# Algorithmic Learning Theory - Exercise sheet 4

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# Exercise 1

Let  $\mathcal{X} = \{0, 1\}^n$ ,  $I \subseteq \{1, 2, ..., n\}$ , and  $h_I = (\sum_{i \in I}) \mod 2$ .

We show  $VCdim(\mathcal{H}_{n-parity}) \geq n$ .

We define the set  $S_n$  as the set of all n unit vectors  $e_j = (e_{j,1}, ..., e_{j,n})$ . If  $e_j$  was labeled positively, we add i to the set I if  $e_{j,i} = 1$ . In the end, we obtain a set of indices in combination with  $h_I$  which fulfills the labeling. With this setup, we can construct all necessary hypotheses. So,  $\mathcal{H}_{n-parity}$  shatters a set of at least size n.

We show  $VCdim(\mathcal{H}_{n-parity}) \leq n$ .

We know that it holds:  $VCdim(\mathcal{H}_{n-parity}) \leq log_2(|\mathcal{H}_{n-parity}|)$ .

Since we only have  $2^n$  parity functions  $(2^n \text{ subsets of } \{1, 2, ..., n\})$  it holds:

 $VCdim(\mathcal{H}_{n-parity}) \le log_2(2^n) = n$ 

# Exercise 2

a.

Chose  $k \geq d \geq 3$  such that there is a set S with size k that can be shattered by  $\bigcup_{i=1}^r \mathcal{H}_i$ . This k exists, because  $\max_i VCdim(\mathcal{H}_i) \geq 3$ , so  $VCdim(\bigcup_{i=1}^r \mathcal{H}_i) \geq 3$  as well.

By Sauer's lemma, we know (assuming  $d \ge m$ ):

$$\tau_{\bigcup_{i=1}^{r} \mathcal{H}_{i}}(m) \leq \sum_{i=1}^{r} \tau_{H_{i}}(m)$$

$$\stackrel{(1)}{\leq} r \cdot \tau_{H_{j}}(m)$$

$$\stackrel{(2)}{\leq} r \cdot \sum_{i=0}^{d} \binom{m}{i}$$

$$\stackrel{(3)}{\leq} r \cdot \sum_{i=0}^{d} \binom{d}{i}$$

$$= r \cdot 2^{d}$$

$$\stackrel{(4)}{\leq} r \cdot k^{d}$$

- (1) Chose  $j = \operatorname{argmax}_{j \in [r]} \{ \tau_{\mathcal{H}_j}(m) \}$
- (2)  $VCdim(\bigcup_{i=1}^r \mathcal{H}_i) < d \Rightarrow VCdim(H_j) < d$ , apply Sauer's lemma
- (3)  $d \ge m$  (Assumption!)
- (4)  $k \geq 3$

Because S can be shattered,  $\tau_{\bigcup_{i=1}^r \mathcal{H}_i}(k) \geq 2^k$ , so we know that  $2^k \leq rk^d$ 

#### b.

Proof that  $VCdim(\mathcal{H}_1 \cup \mathcal{H}_2) \leq 2d + 1$ :

- Let  $VCdim(\mathcal{H}_1 \cup \mathcal{H}_2) \geq 2d + 2$  with  $VCdim(\mathcal{H}_1) = d_1$  and  $VCdim(\mathcal{H}_2) = d_2$  and  $d = max\{d_1, d_2\}$
- W.l.o.g. let  $\mathcal{H}_1$  shatter  $X = \{x_1, \dots, x_{d_1}\}$  and  $\mathcal{H}_2$  shatter  $Y = \{y_1, \dots, y_{d_2}\}$
- Let  $Z = X \bigcup Y \bigcup \{z_1, z_2\} \subseteq \mathcal{X}$  be a set that is shattered by  $\mathcal{H}_u := \mathcal{H}_1 \cup \mathcal{H}_2$
- Let  $h_I$  be the hypothesis that accounts for  $I \subseteq Z$
- Any function  $h_{I'}$  for I' with  $I' \setminus Z \neq \emptyset$ , we can find an equivalent function  $h_{I' \cap Z}$  which accounts for the same subsets of Z
- $\mathcal{H}_1$  must contain  $h_{I_x}$  for all  $I_x \subseteq X$
- $\mathcal{H}_2$  must contain  $h_{I_y}$  for all  $I_y \subseteq Y$
- The  $h_{I_z}$  for all  $I_z \in (\mathcal{P}(Z) \backslash \mathcal{P}(X)) \backslash \mathcal{P}(Y)$  are distributed among  $\mathcal{H}_1$  and  $\mathcal{H}_2$
- Claim: one of these holds:  $\mathcal{H}_1$  shatters  $X \cup \{z_1\}$  or  $X \cup \{z_2\}$  or  $\mathcal{H}_2$  shatters  $Y \cup \{z_1\}$  or  $Y \cup \{z_2\}$  Proof:
  - For simplicity (w.l.o.g) we assume that  $\mathcal{H}_1$  does not shatter  $X \cup \{z_1\}$
  - Then there is one  $I = \{x_{i_1}, \dots, x_{i_n}, z_1\} \subseteq X \cup \{z_1\}$ s.t.  $h_J \notin \mathcal{H}_1$  for  $J = I \cup I'$  for all  $I' \in \mathcal{P}(Y \cup \{z_2\})$ i.e. exclude all hypotheses that could somehow account for I
  - Otherwise any of the  $h_J$  could account for I
  - Then  $h_J \in \mathcal{H}_2$  for all J.
  - This implies that  $\mathcal{H}_2$  shatters  $Y \cup \{z_2\}$  since the  $h_J$ 's can account for all  $I' \in \mathcal{P}(Y \cup \{z_2\})$
- Consequently we have  $VCdim(\mathcal{H}_1) > d_1$  or  $VCdim(\mathcal{H}_2) > d_2$  $\Rightarrow$  Contradiction

