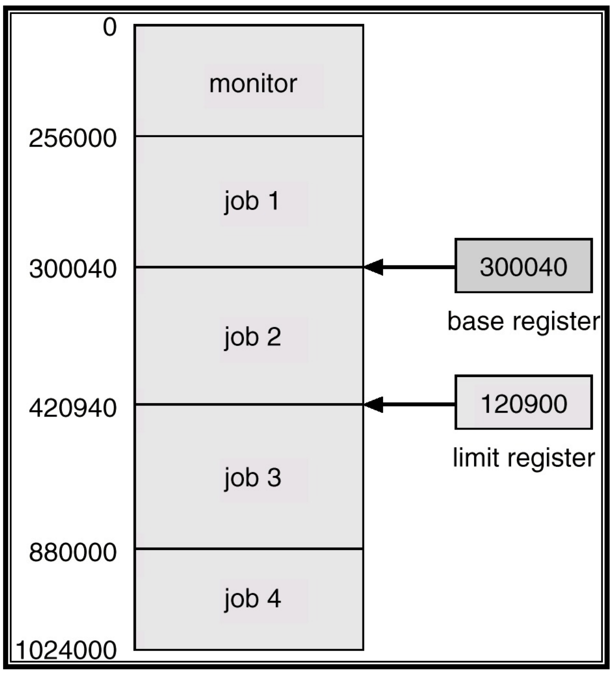
	Computer time	Human time
Processor cycle time (for a 2GHz frequency)	0.5 ns	1 s
L2 cache access*	10 ns	20 s
Local DRAM*	80 ns	160 s (2.7min)
Magnetic disk access time*	8,000,000 ns (8 ms)	16,000,000 s (≈185 days)
CD-ROM access time*	200 ms	400,000,000 s (≈12.7 years)

*Values change depending on the specific HW used

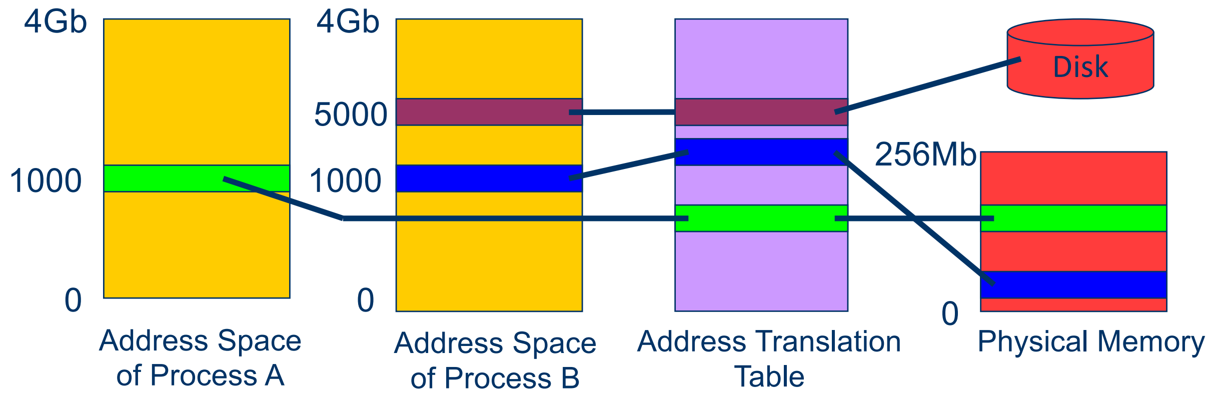
Memory Protection

- Each program must only be able to access its own memory
  - Segmentation
  - Virtual Memory
- Segmentation is based on having a base pointer and a limit for each process running
  - Any access made outside of its “space” generates a trap and normally leads to the process being killed
- Virtual Memory is based on having a table that translates “virtual addresses” into real ones
  - Normally, there is such a table for each process
  - Each process sees all the address space
  - An access made to an non existing page generates a trap and normally leads to the process being killed

