ICC4101 - Algorithms and Competitive Programing

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Dynamic Programing

A generic approach to solve dynamic programing problems is as follows:

- 1. Find sub-problems
- 2. "Guess" (partial) solution
- 3. Relate subproblems
- 4. Recursion and memoization, or use a bottom-up approach to build solutions (avoid repeating yourself)
- 5. Find solution to the original problem

Types of parameters for Dynamic Programing problems

Most DP sub-problems will have parameters (or a combination) of these types:

```
1. Index i of an array [a_0, a_1, ..., a_i, ...].
```

The steps of guessing and relating sub-problmes implies including more elements of the array. For example, the Longest Increasing Subsequence (LIS) problem. 2. Indices (i, j) of two arrays [a_0, a_1, ..., a_i, ...] and [b_0, b_1, ..., b_j, ...]. The steps of guessing and relating sub-problmes implies incresing the index i, j or both. For example, calculating the edit distance.

```
3. Sub-indices (i, j) of an array [a_0, a_1, ..., a_i, a_i+1}, ..., a_j, ...].
```

The steps of guessing and relating sub-problmes implies dividing the problem such as (i, i + k) and (i + k, j). 4. Vertex position in a DAG (usually an implicit graph). Process the neighbors (next or previous vertex).

5. Knapsack-stype parameter.

The step of guessing and realting sub-problem implies incresing (or decreasing) a value (or weight?) until the maximum value or threshold. 6. Small subset, usually modeled with bitmask (for example if only element 1 and 3 are part of the set, this can me modeled with the bitmask 010100...). The step of guessing and relating sub-problems implies changing one or more 0s to 1s.

Previously...

1. Undirectional TSP

Technically a Graph problem. Solution similar to Bellmand-Ford or Floy-Warshall but since it is a DAG (Direct **Acyclic** Graph) a simple solution exists:

- path(i, j) = (c, p) minimum cost c and path p to reach the last row from row i and column j.
- Guess if the next step (j+1) goes through th row i-1, i or i+1.
- Minimize!

```
import sys
from itertools import chain
from functools import lru_cache
def numbers():
    yield from chain.from_iterable(map(int, line.split()) for line
in sys.stdin.readlines())
def main():
    nums = numbers()
    while True:
        try:
            n, m = next(nums), next(nums)
            M = dict()
            for i in range(n):
                for j in range(m):
                    M[i, j] = next(nums)
        except StopIteration:
            return
```

```
@lru_cache(None)
    def path(i, j):
        if i < 0:
            i = n + i
        if i == n:
            i = 0
        if j == m:
            return 0, []

        min_cost, min_path = min(path(i-1, j+1), path(i, j+1),
        path(i+1, j+1))
        return M[i,j] + min_cost, [i+1] + min_path

        min_cost, min_path = min(path(i, 0) for i in range(n))
        print(" ".join(map(str, min_path)))
        print(min_cost)

main()</pre>
```

2. Testing the CATCHER

from sys import stdin

Longest **Decreasing** Sub-sequence

- LIS(i) longest decreasing sub-sequence starting at index i.
- Guess: the next step k of the longest sub-sequence starting at index i

```
from itertools import takewhile, count

for k in count():
    heights = list(takewhile(
        lambda x: x != -1,
        map(int, (stdin.readline() for _ in count()))
    ))
    N = len(heights)
    if N == 0:
        break
```

3. Jill Rides Again

Maximum sum between an interval.

```
1. nice(i, j) nice route value between indices i and j
```

- 2. No Guessing
- 3. nice(i, j) = nice(i, j-1) + route[j-1]

4. Compromise

Longest Common Sub-sequence

- 1. LCS(i, j) is the longest common subsequence starting at index i at the first sequence and index j at the second sub-sequence
- 2. Guess the next word of the LCS

```
@lru_cache(None)
def lcs(i, j):
    if i == len(words1) or j == len(words2):
        return []
    if words1[i] == words2[j]:
        return [words1[i]] + lcs(i+1, j+1)
    return max([lcs(i+1, j), lcs(i, j+1), lcs(i+1, j+1)], key=len)
```

5. Coin Change

- 1. ways (v, i) ways of giving change to v cents using coins with index i or above (to avoid repetitions)
- 2. No quessing
- 3. ways(v, i) = sum(ways(v coins[k], k) for k in range(i, N))

Why use index i?

```
COINS = [50, 25, 10, 5, 1]
@lru_cache(maxsize=None)
def ways(value, i=0):
    if value == 0:
        return 1
    elif value < 0:</pre>
        return None
    ways = 0
    for k, c in enumerate(COINS):
        if k < i:
            continue
        v = ways(value - c, k)
        if v is None:
            continue
        ways_ += v
    return ways_
```

6. Diving for gold

- 1. value(t, i) maximum value of treasures to get with t remainding seconds and only considering treasures with index i and above.
- 2. Guess which is the optimal treasure to get with t remaining seconds
- 3. value(t, i) = max(value(t, id + 1), VAL[id] + value(t 3 times[id], id + 1))

```
# Take item treasure
    if TIME[treasure] <= time:</pre>
        take_value, take_times, take_treasures = value(
            time - TIME[treasure],
            treasure + 1
        take value += GOLD[treasure]
        if take_value > skip_value:
            take_times_ = [time] + take_times
            take_treasures_ = [treasure] + take_treasures
            return take_value, take_times_, take_treasures_
    return skip value, skip times, skip times
while True:
    value.cache_clear()
    T, w = map(int, stdin.readline().split())
    N = int(stdin.readline())
    TIME = N*[0]
    GOLD = N*[0]
    for i in range(N):
        d, GOLD[i] = map(int, stdin.readline().split())
        TIME[i] = 3*w*d
    val, times, treasures = value(T, 0)
    print(val)
    print(len(treasures))
    for t, T in zip(times, treasures):
        print(t, GOLD[t])
    print()
```

For this week!

Explain with comments (one or two sentences) the optimal sub-stucture of your dynamic programming solution!

In []: