

# Uninformed Search

**Chapter 3.1 – 3.4**

# Many AI Tasks can be Formulated as Search Problems

Goal is to find a *sequence of actions*

- **Puzzles**
- **Games**
- **Navigation**
- **Assignment**
- **Motion planning**
- **Scheduling**
- **Routing**

# Models To Be Studied in CS 540

## State-based Models

- **Model task as a graph of all possible states**
- A state captures all the relevant information about the past in order to act (optimally) in the future
- Actions correspond to transitions from one state to another
- Solutions are defined as a sequence of steps/actions (i.e., a path in the graph)
- State-space graphs

# Search Example: Route Finding

The screenshot shows a Google Maps interface in a Mozilla Firefox browser window. The title bar reads "Google Maps - From: 1210 W Dayton St, Madison, WI 53706 to: State St, Madison, WI 53703". The browser's address bar shows the Google Maps URL. The Google Maps logo is visible, along with links for "Maps", "Local Search", and "Directions". The "Directions" tab is active, showing the start address "1210 W Dayton St, Madison, WI 53706" and the end address "State St, Madison, WI 53703". A search button and a "Help" link are present. Below the address fields, the word "Maps" is displayed. The main map area shows a satellite view of a city street grid. A blue line indicates the route from the start address to the end address. The route starts on W Dayton St, goes east, then turns left onto N Frances St, then right onto W Gilman St, then right onto N Henry St, and finally right onto W Gorham St. The map includes a vertical scale bar on the left and a "Map" / "Satellite" toggle on the right. On the right side of the map, there are links for "Print", "Email", and "Link to this page". Below these links, the start and end addresses are repeated, along with the distance "1.2 mi (about 2 mins)". A section titled "Reverse directions" lists five steps: 1. Head east from W Dayton St - go 0.5 mi; 2. Turn left at N Frances St - go 0.2 mi; 3. Turn right at W Gilman St - go 0.3 mi; 4. Turn right at N Henry St - go 0.1 mi; 5. Turn right at W Gorham St - go 0.1 mi. A disclaimer states: "These directions are for planning purposes only. You may find that construction projects, traffic, or other events may cause road conditions to differ from the map results." At the bottom right, it says "Map data ©2005 NAVTEQ™, Tele Atlas".

Google Maps - From: 1210 W Dayton St, Madison, WI 53706 to: State St, Madison, WI 53703 - Mozilla Firefox

File Edit View Go Bookmarks Tools Help

Google Maps

Maps Local Search Directions

1210 W Dayton St, Madison, WI 53706 State St, Madison, WI 53703 Search Help

Start address End address

Maps

Map Satellite

Print Email Link to this page

Start address: 1210 W Dayton St  
Madison, WI 53706

End address: State St  
Madison, WI 53703

Distance: 1.2 mi (about 2 mins)

Reverse directions

1. Head east from W Dayton St - go 0.5 mi
2. Turn left at N Frances St - go 0.2 mi
3. Turn right at W Gilman St - go 0.3 mi
4. Turn right at N Henry St - go 0.1 mi
5. Turn right at W Gorham St - go 0.1 mi

These directions are for planning purposes only. You may find that construction projects, traffic, or other events may cause road conditions to differ from the map results.

Map data ©2005 NAVTEQ™, Tele Atlas

**Actions:** go straight, turn left, turn right

**Goal:** shortest? fastest? most scenic?

# Search Example: River Crossing Problem



**Goal:** All on  
right side of river

## **Rules:**

- 1) Farmer must row the boat
- 2) Only room for one other
- 3) Without the farmer present:
  - Dog bites sheep
  - Sheep eats cabbage

**Actions:**  $F>$ ,  $F<$ ,  
 $FC>$ ,  $FC<$ ,  $FD>$ ,  
 $FD<$ ,  $FS>$ ,  $FS<$

# Search Example: 8-Puzzle

7	2	4
5		6
8	3	1

Start State

	1	2
3	4	5
6	7	8

Goal State

**Actions:** move tiles (e.g., Move2Down)

**Goal:** reach a certain configuration

# Search Example: Water Jugs Problem

Given 4-liter and 3-liter pitchers, how do you get exactly 2 liters into the 4-liter pitcher?



# Search Example: Robot Motion Planning



**Actions:** translate and rotate joints

**Goal:** fastest? most energy efficient? safest?



# Search Example: Natural Language Translation

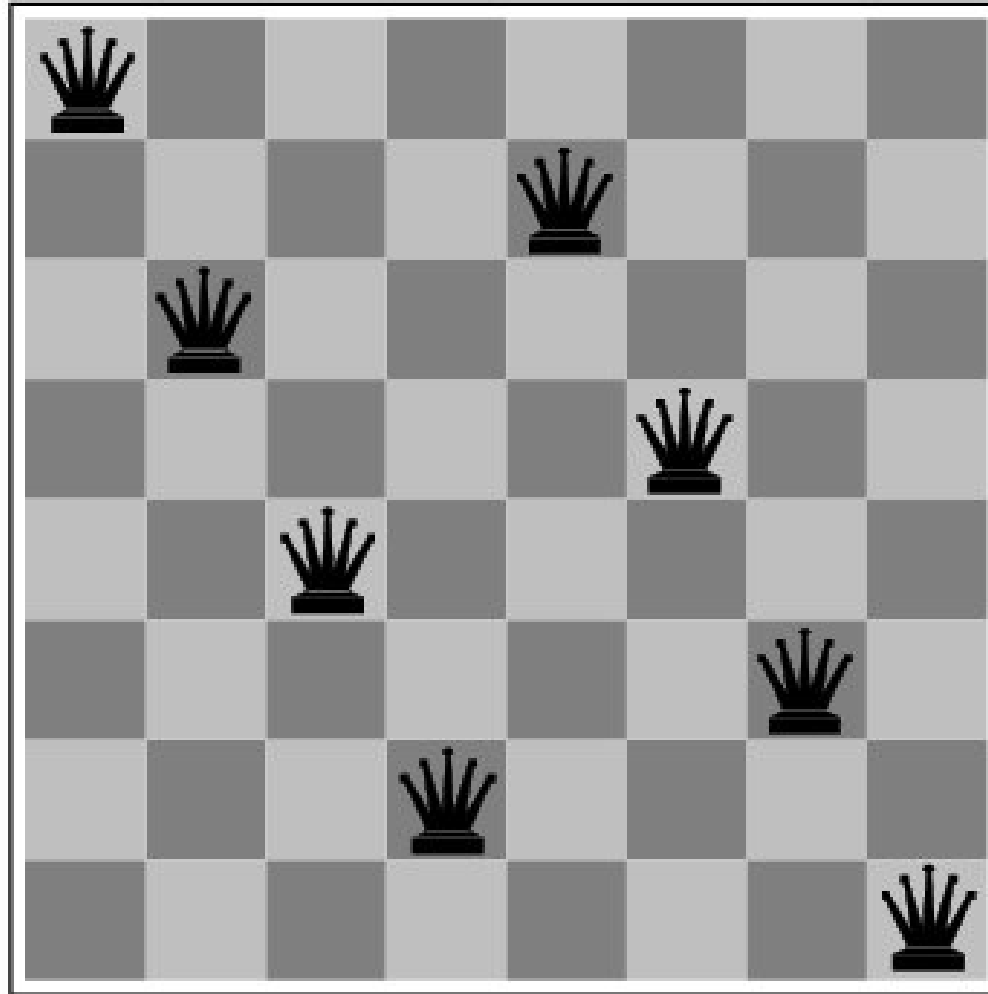
Italian → English:

la casa blu → the blue house

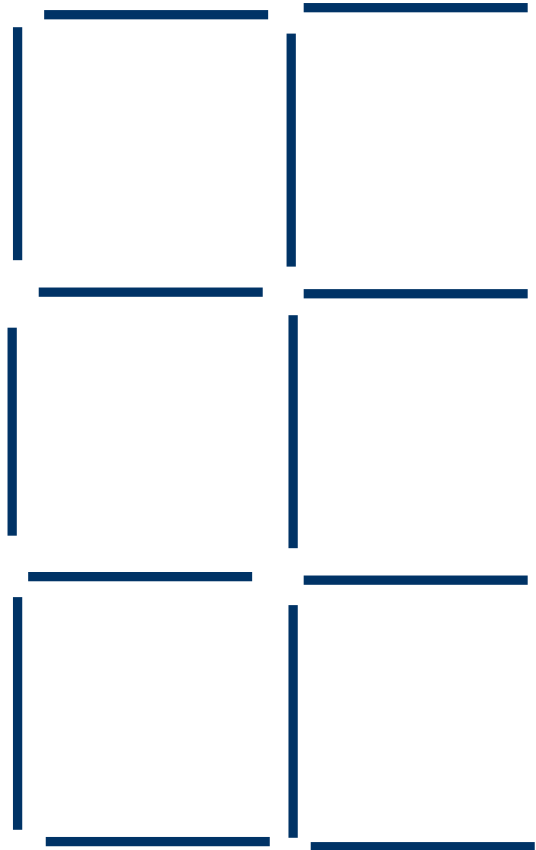
**Actions:** translate single words (e.g., la → the)

**Goal:** fluent English? preserves meaning?

# Search Example: 8-Queens



# Search Example: Remove 5 Sticks Problem



Remove exactly 5  
of the 17 sticks so  
the resulting figure  
forms exactly 3  
squares

# Basic Search Task Assumptions (usually, though not games)

- Fully observable
  - Deterministic
  - Static
  - Discrete
  - Single agent
- 
- Solution is a sequence of actions

# What Knowledge does the Agent Need?

- The information needs to be
  - sufficient to describe all relevant aspects for reaching the goal
  - adequate to describe the world ***state / situation***
- **Fully observable** assumption, also known as the ***closed world assumption***, means
  - *All necessary information about a problem domain is accessible so that each state is a complete description of the world; there is **no missing information** at any point in time*

# How should the Environment be Represented?

- **Knowledge representation problem:**
  - What information from the sensors is relevant?
  - How to represent domain knowledge?
- ***Determining **what** to represent is difficult and is usually left to the system designer to specify***
- Problem **State** = representation of all necessary information about the environment
- **State Space** (aka **Problem Space**) = *all* possible valid configurations of the environment

# What Goal does the Agent want to Achieve?

- **How do you describe the goal?**
  - as a task to be accomplished
  - as a state to be reached
  - as a set of properties to be satisfied
- **How do you know when the goal is reached?**
  - with a **goal test** that defines what it means to have achieved/satisfied the goal
  - or, with a set of **goal states**
- ***Determining the goal is usually left to the system designer or user to specify***

# What Actions does the Agent Need?

- Discrete and Deterministic task assumptions imply
- Given:
  - an ***action*** (aka ***operator*** or ***move***)
  - a description of the current state of the world
- Action completely specifies:
  - if that action *can* be applied (i.e., legal)
  - what the exact state of the world will be after the action is performed in the current state (no "history" information needed to compute the successor state)



# What Actions does the Agent Need?

- **A finite set of actions/operators needs to be**
  - decomposed into atomic steps that are discrete and indivisible, and therefore can be treated as instantaneous
  - sufficient to describe all necessary changes
- ***The number of actions needed depends on how the world states are represented***

# Search Example: 8-Puzzle

7	2	4
5		6
8	3	1

Start State

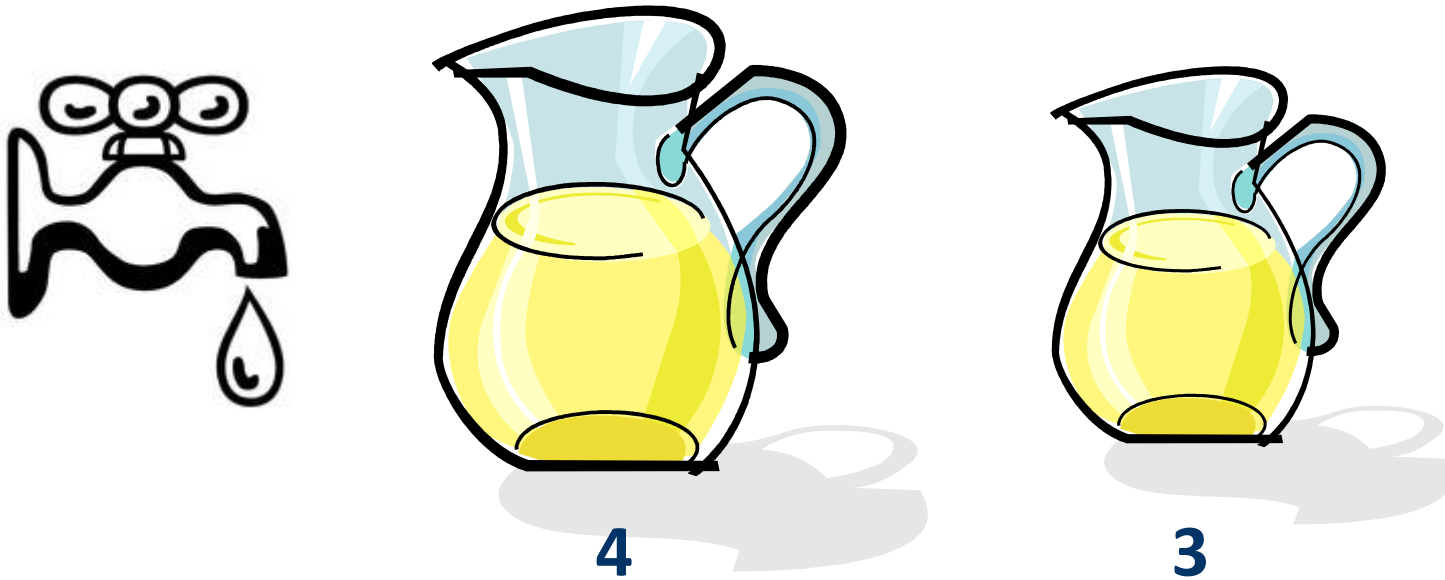
	1	2
3	4	5
6	7	8

Goal State

- **States = configurations**
- **Actions = up to 4 kinds of moves: up, down, left, right**

# Water Jugs Problem

Given 4-liter and 3-liter pitchers, how do you get exactly 2 liters into the 4-liter pitcher?