Memory Protection
Segmentation

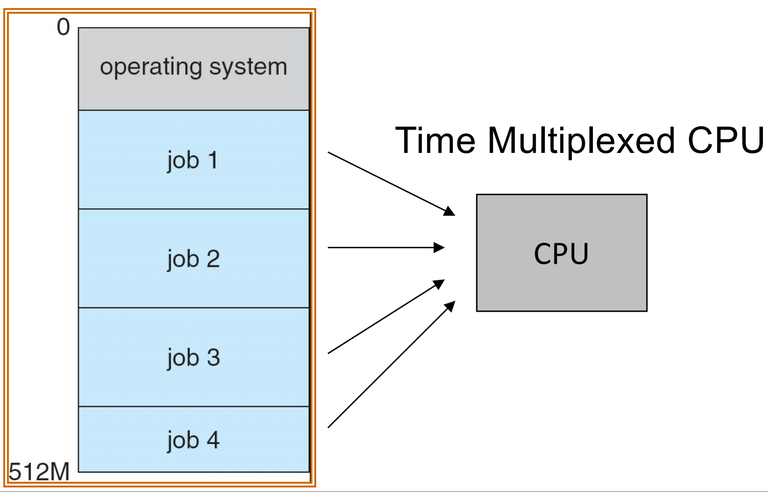


Memory Protection
Virtual Memory



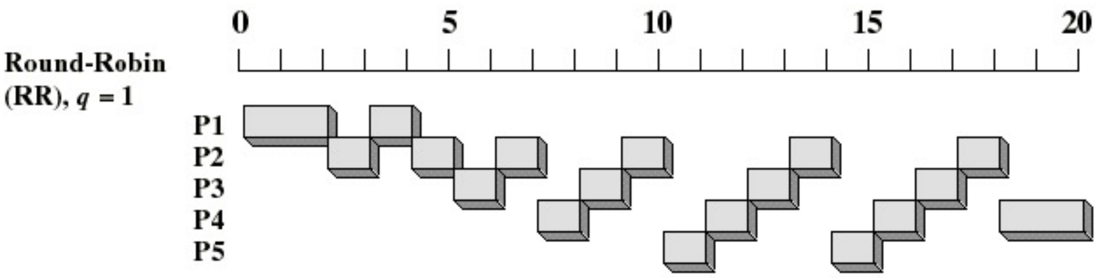
Multiprogramming

- Multiprogramming: maintaining several jobs in memory, active at the same time, so that when the current program can no longer run (for example because it's blocked waiting for I/O), another is available to run
- Optimizes resource utilization: CPU and I/O



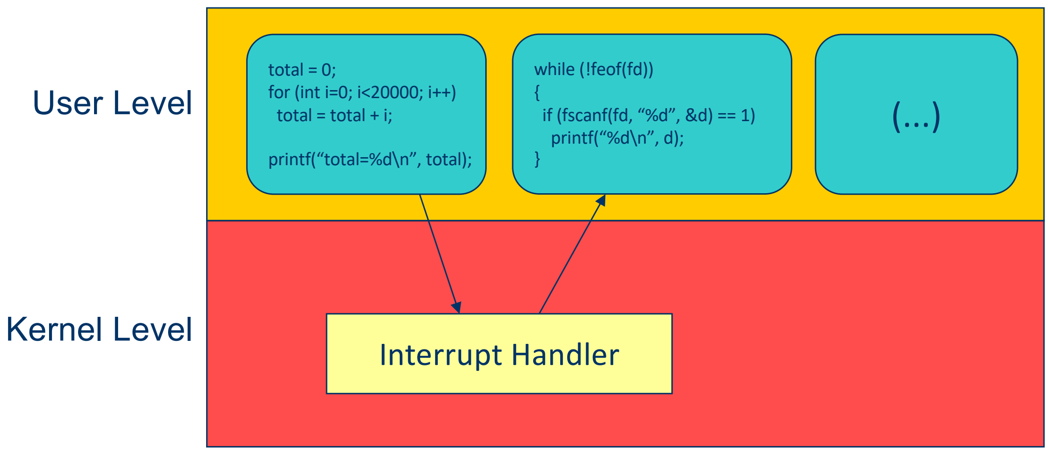
Multitasking

- Multitasking is an extension of multiprogramming in which the execution of the active programs is time-sliced:
  - Each program runs for a short period of time, then another is run.
  - When a program is running and is forcibly replaced by another, typically with a higher priority, it is said to have been preempted, thus the term Preemptive Operating Systems



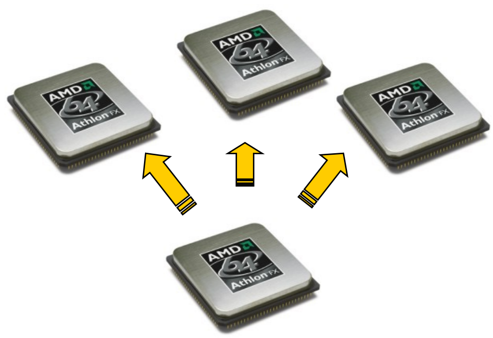
Multitasking

- For implementing multitasking, a non-maskable interrupt must connected to a hardware timer
  - Every few milliseconds (e.g. 100ms), the interrupt causes a task (or process) switch

