#### Procedures and Functions

CMPSC 461
Programming Language Concepts
Penn State University
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#### Parameter Passing

How caller communicates with callee

Formal parameters: names in the declaration of a function

Actual parameters: variables/expressions passed to

a function

```
formal
foo
     (3+5);
        actual
```

#### Parameter Modes

Call-By-Value (CBV)

Call-By-Reference (CBR)

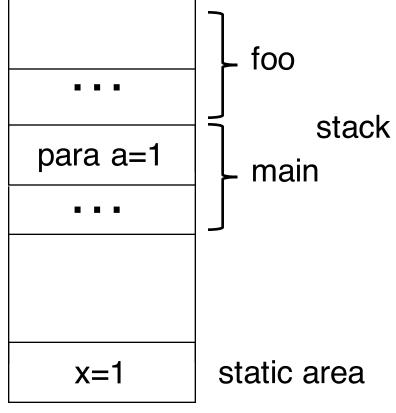
Call-By-Value-Return (CBVR)

Call-By-Name (CBN)

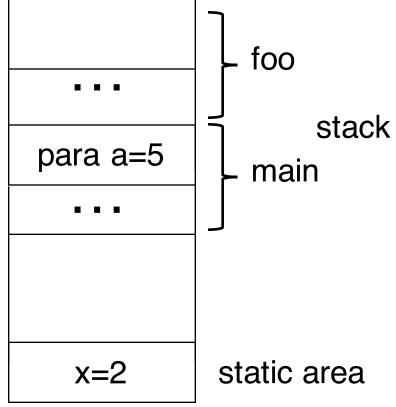
#### Calling mechanism

- Arguments are evaluated to their values
- Memory or registers allocated for arguments on AR
- Argument values copied to AR
- AR destroyed when callee returns

```
int x=1;
int foo (int a) {
    x = 2;
                             para a=1
    a = 5;
    return x+a;
void main() {
    foo(x); //result?
                               x=1
    print(x); //result?
```



```
int x=1;
int foo (int a) {
    x = 2;
    a = 5;
    return x+a;
void main() {
    foo(x); //result?
                               x=2
    print(x); //result?
```



```
int x=1;
int foo (int a) {
    x = 2;
                                              stack
    a = 5;
                                          main
    return x+a;
void main() {
    foo(x); //result?
                                x=2
                                        static area
    print(x); //result?
```

Formal & Actual parameters have separate memory: their values may diverge

#### Characteristics

- Actual parameters may not directly be changed in callee (unless pass in pointers)
- Arguments can be complex expressions
- Simple and intuitive (less error-prone)

#### Call-By-Value: Performance

Primitive types: cost per parameter is small Arrays, records, structures: copying value could be slow

C: programmer has to pass pointers/references to avoid copying cost

Java: only primitive types call-by-value

#### Calling mechanism

- Arguments are evaluated to their values
- Memory or registers allocated for arguments on AR
- Argument address stored in AR
- AR destroyed when callee returns

```
int x=1;
int foo (int a) {
                                               stack
    x = 2;
                               para a
                                           main
    a = 5;
    return x+a;
void main() {
    foo(x); //result?
                                x=1
                                         static area
    print(x); //result?
```

```
int x=1;
int foo (int a) {
                                              stack
    x = 2;
                               para a
                                           main
    a = 5;
    return x+a;
void main() {
    foo(x); //result?
                                x=5
                                         static area
    print(x); //result?
```

```
int x=1;
int foo (int a) {
    x = 2;
                                              stack
    a = 5;
                                          main
    return x+a;
void main() {
    foo(x); //result?
                                        static area
                                x=5
    print(x); //result?
```

Formal parameter is an alias of actual parameter: their values are the same

#### Characteristics

- Actual parameters may directly be changed in callee
- Some language disallows complex expressions as arguments
- Programs are harder to understand (more error-prone)

#### Call-By-Reference: Performance

Avoids the cost of memory copy
Indirect memory access: address → value

C: call-by-value is the default mode Java: constructed types call-by-reference

## Call-By-Value-Result (In-Out)

#### Calling mechanism

- Arguments are evaluated to their values
- Memory or registers allocated for arguments on AR
- Argument values stored in AR
- Before callee returns, AR values copied back to actual arguments
- AR destroyed when callee returns

## Call-By-Value-Result

```
int x=1;
int foo (int a) {
                                               stack
    x = 2;
                              para a=1
                                           main
    a = 5;
    return x+a;
void main() {
    foo(x); //result?
                                         static area
                                x=1
    print(x); //result?
```

## Call-By-Value-Result

```
int x=1;
int foo (int a) {
                                               stack
    x = 2;
                              para a=5
                                           main
    a = 5;
    return x+a;
void main() {
    foo(x); //result?
                                x=2
                                         static area
    print(x); //result?
```

#### Call-By-Value-Result

```
int x=1;
int foo (int a) {
    x = 2;
                                              stack
    a = 5;
                                          main
    return x+a;
void main() {
    foo(x); //result?
                                        static area
                                x=5
    print(x); //result?
```

#### Call-By-Value-Return

#### Characteristics

 Mostly identical to call-by-value, except an extra step of copying values back to actual parameter

# Call-By-Name (Lazy Evaluation)

#### Calling mechanism

- Arguments are not evaluated to their values (like macros)
- Actual parameters replace all formal parameters in body

```
void swap (int a, int b) {
   int t = a;
   a = b;
   b = t;
}
swap(i, A[i]) // value swapped?
```

Still, a useful mode in functional programming, e.g., Haskell

#### Return Values

Fixed size: callee stores values in caller's AR

Size determined at run time: callee stores results in heap, and store their addresses in caller's AR