State:  $(x, y)$  for # liters in 4-liter and 3-liter pitchers, respectively

Actions: empty, fill, pour water between pitchers

Initial state:  $(0, 0)$

Goal state:  $(2, *)$

# Actions / Successor Functions

1.  $(x, y / x < 4) \rightarrow (4, y)$       “Fill 4”
2.  $(x, y / y < 3) \rightarrow (x, 3)$       “Fill 3”
3.  $(x, y / x > 0) \rightarrow (0, y)$       “Empty 4”
4.  $(x, y / y > 0) \rightarrow (x, 0)$       “Empty 3”
5.  $(x, y / x+y \geq 4 \text{ and } y > 0) \longrightarrow (4, y - (4 - x))$   
    “Pour from 3 to 4 until 4 is full”
6.  $(x, y / x+y \geq 3 \text{ and } x > 0) \longrightarrow (x - (3 - y), 3)$   
    “Pour from 4 to 3 until 3 is full”
7.  $(x, y / x+y \leq 4 \text{ and } y > 0) \longrightarrow (x+y, 0)$   
    “Pour all water from 3 to 4”

# Formalizing Search in a State Space

- A **state space** is a directed *graph*:  $(V, E)$ 
  - $V$  is a set of nodes (vertices)
  - $E$  is a set of arcs (edges)
    - each arc is *directed* from one node to another node
- Each **node** is a data structure that contains:
  - a **state** description
  - other information such as:
    - link to parent node
    - name of action that generated this node (from its parent)
    - other bookkeeping data

# Formalizing Search in a State Space

- Each **arc** corresponds to one of the finite number of actions:
  - when the action is applied to the state associated with the arc's source node
  - then the resulting state is the state associated with the arc's destination node
- Each arc has a fixed, positive **cost**:
  - corresponds to the cost of the action

# Formalizing Search in a State Space

- **Each node has a finite set of **successor** nodes:**
  - corresponds to all of the legal actions that can be applied at the source node's state
- **Expanding a node means:**
  - generate *all* of the successor nodes
  - add them and their associated arcs to the state-space search tree

# Formalizing Search in a State Space

- One or more nodes are designated as **start** nodes
- A **goal test** is applied to a node's state to determine if it is a goal node
- A **solution** is a sequence of actions associated with a path in the state space from a start to a goal node:
  - just the goal state (e.g., cryptarithmic)
  - a path from start to goal state (e.g., 8-puzzle)
- The **cost** of a solution is the sum of the arc costs on the solution path



# Search Summary

- **Solution is an ordered sequence of primitive actions (steps)**  
 $f(x) = a_1, a_2, \dots, a_n$  where  $x$  is the input
- **Model task as a graph of all possible states and actions, and a solution as a path**
- **A state captures all the relevant information about the past**

# Sizes of State Spaces

Problem	Nodes	Brute-Force Search Time (10 million nodes/second)
• Tic-Tac-Toe	$3^9$	
• 8 Puzzle	$10^5$	.01 seconds
• $2^3$ Rubik's Cube	$10^6$	.2 seconds
• 15 Puzzle	$10^{13}$	6 days
• $3^3$ Rubik's Cube	$10^{19}$	68,000 years
• 24 Puzzle	$10^{25}$	12 billion years
• Checkers	$10^{40}$	
• Chess	$10^{120}$	

# Formalizing Search



F



C



D



S

A search problem has five components:

$\mathcal{S}$ ,  $I$ ,  $G$ , actions, cost



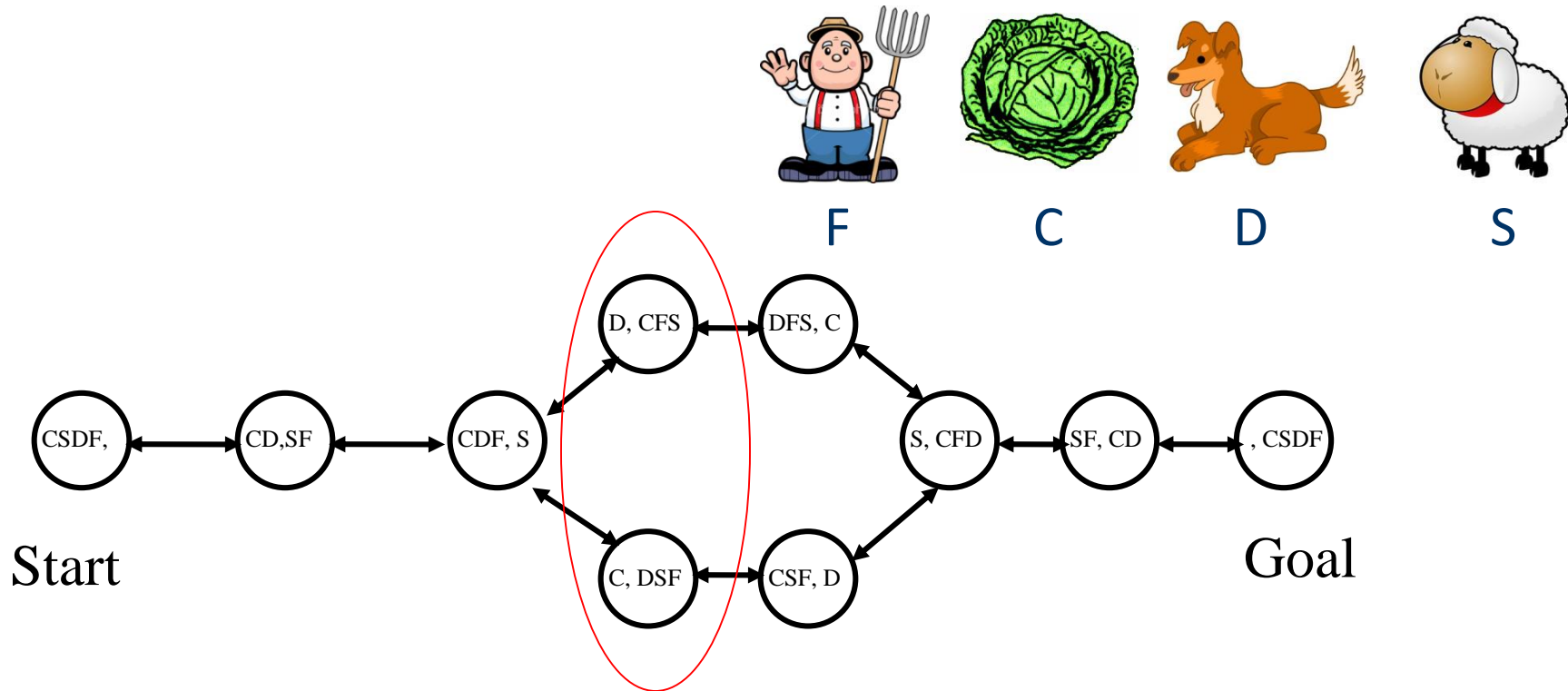
1. **State space**  $\mathcal{S}$ : all valid configurations
2. **Initial states**  $I \subseteq \mathcal{S}$ : a set of start states  $I = \{(FCDS,)\} \subseteq \mathcal{S}$
3. **Goal states**  $G \subseteq \mathcal{S}$ : a set of goal states  $G = \{(:,FCDS)\} \subseteq \mathcal{S}$
4. An **action function**  $successors(s) \subseteq \mathcal{S}$ : states reachable in one step (one arc) from  $s$

$successors((FCDS,)) = \{(CD,FS)\}$

$successors((CDF,S)) = \{(CD,FS), (D,FCS), (C,FSD)\}$

5. A **cost function**  $cost(s, s')$ : The cost of moving from  $s$  to  $s'$
- The goal of search is to find a solution path from a state in  $I$  to a state in  $G$

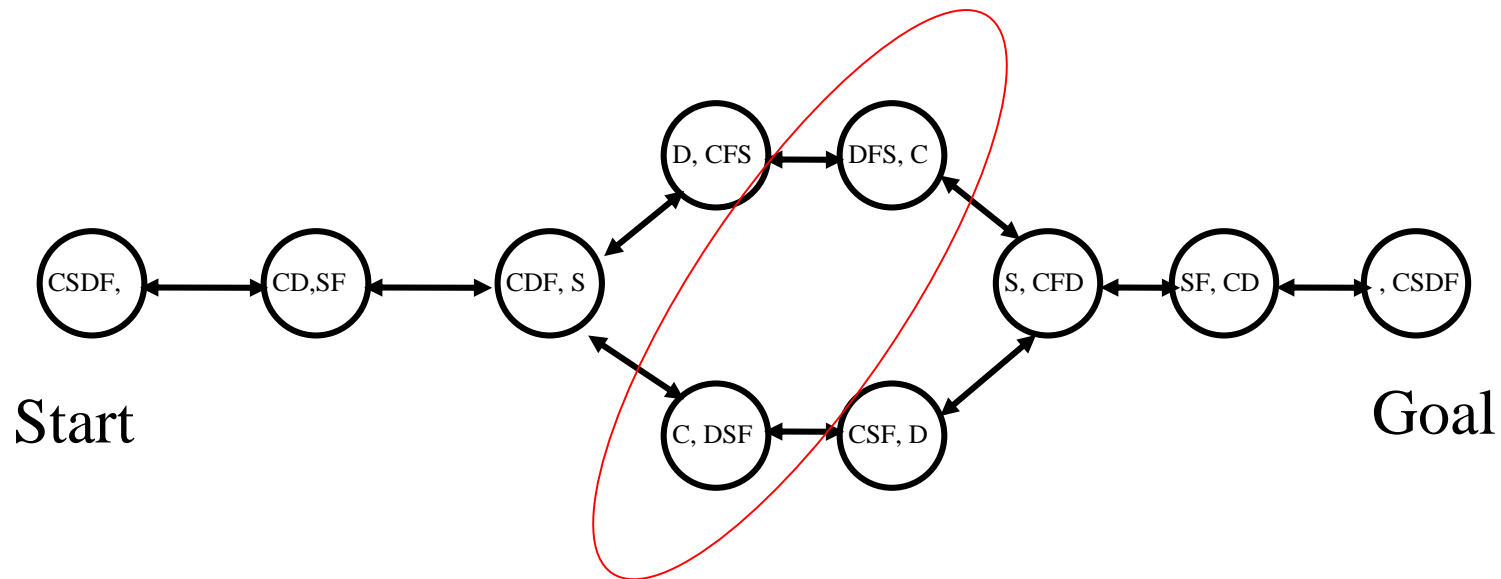
# State Space = A Directed Graph



- In general there will be many generated, but unexpanded, states at any given time
- One has to choose which one to “expand” next

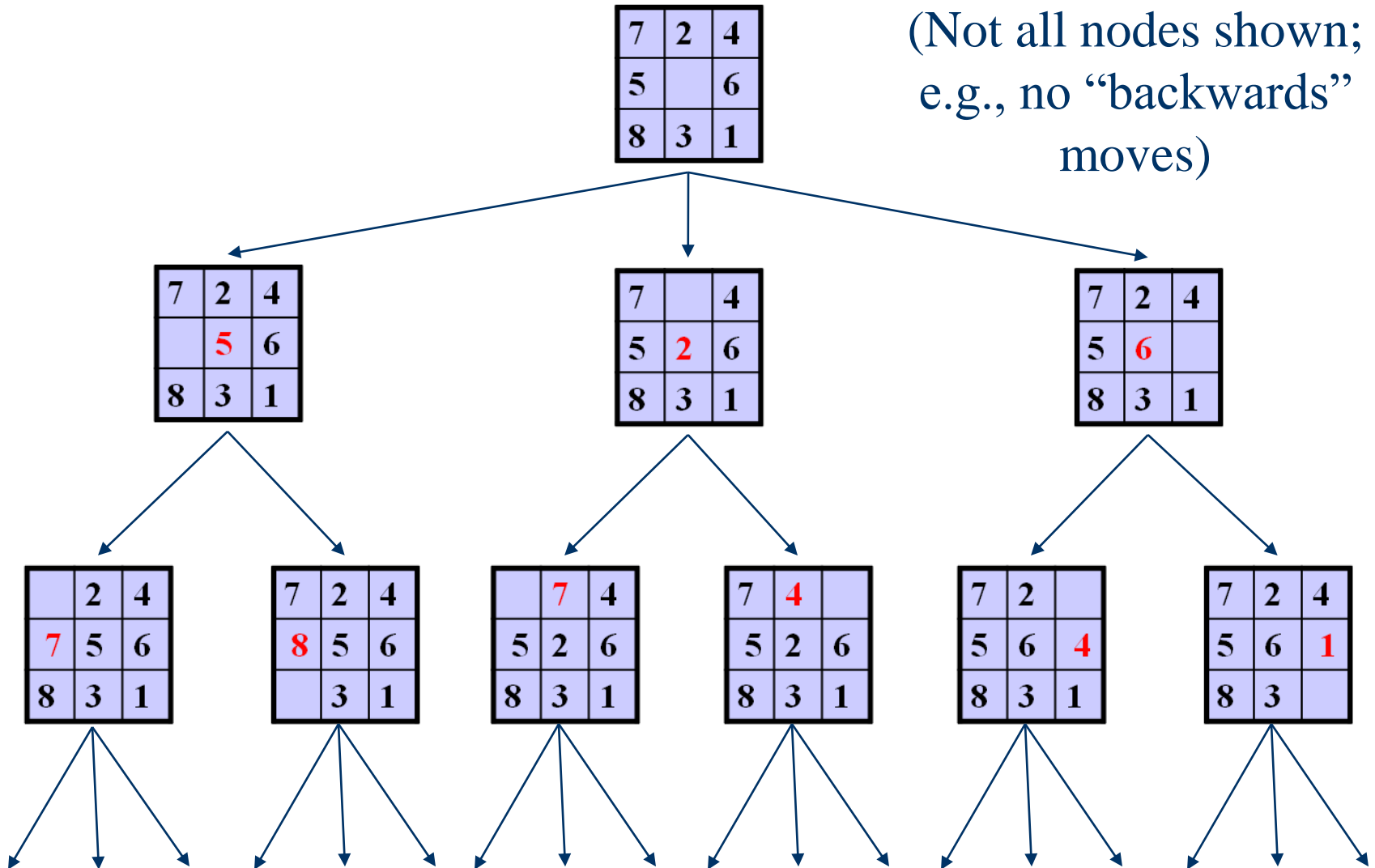
# Different Search Strategies

- The generated, but not yet expanded, states define the *Frontier* (aka *Open* or *Fringe*) set
- The essential difference is, **which one to expand first?**



