Linking

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Today

- Linking
 - Motivation
 - How it works
 - Dynamic linking
- Case study: Library interpositioning

Example C Program

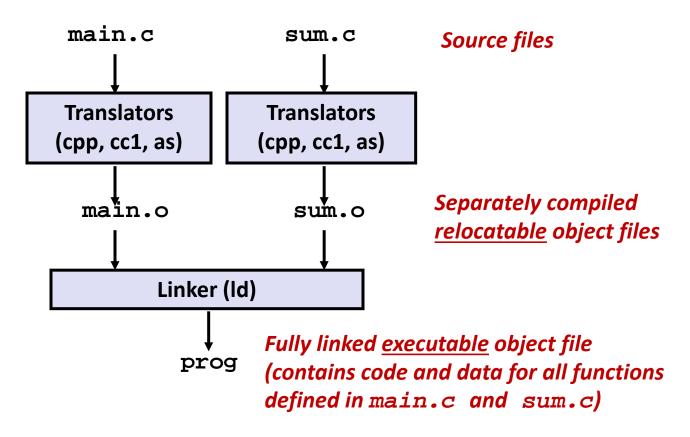
```
int sum(int *a, int n);
int array[2] = {1, 2};
int main(int argc, char** argv)
{
   int val = sum(array, 2);
   return val;
}
```

```
int sum(int *a, int n)
{
   int i, s = 0;

   for (i = 0; i < n; i++) {
       s += a[i];
   }
   return s;
}</pre>
```

Linking

- Programs are translated and linked using a compiler driver:
 - linux> gcc -Og -o prog main.c sum.c
 - linux> ./prog



Why Linkers?

- Reason 1: Modularity
 - Program can be written as a collection of smaller source files, rather than one monolithic mass.
 - Can build libraries of common functions (more on this later)
 - e.g., Math library, standard C library

Why Linkers? (cont)

Reason 2: Efficiency

- Time: Separate compilation
 - Change one source file, compile, and then relink.
 - No need to recompile other source files.
 - Can compile multiple files concurrently.
- Space: Libraries
 - Common functions can be aggregated into a single file...
 - Option 1: Static Linking
 - Executable files and running memory images contain only the library code they actually use
 - Option 2: Dynamic linking
 - Executable files contain no library code
 - During execution, single copy of library code can be shared across all executing processes

What Do Linkers Do?

Step 1: Symbol resolution

Programs define and reference symbols (global variables and functions):

```
void sum() {...} /* define symbol sum */
sum(); /* reference symbol sum */
int *xp = &x; /* define symbol xp, reference x */
```

- Symbol definitions are stored in object file (by assembler) in symbol table.
 - Symbol table is an array of entries
 - Each entry includes name, size, and location of symbol.
- During symbol resolution step, the linker associates each symbol reference with exactly one symbol definition.



Symbols in Example C Program

Definitions

```
int sum(int *a, int n),
int array[2] = {1, 2};
int main(int argc, char** argv)
{
   int val = sum(array, 2);
   return val;
}

main.c
```

```
int sum(int *a, int n)
{
   int i, s = 0;

   for (i = 0; i < n; i++) {
       s += a[i];
   }
   return s;
}</pre>
```

Reference

What Do Linkers Do? (cont)

- Step 2: Relocation
 - Merges separate code and data sections into single sections
 - Relocates symbols from their relative locations in the .o files to their final absolute memory locations in the executable.
 - Updates all references to these symbols to reflect their new positions.

Let's look at these two steps in more detail....

Three Kinds of Object Files (Modules)

Relocatable object file (.o file)

- Contains code and data in a form that can be combined with other relocatable object files to form executable object file.
 - Each .o file is produced from exactly one source (.c) file

Executable object file (a.out file)

 Contains code and data in a form that can be copied directly into memory and then executed.