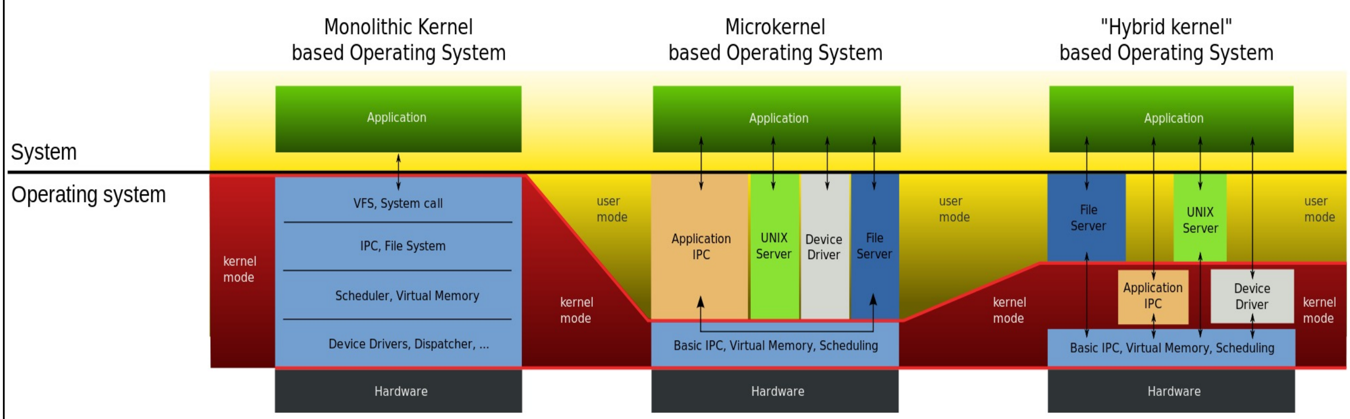


Nowadays…

- Since Windows XP, Linux 2.6.xx:
  - Preemptive multitasking operating systems using virtual memory
  - Each process thinks it has the whole computer for itself
    → “Virtual Machine”



Operating System Architecture and Structures



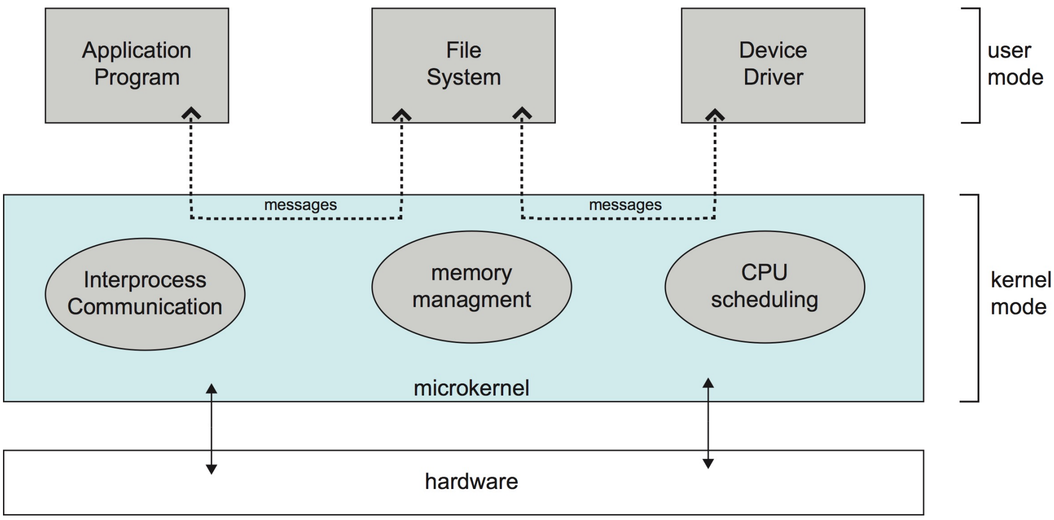
Source: http://en.wikipedia.org/wiki/File:OS-structure2.svg