# 8-Puzzle State-Space Search Tree

(Not all nodes shown;  
e.g., no “backwards”  
moves)



# Uninformed Search Strategies

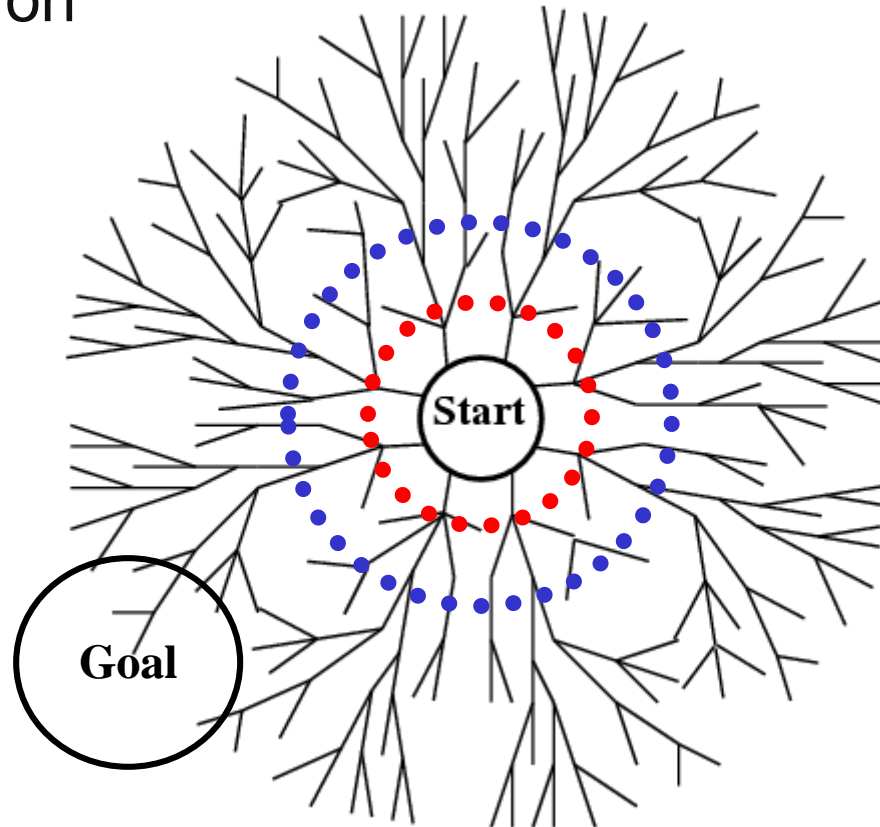
**Uninformed Search:** strategies that order nodes *without* using any domain specific information, i.e., don't use any information stored in a state

- **BFS: breadth-first search**
  - Queue (*FIFO*) used for the Frontier
  - remove from front, add to **back**
- **DFS: depth-first search**
  - Stack (*LIFO*) used for the Frontier
  - remove from front, add to **front**

# Breadth-First Search (BFS)

Expand the shallowest node first:

1. Examine states **one** step away from the initial states
2. Examine states **two** steps away from the initial states
3. and so on



