

MIDTERM REVIEW
Wednesday, March 18, 2015
Regular class

MIDTERM EXAM
Open book and notes (not Internet)
Saturday, March 21, 2015
8:00 am - 12:00 pm (4 hours)
Room LW-844

Assignment 6. Due: 03/09/15

- 14. Consider a modified grammar of Zones. The only difference is the definition of function *ALPHA*:**

$$D(ALPHA) = X \times P \times L_t^{I_0}(S) \times Z_+$$

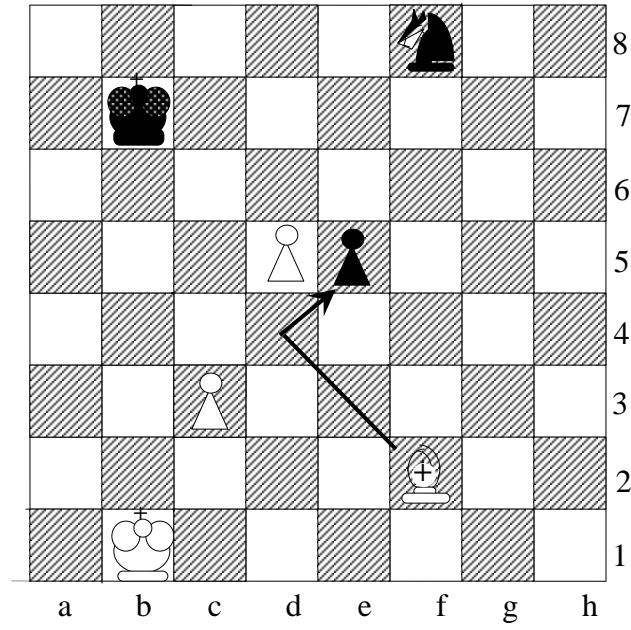
$$ALPHA(x, p_o, t_o, k) = \begin{cases} \min(NEXTTIME(x), k), & \text{if } DIST(x, p_o, t_o) < 2n, \\ NEXTTIME(x), & \text{if } DIST(x, p_o, t_o) = 2n. \end{cases}$$

What is the impact of this new definition on the Zones to be generated by this grammar? Show examples of such Zones. Explain.

Extra Credit.

Do we have to change function *timer* in this case? See “Translations of Languages” in the textbook on LG. Explain.

Second Negation



1	Q_1	$S(u, v, w) \rightarrow A(u, v, w)$		two	\emptyset
2 _i	Q_2	$A(u, v, w) \rightarrow t(h_i^0(u), l_0+1)$ $A((0, 0, 0), g(h_i^0(u), w), zero)$	$TIME(z)=DIST(z, h_i^0(u))$	3	\emptyset
3	Q_3	$A(u, v, w) \rightarrow A(f(u, v), v, w)$	$NEXTTIME(z)=$ $init(u, NEXTTIME(z))$	four	5
4 _j	Q_4	$A(u, v, w) \rightarrow t(h_j(u), TIME(y))$ $A(u, v, g(h_j(u), w))$	$NEXTTIME(z)=$ $ALPHA(z, h_j(u), TIME(y) - l_j+1)$	3	3
5	Q_5	$A(u, v, w) \rightarrow A((0, 0, 0), w, zero)$	$TIME(z) = NEXTTIME(z)$	3	6
6	Q_6	$A(u, v, w) \rightarrow e$		\emptyset	\emptyset

$Q_1(u) = (ON(p_0) = x) \wedge (MAP_{x,p_0}(y) \leq l \leq l_0) \wedge (\exists q ((ON(q) = y) \wedge (OPPOSE(p_0, q))))$

$Q_2(u) = T$; $Q_3(u) = (x \neq n) \vee (y \neq n)$

$Q_4(u) = (\exists p ((ON(p) = x) \wedge (l > 0) \wedge (x \neq x_0) \wedge (x \neq y_0)) \wedge ((\neg OPPOSE(p_0, p) \wedge$

$(MAP_{x,p}(y) = 1)) \vee (OPPOSE(p_0, p) \wedge (MAP_{x,p}(y) \leq l)))$ $Q_5(w) = (w \neq zero); Q_6 = T$

$init(u, r) = \begin{cases} 2n, & \text{if } u = (0, 0, 0), \\ r, & \text{if } u \neq (0, 0, 0). \end{cases}$

$f(u, v) = \begin{cases} (x+1, y, l), & \text{if } ((x \neq n) \wedge (l > 0)) \vee ((y = n) \wedge (l \leq 0)) \\ (1, y+1, TIME(y+1) \times v_{y+1}), & \text{if } (x = n) \vee ((l \leq 0) \wedge (y \neq n)). \end{cases}$

Let $t_0 \in L_t^{l_0}(S)$, $t_0 = a(z_0)a(z_1)...a(z_m)$, $t_0 \in t_{p_0}(z_0, z_m, m)$;

If $((z_m = y_0) \wedge (p = p_0) \wedge (\exists k (1 \leq k \leq m) \wedge (x = z_k))) \vee$
 $((z_m \neq y_0) \vee (p \neq p_0)) \wedge (\exists k (1 \leq k \leq m-1) \wedge (x = z_k))$

then $DIST(x, p_0, t_0) = k+1$ else $DIST(x, p_0, t_0) = 2n$

$ALPHA(x, p_0, t_0, k) = \begin{cases} \max(NEXTTIME(x), k), & \text{if } (DIST(x, p_0, t_0) \neq 2n) \\ & \wedge (NEXTTIME(x) \neq 2n); \\ k, & \text{if } DIST(x, p_0, t_0) \neq 2n \\ & \wedge (NEXTTIME(x) = 2n); \\ NEXTTIME(x), & \text{if } DIST(x, p_0, t_0) = 2n. \end{cases}$

$g_r(p_0, t_0, w) = \begin{cases} 1, & \text{if } DIST(r, p_0, t_0) < 2n, \\ w_r, & \text{if } DIST(r, p_0, t_0) = 2n. \end{cases}$ $TRACKS_{p_0} = \{p_0\} \times (\bigcup_{1 \leq k \leq l} L[G_t^{(2)}(x, y, k, p_0)])$

$1 \leq k \leq l$

If $TRACKS_{p_0} = e$

then $h_i^0(u) = e$

else $TRACKS_{p_0} = \{(p_0, t_1), (p_0, t_2), \dots, (p_0, t_b)\}, (b \leq M)$ and $h_i^0(u) = \begin{cases} (p_0, t_i), & \text{if } i \leq b, \\ (p_0, t_b), & \text{if } i > b. \end{cases}$

$TRACKS = \bigcup_{ON(p)=x} TRACKS_p$, where $TRACKS_p$ is the same as for h_i^0

If $TRACKS = e$

then $h_i(u) = e$

else $TRACKS = \{(p_1, t_1), (p_1, t_2), \dots, (p_m, t_m)\}, (m \leq M)$ and $h_i(u) = \begin{cases} (p_i, t_i), & \text{if } i \leq m, \\ (p_m, t_m), & \text{if } i > m. \end{cases}$

At the beginning : $u = (x_0, y_0, l_0)$, $w = zero$, $v = zero$, $x_0 \in X$, $y_0 \in X$, $l_0 \in \mathbb{Z}_+$,

$p_0 \in P$, and $TIME(z)=2n$, $NEXTTIME(z)=2n$ for all z from X .