### Exercise 3

#### a.

In addition to the all-negative hypothesis,  $3^d$  hypotheses are possible by either choosing for every literal either the positive or negative interpretation, or omitting it from the hypothesis. Hypthoeses, where one literal appears multiple times in the same form are logically equivalent to any hypothesis where this literal appears only once.

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         \Rightarrow |\mathcal{H}^d_{con}| \leq 3^d + 1  We now want to show that VCdim(\mathcal{H}^d_{con}) \leq d \cdot \log_2 3 Proof by contradiction: Assume VCdim(\mathcal{H}^d_{con}) \geq d \cdot \log_2 3 + 1          \Rightarrow \mathcal{H}^d_{con} \text{ shatters } X \text{ with } |X| = \lceil d \cdot log_2 3 + 1 \rceil           \Rightarrow |\mathcal{H}^d_{con}| \geq 2^{|X|} \geq 2^{d \cdot log_2 3 + 1} = 2^{log_2 3^d} \cdot 2 = 3^d \cdot 2 > 3^d + 1 \geq |\mathcal{H}^d_{con}|           \Rightarrow \text{ contradiction}
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#### b.

Let  $S_d = \{e_i | i \leq d\}$  with unit vector  $e_i = (e_{i,1}, ..., e_{i,j}, ..., e_{i,d})$  and  $e_{i,j} = 1$ . So,  $|S_d| = d$ . We construct a family of  $h_i \in \mathcal{H}^d_{con}$  which shatters the set  $S_d$ . The base structure of  $h_i$  is the conjunction where all literals are negated. According to the labeling, we omit a negative literal  $x_j$  if  $e_i$  was labeled positively (Omit the literal where the unit vector has  $e_{i,j} = 1$ ). Consequently, we obtain  $2^d$  different conjunctions/hypotheses and this implies that  $\mathcal{H}^d_{con}$  shatters  $S_d$ .

Assume  $VCdim(\mathcal{H}^d_{con}) \geq d+1$ ). Then there is a set  $C=(c_1,\ldots,c_{d+1})$  of boolean functions with size d+1 and  $c_i \neq c_j \forall i, j \in [d+1] \land i \neq j$  that can be shattered by  $\mathcal{H}^d_{con}$ .

Set  $C_i = C \setminus \{c_i\}$  for all  $i \in [d+1]$ . Because C can be shattered, for every  $C_i$ , there is a hypothesis  $h_i \in \mathcal{H}^d_{con}$  as defined in the hint (it accepts exactly the  $c_j \in C_i$ ).  $l_i$  is defined as in the hint as well, i.e. it's a literal that's false on  $c_i$  and true on all other  $c_j$  with  $j \neq i$ .

Because there are only d variables, but d+1 such literals, at least one variable has to be used at least twice. Let  $l_a, l_b$  be two literals that both use the same variable.

Case 1: Both literals are positive or both are negative. Then both literals are identical because they use the same variable. Then  $h_a = h_b$  and therefore  $S_a = S_b$ . Hence,  $c_a = c_b$ . This is a contradiction, because C was chosen to contain pairwise different elements.

Case 2: One literal is positive, the other negative. W.l.o.g., let  $l_a$  be the positive,  $l_b$  the negative one. From  $h_a$  we know that  $l_a$  is true for all  $c_i$  with  $i \neq a$ . Because  $l_b$  is the negation of  $l_a$ ,  $l_b$  is false for all  $c_i$  with  $i \neq b$ . But this implies that all  $c_i$  with  $i \neq b$  are not in  $h_b$ . This is a contradiction to the choice of  $h_b$ .

