Object-Oriented Programming

CMPSC 461
Programming Language Concepts
Penn State University
Fall 2016

What Makes a Good Language?

Avoid common bugs (e.g., reading a nil pointer)

Ease of understanding

Ease of reuse

What Abstraction Means

Abstraction: omitting or hiding low-level details

Modularity: dividing up a system into components

Encapsulation: building walls around a module

Information hiding: hiding details of a module's implementation from the rest of the system

Separation of concerns: making a feature the responsibility of a single module

Abstract Data Types

Primitive types: values and operation on values User-defined types: records, lists, ...

Focus on values

ADT: defined by a set of operations on a type **Focus on operation**

Stack is a type with new, pop, push, empty ...

Internal representation is less relevant

Classifying Operations

Creators: create new objects of type

Producers: create new objects from old ones

Mutators: change objects, e.g., list.add(n)

Observers: take objects of ADT and return objects with different type, e.g., list.size()

ADT Example

int

Creators: numeric literals 1, 2, 3, ...

Producers: arithmetic operations +, -, *, /, ...

Observers: comparison operators == , !=, <, >

Mutators: none (immutable)

ADT Example

List

Creators: ArrayList, LinkedList, ...

Producers: Collections.unmodifiableList()

Observers: size(), get()

Mutators: add(), remove(), ...

ADT Example

String

Creators: String(), String(char[])

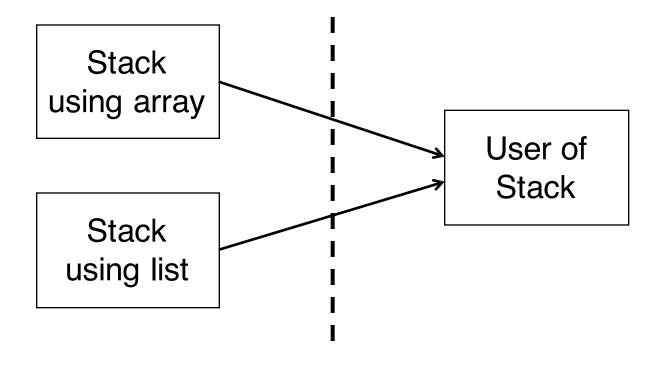
Producers: concat(), substring(),...

Observers: length(), charAt(), ...

Mutators: none (immutable)

Representation Independence

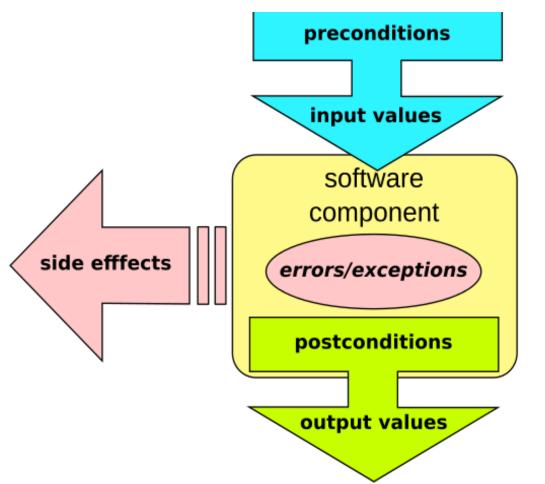
A good ADT should be independent of impl.



Abstraction Boundary

Design by Contract

A good ADT should define a contract:



Contract

Precondition: assumptions on inputs

Side effects: changes to the value

Postcondition: functionality of the operation

Invariants: a property that is always true throughout the operation

```
public class Rational {
   private int p,q; // represents p/q
   // class invariant: q > 0, gcd(p,q) = 1
   //
   Note: gcd(0,x) = x
```

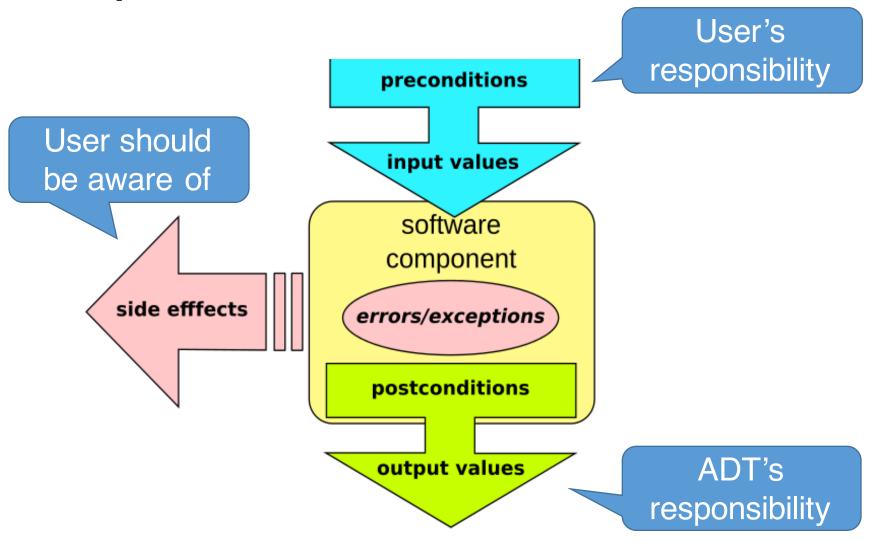
```
/** Create num/den.
                            Precondition
    Requires: den != 0.
public Rational (int num, int den) {
    if (den < 0) {
                                Check den!=0?
         num = -num;
         den = -den;
    int g = \overline{gcd(num, den)};
    p = num/q;
    q = den/q;
```

Side effects

Postcondition

```
/** Returns x+y. */
public Rational plus(Rational x, Rational y) {
   Rational z = new Rational(x.p, x.q);
   z.add(y);
   return z;
}
```

Separation of Concerns



Enforcing Contract

How can we gain confidence that these contracts are all being obeyed?

Assertions: run-time condition check

```
Stops the program if the condition is false
```

```
public class Rational {
    private int p,q; // represents p/q
    // class invariant: q > 0, gcd(p,q) = 1
    // Note: gcd(0,x) = x
    boolean classInv() {
        return q > 0 && gcd(p, q) == 1;
    }
```

```
/** Create num/den.
    Requires: den != 0.
public Rational (int num, int den) {
    assert (den != 0);
    if (den < 0) {
        num = -num;
        den = -den;
    int g = gcd(num, den);
    p = num/q;
    q = den/q;
    assert ClassInv();
```

```
/** Modifies: this to be this+r. */
public void add(Rational r) {
   int g = gcd(q, r.q);
   assert (g != 0);
   p = r.q/g * p + q/g * r.p;
   q *= r.q/g;
   assert ClassInv();
}
```

```
/** Returns x+y. */
public Rational plus(Rational x, Rational y) {
    Rational z = new Rational(x.p, x.q);
    z.add(y);
    assert ClassInv();
    return z;
}
```

Assertion

Assertions are powerful weapons in catching bugs! But they have performance overhead

Java: assertions are turn off by default Pass in -ea to JVM to enable all assertions