Shared object file (.so file)

- Special type of relocatable object file that can be loaded into memory and linked dynamically, at either load time or run-time.
- Called Dynamic Link Libraries (DLLs) by Windows

Executable and Linkable Format (ELF)

- Standard binary format for object files
- One unified format for
 - Relocatable object files (.o),
 - Executable object files (a.out)
 - Shared object files (.so)
- Generic name: ELF binaries

ELF Object File Format

- Elf header
 - Word size, byte ordering, file type (.o, exec, .so), machine type, etc.
- Segment header table
 - Page size, virtual addresses memory segments (sections), segment sizes.
- .text section
 - Code
- .rodata section
 - Read only data: jump tables, ...
- data section
 - Initialized global variables
- bss section
 - Uninitialized global variables
 - "Block Started by Symbol"
 - "Better Save Space"
 - Has section header but occupies no space

ELF header Segment header table (required for executables) . text section .rodata section .data section .bss section .symtab section .rel.txt section .rel.data section .debug section Section header table

ELF Object File Format (cont.)

symtab section

- Symbol table
- Procedure and static variable names
- Section names and locations

rel.text section

- Relocation info for .text section
- Addresses of instructions that will need to be modified in the executable
- Instructions for modifying.

.rel.data section

- Relocation info for .data section
- Addresses of pointer data that will need to be modified in the merged executable

debug section

■ Info for symbolic debugging (gcc -g)

Section header table

Offsets and sizes of each section

ELF header
Segment header table (required for executables)
.text section
.rodata section
.data section
.bss section
.symtab section
.rel.txt section
.rel.data section
.debug section
Section header table

Linker Symbols

Global symbols

- Symbols defined by module m that can be referenced by other modules.
- E.g.: non-static C functions and non-static global variables.

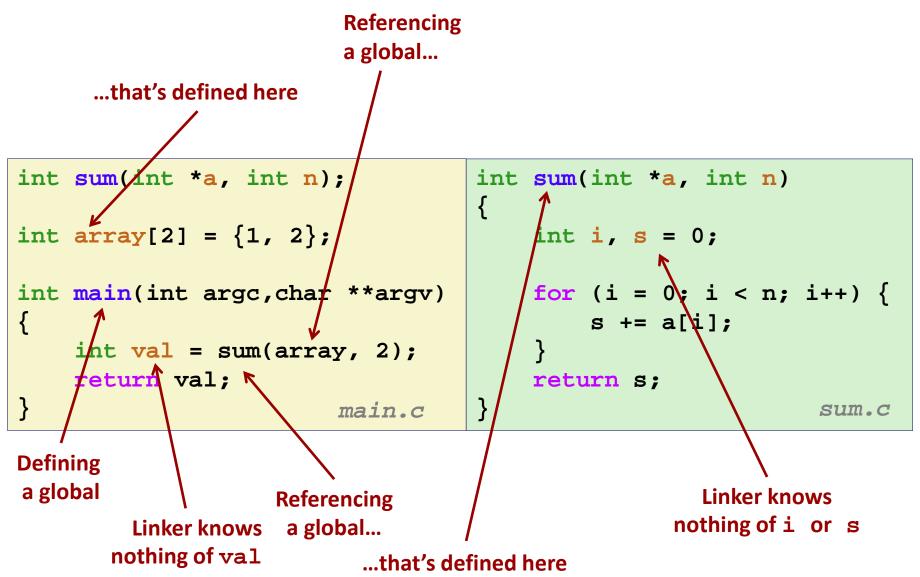
External symbols

 Global symbols that are referenced by module m but defined by some other module.

Local symbols

- Symbols that are defined and referenced exclusively by module m.
- E.g.: C functions and global variables defined with the static attribute.
- Local linker symbols are not local program variables

Step 1: Symbol Resolution





Local Symbols

Local non-static C variables vs. local static C variables

- local non-static C variables: stored on the stack
- local static C variables: stored in either .bss, or .data

```
static int x = 15;
int f() {
    static int x = 17;
    return x++;
int g() {
    static int x = 19;
    return x += 14;
int h() {
    return x += 27;
         static-local.c
```

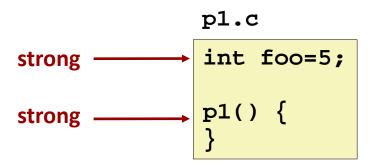
Compiler allocates space in .data for each definition of \mathbf{x}

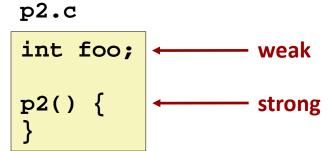
Creates local symbols in the symbol table with unique names, e.g., x, x.1721 and x.1724.