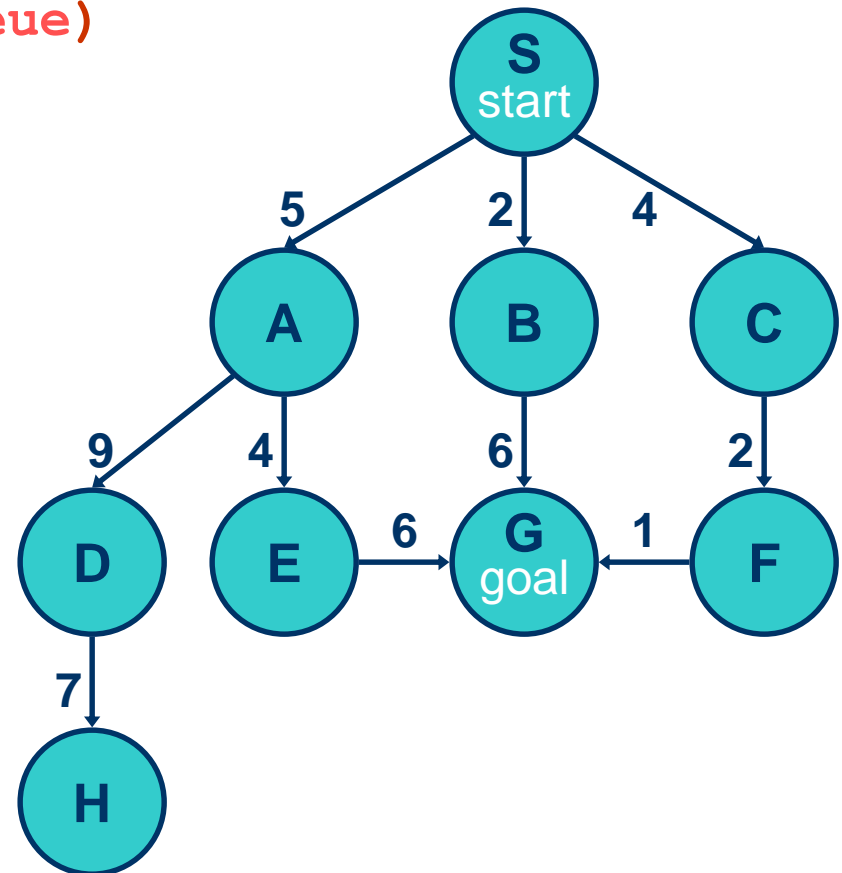


# Breadth-First Search (BFS)

**generalSearch(problem, queue)**

# of nodes tested: 0, expanded: 0

expnd. node	Frontier list
	{S}

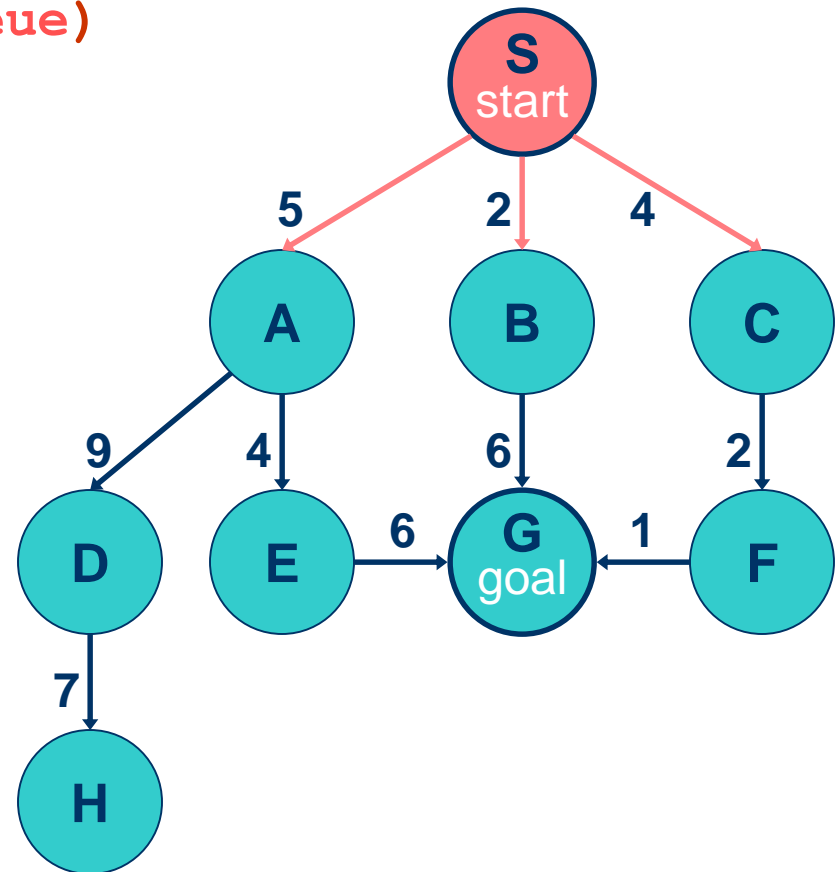


# Breadth-First Search (BFS)

**generalSearch(problem, queue)**

# of nodes tested: 1, expanded: 1

expnd. node	Frontier list
	{S}
S not goal	{A,B,C}

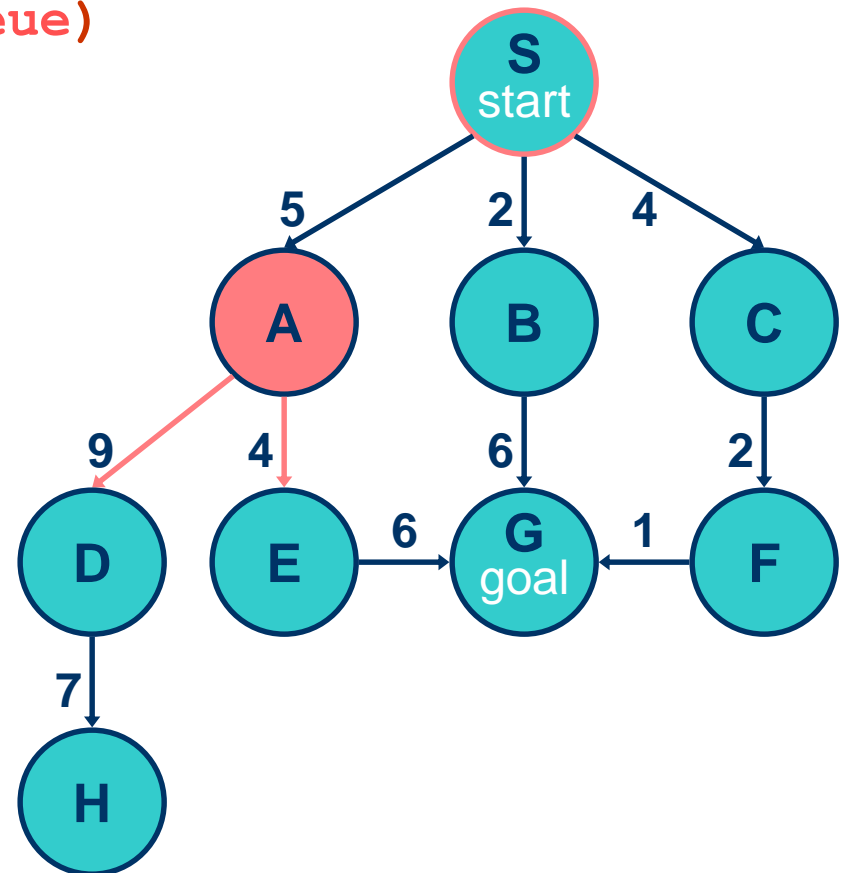


# Breadth-First Search (BFS)

`generalSearch(problem, queue)`

# of nodes tested: 2, expanded: 2

expnd. node	Frontier list
	{S}
S	{A,B,C}
A not goal	{B,C,D,E}

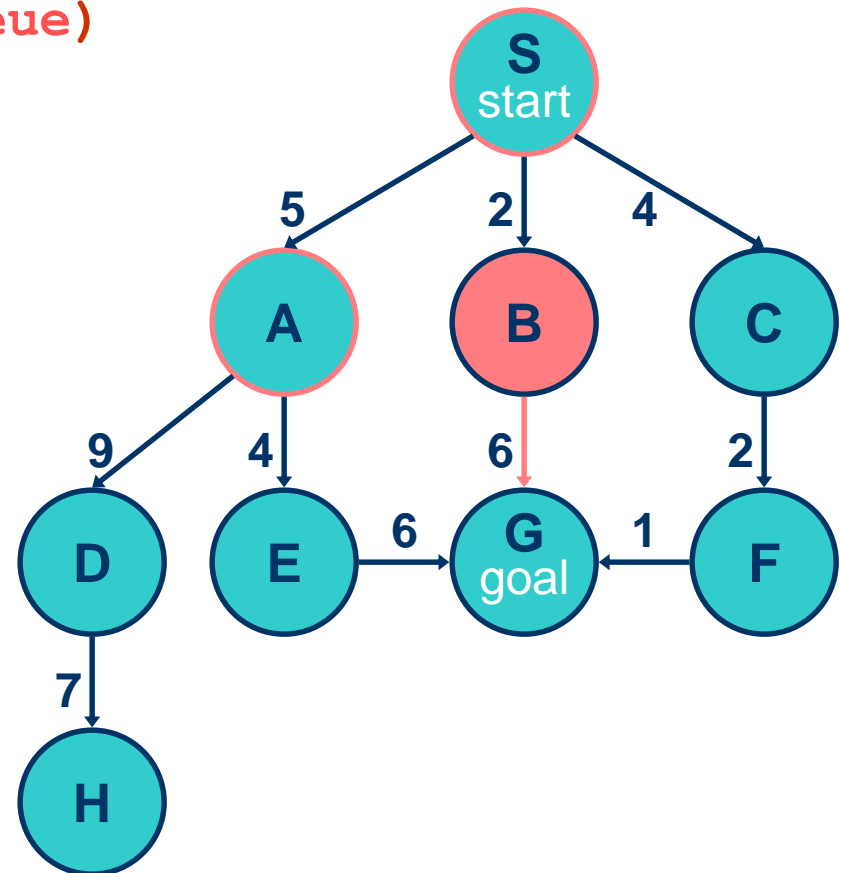


# Breadth-First Search (BFS)

`generalSearch(problem, queue)`

# of nodes tested: 3, expanded: 3

expnd. node	Frontier list
	{S}
S	{A,B,C}
A	{B,C,D,E}
B not goal	{C,D,E,G}

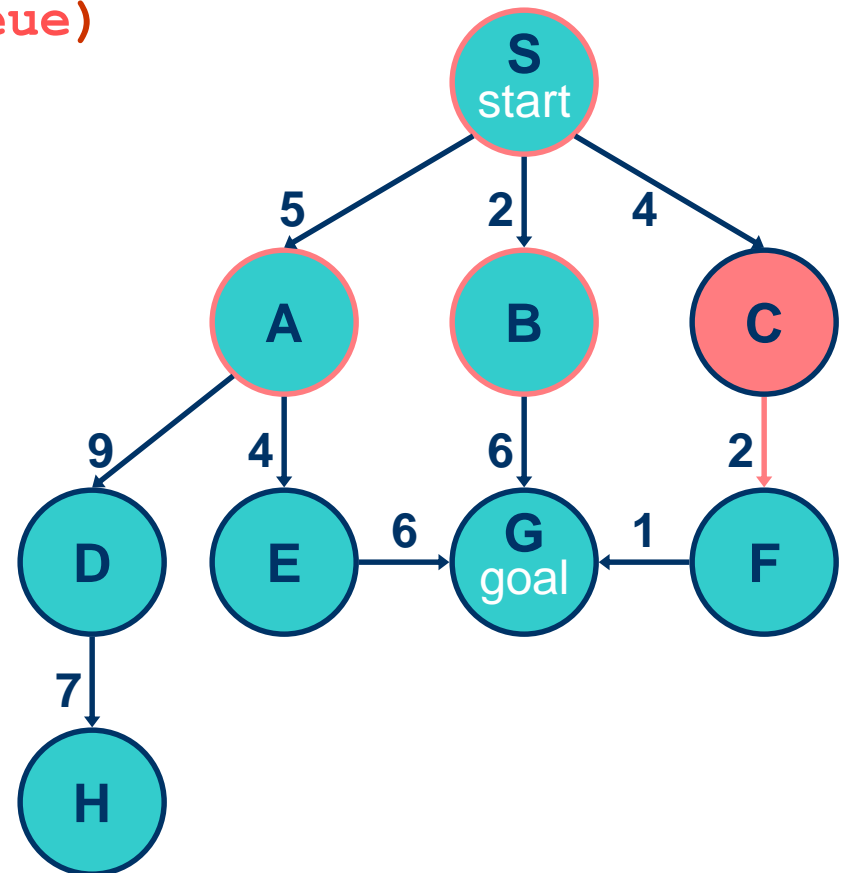


# Breadth-First Search (BFS)

`generalSearch(problem, queue)`

# of nodes tested: 4, expanded: 4

expnd. node	Frontier list
	{S}
S	{A,B,C}
A	{B,C,D,E}
B	{C,D,E,G}
C not goal	{D,E,G,F}

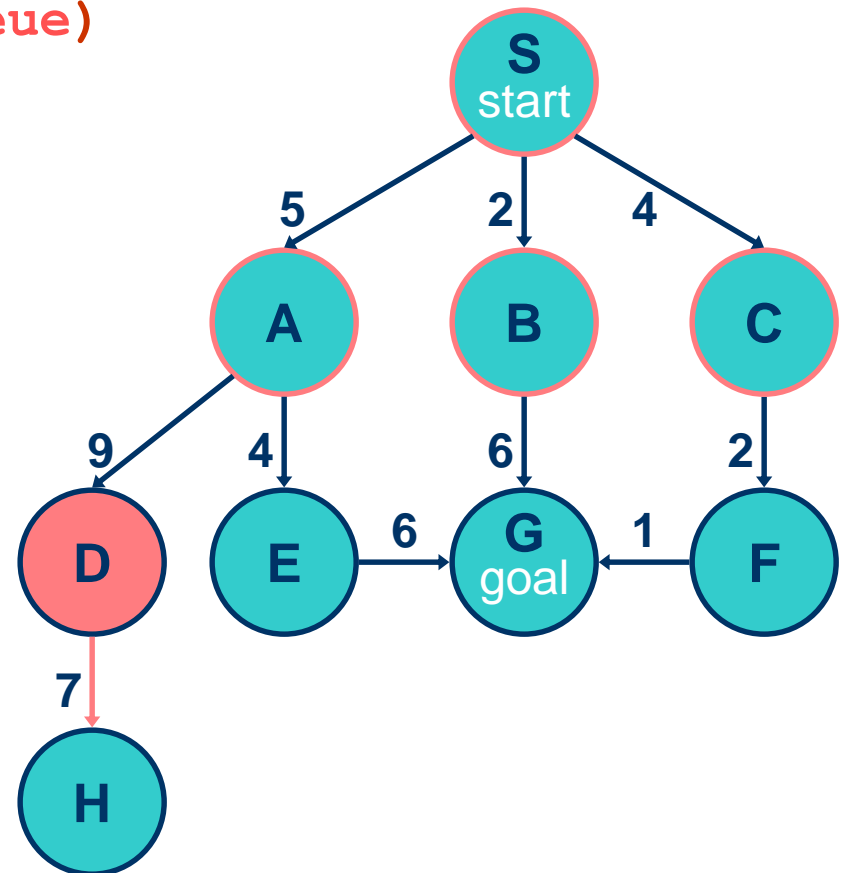


# Breadth-First Search (BFS)

`generalSearch(problem, queue)`

# of nodes tested: 5, expanded: 5

expnd. node	Frontier list
	{S}
S	{A,B,C}
A	{B,C,D,E}
B	{C,D,E,G}
C	{D,E,G,F}
D not goal	{E,G,F,H}

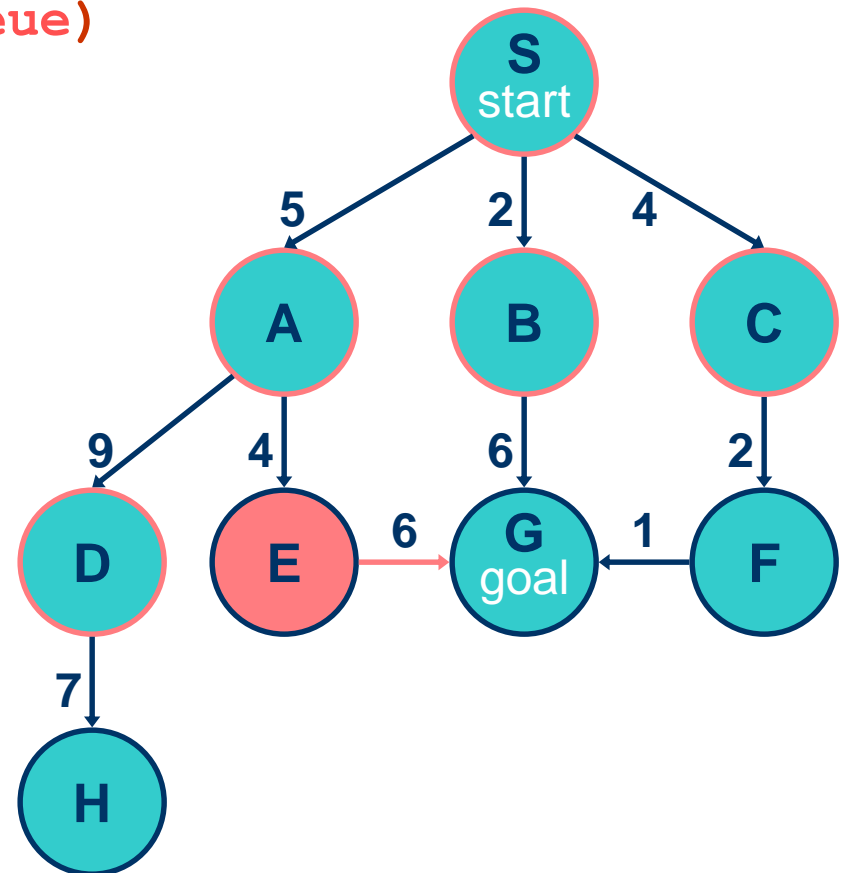


# Breadth-First Search (BFS)

**generalSearch(problem, queue)**

# of nodes tested: 6, expanded: 6

expnd. node	Frontier list
	{S}
S	{A,B,C}
A	{B,C,D,E}
B	{C,D,E,G}
C	{D,E,G,F}
D	{E,G,F,H}
E not goal	{G,F,H,G}

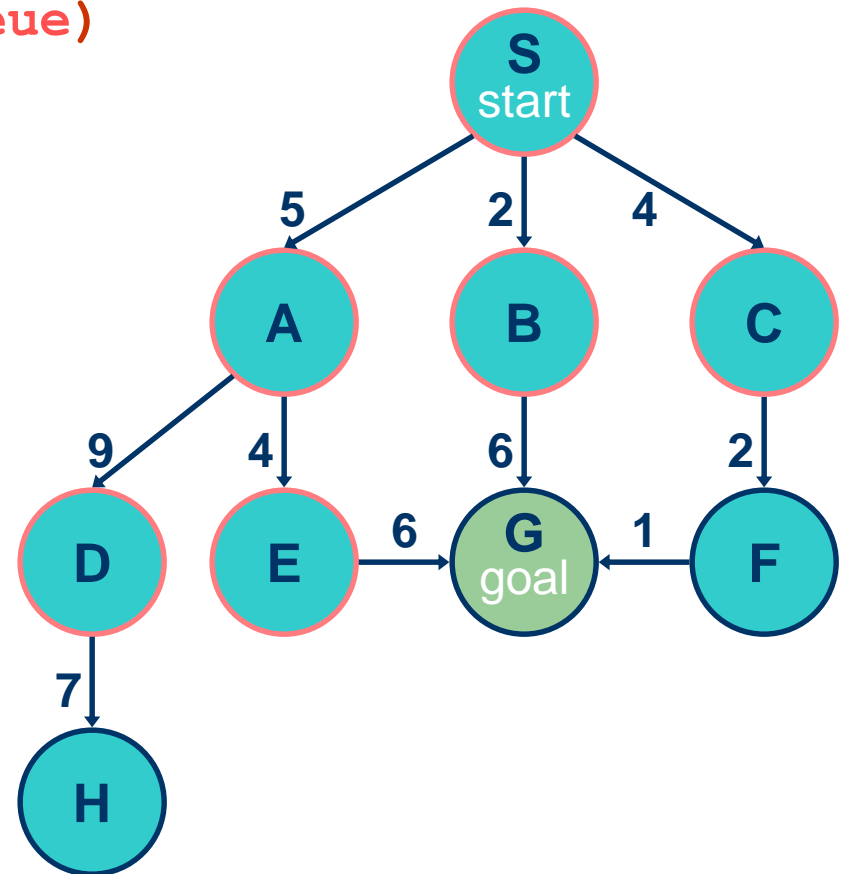


# Breadth-First Search (BFS)

**generalSearch(problem, queue)**

# of nodes tested: 7, expanded: 6

expnd. node	Frontier list
	{S}
S	{A,B,C}
A	{B,C,D,E}
B	{C,D,E,G}
C	{D,E,G,F}
D	{E,G,F,H}
E	{G,F,H,G}
G goal	{F,H,G} no expand



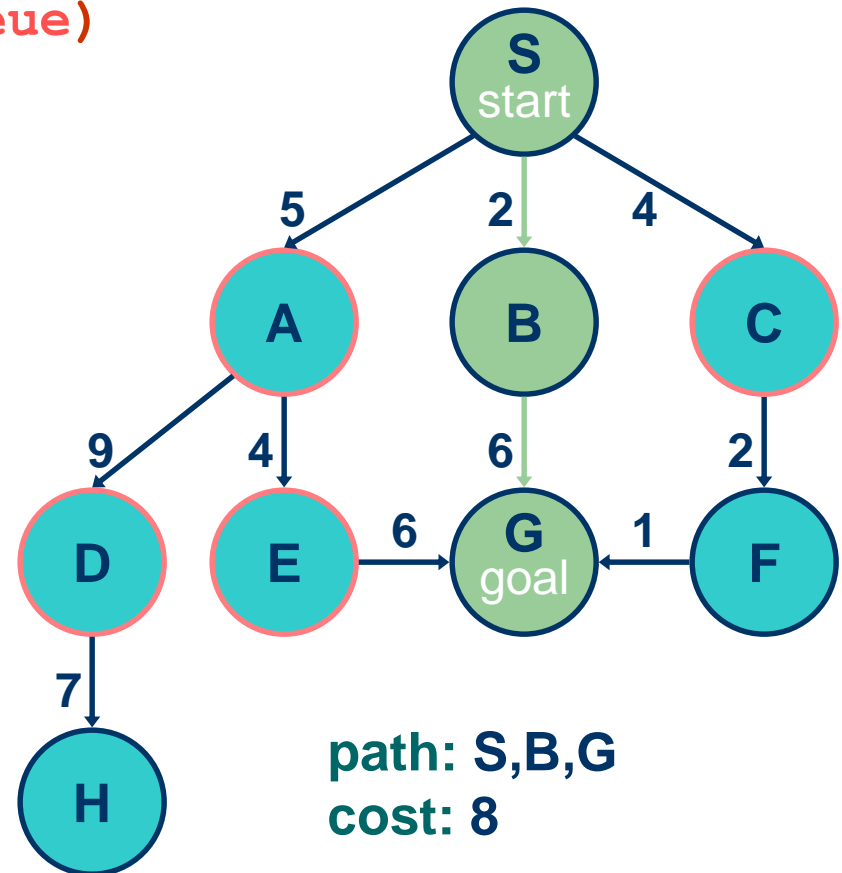


# Breadth-First Search (BFS)

**generalSearch(problem, queue)**

# of nodes tested: 7, expanded: 6

expnd. node	Frontier list
	{S}
S	{A,B,C}
A	{B,C,D,E}
B	{C,D,E,G}
C	{D,E,G,F}
D	{E,G,F,H}
E	{G,F,H,G}
G	{F,H,G}



# Evaluating Search Strategies

- **Completeness**

If a solution exists, will it be found?

