

1. (a) leftmost derivation for the expression “2 * 3 * 6”

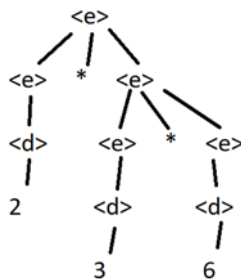
$\langle e \rangle \rightarrow \langle e \rangle * \langle e \rangle \rightarrow \langle d \rangle * \langle e \rangle \rightarrow 2 * \langle e \rangle \rightarrow 2 * \langle e \rangle * \langle e \rangle \rightarrow 2 * \langle d \rangle * \langle e \rangle \rightarrow$
 $2 * 3 * \langle e \rangle \rightarrow 2 * 3 * \langle d \rangle \rightarrow 2 * 3 * 6$

1. (b) rightmost derivation for “2*3*6”

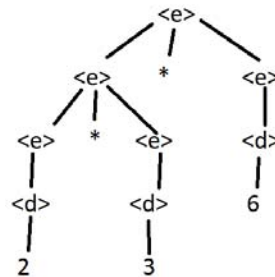
$\langle e \rangle \rightarrow \langle e \rangle * \langle e \rangle \rightarrow \langle e \rangle * \langle d \rangle \rightarrow \langle e \rangle * 6 \rightarrow \langle e \rangle * \langle e \rangle * 6 \rightarrow \langle e \rangle * \langle d \rangle * 6 \rightarrow \langle e \rangle * 3 * 6 \rightarrow \langle d \rangle * 3 * 6 \rightarrow 2 * 3 * 6$

1. (c) two different parse trees for “2*3*6”

For (a):

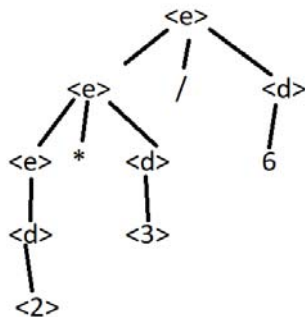


For (b):



1. (d) new grammar – left associative operators

$\langle e \rangle \rightarrow \langle d \rangle \mid \langle e \rangle * \langle d \rangle \mid \langle e \rangle / \langle d \rangle$
 $\langle d \rangle \rightarrow 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$



There are two choices to begin parsing 2*3*6:

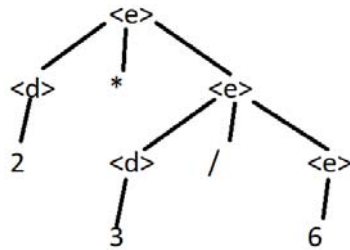
(1) $\langle e \rangle \rightarrow \langle e \rangle * \langle d \rangle$ or (2) $\langle e \rangle \rightarrow \langle e \rangle / \langle d \rangle$

For option (1), it is impossible to get the expression 2*3/6 from $\langle e \rangle \rightarrow \langle e \rangle * \langle d \rangle \rightarrow \langle e \rangle * 6$.

For option (2), it is shown above. Hence, (2) is the only choice.

1. (e) new grammar – right associative operators

$\langle e \rangle \rightarrow \langle d \rangle \mid \langle d \rangle * \langle e \rangle \mid \langle d \rangle / \langle e \rangle$
 $\langle d \rangle \rightarrow 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$



2. An unambiguous BNF grammar

`<e> -> <e>, <n> | <n>`

`<n> -> <n><d> | <d>`

`<d> -> 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9`

Left associativity is enforced on “,” through left recursion; rules of `<n>` also use left recursion. These left-recursive rules remove ambiguity.

3. New grammar for “+” and “-” left-associativity. Precedence of “~” is higher than “+” and “-”

`<e> -> <t> | <e> + <t> | <e> - <t>`

`<t> -> <d> | ~ <t>`

`<d> -> 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9`

4. One BNF grammar is as follows:

`<email> := <account> @ <subdomains> . <topdomain>`

`<account> := <letter> | <account><letter> | <account><digit>`

`<subdomains> := <letOrDigSeq> | <letOrDigSeq> . <subdomains>`

`<letOrDigSeq> := <letter><letOrDigSeq> | <digit><letOrDigSeq>
| <letter> | <digit>`

`<topdomain> := edu | org | com`

`<letter> -> a | b | c | ... | z | A | B | C | ... | Z`

`<digit> -> 0 | 1 | 2 | ... | 9`

5. New grammar

`<canonical-num> -> <digit> | <non-zero-digit> <num>`

`<num> -> <digit> | <digit> <num>`

`<non-zero-digit> -> 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9`

`<digit> -> 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9`

6. equivalent BNF grammar

`<expr> -> -<int> | <int> | -<int>.<int> | <int>.<int>`

`<int> -> <digit> | <digit><int>`

`<digit> -> 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9`