Therefore, no such set C can exist. Because not set of size d+1 can be shattered, not set with size  $\geq d+1$  can be shattered, so  $VCdim(\mathcal{H}^d_{con}) \leq d$ .

#### d.

 $VCdim(\mathcal{H}^d_{mcon}) \leq d$ : The all positive hypothesis is contained in  $\mathcal{H}^d_{con}$ . The all negative hypothesis can be generated by the conjunction  $x_1 \wedge \neg x_1 \in \mathcal{H}^d_{con}$ . All other hypotheses in  $\mathcal{H}^d_{mcon}$  are conjunctions, and as such part of  $\mathcal{H}^d_{con}$ . Therefore,  $\mathcal{H}^d_{mcon} \subseteq \mathcal{H}^d_{con}$ . By reducing the set of available hypotheses, the VC dimension can not be increased. If it is possible to shatter a set of size m using hypotheses in  $\mathcal{H}^d_{mcon}$ , this set could be shattered by  $\mathcal{H}^d_{con}$  as well. In c) we proved that  $VCdim(\mathcal{H}^d_{con}) \leq d$ , so  $VCdim(\mathcal{H}^d_{mcon}) \leq d$  as well.

 $VCdim(\mathcal{H}^d_{con}) \leq d$ , so  $VCdim(\mathcal{H}^d_{mcon}) \leq d$  as well.  $VCdim(\mathcal{H}^d_{mcon}) \geq d$ : Chose  $S = \{s_1, \ldots, s_d\}$  with  $s_{ij} = 0 \Leftrightarrow i = j$  (ie. element  $s_i$  has 1s for every variable except  $x_j$ , there it has a 0).

- The subset  $S' = \emptyset \subseteq S$  can be generated using the all negative hypothesis.
- The subset S' = S can be generated using the all positive hypothesis.
- For every other subset  $S' \subseteq S$ , let  $I = \{i \in [d] | s_i \in S'\}$ . To find a hypothesis h that only accepts elements in S', one can build the conjunction of all variables  $x_j$  with  $j \notin I$ . Due to the choice of S and h, every element not in S' has a negative value for one of the variables used in h. Every element of S' has only positive values for the variables in h. Therefore, h accepts exactly those elements that are in S'. Because h can be build for every  $S' \subseteq S$ , S can be shattered. Hence,  $VCdim(\mathcal{H}^d_{mcon}) \geq d$ .

By combining both proofs, it is shown that  $VCdim(\mathcal{H}^d_{mcon}) = d$ .

#### Exercise 4

Let  $h: \{0,1\}^n \to \{0,1\}$  be a symmetric function and  $\mathcal{H}$  the set of all symmetric functions. First, we show that the  $VCdim(\mathcal{H}) \ge n+1$ .

Therefore, we take a set  $S \subset \{0,1\}^n$  of maximal size such that there are no  $s_1, s_2 \in S$  with the same number of 1's, i.e.  $\{(0,0,...,0), (1,0,...,0), (1,1,0,...,0), ..., (1,1,...,1)\}$ . Consequently, |S| = n+1. Now we show that we can shatter the set S. We define the family of symmetric functions  $h_{1,...,i}$  with  $h_i(x) = 1$  iff x has exactly i ones (e.g.  $h_1((1,0,0)) = 1$ ,  $h_2((1,0,0)) = 0$ ). Since the value of the function  $h_i$  depends on the number of ones, it is symmetric. With this set of functions, we are able to shatter a set of size n+1 by combining  $h_i$  and  $h_{i'}$  ( $i \neq i$ ) to a single function  $h_{i,i'}$  which is also symmetric. Meaning that

$$h_{i_1,\dots,i_n}(x) = \begin{cases} 1 \text{ if number(ones)} = i_1\\ 1 \text{ if number(ones)} = i_2\\ \dots\\ 1 \text{ if number(ones)} = i_n\\ 0 \text{ else} \end{cases}$$

With these symmetric functions, we are able to shatter S by setting the i's according to the labeling.  $\implies VCdim(\mathcal{H}) \ge n+1$ .

Now, we show that  $VCdim(\mathcal{H}) \leq n+1$ :

Assume there is a set |S'| = (n+1)+1 which can be shattered by  $\mathcal{H}$ . Then, there are  $x_1, x_2 \in S'$  with the same number of 1's due to the finite instance space and the pigeonhole principle. This means that for any  $h \in \mathcal{H}$  it always holds  $h(x_1) = h(x_2)$  due to the definition of symmetric functions. Consequently, there is no  $h^*$  with  $h^*(x_1) \neq h^*(x_2)$ . In terms of labeling, there is no different labeling of  $x_1$  and  $x_2$  possible.

 $\implies$  specific labelings are not possible  $\implies$  no shattering of  $S' \implies VCdim(\mathcal{H}) \le n+1$