Third Edition

How Linker Resolves Duplicate Symbol Definitions

- Program symbols are either strong or weak
 - Strong: procedures and initialized globals
 - Weak: uninitialized globals







Linker's Symbol Rules

- Rule 1: Multiple strong symbols are not allowed
 - Each item can be defined only once
 - Otherwise: Linker error
- Rule 2: Given a strong symbol and multiple weak symbols, choose the strong symbol
 - References to the weak symbol resolve to the strong symbol
- Rule 3: If there are multiple weak symbols, pick an arbitrary one
 - Can override this with gcc -fno-common
- Puzzles on the next slide

Linker Puzzles

```
int x; p1() {}
```

Link time error: two strong symbols (p1)

```
int x; p1() {}
```

References to **x** will refer to the same uninitialized int. Is this what you really want?

```
int x;
int y;
p1() {}
```

Writes to **x** in **p2** might overwrite **y**! Evil!

```
int x=7;
int y=5;
p1() {}
```

Writes to **x** in **p2** will overwrite **y**! Nasty!

References to **x** will refer to the same initialized variable.

Nightmare scenario: two identical weak structs, compiled by different compilers with different alignment rules.

Global Variables

Avoid if you can

Otherwise

- Use static if you can
- Initialize if you define a global variable
- Use extern if you reference an external global variable

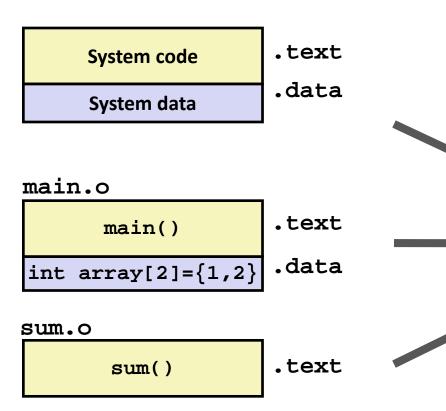
Role of .h Files

```
global.h
c1.c
                             extern int g;
#irclude "global.h"
                             static int init = 0;
int f() {
                             #else
  return g+1;
                               extern int g;
                               static int init = 0;
                             #endif
c2.c
#define INITIALIZE
#include <stdio.h>
                            int g = 23;
#knclude "global.h"
                            static int init = 1;
int main(int argc, char** argv) {
  if (init)
    // do something, e.g., g=31;
  int t = f();
  printf("Calling f yields %d\n", t);
  return 0;
```

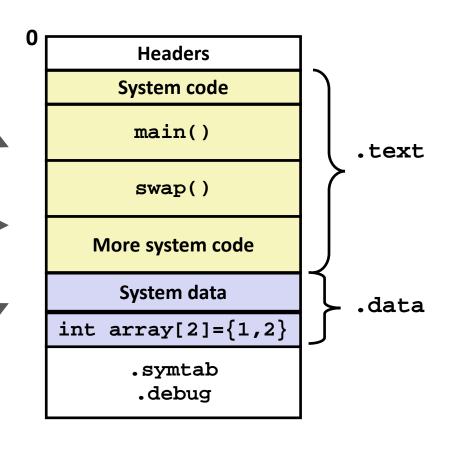


Step 2: Relocation

Relocatable Object Files



Executable Object File





Relocation Entries

```
int array[2] = {1, 2};
int main(int argc, char**
argv)
{
   int val = sum(array, 2);
   return val;
}
```

```
0000000000000000 <main>:
     48 83 ec 08
                                     $0x8,%rsp
  0:
                               sub
  4: be 02 00 00 00
                                     $0x2,%esi
                              mov
  9:
     bf 00 00 00 00
                                     $0x0,%edi
                              mov
                                                  # %edi = &array
                       a: R_X86_64_32 array
                                                    # Relocation entry
       e8 00 00 00 00
                              callq 13 < main + 0x13 > \# sum()
  e:
                       f: R X86 64 PC32 sum-0x4 # Relocation entry
     48 83 c4 08
 13:
                               add
                                     $0x8,%rsp
 17:
       c3
                               retq
                                                                main.o
```

