



Degree Project in Technology

First cycle, 15 credits

# **This is the title in the language of the thesis**

A subtitle in the language of the thesis

**FAKE A. STUDENT**

**FAKE B. STUDENT**



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Bachelor's Programme in Information and Communication Technology

Date: January 29, 2024

Supervisors: A. Busy Supervisor, Another Busy Supervisor, Third Busy Supervisor

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School of Electrical Engineering and Computer Science

Host company: Företaget AB

Swedish title: Detta är den svenska översättningen av titeln

Swedish subtitle: Detta är den svenska översättningen av undertiteln



## 0.1 Inference Rules

### 0.1.1 Extension

#### 0.1.1.1 Typing

$$\text{T-TASK} \frac{x : C; \text{ocap} \vdash t : \tau \quad \Gamma; a \vdash b : Q \triangleright \text{Box}[C]}{\Gamma; a \vdash \text{task}(b)\{x \Rightarrow t\} : Q \triangleright \text{Task}[C]}$$

$$\text{T-ASYNC} \frac{\text{Perm}[Q] \in \Gamma \quad \Gamma \setminus \text{Perm}[Q]; a \vdash s : \sigma \quad \Gamma; a \vdash x : Q \triangleright \text{Task}[C]}{\Gamma; a \vdash \text{async}(x)\{s\} : \perp}$$

$$\text{T-FINISH} \frac{\Gamma; a \vdash t : \tau}{\Gamma; a \vdash \text{finish}\{t\} : \text{null}}$$

#### 0.1.1.2 Evaluation

$$\text{E-TASK} \frac{L(b') = b(o, p)}{H, \{(f, \langle L, \text{let } x = \text{task}(b')\{x \Rightarrow t\} \text{ in } s, P \rangle^l)\} \uplus TS} \\ \rightsquigarrow H, \{(f, \langle L[x \rightarrow \text{task}(b(o, p), t)], s, P \rangle^l)\} \uplus TS$$

$$\text{E-ASYNC} \frac{\begin{array}{l} L(y) = \text{task}(b(o, p), t) \quad p \in P \\ T_1 = (f, \langle L, s, P \rangle^\epsilon) \quad T_2 = (f, \langle [x \rightarrow o], t, \emptyset \rangle^\epsilon) \end{array}}{H, \{(f, \langle L, \text{async}(y)\{s\}, P \rangle^l \circ FS)\} \uplus TS} \\ \rightsquigarrow H, \{T_1, T_2\} \uplus TS$$

$$\text{E-FINISH1} \frac{T = (f', \langle L, t, P \rangle^\epsilon) \quad f' \text{ fresh}}{H, \{(f, \langle L, \text{let } x = \text{finish}\{t\} \text{ in } s, P \rangle^l \circ FS)\} \uplus TS} \\ \rightsquigarrow H, \{(f, \langle \text{FINISH } f' \rangle \circ \langle L[x \rightarrow \text{null}], s, P \rangle^l \circ FS)\} \uplus \{T\} \uplus TS$$

$$\text{E-FINISH2} \frac{\nexists (f', FS) \in TS}{H, \{(f, \langle \text{FINISH } f' \rangle \circ FS)\} \uplus TS} \\ \rightsquigarrow H, \{(f, FS)\} \uplus TS$$

$$\text{E-TASK-DONE} \frac{}{H, \epsilon \uplus TS \rightsquigarrow TS}$$

## 0.1.2 LaCasa

### 0.1.2.1 Well-Formedness

$$\text{WF-VAR} \frac{L(x) = \text{null} \vee L(x) = o \wedge \text{typeof}(H, o) <: \Gamma(x) \vee L(x) = b(o, p) \wedge \Gamma(x) = Q \triangleright \text{Box}[C] \wedge \text{typeof}(H, o) <: C}{H \vdash \Gamma; L; x}$$

$$\text{WF-PERM} \frac{\begin{array}{c} \gamma : \text{permTypes}(\Gamma) \longrightarrow \text{Pinjective} \\ \forall x \in \text{dom}(\Gamma). \\ (\Gamma(x) = Q \triangleright \text{Box}[C] \wedge L(x) = b(o, p) \wedge \text{Perm}[Q] \in \Gamma \vee \\ \Gamma(x) = Q \triangleright \text{Task}[C] \wedge L(x) = \text{task}(b(o, p), t) \wedge \text{Perm}[Q] \in \Gamma) \\ \implies \gamma(Q) = p \end{array}}{\vdash \Gamma; L; P}$$

$$\text{WF-ENV} \frac{\begin{array}{c} \text{dom}(\Gamma) \subseteq \text{dom}(L) \\ \forall x \in \text{dom}(\Gamma). H \vdash \Gamma; L; x \end{array}}{H \vdash \Gamma; L}$$

$$\text{WF-METHOD1} \frac{\Gamma_0, \text{this} : C, x : D; \epsilon \vdash t : E' \quad E' <: E}{C \vdash \text{defm}(x : D) : E = t}$$

$$\text{WF-METHOD2} \frac{\Gamma = \Gamma_0, \text{this} : C, x : Q \triangleright \text{Box}[D], \text{Perm}[Q] \quad Q \text{fresh} \quad \Gamma; \epsilon \vdash t : E' \quad E' <: E}{C \vdash \text{defm}(x : \text{Box}[D]) : E = t}$$

$$\text{WF-PROGRAM} \frac{p \vdash \bar{c}d \quad p \vdash \Gamma_0 \quad \Gamma_0; \epsilon \vdash t : \sigma}{p \vdash \bar{c}d \bar{v}dt}$$

$$\text{WF-CLASS} \frac{\begin{array}{c} C \vdash \bar{m}d \quad D = \text{AnyRef} \vee p \vdash \text{class} D \dots \\ \forall (\text{defm} \dots) \in \bar{m}d.\text{override}(m, C, D) \\ \forall \text{var} f : \sigma \in \bar{f}d.f \notin \text{fields}(D) \end{array}}{p \vdash \text{class} C \text{extends} D \{ \bar{f}d \bar{m}d \}}$$

$$\text{WF-OVERRIDE} \frac{m\text{type}(m, D) \text{notdefined} \vee m\text{type}(m, D) = m\text{type}(m, C)}{\text{override}(m, C, D)}$$