- a complete algorithm will find **a** solution (not all)

- **Optimality / Admissibility**

If a solution is found, is it guaranteed to be optimal?

- an admissible algorithm will find a **solution with minimum cost**

# Evaluating Search Strategies

- **Time Complexity**

How long does it take to find a solution?

- usually measured for worst case
- measured by counting **number of nodes expanded**

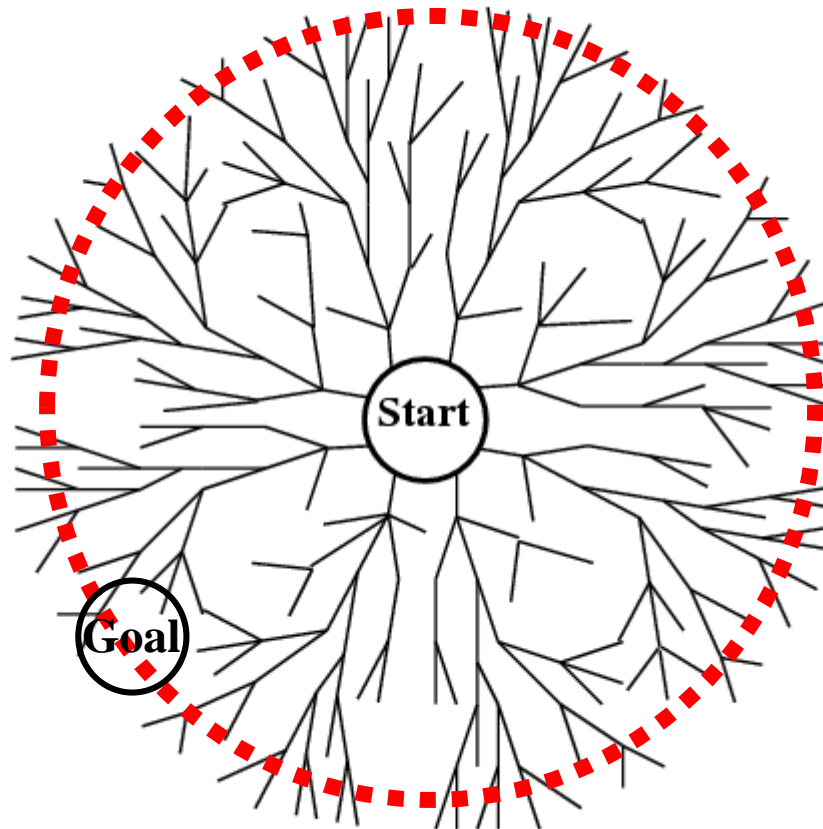
- **Space Complexity**

How much space is used by the algorithm?

- measured in terms of the **maximum size of the *Frontier*** during the search

# What's in the Frontier for BFS?

- If goal is at depth  $d$ , how big is the Frontier (worst case)?



# Breadth-First Search (BFS)

- **Complete**
- **Optimal / Admissible**
  - **Yes**, *if* all operators (i.e., arcs) have the same constant cost, or costs are positive, non-decreasing with depth
  - otherwise, not optimal but *does* guarantee finding solution of shortest *length* (i.e., fewest arcs)

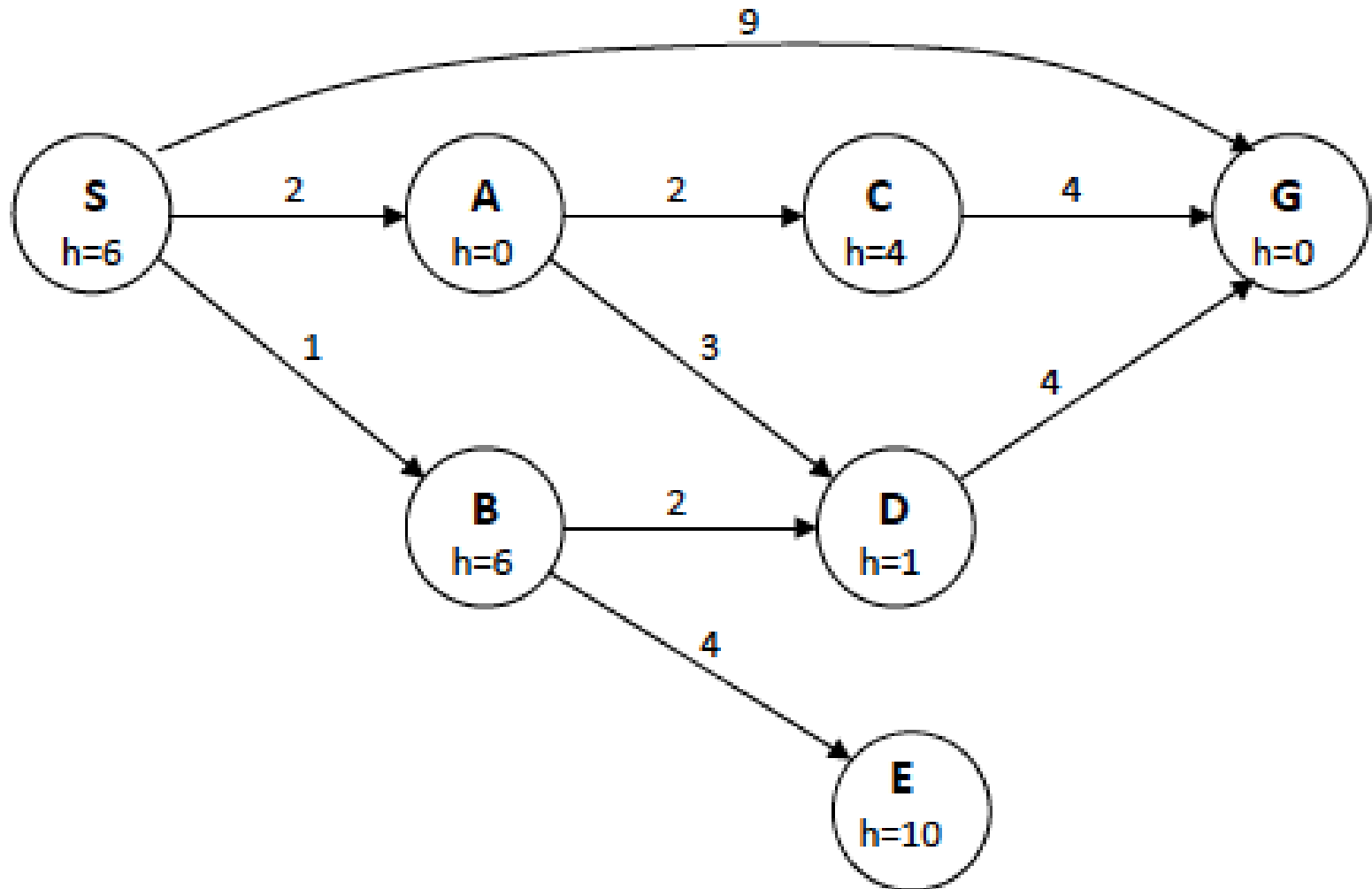
# Breadth-First Search (BFS)

- **Time and space complexity:  $O(b^d)$  (i.e., exponential)**
  - $d$  is the depth of the solution
  - $b$  is the branching factor at each non-leaf node
- Very slow to find solutions with a large number of steps because must look at *all* shorter length possibilities first

# Breadth-First Search (BFS)

- **A complete search tree has a total # of nodes =**  
 $1 + b + b^2 + \dots + b^d = (b^{(d+1)} - 1) / (b-1)$ 
  - $d$ : the tree's depth
  - $b$ : the branching factor at each non-leaf node
- **For example:  $d = 12, b = 10$**   
 $1 + 10 + 100 + \dots + 10^{12} = (10^{13} - 1)/9 = O(10^{12})$ 
  - If BFS expands 1,000 nodes/sec and each node uses 100 bytes of storage, then BFS will take 35 years to run in the worst case, and it will use 111 terabytes of memory!