Relocated .text section

```
00000000004004d0 <main>:
  4004d0:
                48 83 ec 08
                                          $0x8,%rsp
                                  sub
  4004d4:
                be 02 00 00 00
                                          $0x2,%esi
                                  mov
  4004d9:
                bf 18 10 60 00
                                          $0x601018,%edi
                                                         # %edi = &array
                                  mov
  4004de:
                e8 05 00 00 00
                                         4004e8 <sum>
                                                          # sum()
                                  callq
  4004e3:
                48 83 c4 08
                                          $0x8,%rsp
                                  add
  4004e7:
                c3
                                  retq
00000000004004e8 <sum>:
                b8 00 00 00 00
                                                $0x0,%eax
  4004e8:
                                         mov
               ba 00 00 00 00
                                                $0x0,%edx
  4004ed:
                                         mov
                                                4004fd < sum + 0x15 >
  4004f2:
                eb 09
                                         jmp
                48 63 ca
  4004f4:
                                         movslq %edx,%rcx
                                               (%rdi,%rcx,4),%eax
  4004f7:
                03 04 8f
                                         add
                                         add
  4004fa:
               83 c2 01
                                                $0x1,%edx
  4004fd:
                39 f2
                                                %esi,%edx
                                         cmp
  4004ff:
                7c f3
                                         jl
                                                4004f4 < sum + 0xc >
  400501:
                f3 c3
                                         repz retq
```

callq instruction uses PC-relative addressing for sum():

0x4004e8 = 0x4004e3 + 0x5



Loading Executable Object Files

Executable Object File

ELF header	0
Program header table (required for executables)	
.init section	
.text section	
.rodata section	
.data section	
.bss section	
.symtab	
.debug	
.line	
.strtab	
Section header table (required for relocatables)	

Memory invisible to **Kernel virtual memory** user code User stack (created at runtime) %rsp (stack pointer) Memory-mapped region for shared libraries brk **Run-time heap** (created by malloc) Loaded Read/write data segment from (.data, .bss) the **Read-only code segment** executable (.init,.text,.rodata) file

Unused

0x400000

Quiz Time!

K'm-Pee-Du-Wee: Song by guitarist Steve Vai

Check out: quiz: day 13: Linking

https://canvas.cmu.edu/courses/3822



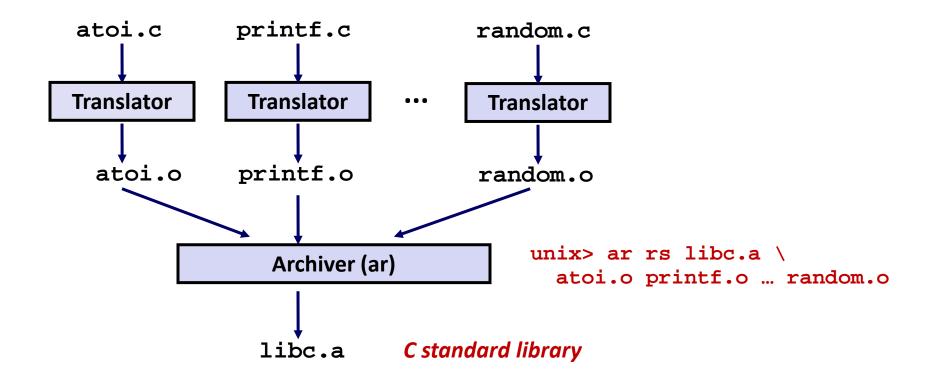
Packaging Commonly Used Functions

- How to package functions commonly used by programmers?
 - Math, I/O, memory management, string manipulation, etc.
- Awkward, given the linker framework so far:
 - Option 1: Put all functions into a single source file
 - Programmers link big object file into their programs
 - Space and time inefficient
 - Option 2: Put each function in a separate source file
 - Programmers explicitly link appropriate binaries into their programs
 - More efficient, but burdensome on the programmer

Old-fashioned Solution: Static Libraries

- Static libraries (.a archive files)
 - Concatenate related relocatable object files into a single file with an index (called an archive).
 - Enhance linker so that it tries to resolve unresolved external references by looking for the symbols in one or more archives.
 - If an archive member file resolves reference, link it into the executable.

