

Chapter 11-2 – Faster Sorting Algorithms

1. Merge Sort – Divide and Conquer sorting algorithm

Divide: divide the n-element sequence to be sorted into two subsequences of $n/2$ elements each

Conquer: sort the two subsequences recursively using merge sort. IF the length of a sequence is 1, do nothing since it is already in order.

Combine: merge two sorted subsequences to produce the sorted answer.

```
// Sorts theArray[first..last] by
// 1. Sorting the first half of the array
// 2. Sorting the second half of the array
// 3. Merging the two sorted halves
mergeSort(theArray: ItemArray, first: integer, last: integer)
{
    if (first < last)
    {
        mid = (first + last) / 2      // Get midpoint

        // Sort theArray[first..mid]
        mergeSort(theArray, first, mid)

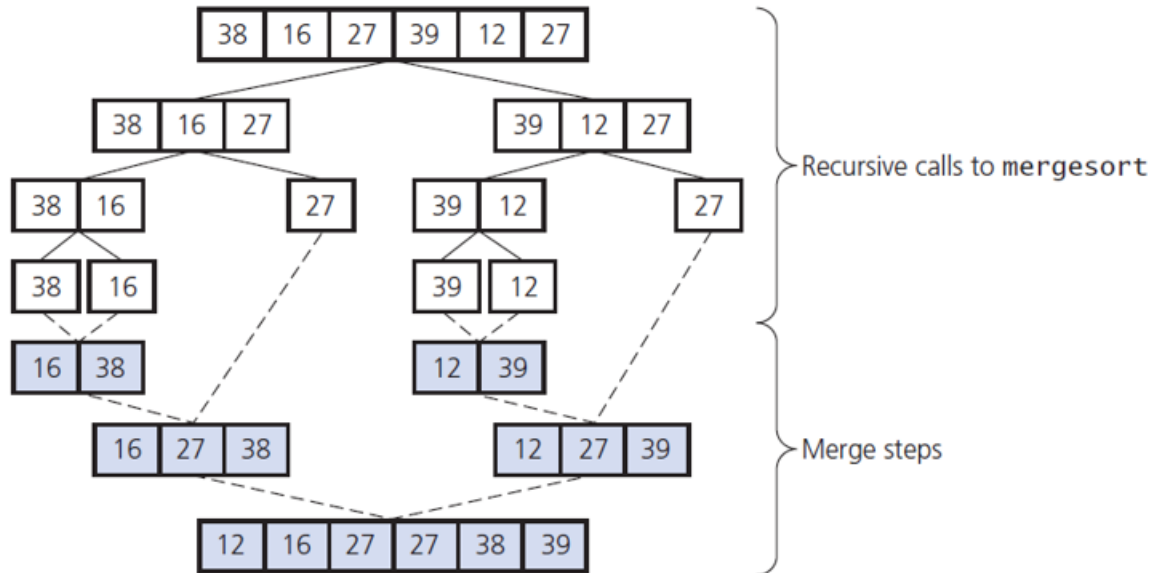
        // Sort theArray[mid+1..last]
        mergeSort(theArray, mid + 1, last)

        // Merge sorted halves theArray[first..mid] and theArray[mid+1..last]
        merge(theArray, first, mid, last)
    }
    // If first >= last, there is nothing to do
}
```

Example 1)

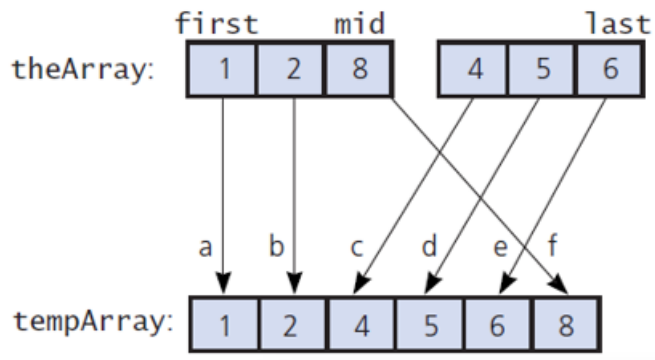
8	1	4	3	2
---	---	---	---	---

Example 2)