## Problem: Given State Space



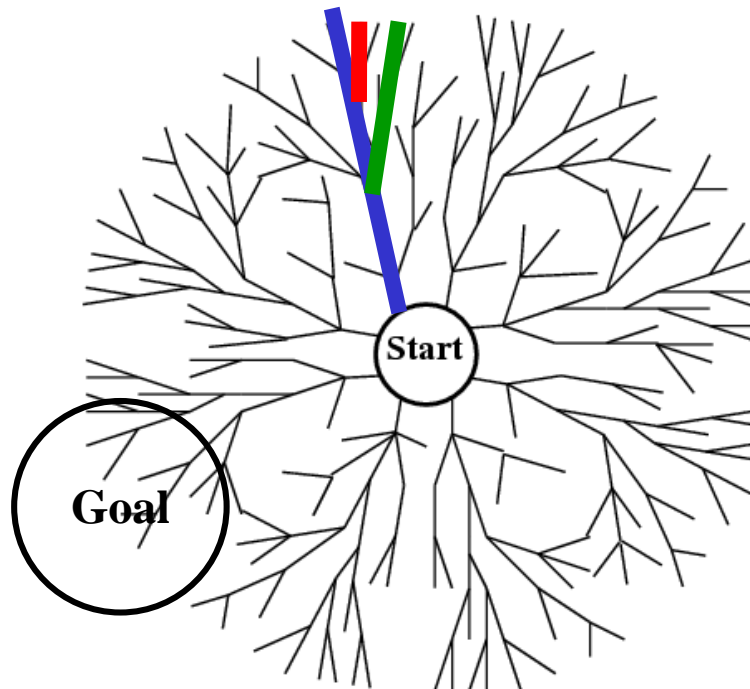


# Depth-First Search

Expand the ***deepest*** node first

1. Select a direction, go deep to the end —
2. Slightly change the end —
3. Slightly change the end some more... —

Use a **Stack** to order nodes on the ***Frontier***

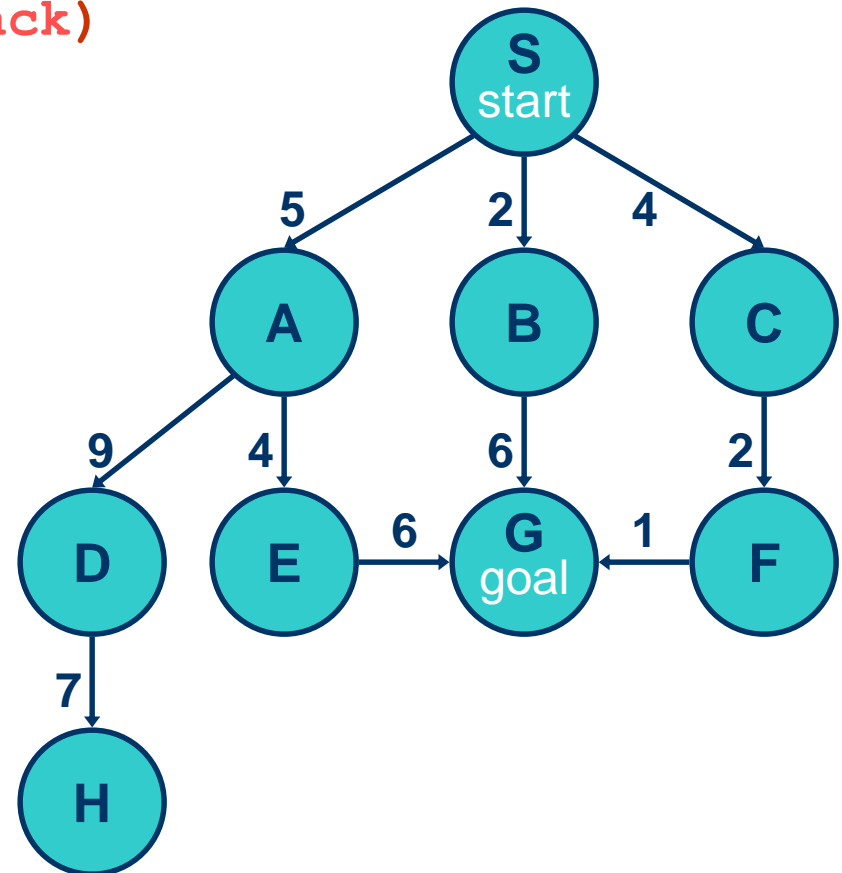


# Depth-First Search (DFS)

`generalSearch(problem, stack)`

# of nodes tested: 0, expanded: 0

expnd. node	Frontier
	{S}

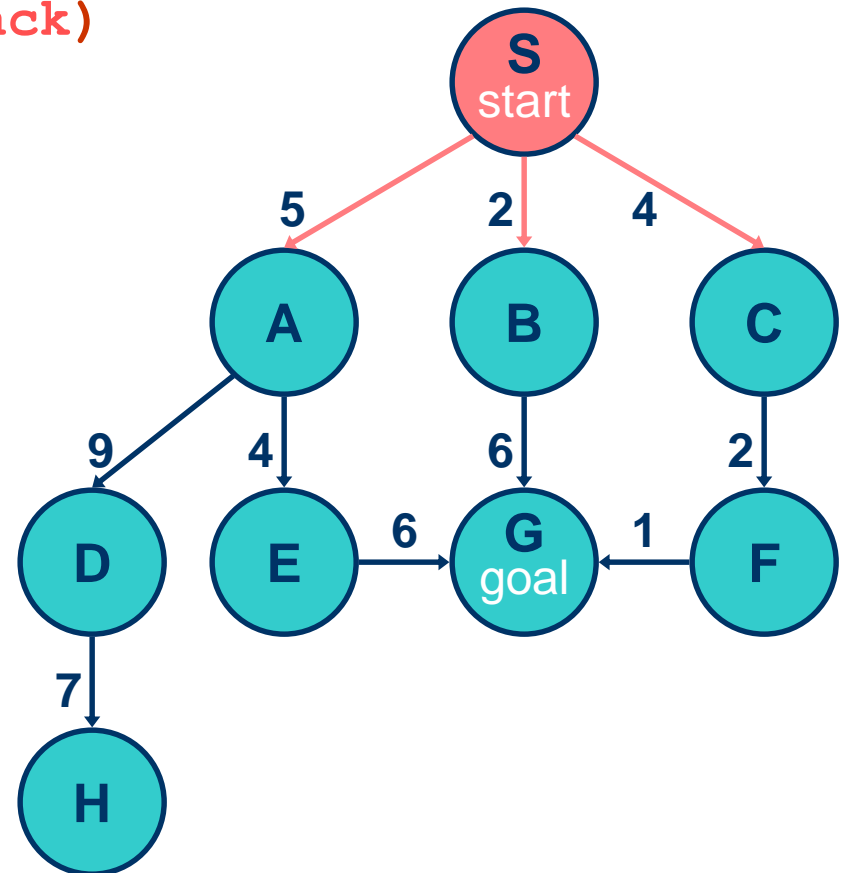


# Depth-First Search (DFS)

`generalSearch(problem, stack)`

# of nodes tested: 1, expanded: 1

expnd. node	Frontier
	{S}
S not goal	{A,B,C}

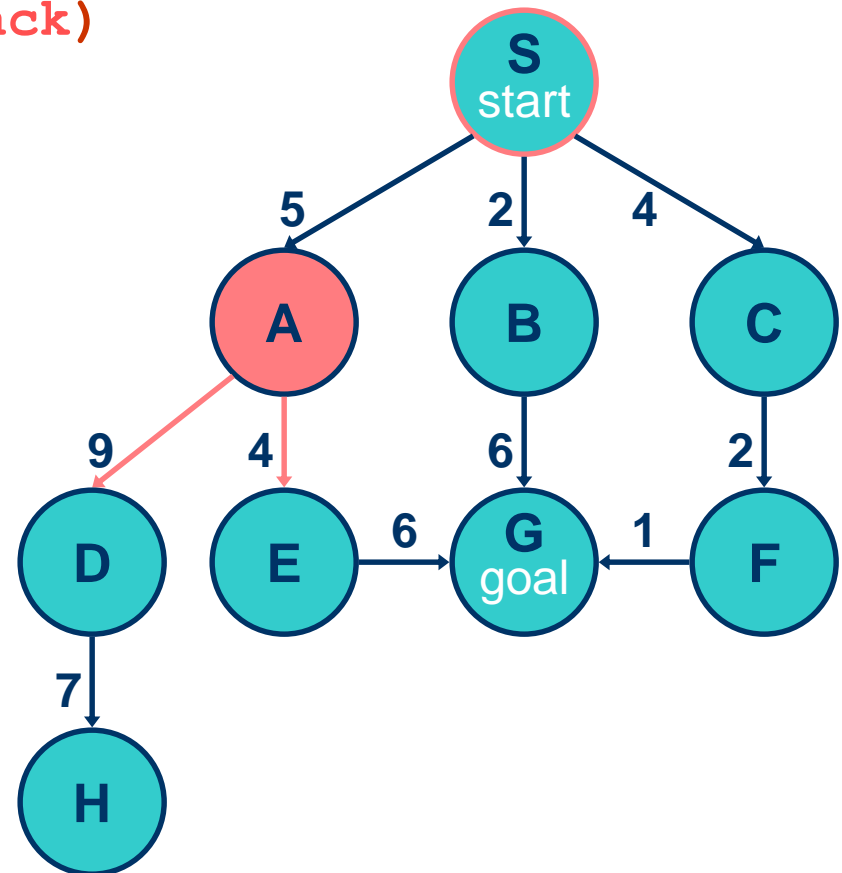


# Depth-First Search (DFS)

`generalSearch(problem, stack)`

# of nodes tested: 2, expanded: 2

expnd. node	Frontier
	{S}
S	{A,B,C}
A not goal	{D,E,B,C}

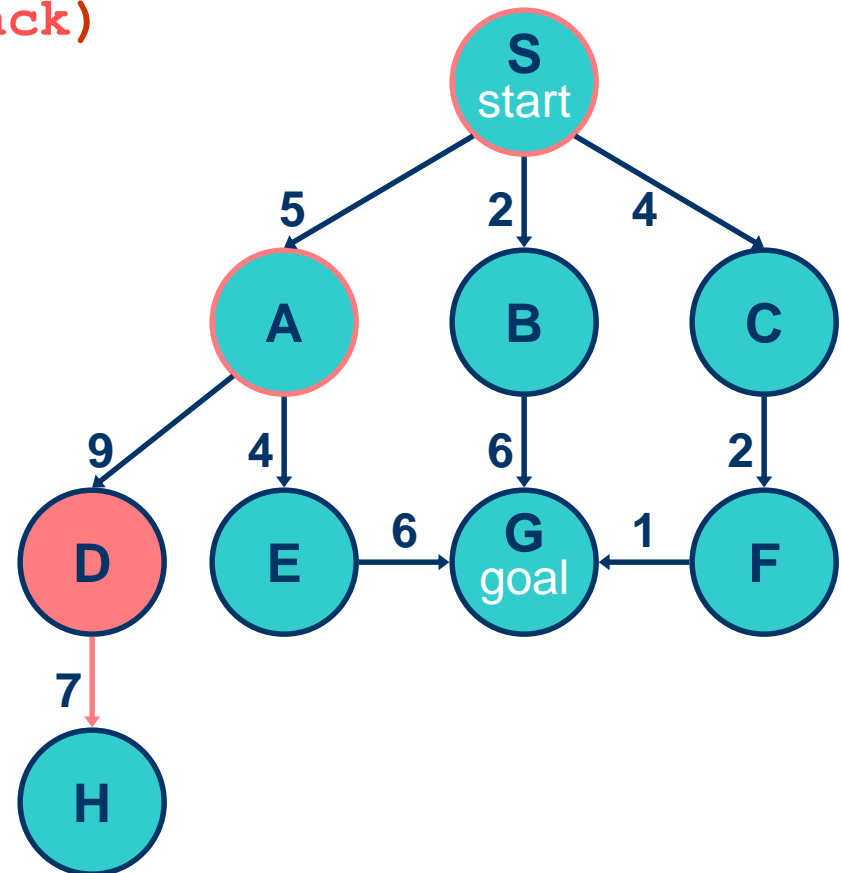


# Depth-First Search (DFS)

`generalSearch(problem, stack)`

# of nodes tested: 3, expanded: 3

expnd. node	Frontier
	{S}
S	{A,B,C}
A	{D,E,B,C}
D not goal	{ <b>H</b> ,E,B,C}

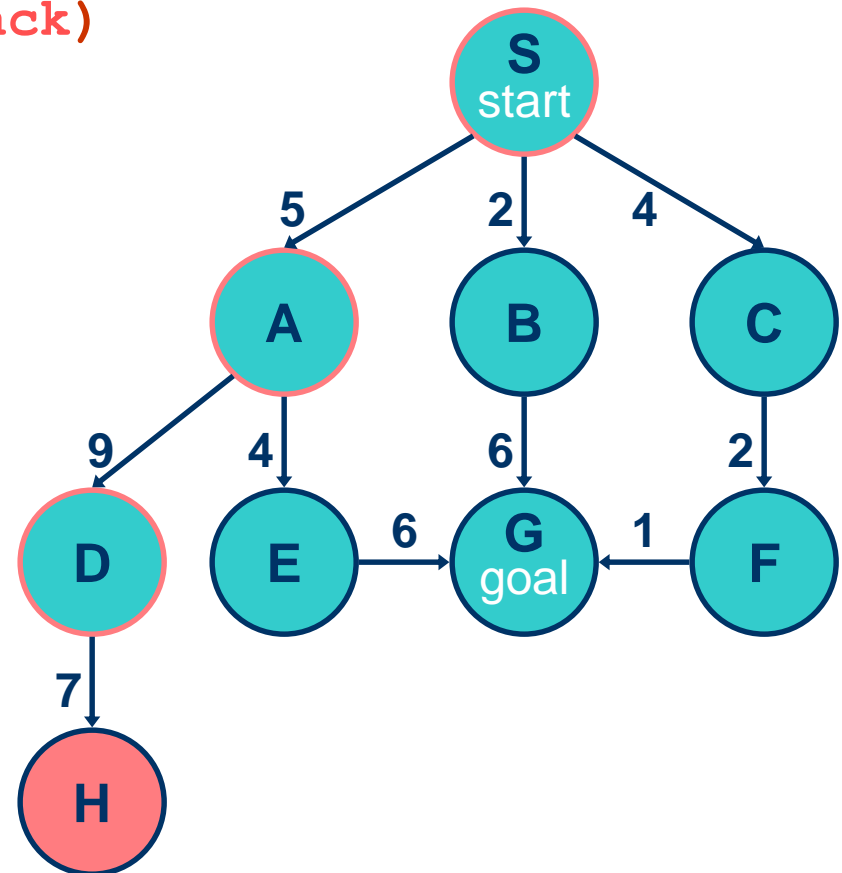


# Depth-First Search (DFS)

`generalSearch(problem, stack)`

# of nodes tested: 4, expanded: 4

expnd. node	Frontier
	{S}
S	{A,B,C}
A	{D,E,B,C}
D	{H,E,B,C}
H not goal	{E,B,C}

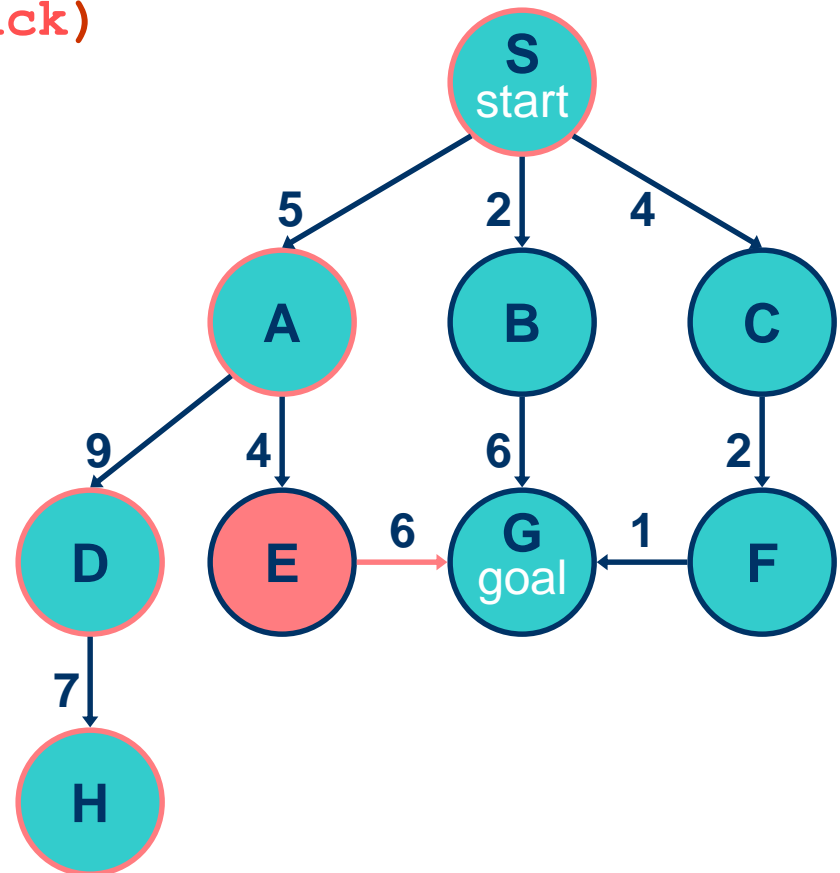


# Depth-First Search (DFS)

`generalSearch(problem, stack)`

# of nodes tested: 5, expanded: 5

expnd. node	Frontier
	{S}
S	{A,B,C}
A	{D,E,B,C}
D	{H,E,B,C}
H	{E,B,C}
E not goal	{G,B,C}

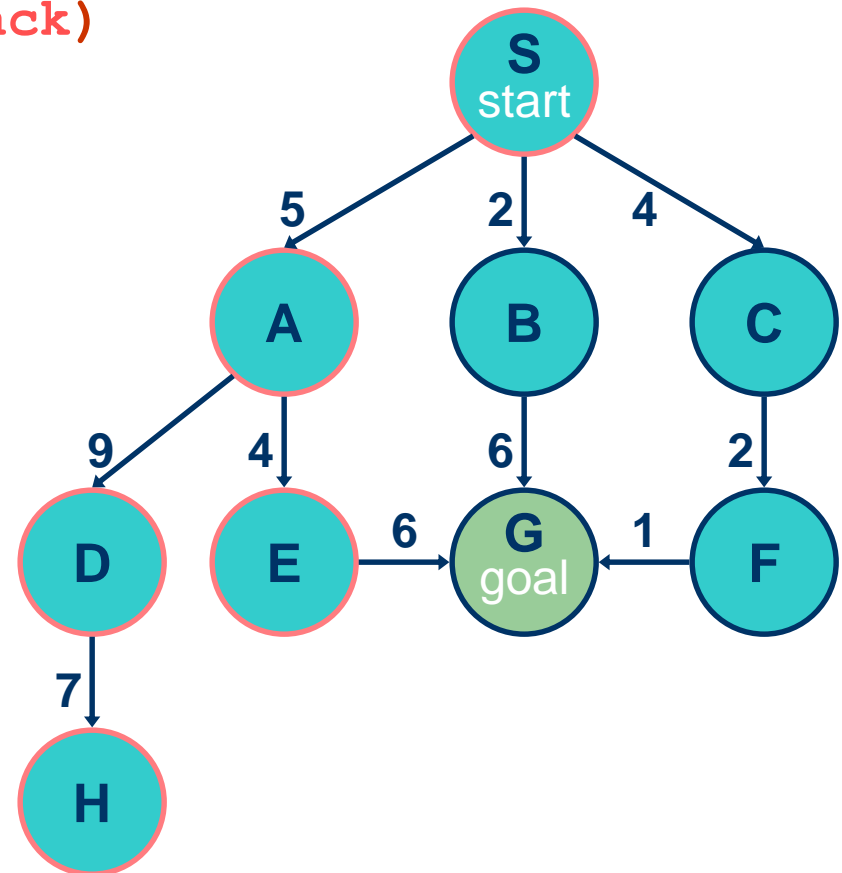


# Depth-First Search (DFS)

`generalSearch(problem, stack)`

# of nodes tested: 6, expanded: 5

expnd. node	Frontier
	{S}
S	{A,B,C}
A	{D,E,B,C}
D	{H,E,B,C}
H	{E,B,C}
E	{G,B,C}
G goal	{B,C} no expand



