

## LP problems

- A standard linear programming (LP) problem is:

$$\begin{aligned} & \min_{z} cz \\ & \text{s.t. } Az \geq b \\ & z \geq 0 \end{aligned} \tag{1}$$

- $z$  is the vector of decision variables, the *solutions*.
  - $c$  is called objective coefficients.
  - Problem (1) is to determine an optimal value of  $z$  (searching for the best  $z$ ), denoted as  $z^*$ , such that  $cz^*$  is the minimum.
  - $Az \geq b$  represent the *constraints* in a matrix form. Each row of  $Az \geq b$  is a *constraint*.
  - $b$  is called the right-hand-side (rhs) vector.
  - $A$  is the left-hand-side (lhs) matrix.
  - $z$  can have upper bounds and lower bounds. Here the lower bounds are 0 and upper bounds are  $\infty$ .
- Gurobi (a *solver*) solves Problem (1). Above list the inputs for Gurobi.
  - (Note) Gurobi does not recommend dense matrix form shown above. Check the comments in the example `dense.c.c`. Gurobi does not store  $A$  in a dense matrix, but in a sparse matrix (CSR format) which ignores 0's.
  - **Task 1** Write a function: given any size of matrices  $A$ ,  $b$  and  $c$ , implement `mip1.c.c` with  $x$ ,  $y$  and  $z$  are continuous and non-negative (as in `dense.c.c`).
  - **Task 2** Read data from a csv file and store as `c` array (matrix).

## DEA problems

- $X = [x_r] = [x_1 \ x_2 \ \cdots \ x_n]$  where  $x_r$  is a  $p \times 1$  vector.
- $Y = [y_r] = [y_1 \ y_2 \ \cdots \ y_n]$  where  $y_r$  is a  $q \times 1$  vector.
- Data envelopment analysis (DEA) problem:

$$\begin{aligned} & \min_{\theta, \lambda} \theta \\ & \text{s.t. } X\lambda \leq x_r\theta \\ & Y\lambda \geq y_r \\ & e\lambda = 1 \\ & \lambda \geq 0 \end{aligned} \tag{2}$$

where is a  $\lambda = [\lambda_1 \ \lambda_2 \ \cdots \ \lambda_n]$  is a  $1 \times n$  vector (or  $n \times 1$  vector).  $x_r$  and  $y_r$  are the  $r$ th column of  $X$  and  $Y$ , respectively.

It is equivalent to:

$$\begin{aligned} & \min_{\theta, \lambda} \theta \\ & \text{s.t. } -X\lambda \geq -x_r\theta \\ & Y\lambda \geq y_r \\ & e\lambda = 1 \\ & \lambda \geq 0 \end{aligned} \tag{3}$$

It can be rearranged as:

$$\begin{aligned}
& \min_{\theta, \lambda} \theta \\
& \text{s.t. } x_r \theta - X \lambda \geq \mathbf{0} \\
& \quad \mathbf{0} \theta + Y \lambda \geq y_r \\
& \quad 0 \theta + e \lambda = 1 \\
& \quad \lambda \geq 0, \theta \geq 0
\end{aligned} \tag{4}$$

#### Solving Problem (4)

- I want to solve Problem (2) (equivalently, Problem (4)), which is an instance of Problem (1). Comparing Problems (1) and (4), we have

$$A = \begin{bmatrix} x_r & -X \\ \mathbf{0} & Y \end{bmatrix}$$

$$b = \begin{bmatrix} \mathbf{0} \\ y_r \end{bmatrix}$$

$x_r$  and  $y_r$  are the  $r$ th column of  $X$  and  $Y$  respectively. Thus here  $A$  is a  $(p+q) \times (1+n)$  matrix, and  $b$  is a  $(p+q) \times 1$  vector.  $\mathbf{0}$  is a  $p \times 1$  vector with zeros.

$$c = [1 \quad \mathbf{0}]$$

$c$  is a  $(1+n) \times 1$  (or  $1 \times (1+n)$ ) vector.

$$z = [\theta \quad \lambda]$$

$z$  is a  $(1+n) \times 1$  (or  $1 \times (1+n)$ ) vector.

- $X$  and  $Y$  are read from a csv file, in which row  $r$  represent column  $[x_r, y_r]$ . In the csv file, the first few columns associates with  $X$ , and follow by  $Y$ . For example, in file 5-2-25k\_10.csv, the first two columns are  $X^T$  (transpose of  $X$ ) and the last three column are  $Y^T$ .
- $0\theta + e\lambda = 1$  in (4) is an *equality constraint*, or *constraint with equality*. It is that you replace  $\leq$  or  $\geq$  by  $=$ . Check the manual for how to set it up. It means that the coefficient matrix (vector) is a  $1 \times (1+n)$  vector  $[0, e]$ , where  $e = [1, 1, 1 \dots, 1]$  is a  $1 \times n$  vector with values 1.
- **Task:** We need to do the following.

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#### Algorithm 1: Naive DEA

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**Input:**  $X, Y$

**Output:** optimal objective values

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1 for  $r = 1, \dots, n$  do
2   | populates  $A, b, c$  defined above;
3   |  $\theta_r^* \leftarrow$  solve Problem (4) ; // get optimal value
4 write  $\theta_r^*$ 's to a csv.
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- Each iteration of Algorithm (1) computes an LP problem.
- Please record the total computation time, but ignore the IO time.