Analysis - Merge step

- worst-case instance of the merge step in a merge sort



Merge the halves:

- compare $1 < 4$, move 1 to the tempArray
- compare $2 < 4$, move 2 to the tempArray
- compare $8 > 4$, move 4 to the tempArray
- compare $8 > 5$, move 5 to the tempArray
- compare $8 > 6$, move 6 to the tempArray
- move 8 to the tempArray
- move all values in the tempArray to the original array

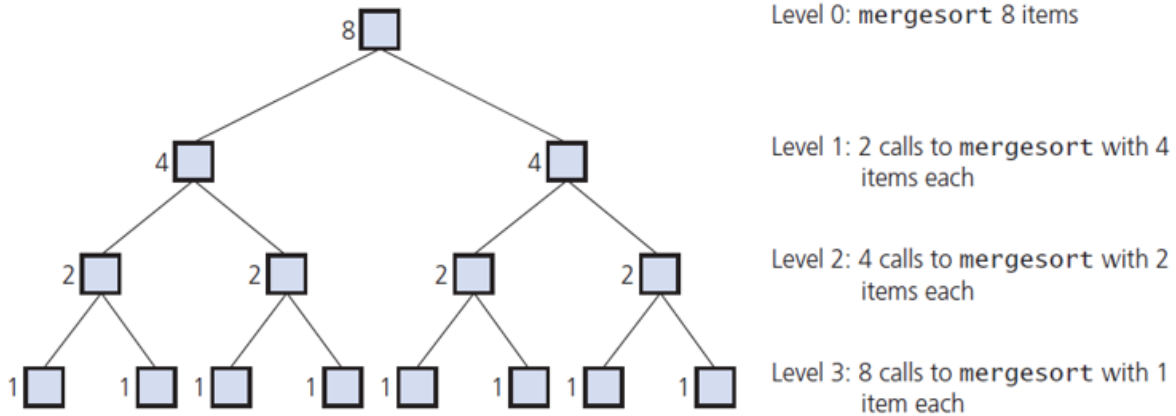
Analysis – mergeSort

```
mergeSort(theArray: ItemArray, first: integer, last: integer)
{
    if (first < last)
    {
        mid = (first + last) / 2      // Get midpoint

        // Sort theArray[first..mid]
        mergeSort(theArray, first, mid)

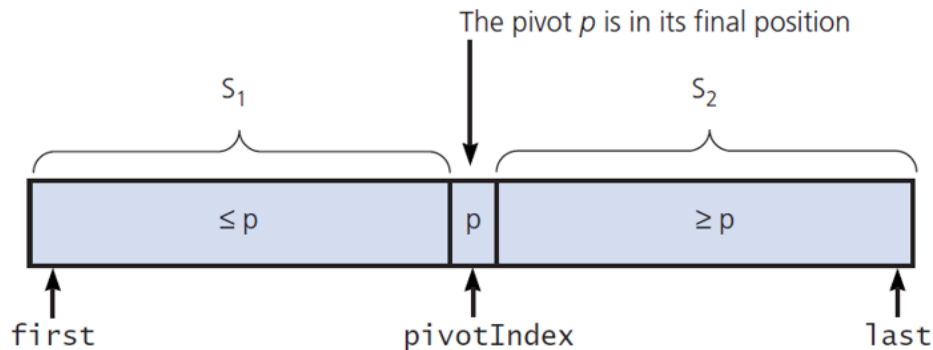
        // Sort theArray[mid+1..last]
        mergeSort(theArray, mid + 1, last)

        // Merge sorted halves theArray[first..mid] and theArray[mid+1..last]
        merge(theArray, first, mid, last)
    }
    // If first >= last, there is nothing to do
}
```