Creating Static Libraries



- Archiver allows incremental updates
- Recompile function that changes and replace .o file in archive.



Commonly Used Libraries

libc.a (the C standard library)

- 4.6 MB archive of 1496 object files.
- I/O, memory allocation, signal handling, string handling, data and time, random numbers, integer math

libm.a (the C math library)

- 2 MB archive of 444 object files.
- floating point math (sin, cos, tan, log, exp, sqrt, ...)

```
% ar -t libc.a | sort
...
fork.o
...
fprintf.o
fpu_control.o
fputc.o
freopen.o
fscanf.o
fseek.o
fstab.o
...
```

```
% ar -t libm.a | sort
...
e_acos.o
e_acosf.o
e_acosh.o
e_acoshf.o
e_acoshl.o
e_acosl.o
e_asin.o
e_asinf.o
e_asinf.o
e_asinf.o
...
```

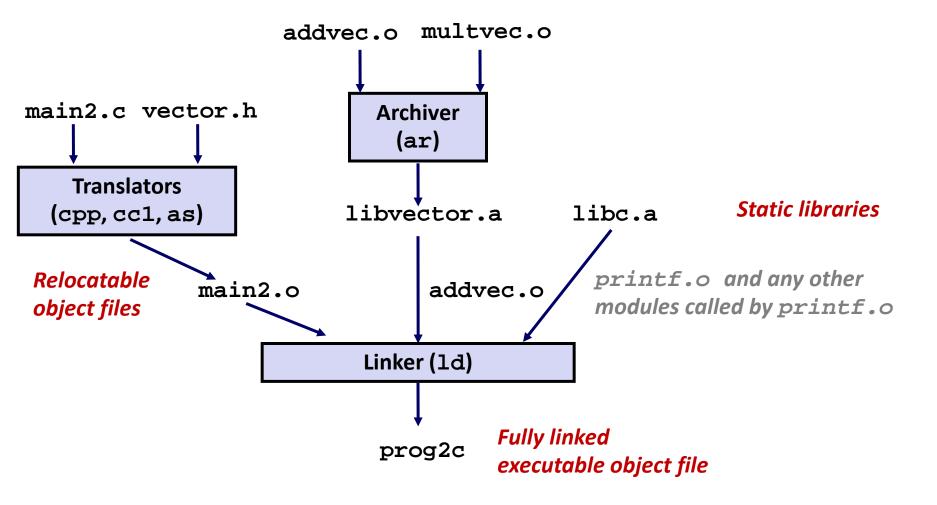
Linking with Static Libraries

```
#include <stdio.h>
#include "vector.h"
int x[2] = \{1, 2\};
int y[2] = \{3, 4\};
int z[2];
int main(int argc, char**
argv)
    addvec(x, y, z, 2);
    printf("z = [%d %d] \n",
           z[0], z[1]);
    return 0;
                    main2.c
```

libvector.a



Linking with Static Libraries



"c" for "compile-time"

Using Static Libraries

Linker's algorithm for resolving external references:

- Scan .o files and .a files in the command line order.
- During the scan, keep a list of the current unresolved references.
- As each new .o or .a file, obj, is encountered, try to resolve each unresolved reference in the list against the symbols defined in obj.
- If any entries in the unresolved list at end of scan, then error.

Problem:

- Command line order matters!
- Moral: put libraries at the end of the command line.

```
unix> gcc -L. libtest.o -lmine
unix> gcc -L. -lmine libtest.o
libtest.o: In function 'main':
libtest.o(.text+0x4): undefined reference to 'libfun'
```

Modern Solution: Shared Libraries

Static libraries have the following disadvantages:

- Duplication in the stored executables (every function needs libc)
- Duplication in the running executables
- Minor bug fixes of system libraries require each application to explicitly relink
 - Rebuild everything with glibc?
 - https://security.googleblog.com/2016/02/cve-2015-7547-glibc-getaddrinfo-stack.html

Modern solution: Shared Libraries

- Object files that contain code and data that are loaded and linked into an application dynamically, at either load-time or run-time
- Also called: dynamic link libraries, DLLs, .so files

Shared Libraries (cont.)

