CS 461

Programming Language Concepts

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The parsing is divided into two steps

- ◆First step: lexical analysis (lexer, scanner)
 - Convert a sequence of chars to a sequence of tokens
 - Token: a logically cohesive sequence of characters
 - · Common tokens
 - Identifiers
 - Literals: 123, 5.67, "hello", true
 - Keywords: bool char ...
 - Operators: + * / ++ ...
 - Punctuation: ; , () { }
- ◆Second step: syntactic analysis (parser)
 - · Convert a sequence of tokens into an AST

Regular Expressions

- ◆Used extensively in languages and tools for pattern matching
 - · E.g., Perl, Ruby, grep
- ◆ Regular expression operations
 - $\bar{\epsilon}$ (pronounced as epsilon) matches the empty string:
 - a, a literal character, matches a single character
 - Alternation: r1 | r2
 - e.g., 0|1|...|9,
 - Concatenation: r1 r2

 - Repetition (zero or more times, Kleene star): r*

Extended Regular Expressions

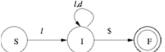
- ◆One or more repetitions
 - r+: digit+ where digit = 0|1|...|9
- ◆Zero or one occurrence: r?
 - E.g., a?
- ◆A set of characters: [aeiou]
- ◆A range of characters in the alphabet
 - a|b|c: [abc]
 - a|b|...|z:[a-z]
 - 0|1|...|9: [0-9]
- ◆Q: How to encode the above constructs using operators in regular expressions?

Lexical Analysis

- ◆Purpose: transform program representation
- ◆Input: a sequence of printable characters
- ♦Output: a sequence of tokens
- ◆Also
 - Discard whitespace and comments
 - · Save source locations (file, line, column) for error messages

Finite State Automata

- ◆A finite set of states
 - · Unique start state
 - One or more final states
 - Drawn in double circles
- ◆Input alphabet + unique end symbol (\$)
- ◆State transition function: T[s,c]
 - Describe how state changes when encountering an input symbol



FSA Execution

◆ An input is *accepted* if, starting with the start state, the automaton consumes all the input and halts in a final state.

```
s = startState;
while( s not in finalState) {
  c = next_input_character;
  s = T[s,c];
}
```

Examples: xx0\$, x12\$; non-examples: 0x\$
The language recognized by an FSA is the set of input strings accepted by the FSA

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Deterministic FSA

- ◆Defn: A finite state automaton is *deterministic* if for each state, there are no two outgoing edges labelled with the same input character
- ◆A deterministic FSA gives a way of recognizing a language
- ◆Theorem: for each RE, we can construct a deterministic FSA that recognizes the language of the RE

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A Running Example for Lexer and Parser

◆A statement language in E-BNF

- <id>= <letter>(<l- <int> = <digit>+

- <float> = <digit>+ - <float> = <digit>+.<digit>+

- punctuation marks: ; , :=, \$

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DFA for the Running Example whitespace whitespace semicolon identifier d d float

Constructing a Lexer: Token.java

```
public class Token {
   public enum TokenType {INT, FLOAT, ID, SEMICOLON,
        ASSIGNMENTOP, EOI, INVALID}

   private TokenType type;
   private String val;

   Token (TokenType t, String v) { type = t; val = v; }
   TokenType getTokenType() {return type;}
   String getTokenValue() {return val;}

   void print () { ... }
   ...
}
```

```
The Structure of Lexer.java
```

Lexer: nextToken(), part I

Lexer: nextToken(), part II

Some Aux. Functions for the Lexer

```
private char nextChar() {
    char ch = stmt.charAt(index); index = index+1;
    return ch;
}

private boolean check (char c) {
    ch = nextChar();
    if (ch == c) {ch = nextChar(); return true;}
    else return false;
}

private String concat (String set) {
    StringBuffer r = new StringBuffer("");
    do { r.append(ch); ch = nextChar(); } while (set.indexOf(ch) >= 0);
    return r.toString();
}
```

An Example of Running the Lexer

```
lexer = new Lexer ("x := 1; y := x $");
tk = lexer.nextToken();
while (tk.getTokenType() != Token.TokenType.EOI) {
    tk.print(); System.out.print(" ");
    tk = lex.nextToken();
}
```

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Recursive descent parsing

- ◆Implementation follows directly the BNF grammar <stmt> -> <assignment> {;<assignment>} <assignment> -> <id> := <exp> <exp> -> <id> | <int> | <float>
- ◆ Each non-terminal comes with a parser method
 - statement(); assignmentStmt(); expression();
 - Usually a parser method returns an object of corresponding class
 - E.g., expression() should return an expression object and statement() should return a statement object
 - The code we show next, however, just prints out the parse tree

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Parser Method for Statements

```
public void statement () {
    System.out.println("<Statement>");
    assignmentStmt();
    while (token.getTokenType() == Token.TokenType.SEMICOLON) {
        System.out.println("\t-Semicolon>;</Semicolon>");
        token = lexer.nextToken();
        assignmentStmt();
    }
    match(Token.TokenType.EOI);
    System.out.println("</Statement>");
}

<stmt> -> <assignment> {;<assignment>}
```

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Parser Method for Assignment

```
public void assignmentStmt () {
    System.out.println("\t<Assignment>");
    String val = match(Token.TokenType.ID);
    System.out.println("\t\t<Identifier>" + val + "</Identifier>");
    match(Token.TokenType.ASSIGNMENTOP);
    System.out.println("\t\t<AssignmentOp>:=</AssignmentOp>");
    expression();
    System.out.println("\t</Assignment>");
}

<assignment> -> <id>:= <exp>
```

Parser Method for Expression

Auxiliary Method for the Parser

```
private String match (Token.TokenType tp) {
    String value = token.getTokenValue();
    if (token.getTokenType() == tp)
        token = lexer.nextToken();
    else error(tp);
    return value;
}
```

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