2. The Quick Sort

- Another divide-and-conquer algorithm
- Partitions an array into items that are
 - Less than or equal to the pivot and
 - Those that are greater than or equal to the pivot
- Partitioning places pivot in its correct position within the array
 - Place chosen pivot in theArray[last] before partitioning



```
// Sorts theArray[first..last].
quickSort(theArray: ItemArray, first: integer, last: integer): void
{
    if (first < last)
    {
        Choose a pivot item  $p$  from theArray[first..last]
        Partition the items of theArray[first..last] about  $p$ 
        // The partition is theArray[first..pivotIndex..last]
        quickSort(theArray, first, pivotIndex - 1) // Sort  $S_1$ 
        quickSort(theArray, pivotIndex + 1, last) // Sort  $S_2$ 
    }
    // If first  $\geq$  last, there is nothing to do
}
```

Example A partitioning of an array during a quick sort

3	5	0	7	6	1	2	4
---	---	---	---	---	---	---	---

```

// Partitions theArray[first..last].
partition(theArray: ItemArray, first: integer, last: integer): integer
{
    // Choose pivot and reposition it
    mid = first + (last - first) / 2
    sortFirstMiddleLast(theArray, first, mid, last)
    Interchange theArray[mid] and theArray[last - 1]
    pivotIndex = last - 1
    pivot = theArray[pivotIndex]

    // Determine the regions  $S_1$  and  $S_2$ 
    indexFromLeft = first + 1
    indexFromRight = last - 2

    done = false
    while (not done)
    {
        // Locate first entry on left that is  $\geq$  pivot
        while (theArray[indexFromLeft] < pivot)
            indexFromLeft = indexFromLeft + 1

        // Locate first entry on right that is  $\leq$  pivot
        while (theArray[indexFromRight] > pivot)
            indexFromRight = indexFromRight - 1

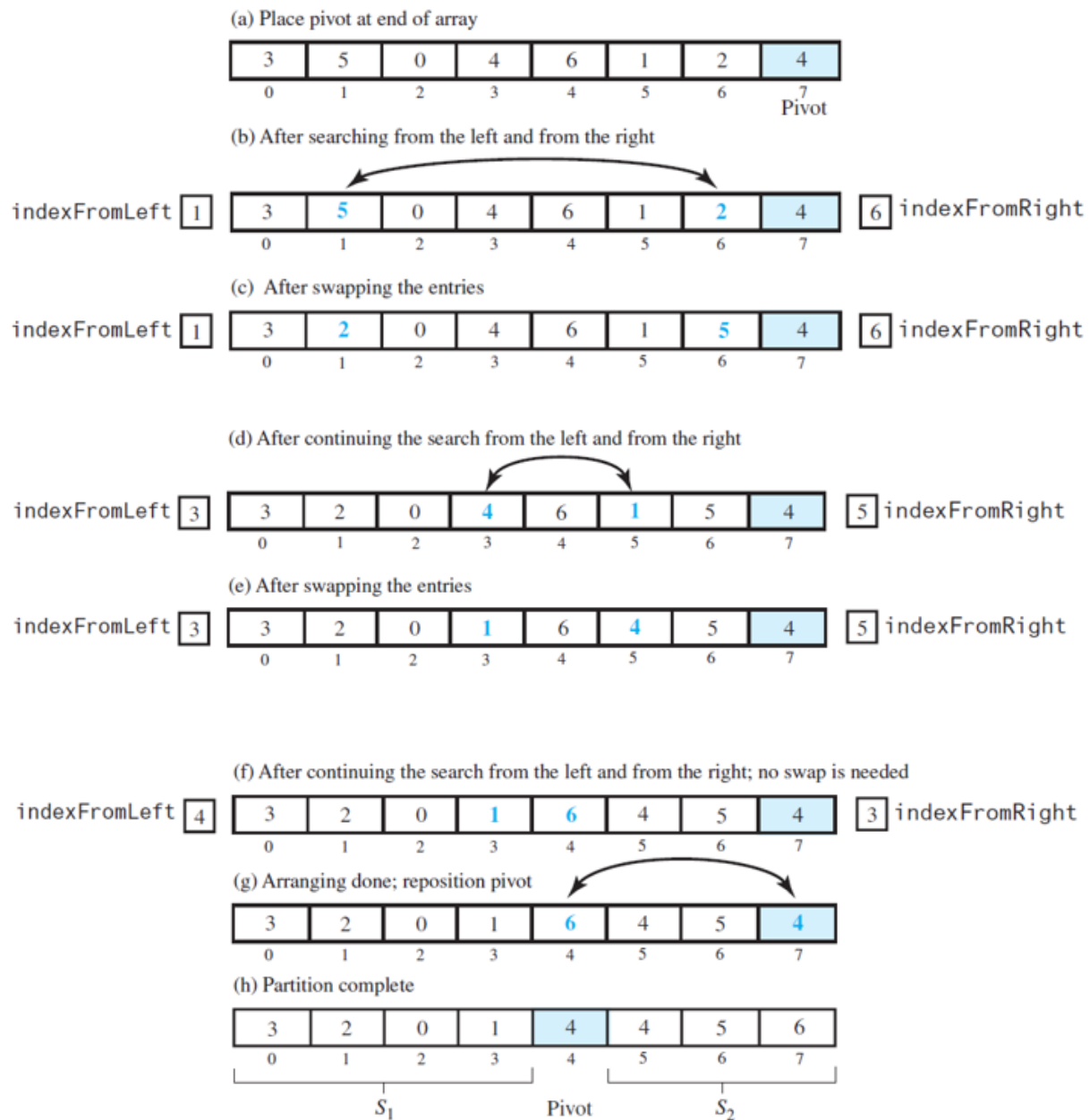
        if (indexFromLeft < indexFromRight)
        {
            Interchange theArray[indexFromLeft] and theArray[indexFromRight]
            indexFromLeft = indexFromLeft + 1
            indexFromRight = indexFromRight - 1
        }
        else
            done = true
    }

    // Place pivot in proper position between  $S_1$  and  $S_2$ , and mark its new location
    Interchange theArray[pivotIndex] and theArray[indexFromLeft]
    pivotIndex = indexFromLeft

    return pivotIndex
}