- Dynamic linking can occur when executable is first loaded and run (load-time linking).
 - Common case for Linux, handled automatically by the dynamic linker (ld-linux.so).
 - Standard C library (libc.so) usually dynamically linked.
- Dynamic linking can also occur after program has begun (run-time linking).
 - In Linux, this is done by calls to the dlopen() interface.
 - Distributing software.
 - High-performance web servers.
 - Runtime library interpositioning.
- Shared library routines can be shared by multiple processes.
 - More on this when we learn about virtual memory

What dynamic libraries are required?

.interp section

Specifies the dynamic linker to use (i.e., ld-linux.so)

.dynamic section

- Specifies the names, etc of the dynamic libraries to use
- Follow an example of csim-ref from cachelab

```
(NEEDED) Shared library: [libm.so.6]
```

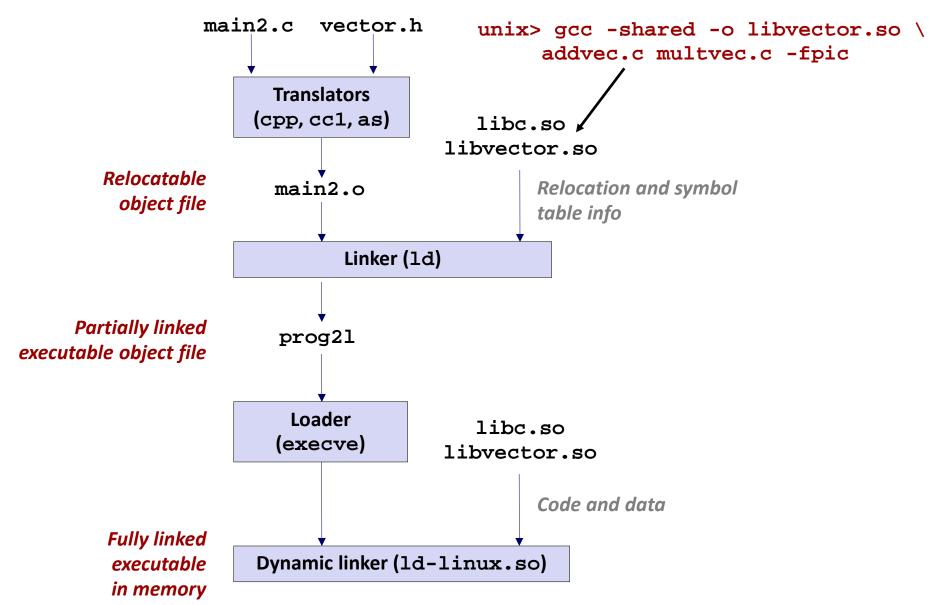
Where are the libraries found?

Use "ldd" to find out:

```
unix> ldd csim-ref
  linux-vdso.so.1 => (0x00007ffc195f5000)
  libc.so.6 => /lib/x86_64-linux-gnu/libc.so.6 (0x00007f345eda6000)
  /lib64/ld-linux-x86-64.so.2 (0x00007f345f181000)
```



