Names, Scopes, Bindings

CMPSC 461
Programming Language Concepts
Penn State University
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Names in Programs

Scope: visibility of names

Storage: memory space associated with names

Lifetime: the time interval a variable is allocated with memory

Binding: a mapping between a name and its property

Scope

The visibility or availability of names

Which code region is a name visible?

When is it determined?

- Static or Dynamic?
- Lexical or Control-Flow?

How & when does it change?

Static (Lexical) Scoping

Scopes can be determined by the compiler

All bindings for identifiers can be resolved by examining the program

Typically, use the closet binding

Used in most compiled languages (C, Java)

Scope Rules in Scheme

```
(let ((x 0)) 
 \underline{(+ x x))} Scope of x
```

```
(let ((x 0))

(let ((y (+ x 1)))

(+ x y))
```

Scope of x

Scope of z

(let
$$((x 17))$$

 $(lambda (z) (+ z x))$

(let*) Scope of x

Scope of inner y

```
(define (f x)
(let ((* +))
(* x x)))
```

 $(q \times x))$

Scope of *

Scope of g

Scope Rules in Scheme

Three "let"s in Scheme

let

let*

$$\frac{(\text{let* }((\text{x 0}) \ \underline{(\text{y x}))}}{(+\text{x y})} \frac{(\text{y x})}{} \text{Scope of x}$$

letrec

Scope of is Even and is Odd

Nested Scope

In Scheme, each (let ...) defines one scope

In Pascal, ALGOL, function can be nested

```
function E (int x) {
  function F (int y) ...
  return F(1)+x; }
// F is not visible here
```

In C, block scopes can be nested

```
function E (int x) {
    { int y=3; x=x+y;}

    // y is not visible here
    return x; }
```

How are scope rules being checked?

Scope Check

Symbol table: a table of names and their bindings

Single scope: a dictionary or hash map

Nested scope: a tree of symbol tables

- Each scope has one symbol table
- Each symbol table may have a parent

Symbol Table

Each entry in the symbol table contains

- The name of an identifier
- Additional information: its kind, its type, if it is constant and so on

Name	Kind	Constant?
X	ld	0

Scope Check Example

```
(let ((x 0) (y 7))
(let ((y (+ x 1))
(x (+ y 1)))
(+ x y)))
```

		(global)
Name	Kind	Constant?
		(first let)
Name	Kind	Constant?
X	ld	0
У	ld	7
		(second let)
Name	Kind	Constant?
у	ld	1
X	ld	8

(alobal)

Name Collision: Identifier with Same Name

(global)

(let ((x 0)	(у	7))	
(let ((y	(+	X	1))	
(X	(+	У	1)))	
(+ x y)))				
which y?				

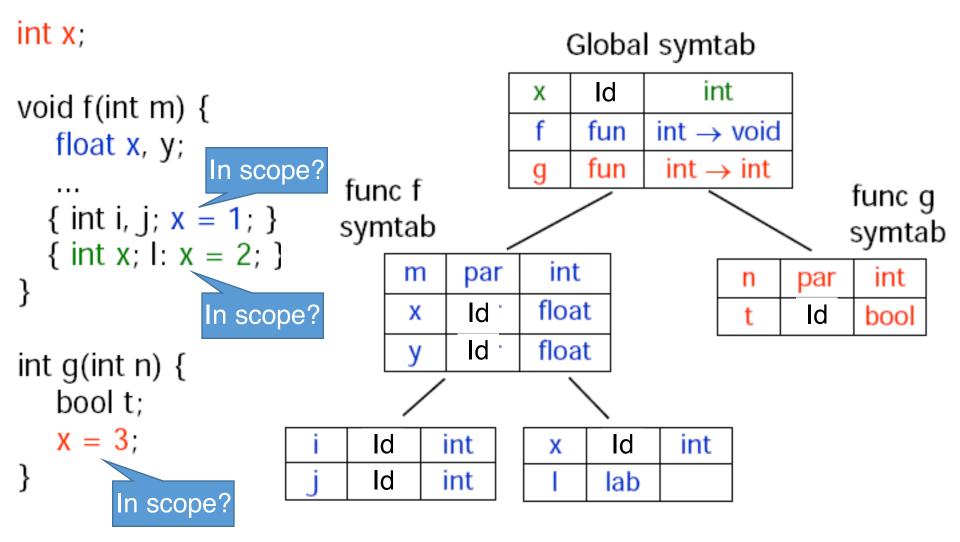
o i taiii		(giobai)		
Name	Kind	Constant?		
		(second let)		
Name	Kind	Constant?		
X	ld	0		
У	ld	7		
		(last line)		
Name	Kind	Constant?		
У	ld	1		
X	ld	8		

To find a binding:

- 1. Start from table of current scope
- 2. Go up until the name is found
- 3. Fail if no table contains the name

Name Search

Scope Check Example



Dynamic Scoping

Static Scoping

- Bindings determined by lexical structure
- Local renaming principle

Dynamic Scoping

- Bindings depend on flow of control
- Always use most recent, active binding
- Names important!
- Meaning of a variable can change

Dynamic Scoping

(global)

```
int n=2;
void first() {
n = 1
void second() {
 int n=0;
 first();
first();
second();
```

Name	Kind	Туре	
n	Id	int	
first	fun	void ->void	
second	fun	Void->void	
		(second)	
Name	Kind	Type	
n	Id	int	
		(first)	
Name	Kind	Type	

Symbol tables changes at run time! Always use most recent, active binding