```

Example A partitioning of an array during a quick sort



Median-of-three pivot selection

```
// Arranges the first, middle, and last entries in an array into ascending order.
sortFirstMiddleLast(theArray: ItemArray, first: integer, mid: integer,
                    last: integer): void
{
    if (theArray[first] > theArray[mid])
        Interchange theArray[first] and theArray[mid]

    if (theArray[mid] > theArray[last])
        Interchange theArray[mid] and theArray[last]

    if (theArray[first] > theArray[mid])
        Interchange theArray[first] and theArray[mid]
}
```

(a) The original array

5	8	6	4	9	3	7	1	2
---	---	---	---	---	---	---	---	---

(b) The array with its first, middle, and last entries sorted

2	8	6	4	5	3	7	1	9
---	---	---	---	---	---	---	---	---

Pivot

(c) The array after positioning the pivot and just before partitioning

2	8	6	4	1	3	7	5	9
---	---	---	---	---	---	---	---	---

indexFromLeft indexFromRight Pivot

2	8	6	4	1	3	7	5	9
---	---	---	---	---	---	---	---	---

Example Trace the quicksort's partitioning algorithm as it partitions the following array

38	16	40	39	12	27
----	----	----	----	----	----

Analysis

Partitioning

Average case

Worst case

A Comparison of sorting algorithms

	<u>Worst case</u>	<u>Average case</u>
Selection sort	n^2	n^2
Bubble sort	n^2	n^2
Insertion sort	n^2	n^2
Merge sort	$n \times \log n$	$n \times \log n$
Quick sort	n^2	$n \times \log n$