### 0.1.2.2 Typing

$$\text{T-FRAME1} \frac{\begin{array}{c} \Gamma; a \vdash t : \sigma \quad l \neq \epsilon \implies \sigma <: C \\ H \vdash \Gamma; L \quad H \vdash \Gamma; L; P \end{array}}{H \vdash \langle L, t, P \rangle^l : \sigma}$$

T-FRAME2	$\frac{\Gamma; x : \tau; a \vdash t : \sigma \quad l \neq \epsilon \implies \sigma <: C \quad \begin{array}{c} H \vdash \Gamma; L \\ H \vdash \Gamma; L; P \end{array}}{H \vdash_x^\tau \langle L, t, P \rangle^l : \sigma}$
—	
T-FRAME-NA	$\frac{H \vdash F^\epsilon : \sigma \quad H \vdash FS}{H \vdash F^\epsilon \circ FS}$
—	
T-FRAME-NA2	$\frac{H \vdash_x^\tau F^\epsilon : \sigma \quad H \vdash FS}{H \vdash_x^\tau F^\epsilon \circ FS}$
—	
T-FRAME-A	$\frac{H \vdash F^x : \sigma \quad H \vdash_x^\sigma FS}{H \vdash F^x \circ FS}$
—	
T-FRAME-A2	$\frac{H \vdash_x^\tau F^y : \sigma \quad H \vdash_y^\sigma FS}{H \vdash_x^\tau F^y \circ FS}$
—	
—	
T-VAR	$\frac{x \in \text{dom}(\Gamma)}{\Gamma; a \vdash x : \Gamma(x)}$
—	
T-INVOKE	$\frac{\begin{array}{c} \Gamma; a \vdash x : C \quad mtype(C, m) = \sigma \rightarrow \tau \\ \Gamma; a \vdash y : \sigma' \quad \sigma' <: \sigma \vee \\ (\sigma = \text{Box}[D] \wedge \sigma' = Q \triangleright \text{Box}[D] \wedge \text{Perm}[Q] \in \Gamma) \end{array}}{\Gamma; a \vdash x.m(y) : \tau}$
—	
T-NEW	$\frac{a = \text{ocap} \implies \text{ocap}(C) \quad \forall \text{var } f : \sigma \in \text{fd}_{\text{within}} \text{above}. \exists D. \sigma = D}{\Gamma; a \vdash \text{new } C : C}$
—	
T-OPEN	$\frac{\Gamma; a \vdash x : Q \triangleright \text{Box}[C] \quad \text{Perm}[Q] \in \Gamma \quad y : C; \text{ocap} \vdash t : \sigma}{\Gamma; a \vdash x.\text{open}\{y \Rightarrow t\} : Q \triangleright \text{Box}[C]}$
—	
T-BOX	$\frac{\text{ocap}(C) \quad Q \text{fresh} \quad \Gamma; x : Q \triangleright \text{Box}[C]; \text{Perm}[Q]; a \vdash t : \text{sigma}}{\Gamma; a \vdash \text{box}[C]\{x \Rightarrow t\} : \perp}$
—	
T-CAPTURE	$\frac{\begin{array}{c} \Gamma; a \vdash x : Q \triangleright \text{Box}[C] \quad \Gamma; a \vdash y : Q' \triangleright \text{Box}[D] \\ \{\text{Perm}[Q], \text{Perm}[Q']\} \subseteq \Gamma \quad D <: \text{ftype}(C, f) \\ \Gamma \{\text{Perm}[Q']\}, z : Q \triangleright \text{Box}[C]; a \vdash t : \sigma \end{array}}{\Gamma; a \vdash \text{capture}(x.f, y)\{z \Rightarrow t\} : \perp}$
—	
T-SWAP	$\frac{\begin{array}{c} \Gamma; a \vdash x : Q \triangleright \text{Box}[C] \quad \Gamma; a \vdash y : Q' \triangleright \text{Box}[D'] \\ \{\text{Perm}[Q], \text{Perm}[Q']\} \subseteq \Gamma \quad \text{ftype}(C, f) = \text{Box}[D] \\ D' <: D \quad R \text{fresh} \\ \Gamma \{\text{Perm}[Q']\}, z : R \triangleright \text{Box}[D], \text{Perm}[R]; a \vdash t : \sigma \end{array}}{\Gamma; a \vdash \text{swap}(x.f, y)\{z \Rightarrow t\} : \perp}$

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### 0.1.2.3 Evaluation

### 0.1.2.4 Other

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$$\text{F-OK} \frac{\begin{array}{c} \text{boxSep}(H, F) \quad \text{boxObjSep}(H, F) \quad \text{boxOcap}(H, F) \\ a = \text{ocap} \implies \text{globalOcapSep}(H, F) \quad \text{fieldUniqueness}(H, F) \end{array}}{H; a \vdash F \text{ok}}$$