Dynamic Linking at Load-time





Dynamic Linking at Run-time

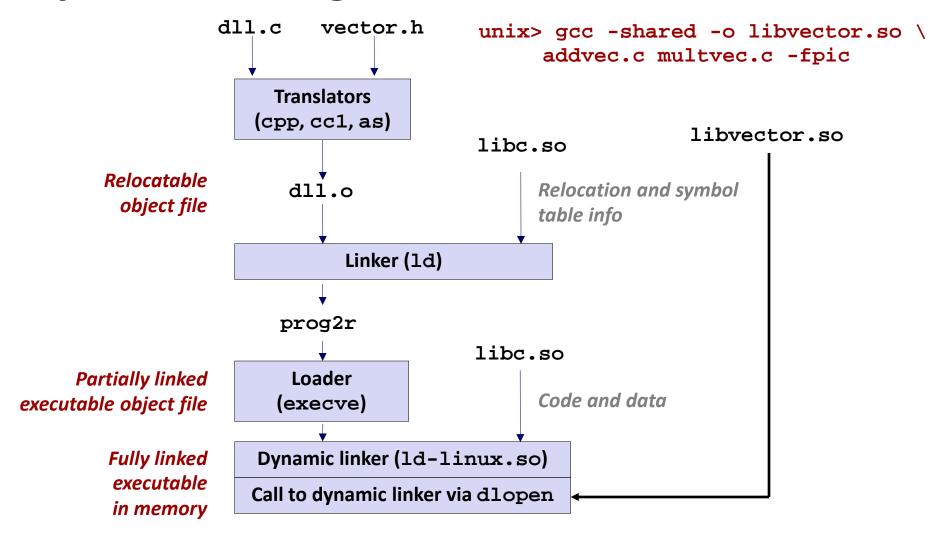
```
#include <stdio.h>
#include <stdlib.h>
#include <dlfcn.h>
int x[2] = \{1, 2\};
int y[2] = \{3, 4\};
int z[2];
int main(int argc, char** argv)
   void *handle:
   void (*addvec)(int *, int *, int *, int);
    char *error:
    /* Dynamically load the shared library that contains addvec() */
   handle = dlopen("./libvector.so", RTLD LAZY);
    if (!handle) {
        fprintf(stderr, "%s\n", dlerror());
       exit(1);
                                                                dll.c
```

Dynamic Linking at Run-time (cont)

```
/* Get a pointer to the addvec() function we just loaded */
addvec = dlsym(handle, "addvec");
if ((error = dlerror()) != NULL) {
    fprintf(stderr, "%s\n", error);
    exit(1);
/* Now we can call addvec() just like any other function */
addvec(x, y, z, 2);
printf("z = [%d %d]\n", z[0], z[1]);
/* Unload the shared library */
if (dlclose(handle) < 0) {</pre>
    fprintf(stderr, "%s\n", dlerror());
    exit(1);
return 0:
                                                        dll.c
```



Dynamic Linking at Run-time



Linking Summary

- Linking is a technique that allows programs to be constructed from multiple object files.
- Linking can happen at different times in a program's lifetime:
 - Compile time (when a program is compiled)
 - Load time (when a program is loaded into memory)
 - Run time (while a program is executing)
- Understanding linking can help you avoid nasty errors and make you a better programmer.

Today

- Linking
- Case study: Library interpositioning

Case Study: Library Interpositioning

- Library interpositioning: powerful linking technique that allows programmers to intercept calls to arbitrary functions
- Interpositioning can occur at:
 - Compile time: When the source code is compiled
 - Link time: When the relocatable object files are statically linked to form an executable object file
 - Load/run time: When an executable object file is loaded into memory, dynamically linked, and then executed.



Some Interpositioning Applications

Monitoring and Profiling

- Count number of calls to functions
- Characterize call sites and arguments to functions
- Malloc tracing
 - Detecting memory leaks
 - Generating address traces



Example program

```
#include <stdio.h>
#include <malloc.h>
#include <stdlib.h>
int main(int argc,
         char *argv[])
  int i;
  for (i = 1; i < argc; i++) {
    void *p =
          malloc(atoi(argv[i]));
    free(p);
  return(0);
                             int.c
```

- Goal: trace the addresses and sizes of the allocated and freed blocks, without breaking the program, and without modifying the source code.
- Three solutions: interpose on the library malloc and free functions at compile time, link time, and load/run time.



Compile-time Interpositioning

```
#ifdef COMPILETIME
#include <stdio.h>
#include <malloc.h>
/* malloc wrapper function */
void *mymalloc(size_t size)
    void *ptr = malloc(size);
    printf("malloc(%d)=%p\n", (int)size, ptr);
    return ptr;
/* free wrapper function */
void myfree(void *ptr)
    free(ptr);
    printf("free(%p)\n", ptr);
#endif
                                                     mymalloc.c
```



Compile-time Interpositioning

```
#define malloc(size) mymalloc(size)
#define free(ptr) myfree(ptr)
void *mymalloc(size t size);
void myfree(void *ptr);
                                                         malloc.h
linux> make intc
qcc -Wall -DCOMPILETIME -c mymalloc.c
gcc -Wall -I. -o intc int.c mymalloc.o
linux> make runc
./intc 10 100 1000
                              Search for <malloc, h > leads to
malloc(10) = 0x1ba7010
                              /usr/include/malloc.h
free(0x1ba7010)
malloc(100) = 0x1ba7030
free(0x1ba7030)
malloc(1000) = 0x1ba70a0
                            Search for <malloc.h > leads to
free(0x1ba70a0)
linux>
```