Operating System Architecture and Structures

- **Microkernel**:
  - Moves as much as possible from the kernel into "user" space
  - The main function of the microkernel is to provide communication between the client program and the various services that are also running in user space. Communication is provided through message passing
  - e.g., Mach, XNU (used in OpenStep and Darwin, part of MAC OS X)
  - Benefits:
    - Easier to extend a microkernel
    - Easier to port the operating system to new architectures
    - More reliable (less code is running in kernel mode)
    - More secure
  - Detriments:
    - SLOW: Performance overhead of user space to kernel space communication

Operating System Architecture and Structures

- **Microkernel**:



Operating System Architecture and Structures

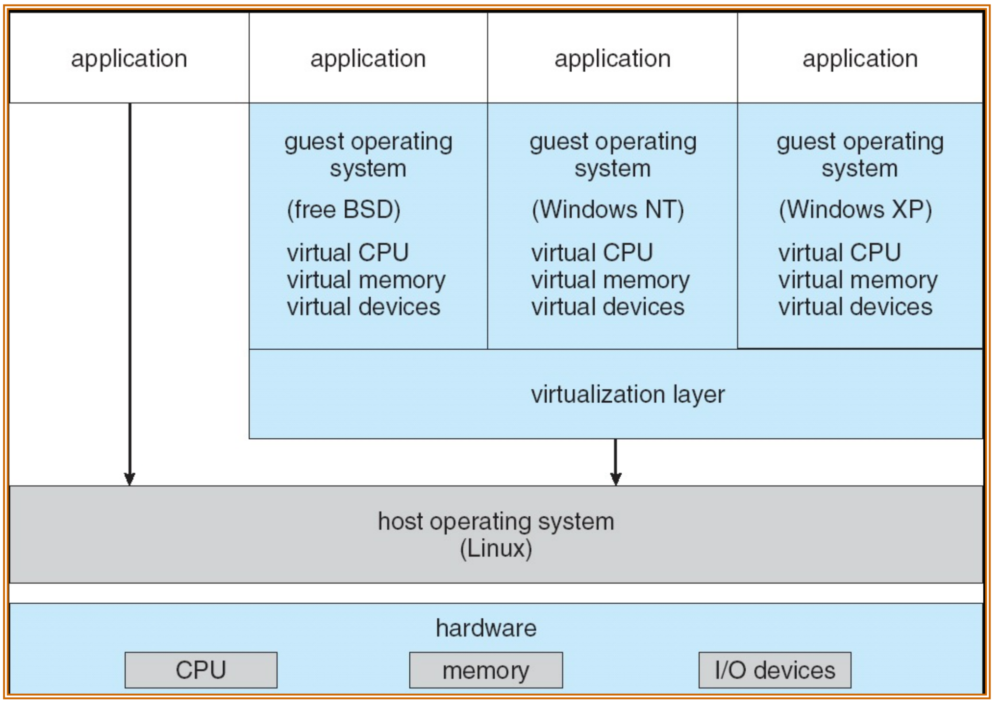
- **Monolithic**:
  - Everything is put below the system-call interface and above the physical hardware
  - Provides the file system, CPU scheduling, memory management, and other operating-system functions; a large number of functions for one level
  - e.g., Unix, Linux, MS-DOS, Windows 9* family
  - Benefits:
    - FAST!
  - Detriments:
    - Less reliable (more code is running in kernel mode)
    - Less secure
    - Not so easy to port to new architectures (… it depends)

Virtual Machines

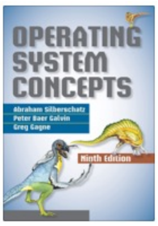
- Process Virtual Machines:
  - “Abstract Machine” that runs bytecode (e.g. Java, MS.NET)
  - Good for portability, allowing running the same unmodified application in different physical architectures
  - Designed to execute computer programs in a platform-independent environment

- System Virtual Machines:
  - Software that allows different operating systems to be run in parallel giving them the illusion of “having a whole machine”
    (e.g. VMWARE, XEN, VirtualBox, …)
  - Some Virtual Machines run on top of the hardware: IBM z/VM
  - Others run on top of the host OS: XEN, VMWARE, VirtualBox, etc
  - Allows for server consolidation, software testing, software fault tolerance, etc.

Virtual Machines
VMWARE Architecture



References



- [Silberschatz13]
  - Chapter 1: Introduction
    - 1.1, 1.2, 1.4, 1.5, 1.6, 1.7, 1.8, 1.9
  - Chapter 2: Operating-System Structures
    - All chapter 2!

Thank you! Questions?

I keep six honest serving men. They taught me all I knew. Their names are What and Why and When and How and Where and Who.
—Rudyard Kipling