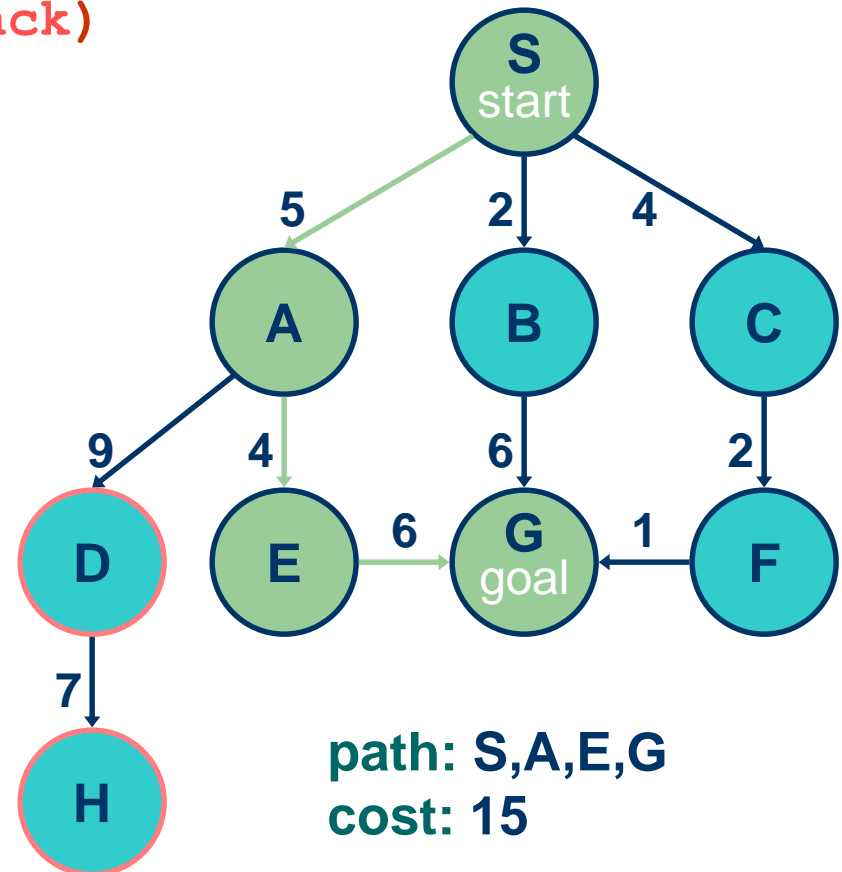


# Depth-First Search (DFS)

`generalSearch(problem, stack)`

# of nodes tested: 6, expanded: 5

expnd. node	Frontier
	{S}
S	{A,B,C}
A	{D,E,B,C}
D	{H,E,B,C}
H	{E,B,C}
E	{G,B,C}
G	{B,C}



# Depth-First Search (DFS)

- May not terminate without a **depth bound**  
i.e., cutting off search below a fixed depth,  $D$
- **Not complete**
  - with or without cycle detection
  - and, with or without a depth cutoff
- **Not optimal / admissible**
- *Can find long solutions quickly if lucky*

# Depth-First Search (DFS)

- **Time complexity:**  $O(b^d)$  exponential  
**Space complexity:**  $O(bd)$  linear
  - $d$  is the depth of the solution
  - $b$  is the branching factor at each non-leaf node
- Performs “**chronological backtracking**”
  - i.e., when search hits a dead end, backs up *one* level at a time
  - problematic if the mistake occurs because of a bad action choice near the top of search tree

# Uniform-Cost Search (UCS)

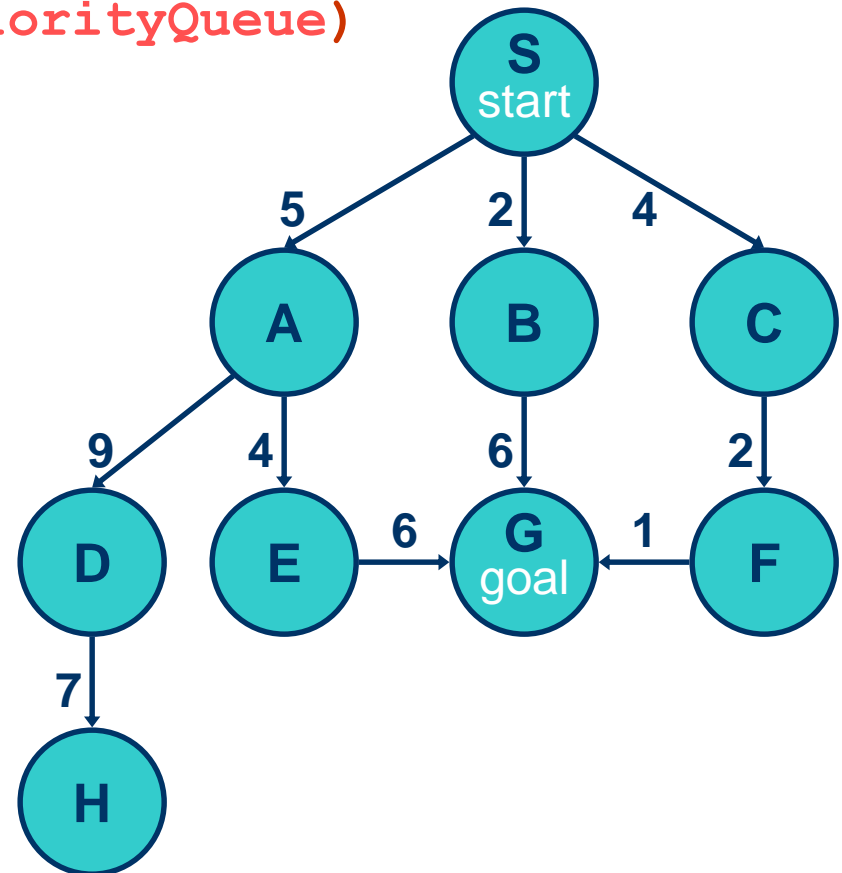
- Use a “**Priority Queue**” to order nodes on the *Frontier* list, sorted by path cost
- Let  $g(n)$  = cost of path from start node  $s$  to current node  $n$
- Sort nodes by increasing value of  $g$

# Uniform-Cost Search (UCS)

`generalSearch(problem, priorityQueue)`

# of nodes tested: 0, expanded: 0

expnd. node	Frontier list
	{S}

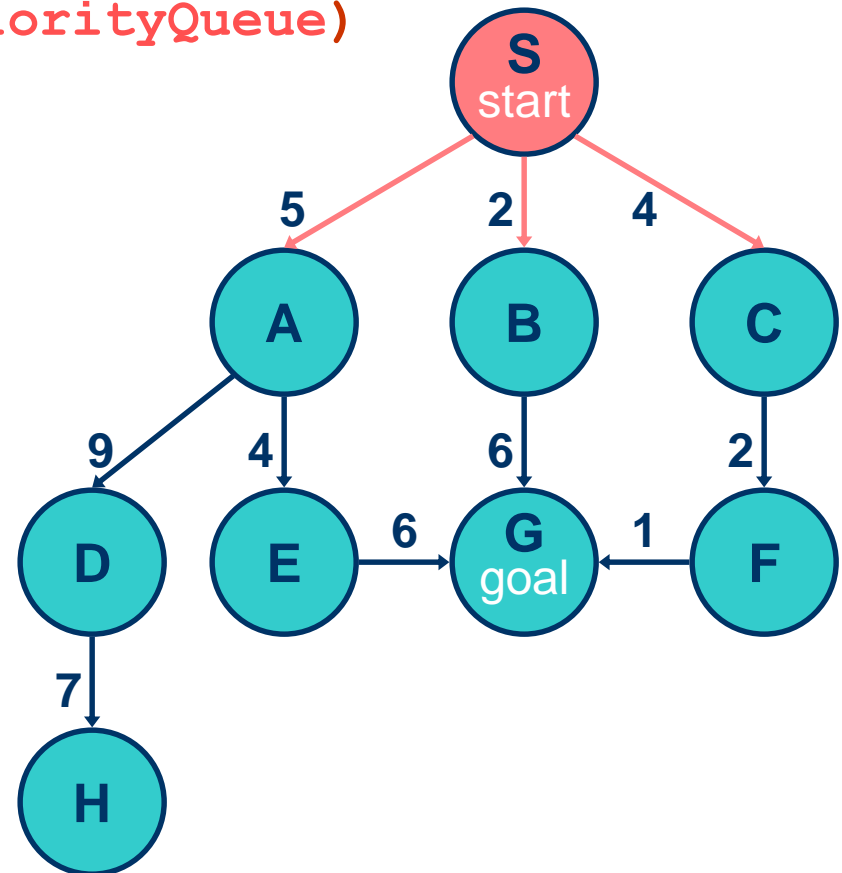


# Uniform-Cost Search (UCS)

`generalSearch(problem, priorityQueue)`

# of nodes tested: 1, expanded: 1

expnd. node	Frontier list
	{S:0}
S not goal	{B:2,C:4,A:5}

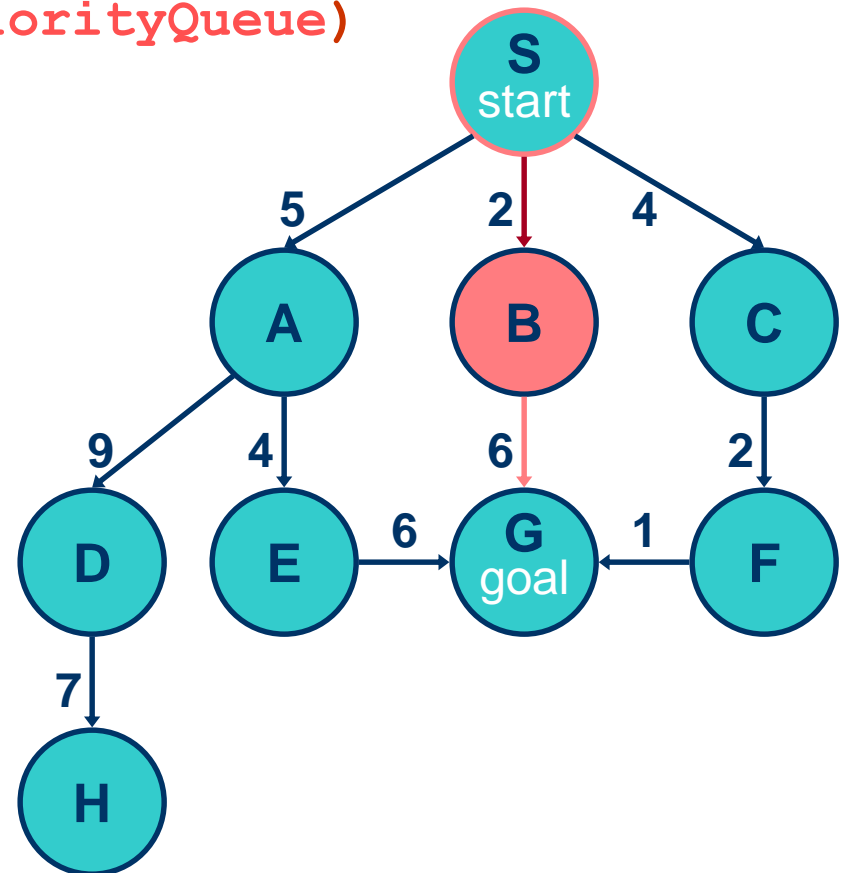


# Uniform-Cost Search (UCS)

`generalSearch(problem, priorityQueue)`

# of nodes tested: 2, expanded: 2

expnd. node	Frontier list
	{S}
S	{B:2,C:4,A:5}
B not goal	{C:4,A:5,G:2+6}

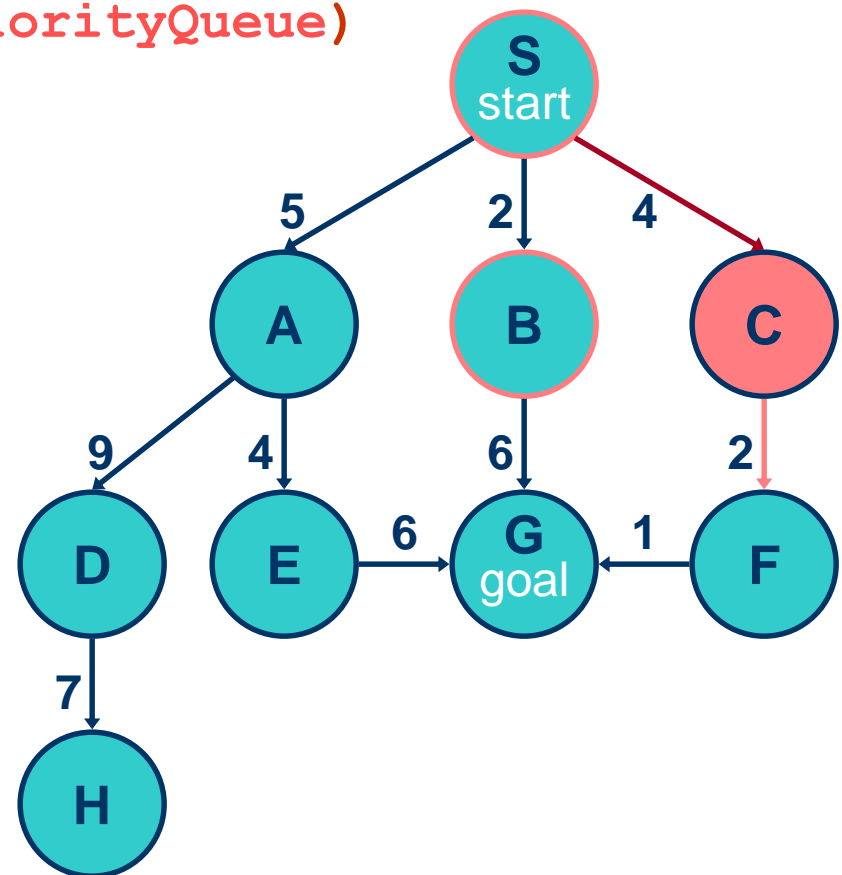


# Uniform-Cost Search (UCS)

`generalSearch(problem, priorityQueue)`

# of nodes tested: 3, expanded: 3

expnd. node	Frontier list
	{S}
S	{B:2,C:4,A:5}
B	{C:4,A:5,G:8}
C not goal	{A:5,F:4+2,G:8}