Link-time Interpositioning

```
#ifdef LINKTIME
#include <stdio.h>
void * real malloc(size t size);
void __real_free(void *ptr);
/* malloc wrapper function */
void * wrap malloc(size t size)
   void *ptr = __real_malloc(size); /* Call libc malloc */
    printf("malloc(%d) = %p\n", (int)size, ptr);
    return ptr;
/* free wrapper function */
void __wrap_free(void *ptr)
    real free(ptr); /* Call libc free */
   printf("free(%p)\n", ptr);
#endif
```



Link-time Interpositioning

```
linux> make intl
gcc -Wall -DLINKTIME -c mymalloc.c
gcc -Wall -c int.c
gcc -Wall -Wl,--wrap,malloc -Wl,--wrap,free -o intl \
   int.o mymalloc.o
linux> make runl
./intl 10 100 1000
malloc(10) = 0x91a010
free(0x91a010)
. . .
```

- The "-W1" flag passes argument to linker, replacing each comma with a space.
- The "--wrap, malloc" arg instructs linker to resolve references in a special way:
 - Refs to malloc should be resolved as __wrap_malloc
 - Refs to ___real_malloc should be resolved as malloc



Load/Run-time Interpositioning

```
#ifdef RUNTIME
                                           Interpositioning
#define GNU SOURCE
#include <stdio.h>
#include <stdlib.h>
                            Observe that DON'T have
#include <dlfcn.h>
                            #include <malloc.h>
/* malloc wrapper function */
void *malloc(size_t size)
   void *(*mallocp)(size t size);
    char *error:
   mallocp = dlsym(RTLD_NEXT, "malloc"); /* Get addr of libc malloc */
    if ((error = dlerror()) != NULL) {
        fputs(error, stderr);
       exit(1);
    char *ptr = mallocp(size); /* Call libc malloc */
   printf("malloc(%d) = %p\n", (int)size, ptr);
    return ptr;
                                                            mymalloc.c
```



Load/Run-time Interpositioning

```
/* free wrapper function */
void free(void *ptr)
   void (*freep)(void *) = NULL;
    char *error;
    if (!ptr)
        return:
    freep = dlsym(RTLD NEXT, "free"); /* Get address of libc free */
    if ((error = dlerror()) != NULL) {
        fputs(error, stderr);
        exit(1);
    freep(ptr); /* Call libc free */
   printf("free(%p)\n", ptr);
#endif
```

mymalloc.c



Load/Run-time Interpositioning

```
linux> make intr
gcc -Wall -DRUNTIME -shared -fpic -o mymalloc.so mymalloc.c -ldl
gcc -Wall -o intr int.c
linux> make runr
(LD_PRELOAD="./mymalloc.so" ./intr 10 100 1000)
malloc(10) = 0x91a010
free(0x91a010)
. . .
linux>
Search for <malloc.h> leads to
/usr/include/malloc.h
```

- The LD_PRELOAD environment variable tells the dynamic linker to resolve unresolved refs (e.g., to malloc) by looking in mymalloc.so first.
- Type into (some) shells as:

```
(setenv LD_PRELOAD "./mymalloc.so"; ./intr 10 100 1000)
```



Interpositioning Recap

Compile Time

- Apparent calls to malloc/free get macro-expanded into calls to mymalloc/myfree
- Simple approach. Must have access to source & recompile

Link Time

- Use linker trick to have special name resolutions
 - malloc → __wrap_malloc
 - __real_malloc → malloc

Load/Run Time

- Implement custom version of malloc/free that use dynamic linking to load library malloc/free under different names
- Can use with ANY dynamically linked binary

```
(setenv LD_PRELOAD "./mymalloc.so"; gcc -c int.c)
```

Linking Recap

- Usually: Just happens, no big deal
- Sometimes: Strange errors
 - Bad symbol resolution
 - Ordering dependence of linked .o, .a, and .so files
- For power users:
 - Interpositioning to trace programs with & without source