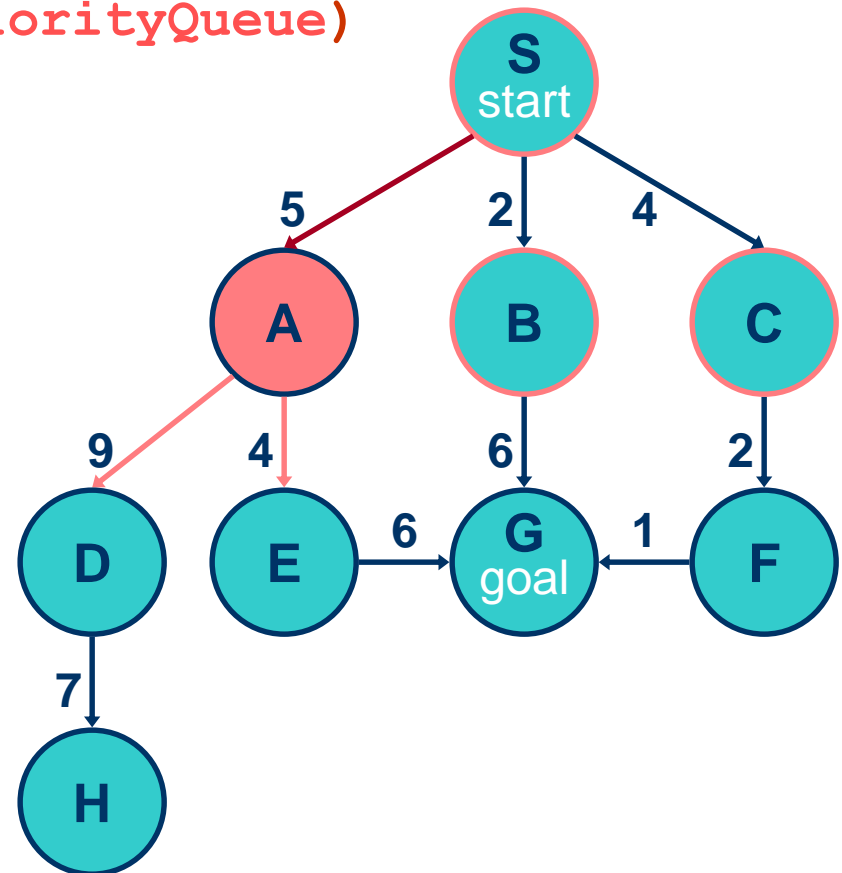


# Uniform-Cost Search (UCS)

`generalSearch(problem, priorityQueue)`

# of nodes tested: 4, expanded: 4

expnd. node	Frontier list
	{S}
S	{B:2,C:4,A:5}
B	{C:4,A:5,G:8}
C	{A:5,F:6,G:8}
A not goal	{F:6,G:8,E:5+4, D:5+9}

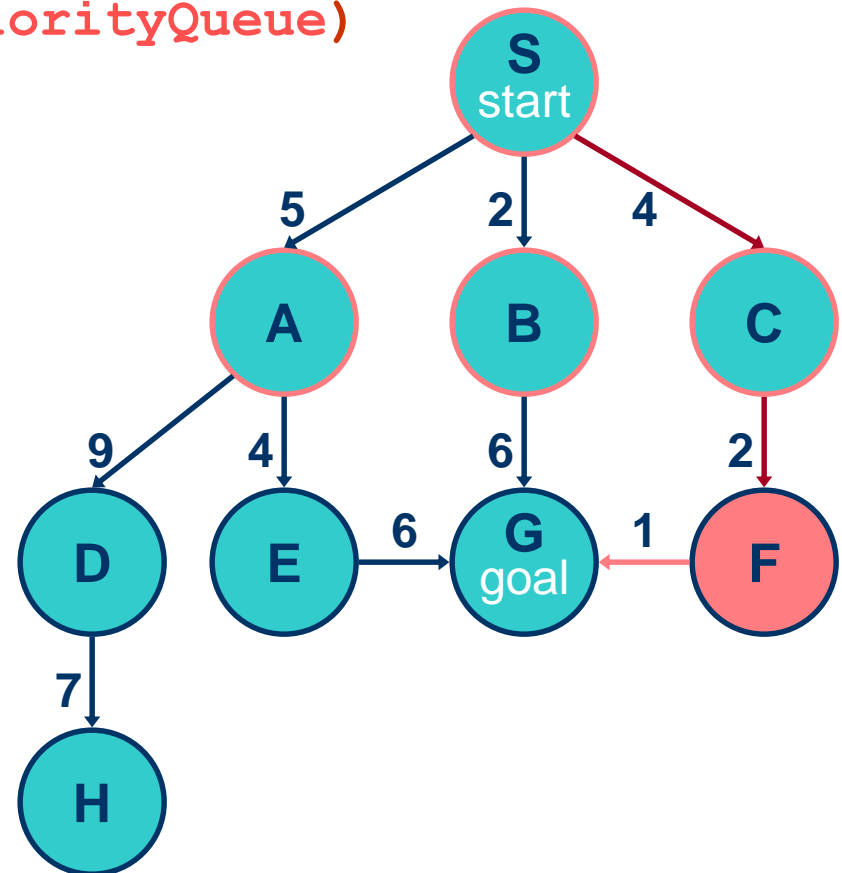


# Uniform-Cost Search (UCS)

`generalSearch(problem, priorityQueue)`

# of nodes tested: 5, expanded: 5

expnd. node	Frontier list
	{S}
S	{B:2,C:4,A:5}
B	{C:4,A:5,G:8}
C	{A:5,F:6,G:8}
A	{F:6,G:8,E:9,D:14}
F not goal	{G:4+2+1, G:8, E:9, D:14}

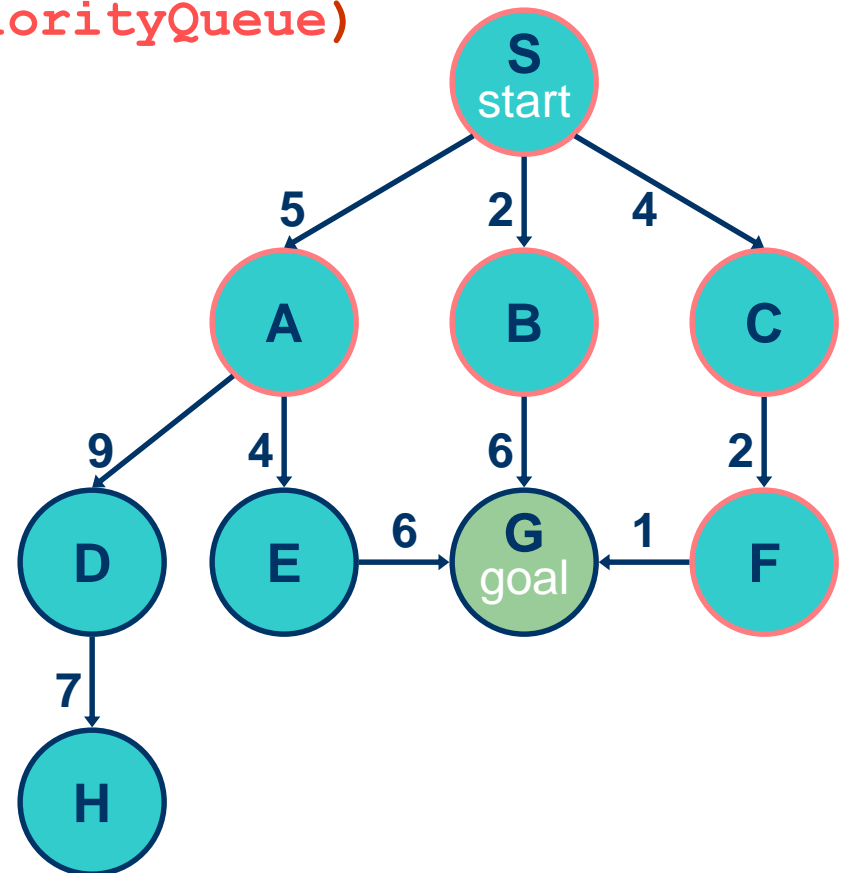


# Uniform-Cost Search (UCS)

`generalSearch(problem, priorityQueue)`

# of nodes tested: 6, expanded: 5

expnd. node	Frontier list
	{S}
S	{B:2,C:4,A:5}
B	{C:4,A:5,G:8}
C	{A:5,F:6,G:8}
A	{F:6,G:8,E:9,D:14}
F	{G:7,G:8,E:9,D:14}
G goal	{G:8,E:9,D:14} no expand

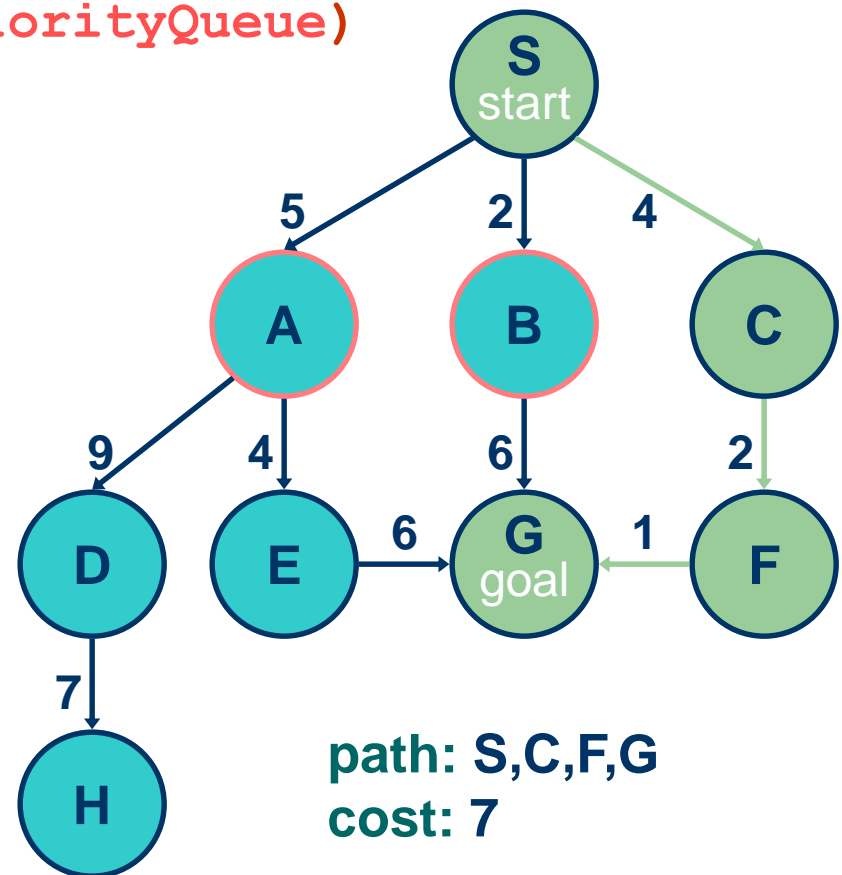


# Uniform-Cost Search (UCS)

`generalSearch(problem, priorityQueue)`

# of nodes tested: 6, expanded: 5

expnd. node	Frontier list
	{S}
S	{B:2,C:4,A:5}
B	{C:4,A:5,G:8}
C	{A:5,F:6,G:8}
A	{F:6,G:8,E:9,D:14}
F	{G:7,G:8,E:9,D:14}
G	{G:8,E:9,D:14}



# Uniform-Cost Search (UCS)

- Called *Dijkstra's Algorithm* in the algorithms literature
- Similar to *Branch and Bound Algorithm* in Operations Research literature
- Complete
- Optimal / Admissible
  - requires that the goal test is done when a node is ***removed*** from the *Frontier* rather than when the node is generated by its parent node

# Uniform-Cost Search (UCS)

- **Time and space complexity:  $O(b^d)$  (i.e., exponential)**
  - $d$  is the depth of the solution
  - $b$  is the branching factor at each non-leaf node
- **More precisely, time and space complexity is  $O(b^{C^*/\epsilon})$  where all edge costs  $\epsilon \geq \epsilon_{\min} > 0$ , and  $C^*$  is the best goal path cost**